



# Lecture 2 Relational Database Language Part 2

关系数据库语言 第2部分

### Outline

More SQL 更多SQL知识

• Subquery 子查询

Group and Aggregation 分组与聚合

Modification 更新

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#### Movie Database

Movies (title, year, length, genre, studioName, producerC#)

MovieStar (name, address, gender, birthdate)

StarsIn (movieTitle, movieYear, starName)

MovieExec (name, address, cert#, netWorth)

Studio (name, address, presC#)

#### SELECT-FROM-WHERE

**SELECT** \*

**FROM** Movies

WHERE studioName='Disney' AND year=1990;

#### FROM

Relation(s) to which the query refers

#### WHERE

- A condition
- Tuples must satisfy the condition in order to match the query

#### SELECT

 Tells which attributes of the tuples matching the condition are produced as part of the answer

## SQL Queries and Relational Algebra

SELECT 
$$L$$
FROM  $R$ 

$$\pi_L(\sigma_C(R))$$
WHERE  $C$ 

- -L is a list of expressions
- -R is a relation
- C is a condition

## Pattern Matching in SQL (Cont'd)

```
SELECT title
FROM Movies
WHERE title LIKE 'Star ____';
```

- We remember a movie "Star something",
- And we remember that the something has four letters.
- What could this movie be?

## Pattern Matching in SQL (Cont'd)

```
SELECT title
FROM Movies
WHERE title LIKE '%"s%';
```

- Search for all movies with a possessive ('s) in their titles
- Any title with 's as a substring will match the pattern

## Pattern Matching in SQL (Cont'd)

Escape characters in LIKE expressions

```
s LIKE 'x%%x%' ESCAPE 'x'
```

- What if the pattern we wish to use in a LIKE expression involves the character % or \_?
- SQL allows us to specify any one character we like as the escape character for a single pattern
- Matches any string that begins and ends with the character
   %

## Dates and Times 不同DBMS语法差异较大

- Date constant
  - DATE '1984-05-14'
- Time constant
  - TIME '15:00:02.5'
  - TIME '12:00:00-8:00'
    - Noon in Pacific Standard Time, eight hours behind GMT
- Combine dates and times
  - TIMESTAMP '1984-05-14 12:00:00'
- Compare dates and times using <</li>

#### **Null Values**

- Null values NULL
  - Value unknown
    - An unknown birthdate
  - Value inapplicable
    - If we had a **spouse** attribute for the **MovieStar** relation, then an unmarried star might have **NULL**
  - Value withheld
    - An unlisted phone number might appear as **NULL** in the component for a **phone** attribute

# Comparisons Involving NULL

#### Two rules

- Operate on a NULL and any value, using an arithmetic operator like \* or +, the result is NULL
- Compare a NULL and any value, using a comparison operator like = or >, the result is UNKNOWN
- NULL is not a constant
  - Cannot use NULL as an operand

## Comparisons Involving NULL (Cont'd)

- Let x have the value NULL
- The value of x + 3 is also NULL
- NULL + 3 is not a legal SQL expression
- The value of x = 3 is UNKNOWN
- NULL = 3 is not correct SQL

## Comparisons Involving NULL (Cont'd)

- IS NULL
  - The correct way to ask if x has the value NULL
     x IS NULL
    - TRUE if x has the value NULL, FALSE otherwise
- IS NOT NULL
  - x IS NOT NULL
    - TRUE unless the value of x is NULL

#### The Truth-Value UNKNOWN

- Three truth-values
  - TRUE as 1, FALSE as 0, and UNKNOWN as ½
  - The rules
    - 1. The AND of two truth-values is the minimum of those values
    - 2. The OR of two truth-values is the maximum of those values
    - 3. The negation of truth-value v is 1 v

### The Truth-Value UNKNOWN (Cont'd)

#### Truth table for three-valued logic

X	у	x AND y	x OR y	NOT x
TRUE	TRUE	TRUE	TRUE	FALSE
TRUE	UNKNOWN	UNKNOWN	TRUE	FALSE
TRUE	FALSE	FALSE	TRUE	FALSE
UNKNOWN	TRUE	UNKNOWN	TRUE	UNKNOWN
UNKNOWN	UNKNOWN	UNKNOWN	UNKNOWN	UNKNOWN
UNKNOWN	FALSE	FALSE	UNKNOWN	UNKNOWN
FALSE	TRUE	FALSE	TRUE	TRUE
FALSE	UNKNOWN	FALSE	UNKNOWN	TRUE
FALSE	FALSE	FALSE	FALSE	TRUE

### The Truth-Value UNKNOWN (Cont'd)

#### Surprising behavior

Movies (<u>title, year,</u> length, genre, studioName, producerC#)

```
SELECT *
FROM Movies
WHERE length <= 120 OR length > 120;
```

- Movies tuples with NULL in the length component
- length <= 120 OR length > 120 evaluate to UNKNOWN
- Such a tuple is not returned as part of the answer to the query
- "Find all the Movies tuples with non-NULL lengths"

# Ordering the Output

The ORDER BY clause

ORDER BY < list of attributes >

- The order is by default ascending
- The keyword DESC for descending
- The keyword ASC for ascending (unnecessary)
  - The ordering is performed on the result of the FROM, WHERE, and other clauses, just before we apply the SELECT clause

# Ordering the Output (Cont'd)

Example

```
SELECT *
FROM Movies
WHERE studioName='Disney' AND year=1990
ORDER BY length, title;
```

• All the attributes of **Movies** are available at the time of sorting, even if they are not part of the **SELECT** clause

#### Products and Joins in SQL

#### Example

- Want to know the name of the producer of *Star Wars*
- Relations we need

```
Movies (<u>title, year,</u> length, genre, studioName, producerC#)
MovieExec (name, address, <u>cert#</u>, netWorth)
```

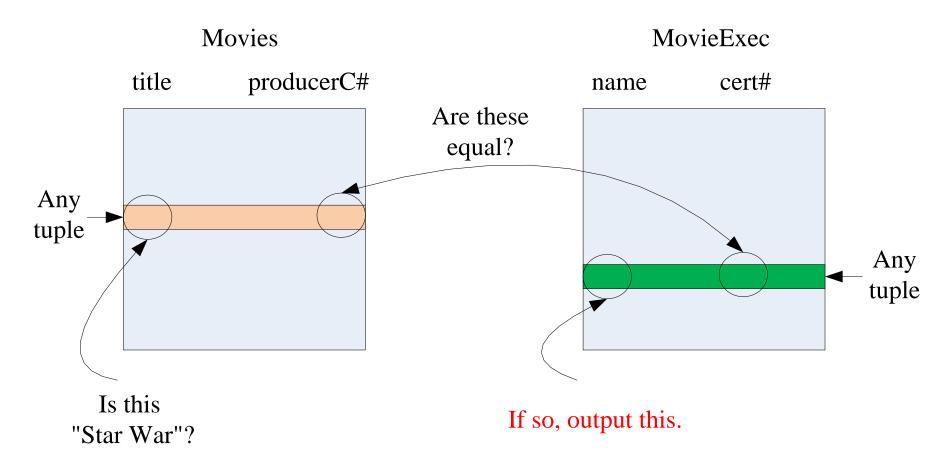
The query

**SELECT** name

FROM Movies, MovieExec

WHERE title= 'Star Wars' AND producerC# = cert#;

## Products and Joins in SQL (Cont'd)



• To pair every tuple of Movies with every tuple of MovieExec and test two conditions

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## **Tuple Variables**

#### Tuple variables

Disambiguate more than one occurrences of the same relation

#### Example

Want to know about two stars who share an address

```
SELECT Star1.name, Star2.name
FROM MovieStar Star1, MovieStar Star2
WHERE Star1.address = Star2.address
AND Star1.name < Star2.name</pre>
```

#### Interpreting Multirelation Queries (Cont'd)

- Conversion to Relational Algebra
  - FROM clause → Cartesian product
  - WHERE clause → selection
  - SELECT clause → projection

#### Interpreting Multirelation Queries (Cont'd)

- Conversion to Relational Algebra
  - Example

```
SELECT Star1.name, Star2.name
```

FROM MovieStar Star1, MovieStar Star2

WHERE Star1.address = Star2.address

AND Star1.name < Star2.name

$$\pi_{A1,A5}(\sigma_{A2=A6~\text{AND}~A1< A5}(\rho_{M(A1,A2,A3,A4)}(MovieStar) \times \\ \rho_{N(A5,A6,A7,A8)}(MovieStar)))$$

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# Subquery

#### Subquery

- A query that is part of another
- Subqueries can have subqueries, and so on

#### Subquery

- Can return a single constant, used in WHERE
- Can return relations, used in WHERE
- Can appear in FROM, followed by a tuple variable

## Subqueries that Produce Scalar Values

#### Scalar

An atomic value that can appear as one component of a tuple

#### Example

```
Movies (<u>title</u>, <u>year</u>, length, genre, studioName, producerC#)
MovieExec (name, address, <u>cert#</u>, netWorth)
```

SELECT name
FROM Movies, MovieExec

WHERE title= 'Star Wars' AND producerC# = cert#;

```
SELECT name
FROM MovieExec
WHERE cert# =

(SELECT producerC#
FROM Movies
WHERE title = 'Star Wars'
);
```

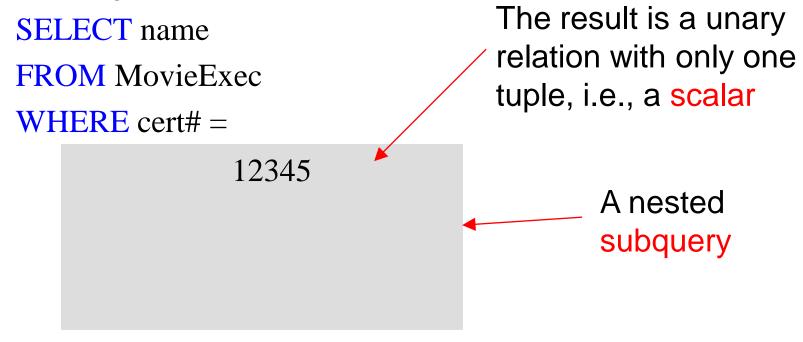
```
FROM MovieExec

WHERE cert# =

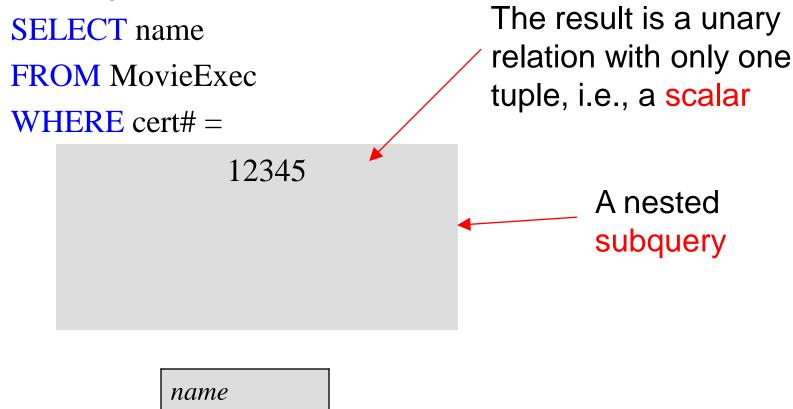
(SELECT producerC#
FROM Movies
WHERE title = 'Star Wars'
);

The result is a unary relation with only one tuple, i.e., a scalar

A nested subquery
```



#### Example



George Lucas

#### SQL operators

- Apply to a relation R
- Produce a Boolean result
- -R must be expressed as a subquery

#### 1. EXISTS R

is a condition that is true if and only if *R* is not empty

#### SQL operators

- Apply to a relation R
- Produce a Boolean result
- -R must be expressed as a subquery

#### 2. s IN R

is true if and only if *s* is equal to one of the values in *R*. Assume *R* is a unary relation

#### SQL operators

- Apply to a relation R
- Produce a Boolean result
- -R must be expressed as a subquery

#### 3. s > ALL R

is true if and only if *s* is greater than every value in unary relation *R*.

$$s < ALL R$$
,  $s < = ALL R$ ,  $s > = ALL R$ ,  $s = ALL R$ ,  $s < ALL R$  is the same as  $s NOT IN R$ 

#### SQL operators

- Apply to a relation R
- Produce a Boolean result
- -R must be expressed as a subquery

#### 4. s > ANY R

is true if and only if s is greater than at least one value in unary relation R.

$$s < ANY R$$
,  $s < = ANY R$ ,  $s > = ANY R$ ,  $s < > ANY R$ ,  $s < > ANY R$ ,  $s < ANY$ 

#### SQL operators

 EXISTS, ALL and ANY can be negated by putting NOT in front of the entire expression.

**NOT EXISTS** is true iff *R* is empty

NOT s >= ALL R is true iff s is not the maximum value in R

**NOT** s > ANY R is true iff s is the minimum value in R

## ANY或ALL谓词

- 等价关系
  - ANY和ALL谓词、聚合函数、IN谓词

	=	<b>&lt;&gt;</b>	<	<=	^	>=
ANY	IN		<max< th=""><th>&lt;=MAX</th><th>&gt;MIN</th><th>&gt;=MIN</th></max<>	<=MAX	>MIN	>=MIN
ALL		NOT IN	<min< th=""><th>&lt;=MIN</th><th>&gt;MAX</th><th>&gt;=MAX</th></min<>	<=MIN	>MAX	>=MAX

### Example

```
SELECT name
FROM MovieExec
WHERE cert# IN
    (SELECT producerC#
    FROM Movies
    WHERE (title, year) IN
         (SELECT movieTitle, movieYear
         FROM StarsIn
         WHERE starName = 'Harrison Ford'
```

Main query Example **SELECT** name Subquery FROM MovieExec WHERE cert# IN (SELECT producerC# A third **FROM** Movies subquery WHERE (title, year) IN (SELECT movieTitle, movieYear FROM StarsIn WHERE starName = 'Harrison Ford' 39

#### Example

```
SELECT name
FROM MovieExec
WHERE cert# IN

(SELECT producerC#
FROM Movies
WHERE (title, year) IN
```

title	year
Star Wars	1977
Raiders of the Lost Ark	1981
The Fugitive	1993

);

#### Example

**SELECT** name

FROM MovieExec

WHERE cert# IN	producerC#		
	12345		
	23456		
	34567		
	•••		

title	year
Star Wars	1977
Raiders of the Lost Ark	1981
The Fugitive	1993

## Conditions Involving Tuples (Cont'd)

#### Example

A single select-from-where expression

```
FROM MovieExec, Movies, StarsIn

WHERE cert# = producerC# AND

title = movieTitle AND

year = movieYear AND

starName = 'Harrison Ford';
```

## Correlated Subqueries

- Correlated subquery 相关子查询
  - Requires the subquery to be evaluated many times
    - Once for each assignment of a value to some term in the subquery that comes from a tuple variable outside the subquery

## Correlated Subqueries: Example

#### Example

Find the titles that have been used for two or more movies

```
SELECT title
FROM Movies Old
WHERE year < ANY
(SELECT year
FROM Movies
WHERE title = Old.title
```

A movie made twice will be listed once, a movie made three times will be listed twice, and so on.

For each such tuple, we ask in the subquery whether there
 is a movie with the same title and a greater year

# Subqueries in FROM Clauses

#### Example

```
SELECT name

FROM MovieExec, (SELECT producerC#

FROM Movies, StarsIn

WHERE title = movieTitle AND

year = movieYear AND

starName = 'Harrison Ford'

) Prod

WHERE cert# = Prod.producerC#;
```

# Subqueries in FROM Clauses

• Example

SELECT name

FROM MovieExec,

(SELECT producerC#

FROM Movies, StarsIn

WHERE title = movieTitle AND

year = movieYear AND

starName = 'Harrison Ford'

) Prod

WHERE cert# = Prod.producerC#;

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# **Full-Relation Operations**

Eliminating Duplicates: DISTINCT

Grouping

Aggregation

HAVING Clauses

# Grouping and Aggregation

#### Grouping

- Partition the tuples of a relation into "groups",
- based on the values of tuples in one or more attributes

#### Aggregation

- Aggregate certain other columns of the relation by applying "aggregation" operators to those columns
- The aggregation is done separately for each group

# Aggregation Operators

Five aggregation operators in SQL

Sum of a column with numerical values

AVG Average of a column with numerical values

MIN Smallest value of a column with numerical values

MAX Largest value of a column with numerical values

COUNT Number of (not necessarily distinct) values in a column

# **Aggregation Operators**

- Five aggregation operators in SQL
  - NULL value

```
遇到空值

•处理

   COUNT(*)

•忽略

   COUNT、SUM、AVG、MAX、MIN
```

## Aggregation Operators: Example

#### Example

Finds the average net worth of all movie executives

MovieExec (name, address, cert#, netWorth)

SELECT AVG(netWorth)

FROM MovieExec;

## Aggregation Operators: Example

#### Example

Counts the number of tuples in the StarsIn
 StarsIn (movieTitle, movieYear, starName)

SELECT COUNT(\*)
FROM StarsIn;

The similar query

Do these two queries have the same result?

SELECT COUNT(starName)
FROM StarsIn;

## Aggregation Operators: Example

#### Example

Counts the number of tuples in the StarsIn
 StarsIn (movieTitle, movieYear, starName)

SELECT COUNT(\*)
FROM StarsIn;

Duplicate values are eliminated before we count.

The similar query

SELECT COUNT(DISTINCT starName)
FROM StarsIn;

# Grouping

- The GROUP BY clause
  - Followed by a list of grouping attributes
  - Has tuples grouped according to their values in the grouping attributes
  - Whatever aggregation operators are used in the SELECT clause are applied only within groups

# Grouping: Example

#### Example

Movies (<u>title, year,</u> length, genre, studioName, producerC#)

- Finds the sum of the lengths of all movies for each studio

**SELECT** studioName, **SUM**(length)

**FROM** Movies

**GROUP BY** studioName;

# Grouping: Example

#### Example

**SELECT** studioName, **SUM**(length)

**FROM** Movies

**GROUP BY** studioName;

		studioName	
		Disney	
		Disney	
		Disney	
		MGM	
		MGM	
		0	
		o	

# Grouping: The SELECT Clause

**SELECT** studioName, **SUM**(length)

**FROM** Movies

**GROUP BY** studioName;

- Two kinds of terms in the SELECT clause
  - Aggregations
    - An aggregate operator is applied to an attribute
    - These terms are evaluated on a per-group basis
  - Attributes in the GROUP BY clause
    - Only those attributes may appear unaggregated in the SELECT clause

# Grouping: The SELECT Clause

**SELECT** studioName, **SUM**(length)

**FROM** Movies

**GROUP BY** studioName;

Two kinds of terms in the SELECT clause

#### 规律

- •出现在SELECT后面的列,
  - · 要么是GROUP BY后面的分组列,
  - 要么是对其他列应用聚合函数

# Grouping (Cont'd)

No aggregations

SELECT studioName
FROM Movies
GROUP BY studioName;

Has the same effect as

SELECT DISTINCT studioName FROM Movies;

### **HAVING** Clauses

- WHERE clause
  - Restrict the tuples prior to grouping
    - Only wanted the total length of movies for producers with a net worth of more than \$10,000,000

```
SELECT name, SUM(length)
FROM MovieExec, Movies
WHERE producerC# = cert# AND netWorth > 10000000
GROUP BY name;
```

# **HAVING** Clauses (Cont'd)

- HAVING clause
  - We want to choose groups based on some
     aggregrate property of the group itself

```
GROUP BY < grouping attributes>
HAVING < condition about the group>
```

# HAVING Clauses (Cont'd)

#### Example

 Print the total film length for only those producers who made at least one film prior to 1930

SELECT name, SUM(length)
FROM MovieExec, Movies
WHERE producerC# = cert#
GROUP BY name
HAVING MIN(year) < 1930

The resulting query would remove from the grouped relation all those groups in which every tuple had a year component 1930 or higher.

# **HAVING** Clauses (Cont'd)

- Two rules about HAVING clauses
  - 1. An aggregation in a HAVING clause applies only to the tuples of the group being tested
  - 2. Any attribute of relations in the FROM clause may be aggregated in the HAVING clause, but only those attributes that are in the GROUP BY list may appear unaggregated in the HAVING clause
    - The same rule as for the **SELECT** clause

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### **Database Modifications**

- Three types
  - 1. Insert tuples into a relation

- 2. Delete certain tuples from a relation
- 3. Update values of certain components of certain existing tuples

### Insertion

#### Basic form

INSERT INTO 
$$R(A_1, ..., A_n)$$
 VALUES  $(v_1, ..., v_n)$ ;

- A tuple is created using the value  $v_i$  for attribute  $A_i$ , for i = 1, 2, ..., n
  - If the list of attributes does not include all attributes of the relation *R*, then the tuple created has default values for all missing attributes

# Insertion: Example

#### Example

```
INSERT INTO StarsIn(movieTitle, movieYear, starName) VALUES('The Maltese Falcon', 1942, 'Sydney Greenstreet');
```

 If we provide values for all attributes of the relation, then we may omit the list of attributes that follows the relation name.

```
INSERT INTO StarsIn VALUES('The Maltese Falcon', 1942, 'Sydney Greenstreet');
```

# Insertion: Subquery

#### Example

 Add to the relation **Studio** all movie studios that are mentioned in the relation **Movies**, but do not appear in **Studio**

```
INSERT INTO Studio(name)

SELECT DISTINCT studioName
FROM Movies

WHERE studioName NOT IN

(SELECT name
FROM Studio);
```

### Deletion

Form

DELETE FROM *R* WHERE *<condition>*;

 Every tuple satisfying the condition will be deleted from relation R

# Deletion: Example

#### Example

StarsIn (movieTitle, movieYear, starName)

Delete the fact that Sydney Greenstreet was a star in The Maltese Falcon

```
DELETE FROM StarsIn

WHERE movieTitle = 'The Maltese Falcon' AND

movieYear = 1942 AND

starName = 'Sydney Greenstreet';
```

# Deletion: Example

Example

StarsIn (movieTitle, movieYear, starName)

- Delete all tuples in the relation StarsIn
  - Make the relation **StarsIn** empty

**DELETE FROM StarsIn**;

## Updates

- Updates in SQL
  - One or more tuples that already exist in the database have some of their components changed
- Form

```
UPDATE R SET <new-value assignments>
WHERE <condition>;
```

## Updates: Example

#### Example

Modify the relation

MovieExec (name, address, <u>cert#</u>, netWorth)

• by attaching the title Pres. In front of the name of every movie executive who is the president of a studio

```
UPDATE MovieExec
SET name = 'Pres. ' || name
WHERE cert# IN (SELECT presC# FROM Studio);
```





## ...The End of This Lecture...