VFS and filesystems in Linux

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VFS layer

- Linux since long time has supported many various filesystems (both physical and virtual), now over 30
- The natural way of handling all filesystems was to create separate layer with common functions
- This layer is called Virtual File System (VFS) layer

Types of filesystems handled by VFS

- Physical (backed by physical device, like hard disk, flash or CD) – e.g. ext3, reiserfs, jfs, xfs, jffs, iso9660
- Network filesystems (including cluster filesystems) – e.g. nfs, coda, smbfs, gfs, ocfs
- Pseudofilesystems provide various information based on some metainfo or allowing some actions – e.g. proc, sysfs, pipefs, futexfs, usbfs
- Special-purpose filesystems tmpfs, ramfs, rootfs

VFS architecture

- Filesystems handling consists of:
 - generic operations handled by VFS
 - structures representing files and other filesystem objects (handled by VFS)
 - caches (handled by VFS)
 - structures specific to filesystem (defined by filesystem)
 - operations on structures specific to filesystem (defined by filesystem)

VFS and page cache

- VFS is integrated with page cache
- If allows to provide zero-copy semantics (file contents is served from page cache, no copying)
- It simplifies VFS in some cases all operations (e.g. read) are done on pages
- It also simplifies FS implementation filesystem can provide page-level operations and page-cache layer together with VFS will handle most work (like mmap and AIO)

VFS abstract structures

- Structures which represent generic (abstract) objects:
 - inode represents inode on disk
 - dentry represents directory entry
 - file represents open file
 - super_block represents filesystem instance
 - file_system_type represents filesystem type handling
- Can contain filesystem-specific extra information

VFS caches

- inode cache caches read inodes and performs delayed writeback of inodes
- dentry cache caches directory entries to speed up path lookup (dentries are not written to disk)

Generic operations handled by VFS

- VFS code handles:
 - sys_open() and related functions path lookup, crossing mounts, etc.
 - sys_read(), sys_write(), sys_mmap(), aio, and all other system-call interface functions
 - file descriptor operations (closing, duplicating, etc.)
 - mounting filesystems
- Most operations are done by doing generic work while delegating filesystemspecific functions to filesystem defined hooks (file/inode/super_block/dentry/...)

Example - opening file

- Function do_sys_open():
 - Get unused file descriptor
 - Perform path lookup using namei/ path_lookup functions (open_namei() and do_path_lookup())
 - get root
 - get directory entry read using FS-specific function if not present in cache, calls iget in superblock if inode not present in cache
 - if directory lookup name in directory
 - cross mount points, follow symlinks, ...
 - Call filesystem-specific open routine from file_operations (__dentry_open())

Example – filesystem mounting

- Code: do_mount()
 - Lookup directory to be mounted
 - Load filesystem handling module if not present (which registers its filesystem_type)
 - Get filesystem superblock using filesystem_type->get_sb() (vfs_kern_mount())
 - Add mount point to namespace

struct filesystem_type

- Represents filesystem type and is registered by module which handles such filesystem type
- Important fields (struct file_system_type):
 - get_sb() mounts new instance of filesystem on given device and directory (e.g. ext3_get_sb()) and creates and fills in super_block
 - kill_sb() unmounts filesystem instance
 - name name of filesystem type (e.g. "ext3")
 - fs_flags various flags which tell how to handle this filesystem type (e.g. FS_REQUIRES_DEV)

struct filesystem_type (2)

- VFS provides convenience functions for filesystem type operations:
 - get_sb_bdev() fill superblock for block device-based filesystems
 - get_sb_nodev() for non-physical filesystems
 - get_sb_single() for filesystems with single instance
 - kill_block_super() unmounts block-device based filesystem
 - kill_litter_super()/kill_anon_super() unmounts non-physical filesystems

struct super_block

- Represents mounted filesystem instance
- Most important fields (struct super_block)
 - struct super_operations *s_op operations on filesystem instance
 - struct file_system_type *s_type FS type
 - struct export_operations *s_export_op NFS export operations
 - s_flags mounted filesystem instance flags
 - struct dentry *s_root dentry of root

struct super_operations

- Keeps FS-specific operations performed by filesystem instance
- Important fields (struct super_operations)
 - read_inode reads inode from disk to cache
 - dirty_inode marks inode as dirty
 - write_inode writes inode to disk
 - put_inode inode is removed from cache
 - delete_inode remove inode from disk (last hardlink was removed and inode is not used)

struct super_operations (2)

- Superblock important fields (filesystem wide):
 - put_super filesystem is unmounted
 - write_super superblock should be written to disk
 - sync_fs write all state to disk
 - statfs report statistics about filesystem (used and available space, etc.)

struct inode

- Keeps information about inode in memory
- Important fields (struct inode)
 - i_ino unique inode number (inode identifier)
 - i_nlink number of references to inode on disk (hardlinks) and by processes
 - struct inode_operations *i_op filesystemspecific inode operations
 - struct file_operations *i_fop default filesystem-specific file operations
 - struct super_block *i_sb pointer to super block of filesystem

struct inode (2)

- address_space *i_mapping pagecache operations
- struct pipe_inode_info *i_pipe inode information specific to pipes
- struct block_device *i_bdev inode information specific to block device files
- struct cdev i_cdev inode information specific to character device files
- i_generation generation (version) of inode, used mainly for NFS exporting
- i_private private data for filesystem kept with inode

struct inode_operations

- Important fields (struct inode_operations, e.g. ext3_dir_inode_operations, ext3_file_inode_operations, docs)
 - truncate truncate file
 - create create file in directory
 - mkdir create directory in directory
 - lookup lookup file name in directory inode
 - link create hardlink
 - unlink remove link to inode (and possibly delete if number of references drops to 0)
 - setattr set file attributes

struct file

- Represents open file
- Most important fields (struct file)
 - f_dentry dentry associated with file
 - f_vfsmnt mount point of FS for file
 - struct file_operations *f_op file operations
 - f_pos file position
 - f_ra readahead state
 - private_data filesystem private data for file
 - struct address_space *f_mapping address space object for file

struct file_operations

- Important fields (struct file_operations)
 - open open file
 - read/write standard synchronous read/write
 - aio_read/aio_write asynchronous read/write
 - poll descriptor status polling
 - mmap memory mapping of file
 - flush close file descriptor
 - release close last reference to file
 - fsync/fasync sync file contents with disk
 - readv/writev scatter/gather operations

struct file_operations (2)

- sendfile stream file contents to socket
- lock/setlease file locking
- splice_read/splice_write transfer file contents directly between descriptors (no userspace copy using splice() system call)

struct dentry

- Represents directory entry (file name in directory) which points to inode
- Dentries are only cache and not synced to disk
- Important fields (struct dentry)
 - struct inode *d_inode inode associated with dentry
 - struct dentry *d_parent dentry of parent directory
 - struct qstr d_name file name in dentry
 - struct dentry operations *d op operations

struct dentry_operations

- Most important fields (struct dentry_operations)
 - d_compare comparison function for dentries (for example for case-insensitive comparison)
 - d_delete all references to dentry have been removed, but it stays in cache
 - d_release free dentry
 - d_dname dynamic creation of file name in dentry
- dentry_operations are not so often overridden, for most filesystems it is not necessary

struct address_space (reminder from VM)

- address_space object represents inode which hosts data of some pages and defines how to read/write them on disk
- Page in page cache is represented as pointer to address_space object and index – i.e. offset in page units in file
- As reading/writing files is unified using page cache, filesystems define only address_space_operations (e.g. readpage, writepage) and they do not care if file is written using read/write or using mmap – this is handled by pagecache layer

struct address_space (2)

 Generic VFS functions use address_space object, so if filesystem uses generic functions, it gets much functionality for free (mmap, async IO), it just has to define address_space operations

Generic functions provided by VFS

- VFS provides also generic functions which can be used by filesystems (but they can implement them on their own)
- It is recommended to use VFS-provided functions as they are updated when VFS architecture changes and cause less code duplication (which avoids errors)
- Using generic functions you also get much functionality for free (like mmap, AIO)
- Some functions can be put directly in structure or can be called from filesystemspecific function after FS-specific job

Generic file functions provided by VFS

- Examples of generic file operations
 - generic_file_llseek() performs llseek()
 operation as expected in standard systems,
 using struct file fields
 - generic file open() opens the file
 - do_sync_read()/write() standard, sync ops
 - generic_file_aio_read/write() AIO
 - generic_file_mmap() mmap handling
 - generic_file_splice_read()/write() splice
- Example: struct ext3_file_operations

Generic inode operations provided by VFS

Generic inode operations

- generic_readlink() generic symlink read
- generic_setxattr()/getxattr/listxattr/removexa
 ttr extended attributes routines
- generic_fillattr() fills inode attributes
- generic_permission() checks permissions for inode
- generic_delete_inode() deletes inode
- generic_drop_inode() generic drop inode

Generic address_space ops provided by VFS

- Generic address_space operations:
 - mpage_readpages(),mpage_readpage reading page using FS-provided get block function
 - mpage_writepages/mpage_writepage page write routine using get block function
 - block_sync_page() for sync_page
 - generic_block_bmap() for bmap

Creating simple FS using generic functions

- So, to create the core of simple devicebased filesystem, it is enough to provide:
 - get_block routine, used by generic address_space operations, used by generic file operations
 - inode read operation
- Remaining part can be done using generic functions, some default functions are transparently used by VFS is FSspecific function is not provided

Namespaces

- Linux VFS supports namespaces view of filesystem structure can be different for processes
- Used to implement chroot(), but has more power
- It is possible to do tricks like bind mounts (mount single directory/file in 2 places at the same time)
- All dirty work is done by VFS

Special case - NFS exporting

- Standard VFS interface does not suffice to export filesystem using NFS
- Filesystem must provide struct export_operations in s_export_op of its super_block if it wants to be exported by NFS

Other VFS features

- Linux VFS is flexible enough to create, for example, stacked filesystems
- Linux VFS supports inotify/dnotify notification of filesystem changes
- VFS supports extended attributes –
 extended attributes operations can be
 provided in file_operations or generic
 functions can be used, which use s_xattr
 handler defined in super_block

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