

SPRINTS

# Healthcare System

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System design case study using RTOS

Submitted by:

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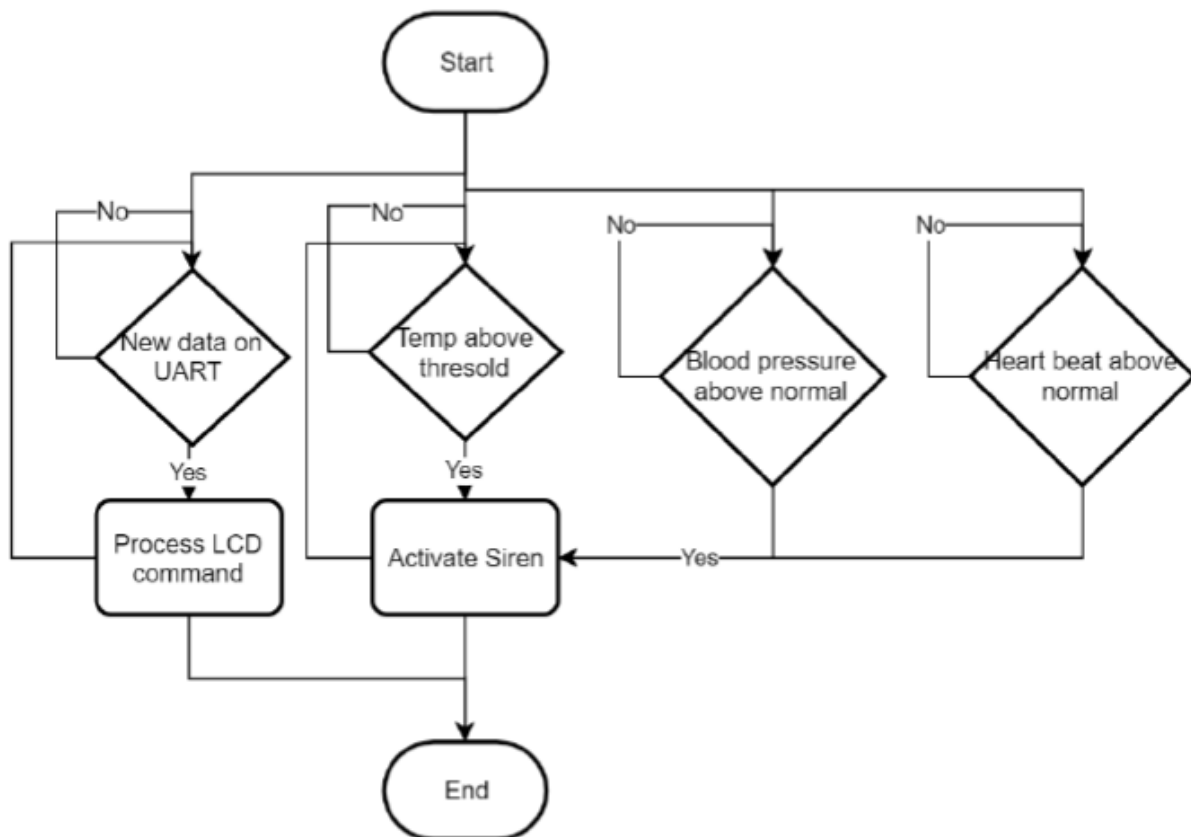
## - Contents:

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- Calculating system tick rate.
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## - Problem Statement and system diagram:

**Task:** Design a healthcare system using RTOS with the following requirements:

- A touch LCD as input that can control the system and give commands. Every LCD command is represented in 4 bytes. LCD is connected to the micro-controller through UART with speed 9600 bps [Bit per second]. (Reading 4 bytes and processing the command takes 2 ms)
- Blood pressure sensor with new data every 25ms. (Reading the sensor and processing its data takes 3 ms)
- Heart beat detector with new data every 100ms. (Reading the sensor and processing its data takes 1.5 ms)
- Temperature sensor with new data every 10ms. (Reading the sensor and processing its data takes 2.5 ms)
- Alert siren. (Activate or Deactivate the siren takes 1 ms)



- Breaking down tasks:

**Five tasks, they are:**

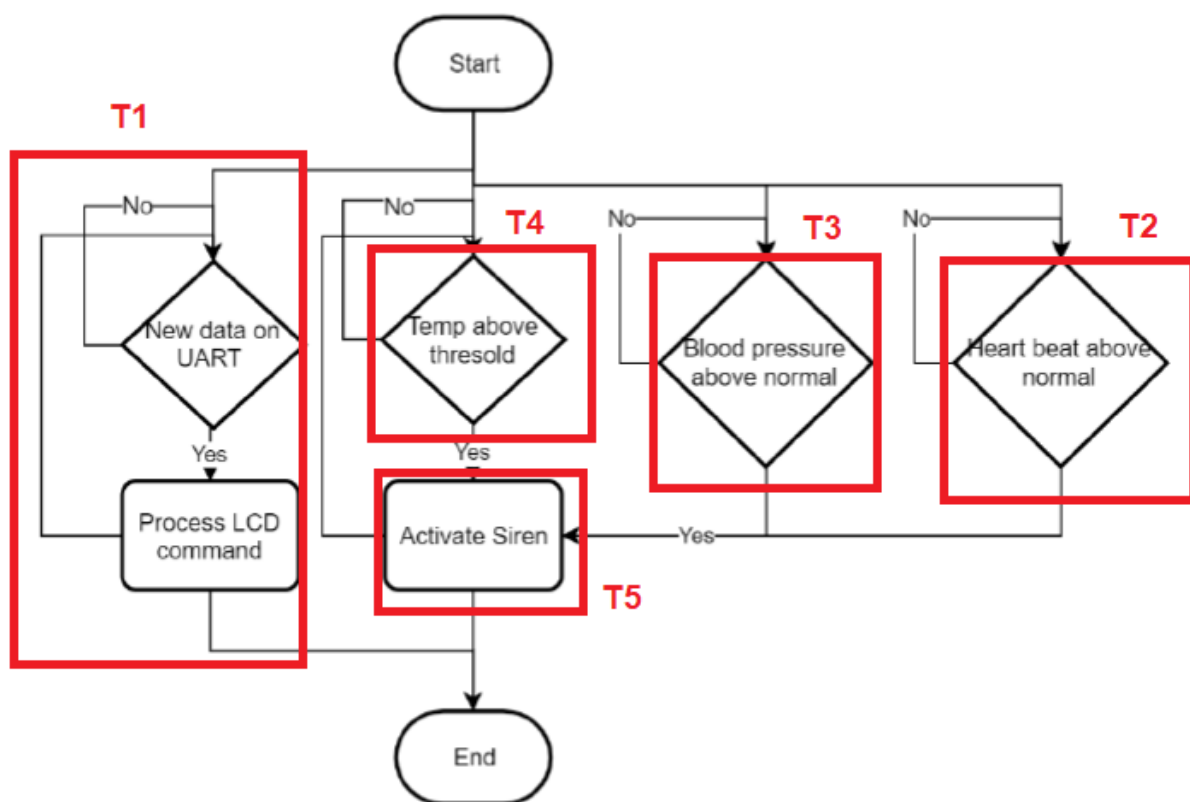
**T1: UART/LCD**

**T2: Heart Beat Sensor**

**T3: Blood Pressure Detector**

**T4: Temperature Sensor**

**T5: Alert Siren**



- Finding task parameters (Priority - Periodicity - Deadline):

**T1: UART/LCD (P = D = 20 ms, E = 2 ms)**

**T2: Heart Beat Sensor (P = D = 100 ms, E = 1.5 ms)**

**T3: Blood Pressure Detector (P = D = 25 ms, E = 3 ms)**

**T4: Temperature Sensor (P = D = 10 ms, E = 2.5 ms)**

**T5: Alert Siren (P = D = 10 ms, E = 1 ms)**

Assumptions:

- UART/LCD task runs every 20 ms to ensure debouncing and high-precision input
- Alert Siren task runs every 10 ms as it has to immediately respond to any enabling or disabling as fast as the temperature sensing.
- Deadline is the same as periodicity.

- Calculating system tick rate:

**Tick rate = Highest common factor of all periodicities in the system**

**H.C.F. (20, 100, 25, 10, 10) = 5 ms**

**20 = 5 x 2 x 2**

**100 = 5 x 5 x 2 x 2 (Using prime factorization)**

**25 = 5 x 5**

**10 = 5 x 2**

- Calculating hyper period:

**Hyperperiod = Lowest common multiple of all periodicities in the system**

**L.C.M. (20, 100, 25, 10, 10) = 100 ms**

$$20 = 5 \times 2 \times 2$$

$$100 = 5 \times 5 \times 2 \times 2$$

$$25 = 5 \times 5$$

$$10 = 5 \times 2$$

$$\text{LCM} = 5 \times 5 \times 2 \times 2 = 100$$

(Multiply max occurrence of prime factors)

- Calculating CPU load:

$$\text{CPU load} = \frac{\sum E_i \times \frac{H}{P_i}}{H} = \sum_{i=1}^5 \frac{E_i}{P_i}$$

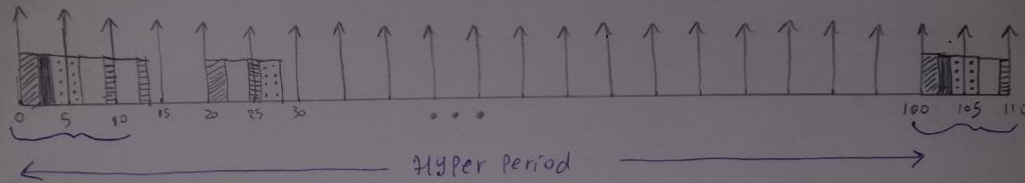
$$\begin{aligned} \therefore \text{CPU load} &= \frac{2}{20} + \frac{1.5}{100} + \frac{3}{25} + \frac{2.5}{10} + \frac{1}{10} \\ &= 0.1 + 0.015 + 0.12 + 0.25 + 0.1 \\ &= 0.585 \times 100\% = \boxed{58.5\%} \end{aligned}$$

(Analytically)  
system is healthy

## - Drawing system manually and analyzing schedulability:

Finding schedulability

Case 1: No Priority & No Preemption (Tasks are running in order  $T_1 \rightarrow T_5$ )



$\therefore \sum E_i = 10$  and system is non-preemptive  
    ↳ which is the least periodicity  
 $\therefore$  HyperPeriod on graph matches the analytical value

Tasks will not miss their deadlines  
& system is schedulable

Case 2: Priority & Preemption (Drawn on simso since hand drawing is more complex)

same results

- system is healthy
- Tasks will not miss their deadlines
- system is schedulable

Are they really the same? Let's check more details

## - Modeling the system on simso and verifying schedulability:

### Case 1: No priority

The Qt Model data window displays the configuration for five tasks. All tasks are periodic with a priority of 0. The 'Followed by' column is currently empty for all tasks.

id	Name	Task type	Abort on miss	Act. Date (ms)	Period (ms)	List of Act. dates (ms)	Deadline (ms)	WCET (ms)	Followed by	priority
1	TASK T1	Periodic	<input type="checkbox"/> No	0.0	20.0	-	20.0	2.0		0
2	TASK T2	Periodic	<input type="checkbox"/> No	0.0	100.0	-	100.0	1.5		0
3	TASK T3	Periodic	<input type="checkbox"/> No	0.0	25.0	-	25.0	3.0		0
4	TASK T4	Periodic	<input type="checkbox"/> No	0.0	10.0	-	10.0	2.5		0
5	TASK T5	Periodic	<input type="checkbox"/> No	0.0	10.0	-	10.0	1.0		0

Buttons: Remove selected task(s), Add task, Generate Task Set

### Result:

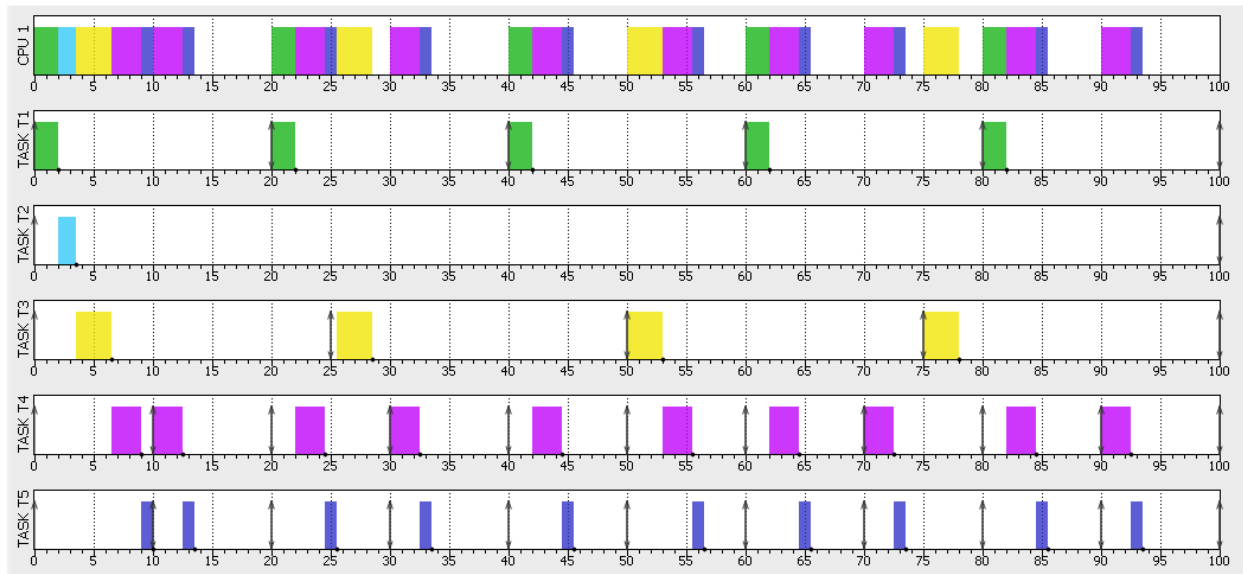
The Qt Results window displays the simulation results. The observation window is set from 0.00 to 100.00 ms. The results table shows a total load of 0.5850, a payload of 0.5850, and a system load of 0.0000 for both CPU 1 and the average.

	Total load	Payload	System load
CPU 1	0.5850	0.5850	0.0000
Average	0.5850	0.5850	0.0000

**CPU Load = 58.5 %**  
**(Matches the analytical method)**



### Gantt chart:



### Case 2: Multiple priority levels based on periodicity

Qt Model data

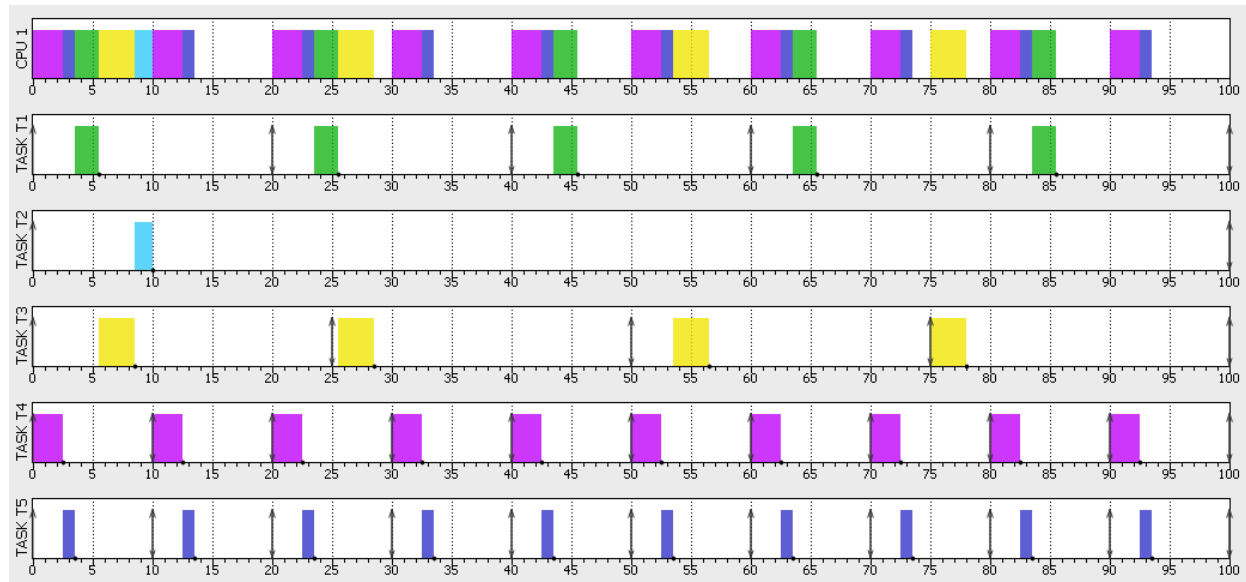
id	Name	Task type	Abort on miss	Act. Date (ms)	Period (ms)	List of Act. dates (ms)	Deadline (ms)	WCET (ms)	Followed by	priority
1	TASK T1	Periodic	<input type="checkbox"/> No	0.0	20.0	-	20.0	2.0	2	
2	TASK T2	Periodic	<input type="checkbox"/> No	0.0	100.0	-	100.0	1.5	0	
3	TASK T3	Periodic	<input type="checkbox"/> No	0.0	25.0	-	25.0	3.0	1	
4	TASK T4	Periodic	<input type="checkbox"/> No	0.0	10.0	-	10.0	2.5	3	
5	TASK T5	Periodic	<input type="checkbox"/> No	0.0	10.0	-	10.0	1.0	3	

Edit data fields...

Remove selected task(s) Add task Generate Task Set

**Result: the same regardless the case**

### Gantt chart:



Analytical results are verified!

- Introducing new tasks:

**Introducing a new task to the system to load the system:**

**T6 ( $P = D = 5$  ms,  $E = 2$  ms)**

## Case 1: No priority

Qt Model data

General Scheduler Processors Tasks

id	Name	Task type	Abort on miss	Act. Date (ms)	Period (ms)	List of Act. dates (ms)	Deadline (ms)	WCET (ms)	Followed by	priority
1	TASK T1	Periodic	<input type="checkbox"/> No	0.0	20.0	-	20.0	2.0	▼	0
2	TASK T2	Periodic	<input type="checkbox"/> No	0.0	100.0	-	100.0	1.5	▼	0
3	TASK T3	Periodic	<input type="checkbox"/> No	0.0	25.0	-	25.0	3.0	▼	0
4	TASK T4	Periodic	<input type="checkbox"/> No	0.0	10.0	-	10.0	2.5	▼	0
5	TASK T5	Periodic	<input type="checkbox"/> No	0.0	10.0	-	10.0	1.0	▼	0
6	TASK T6	Periodic	<input type="checkbox"/> No	0.0	5	-	5	2	▼	0

Edit data fields...

## Result:

Qt Results

General Logs Tasks Scheduler Processors

Observation Window:

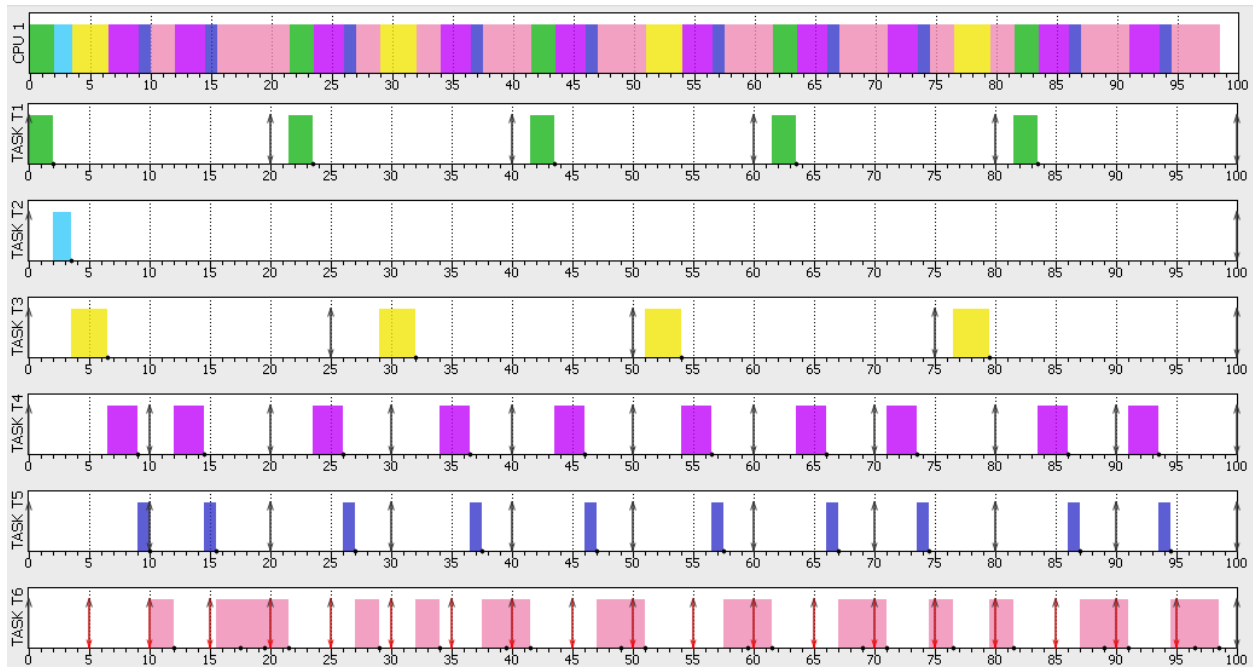
from 0.00 to 100.00 ms

	Total load	Payload	System load
CPU 1	0.9850	0.9850	0.0000
Average	0.9850	0.9850	0.0000

**CPU Load = 98.5 %**

**(Heavy loaded)**

## Gantt chart:



## Case 2: Multiple priority levels based on periodicity

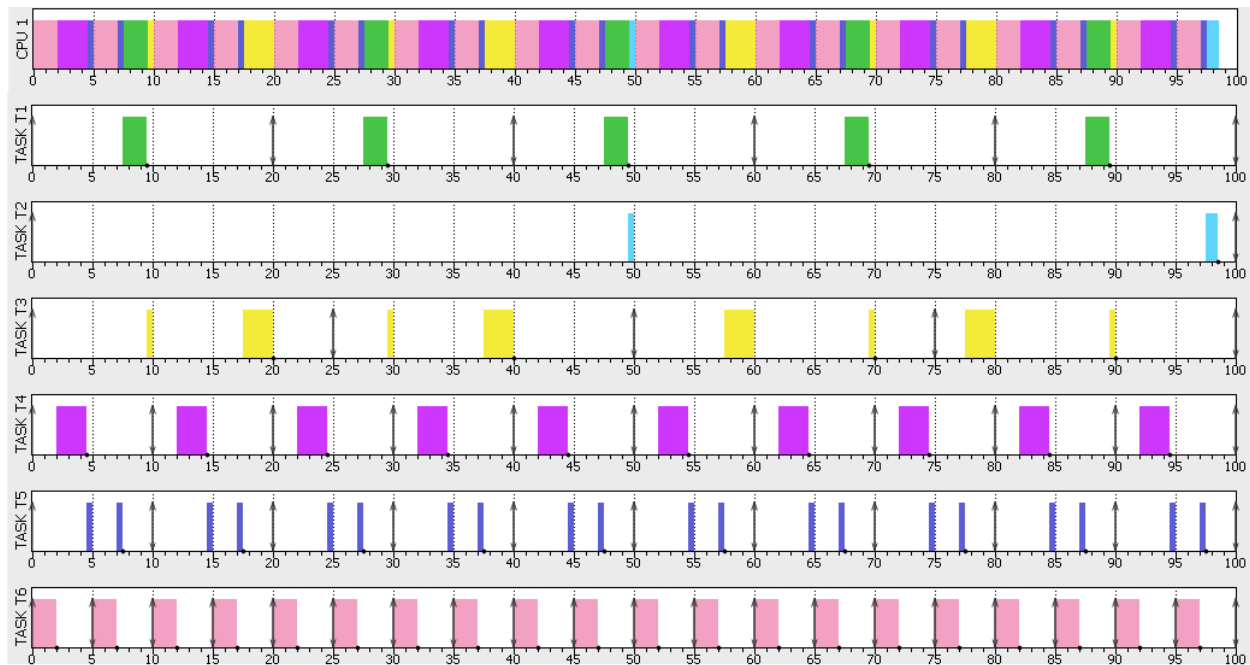
Qt Model data										
General Scheduler Processors Tasks										
id	Name	Task type	Abort on miss	Act. Date (ms)	Period (ms)	List of Act. dates (ms)	Deadline (ms)	WCET (ms)	Followed by	priority
1	TASK T1	Periodic	<input type="checkbox"/> No	0.0	20.0	-	20.0	2.0	▼ 2	
2	TASK T2	Periodic	<input type="checkbox"/> No	0.0	100.0	-	100.0	1.5	▼ 0	
3	TASK T3	Periodic	<input type="checkbox"/> No	0.0	25.0	-	25.0	3.0	▼ 1	
4	TASK T4	Periodic	<input type="checkbox"/> No	0.0	10.0	-	10.0	2.5	▼ 3	
5	TASK T5	Periodic	<input type="checkbox"/> No	0.0	10.0	-	10.0	1.0	▼ 3	
6	TASK T6	Periodic	<input type="checkbox"/> No	0.0	5	-	5	2	▼ 4	

Edit data fields...

Remove selected task(s) Add task Generate Task Set

**Result: The same regardless the case**

## Gantt chart:



## - Final comments:

**To conclude these experiments:**

- 1) Our system is healthy and schedulable regardless the usage of priority.**
- 2) When further loading the system, it's wise to start adding priority to the system so the more frequent tasks don't miss their deadline.**
- 3) Overall, preemptive system is more preferable if the system is eligible for upgrading (by adding more functionalities to the system).**