## Search-Based Test Data Generation for SQL Queries

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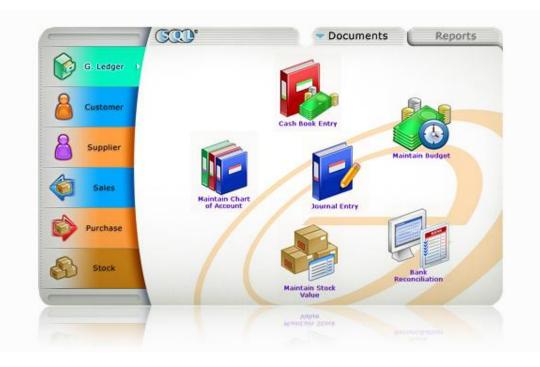
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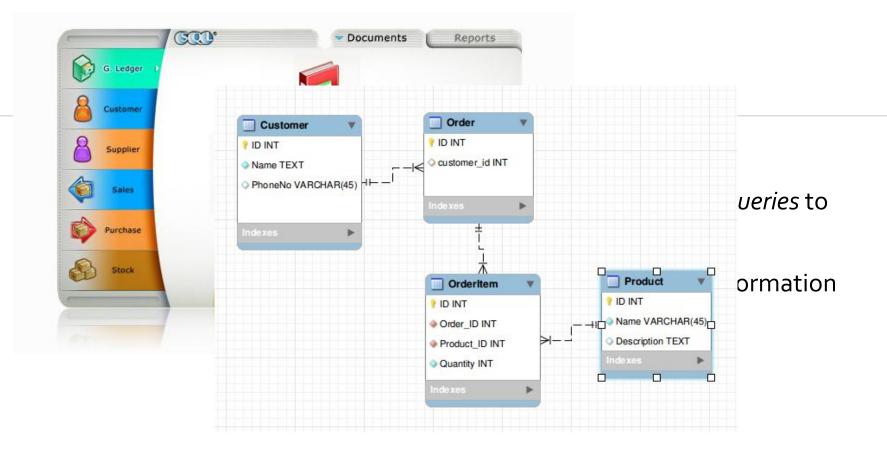
## SQL is everywhere!

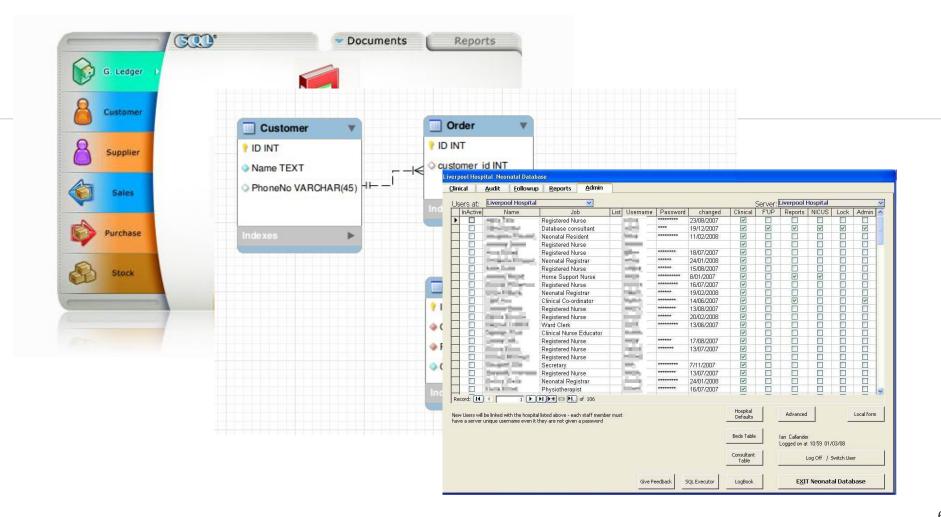
- Database-centric systems strongly rely on *SQL Queries* to manage and manipulate their data.
- Accounting, Customer management, Medical Information etc.



/ rely on *SQL Queries* to a.

ent, Medical Information







## **Related Work**

- 5 similar tools in total
- Problems:
  - Constraint Solver
  - Restriction to certain constraints
    - Subqueries
    - String predicates



# Goal: Generate Data for SQL-Queries in a Search-Based Manner

- Limits of constraints solvers:
  - No support for strings
  - Mapping from SQL to Constraints highly demanding (so no advanced SQL)
- Solution:
  - Use a search based method to move through solution space
  - Use fully functioning SQL database engine → support of whole standard SQL syntax



### **Problem Definition**

#### Given

- Query
- Database Scheme
- Time Budget

#### Use

- Random (as a baseline)
- Biased Random
- Genetic Algorithm

Database that covers the query sufficiently



## **Problem Space**

- ullet Assume a database scheme DS
- ullet Given a query Q
- We want to cover all of its coverage targets (by SQLFpc\*)

$$CT = \{ct_1, ct_2, \dots, ct_n\}$$

- Example:  $DS = \{T_1 = \{c_1 = \mathrm{ID}, c_2 = \mathrm{Category}\}\}$ 
  - Q = "SELECT \* FROM Product WHERE (Category = 'Toy')"
  - CT = {"SELECT \* FROM Product WHERE (Category = 'Toy')",

    "SELECT \* FROM Product WHERE NOT (Category = 'Toy')"}



## **Solution Space**

- Solution  $S = \{T_1, T_2, \dots, T_m\}$
- Table  $T_i = \{R_1, R_2, \dots, R_k\}$
- Row  $R_i = < V_1, V_2, \ldots, V_c >$ 
  - c is the number of columns given by the DB scheme
- Example  $DS=\{T_1=\{c_1=\mathrm{ID},c_2=\mathrm{Category}\}\}$  ct, = "SELECT \* FROM Product WHERE (Category = 'Toy')"

$$S = \{T_1 = \{<1, ``Toy">\}\}$$

id	category
1	Toy



#### **Fitness Function**

(Query Execution Plan)

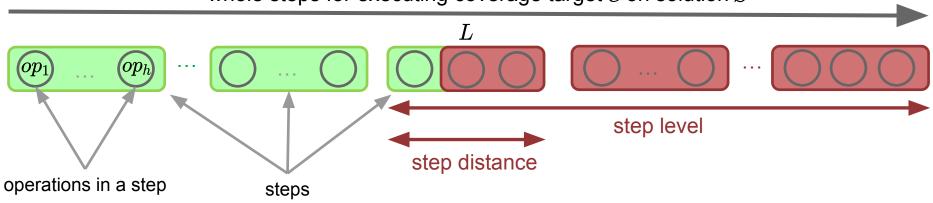
```
FROM Cars Step 1
JOIN Tires
ON Cars.tire_id = Tires.id
WHERE model = 'Ferrari' and color = 'red
Step 2
```



### **Fitness Function**

(High level view)





$$F(c,S) = step\_level(c,S) + step\_distance(S,L)$$



#### **Fitness Function**

#### (step\_distance)

- step\_distance  $(S,L)=\phi\left[\sum_{i=1}^h\phi\left[dist(S,op_i)
  ight]
  ight]$  where  $\phi(x)=rac{x}{1+x}$
- Various distance operators  $op_i$  possible:
  - Comparison operators (=, <>, >, >=, < and <=)</li>
  - Logical operators (e.g. AND)
  - SQL operators (BETWEEN, IS (NOT) NULL, IN,
     LIKE and EXISTS)



## Genetic Algorithm

- Generate population (randomly generated solutions)
- Tournament selection based on fitness function for next generation

  • Uniform crossover with two solutions

  • Mutation on one table when a solution is mutated

- > Selection-reproduction cycle ends if one achieves the coverage target



- Delete
- Insert
  - duplicate one existing row
  - o add a newly generated row
- Change
  - o integers: delta mutation
  - booleans: flipping values
  - If the column is nullable, one of the value is set to NULL with some probability

id	category		
1	Тоу		
2	Food		
3	Car		
<b>\</b>			
id	category		
1	Toy		

2

Food



- Delete
- Insert
  - duplicate one existing row
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- Change
  - o integers: delta mutation
  - o booleans: flipping values
  - If the column is nullable, one of the value is set to NULL with some probability

id	category
1	Toy
2	Food



id	category
1	Toy
2	Food
2	Food



- Delete
- Insert
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id	category
1	Toy
2	Food



id	category
1	Тоу
2	Food
3	Car



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id	category		
1	Тоу		
2	Food		
3	Car		
<b>\</b>			

id	category
1	Toy
2	Food
11	Car

## **Seeding Strategy**

- Insert values from a seeding pool into population with some probability
- Seeding pool contains all constants in the query

SELECT \* FROM Cars WHERE model = 'Ferrari' AND id = 10

#### <Random Search>

generate random solutions and repeat until the coverage target is satisfied

#### <Biased Random Search>

random search with seeding strategy



#### Input

- Query
- Database Scheme
- Time Budget

#### Output

Test data for SQLFpc\* coverage target



## **Collected queries**

totally 2,135 queries (Alura, EspCRM, SuiteCRM, ERPNext)

494 149	258 40	249 40
140	40	40
14/	10	40
704	280	279
17,761	1,631	1,567
19,108	2,209	2,135
	VC (1003   <b>V</b> (1) (1004   1004	A 00000 200 0000000 0000000000000000000

Table 1: Queries collected per application



## Configuration of the Search Parameters

The exercised configurations (in a total of 108 different combinations)

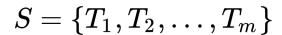
- (1) NULL mutation  $(p_{null}) = \{0.01, 0.10, 0.50\},$
- (2) Inserting, deleting mutation  $(p_m) = \{\frac{1}{3}, \frac{1}{8}\}$
- (3) Changing mutation  $(pc) = \{1/n, 1\}$ , where n is the number of rows in the mutated table.
- (4) seeding = {0.01, 0.10, 0.50}
- (5) crossover =  $\{0.0, 0.6, 0.75\}$

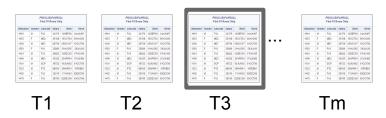
We selected the configuration with the *smallest execution time*.



## **Configuration of the Search Parameters**

- Population size = 50
- Tournament size = 4
- NULL mutation  $(p_{null}) = 0.1$
- Inserting, deleting mutation  $(p_m) = \frac{1}{3}$
- Changing mutation  $(p_c) = 1$
- Seeding = 0.5
- Crossover = 0.75
- Cloning from previous target population = 0.6





$$p_{table} = 1/K$$
, ( $K tables$ )



## **Experimental Procedure**

- executed the three approaches(Random, Biased, GA) across the
   10 executions.
- manually removed infeasible coverage targets("A > 10 and A < 10"?)</li>
   (127 out of 12,991)
- set time budget as 30 minutes for one query



## RQ1: What is the coverage achieved?

	fully covered	less than 10 cover. target, fully covered	10~20 cover. target, fully covered	more than 20 cover. target, fully covered	achieved 0% coverage
Random Search	6.5%	7.3%	0%	0%	28.33%
Biased search	90%	96.54%	49.36%	7.04%	0.05%
GA	98.64%	99.999%	94.93%	74.64%	0%

fully covered: covered in all 10 executions



## **RQ2**: What is the performance(speed)?

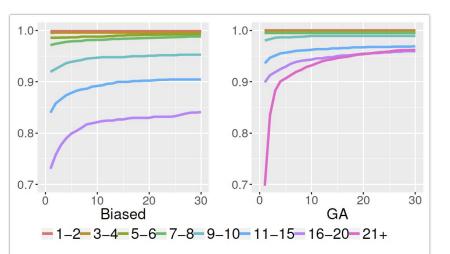


Figure 3: The average coverage of a SQL query (Y axis) in a given time budget (X axis, in minutes) controlled by the number of coverage targets. Figure better visualized in colors.

	1st quantile Media		an	3rd qu	antile	
	Biased	GA	Biased	GA	Biased	GA
1-2	0.03	0.15	0.05	0.22	0.07	0.37
3-4	0.09	0.2	0.15	0.28	0.36	0.46
5-6	0.17	0.3	0.34	0.41	0.56	0.58
7-8	1.23	0.59	2.25	0.88	7.64	1.45
9-10	1.26	0.9	5.95	1.48	100.80	2.49
11-15	7.47	2.21	74.04	3.65	844.10	8.80
16-20	83.1	7.33	844.00	15.69	1539.00	41.79
21+	1027	81.36	1688.00	155.90	1800.00	581.90

Table 4: Descriptive statistics of the biased and GA approaches' average runtime (in seconds).



## RQ3: Why not 100% coverage?

- ✓ How?
  - generated a decision tree model
  - classified whether coverage target is likely to be covered or not
- Results (accuracy random: 85.93% / BS: 90.27% / GA: 91.88%)
  - The random search can not deal with JOINs and strings.
  - The biased search suffers from queries with many predicates.
  - The size of the query impacts the GA.



## Analysis: BS vs GA

✓ Limit on Biased Random Search : complex string manipulation functions

```
    SELECT * FROM product WHERE length(name) = 12 AND left(name, 5) = 'REFRI' AND right(name, 7) = 'GERATOR'
    SELECT * FROM product WHERE reverse(name) = 'ROTAREGIRFER'
```

Additional steps in for GA are an overhead for simple queries



## Threats to Validity

- Limited diversity of queries
- Manual removal of infeasible targets
- Empirical selection of configurations for GA\*



## Threats to Validity

15 16 - 2 27 1	
27 1	
21	16 6
520	1 1
1	
12	
20 1	13 20
.07 5	51 71
	1 12 20

- ✓ Limited diversity of queries (not a threat)
- $\checkmark$  Manual removal of infeasible targets  $\rightarrow$  part of our project
- Empirical selection of configurations for GA
  - → Hyperparameter optimization (e.g. Bayesian Optimization)



- ✓ This paper, modeled the problem of test data generation for SQL queries as a **search-based problem**.
- Implemented three different search approaches in tool, named *EvoSQL* (random search, biased random search, genetic algorithm)
- Executed on 2,135 queries extracted from 4 software systems
- Achieved 98.64% Coverage with GA

## Backups





# Distance Operators (Comparison Operator)

Datatype	Distance
Numbers, Boolean	Standard Branch Distance
Strings	Enhanced Edit Distance (Standard + Character Distance
Dates	Sum of differences for each numerical calendar part



## Distance Operators (SQL Operators 1/2)

Operator	Distance
BETWEEN	$egin{aligned} lb <= v <= ub \  ightarrow lb <= v  ext{ AND } v <= ub \end{aligned}$
IS NULL	eq. to boolean $v={ m NULL}$
IN	$dist(S,op_i) = \min_{e \in  ext{list}}  ext{abs} (v-e)$
LIKE	Branch distance for pattern matching



## Distance Operators (SQL Operators 2/2)

Operator	Distance
EXISTS	Operation distance of subquery
JOIN	Branch distance for equality operator between columns



#### Change operation

- Floating-point numbers → polynomial mutation
- Integers -> delta mutation
- Date→ delta mutation to all its calendar parts
- Strings→ adding, removing or replacing characters
- Booleans→ flipping values
- If the column to mutae is nullable, one of the values is set to NULL with some probability



#### RQ3: Why not 100% coverage?

- √ J48 decision tree
  - classifies whether a coverage target is likely to be covered or not
  - features to the model
    - #of tables
    - SQL query's total # of coverage targets
    - # of SQL predicates or functions



### Search-based Test Data Generation for SQL Queries

Jeroen Castelein et al. ICSE 18'

# **Problem** Let $R = \{r_1, ..., r_m\}$ be the set of coverage targets for a query Q according to a coverage criterion. Find test data S that satisfies all coverage targets in R.

## **Solution** $T = \{T1, ..., Tn\}$ : set of tables each table $Ti = \{R1, ..., Rk\}$ : composed of references

each table  $Ti = \{R1, ..., Rk\}$ : composed of rows each row  $Rj = \{V1, ..., Vc\}$ 

(c = number of columns in Ti)

```
SELECT items .* FROM invoice
  JOIN items ON invoice.id =
        items.invoiceid
WHERE amount > 1000 OR taxFree =
        true
```

SELECT items .\* FROM invoice
JOIN items ON invoice.id =
 items.invoiceid
WHERE amount > 1000 OR taxFree =
 true

## How should we test this query?

66



#### **Fitness Function**

```
FROM Cars
JOIN Tires
ON Cars.tire_id = Tires.id
WHERE model = 'Ferrari'
Step 2
```

If Step 1 doesn't produce an output, database stops its execution before proceeding to Step 2



#### **Fitness Function**

- ✓ Satisfying coverage target means solution T yields a non-empty result. ( $T = \{T1, ..., Tn\}$  : set of tables)
- ✓ If T does not cover the given target, measure the distance of T to cover the target
  - Counting the number of yet unsatisfied steps (step level)
  - Measuring how far T is to satisfy the step where execution stopped (step distance)



#### Genetic Algorithm

- ✓ For single coverage target from a query, implement standard GA(with Crossover, Mutation).
- ✓ Using <u>seeding strategies</u>, it uses a seeding pool containing all constants appearing in the query.
- ✓ Post Processing for readability of the generating data.

There are three approaches, Random Search, Biased Random Search(seeding strategy) and, GA.

## What can we develop?

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## Inefficiency of GA in EvoSQL

calculating the fitnesses & applying the search operators

## Inefficiency of GA in EvoSQL

calculating the fitnesses & applying the search operators



Re-executed multiple times, once for each coverage target

## Problem: Inefficiency of single-target strategy

1. The order of each coverage target is not optimized

2. Inefficient allocation of the budget might happen, such as infeasible coverage target.

# **Problem** Let $R = \{r_1, ..., r_m\}$ be the set of coverage targets for a query Q according to a coverage criterion. Find test data S that satisfies all coverage targets in R.

## **Solution** $T = \{T1, ..., Tn\}$ : set of tables each table $Ti = \{R1, ..., Rk\}$ : composed of respectively.

each table  $Ti = \{R1, ..., Rk\}$ : composed of rows each row  $Rj = \{V1, ..., Vc\}$ 

(c = number of columns in Ti)



#### **Fitness Function**

#### **Distance**

Operations required to process the query and the order by which they needed to be performed

#### **Function**

Operations required to process the query and the order by which they needed to be performed



#### **Testing SQL Queries**

- Database-centric systems strongly rely on SQL Queries to manage and manipulate their data.
- How developers test SQL queries?
  - => Actually, they need to create test data by hand.

But, this can be challenging and time-consuming for complex queries.

#### Solution

Multi-Objective Optimization!





SELECT \* FROM Product WHERE (Category = 'Toy')



SELECT \* FROM Product WHERE (Category = 'Toy')

id	category		
1	Toy		
2	Food		
3	Car		

Product



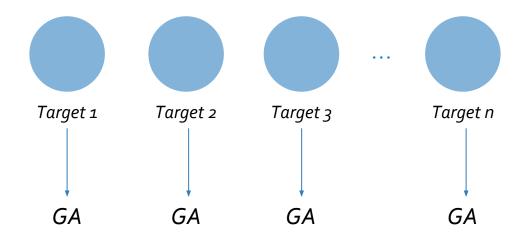
SELECT \* FROM Product WHERE (Category != 'Toy')

id	category		
1	Toy		
2	Food		
3	Car		

**Product** 

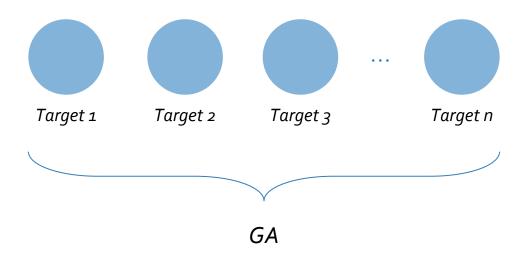


Challenge 1. The order of each coverage target being executed is not optimized



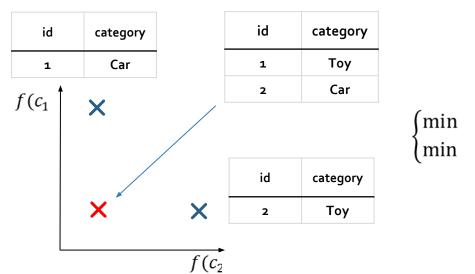


Challenge 1. The order of each coverage target being executed is not optimized





#### SELECT \* FROM Product WHERE (Category = 'Toy')



 $\begin{cases} \min & f(c_1,t) & c_1 : WHERE \ Category = 'Toy' \\ \min & f(c_2,t) & c_2 : WHERE \ Category ! = 'Toy' \end{cases}$ 



Challenge 2. Inefficient allocation of the budget might happen, such as infeasible coverage target.

"A > 10 and A < 10"?  $\rightarrow$  infeasible, no solution from GA

$$\begin{cases} \min & f(c_1,t) & f(c_1,t) => feasible \\ \min & f(c_2,t) & f(c_2,t) => Infeasible \\ \dots & \dots \\ \min & f(c_m,t) & f(c_m t) => feasible \end{cases}$$

Solution from EvoSQL is to manually remove infeasible coverage targets.



#### RQ1: What is the coverage achieved?

	all targets, fully covered	less than 10 predicates, fully covered	more than 20 predicates, fully covered	all targets, achieved 0% coverage	partial median coverage
Random Search	6.5%			28.33%	33.75%
Biased search	90%	96.54%	7.04%	0.05%	79.76%
GA	98.64%	99.999%	74.64%	0%	86.66%

fully covered: covered in all 10 executions