INTRODUCTION TO COMPUTER SYSTEMS

Tanzir Ahmed CSCE 313 Spring 2020

Lecture Outline

- Part1
 - Course Objectives and Outcome
 - Logistics
- Part2
 - Course Context

Essence of the Course

- To learn the software interface to the Operating Systems with a view to becoming more educated and efficient programmers
- This is, NOT an Operating System (OS) course
 - Can take CSCE 410: Operating Systems
- However, the content will discuss OS quite a bit from a programmer's perspective, because:
 - We want to <u>harness OS basics</u> to write high-performance, inter-dependent, and hardware and OS-aware applications
 - We can also <u>borrow techniques</u> used in OS as a huge software and apply those to everyday software development practices
 - For instance, the concept of timers in OS is crucial for Asynchronous programming required in real-time systems

Why do we need OS in 313?

- Some of you will design and build OS's or components of them
 - Perhaps more now than ever, because there are many different OS's targeted for specific devices (e.g., robots, servers, signal lights, cameras, vehicle infotainment systems)
- Many of you will create systems that utilize the core concepts in operating systems
 - The concept of "hierarchy" and "abstraction" are all too imp
 - Whether you build software or hardware
 - The concepts and design patterns appear at many levels
- All of you will build applications, etc. that utilize operating systems
 - The better you understand their design and implementation, the better use you'll make of them.

Learning Outcome

- In this course, you will learn how to use different OS services and features:
 - What is an operating system; its components; how it works
 - Execution of a program; function calls; interrupts; system calls; process control
 - File system
 - Concurrency and Synchronization
 - Inter-Process Communication (IPC)
 - Network/socket programming
 - Network Threats/Security Basics

CSCE 313 Instruction Team

- Instructor: Dr. Sarker Tanzir Ahmed
 - Instructional Assistant Professor
 - Background: PhD from TAMU CSE in High Performance Computing, Big Data, Web Crawling
- Teaching Assistant:
 - Minhwan Oh (501-4)
 - Di Xiao (505-8)
 - Ananya Bothra (509-12)

Textbook, Reference Books

■ Text:

- Operating Systems: Principles and Practice, Second Edition,
 Thomas Anderson and Michael Dahlin, Recursive Books, 2014.
- Operating Systems: Three Easy Pieces, Remzi H. Arpaci-Dusseau and Andrea C. Arpaci-Dusseau, available online at: http://pages.cs.wisc.edu/~remzi/OSTEP/

■ Reference:

- Main: Operating Systems Internals and Design Principles,
 William Stallings, Ninth Edition
- Secondary: Understanding Unix/Linux Programming A Guide to Theory and Practice, Bruce Molay, Pearson Education Inc., 2003

Other Interesting Readings

 Computer Systems: A Programmer's Perspective, Randal E. Bryant and David R. O'Hallaron, Prentice Hall, 2011

How Success will be Measured

- The course will have several quizzes/mini assignments, two exams, and a series of Programming Assignments (PA). The grade allocation is as follows:
- Total 100 points
 - Quizzes (4-5): 15 points
 - 2 Exams: 35 points
 - Programming Assignments (5-6): 48 points
 - Attendance (lab + lecture): 2 points
- The grading scale is as follows:
 - 90 100: A
 - 80 89: B
 - 70 79: C
 - 60 69: D
 - 59 or below: F

Exams

- Two exams: midterm 15%, final 20%
- Exams are closed book
- Cheat-sheets are allowed
 - A4 page, double-sided, hand-written
- Final Exam is NOT cumulative
 - Will be based on material covered after the midterm
- The exams are relatively "hard"
 - Sample exams from previous semesters will ease them significantly

Programming Assignments

- There will be roughly 7 (+/- 1) Programming Assignments (PA) assigned weekly or bi-weekly basis.
- Each PA will be individual effort
 - You are not allowed to work in teams
 - You are not supposed to collaborate with other students on any peic
- Worth 48% of your grade
- PAs are posted and submitted through Ecampus

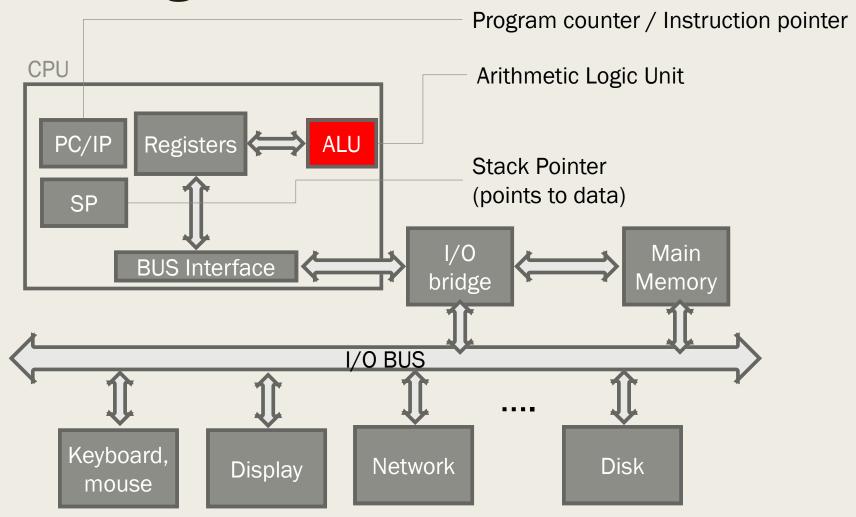
Late Policy for Programming Assignments

- Penalty: Unless stated otherwise, lateness is penalized as a simple linear function of hours late
 - 0.5% /hour of your score
- As a result on late submissions you loose:
 - 12% / day
 - 100% if more than 8 days late

Lecture Outline

- Part1
 - Course Objectives and Outcome
 - Logistics
- Part2
 - Introduction to Computer Systems
 - Background What we know already
 - What we are going to learn

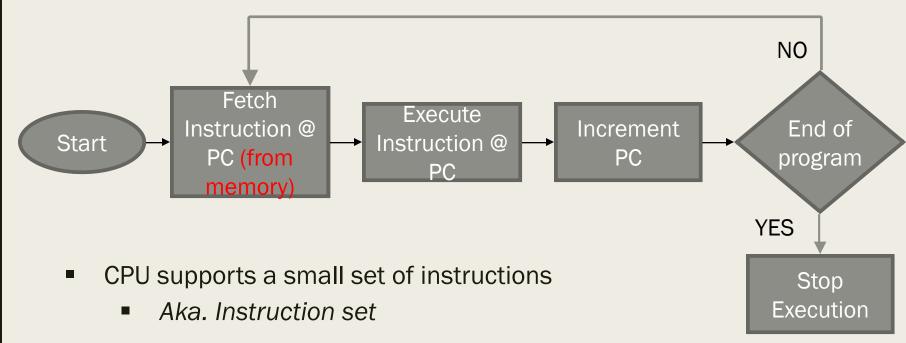
Background – From 312



Simplified hardware organization of a typical system

Instruction Execution

Simple Instruction execution model



 Now, we must think about how to run a program using this model

Before Even That: Preparing a Program to Run

- Write the program in a high-level language e.g., C
- Convert the human-readable C program to an assembly program
 (i.e., hello.as) by Compilation
 - Assembly program still human-readable
- Assemble the program into relocatable machine code (hello.o)
- Link the program to other required libraries that are precompiled and make a single executable (hello)

printf.o Sequence Humanof machine readable hello.s hello.c hello.o hello instructions file stored stored in a in disk disk file compile assemble link

Background – Loading the hello program

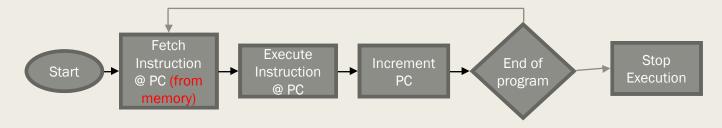
- A program although stored in the disk, cannot run while in the disk
 - Why?? Disks are million times slower than a CPU speed
 - It must be loaded to the main memory before running
- Load the executable program hello (its code and data) into main memory
- Set the PC to the first byte of the first instruction of hello
- Let the CPU do what it does execute an instruction, proceed to the next one and repeat

Background – A Few Questions

- Who loads a program into main memory?
 - A part of the Operating System called loader
 - But what if you like another faster loader from another vendor
- Alternatives??
 - One solution, each programmer write his own loader. But is that feasible?
 - How about compiler, linker, assembler etc.?
 - How about other common utilities (e.g., file explorer, shell)
- Seems like all would be best IF these utilities can be obtained from third party easily, but they go through a supervisory program called the Operating System (OS)
 - The OS makes sure that everything is working safely and securely
- That is the whole point of an OS

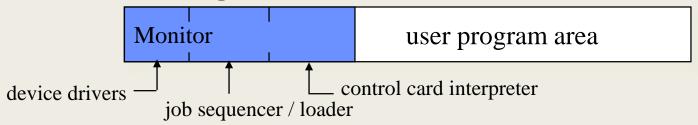
A Short Historical Tour

- **First Generation** Computer Systems (1949-1956):
 - Single user: writes program, operates computer through console or card reader / printer
 - Absolute machine language
 - No concept of an Operating Systems
 - Hardware simple, but not the programs written on it
 - Programmer must compile, load, and run program repeat



Second-Generation Computers (1956-1963)

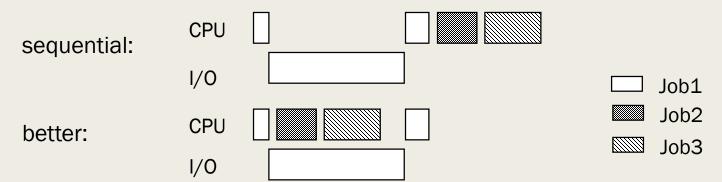
- Automation of Load/Translate/Load/Execute
 - Batch systems
 - Monitor programs



- Problem: Resource management and I/O still under control of programmer. Issues??
 - Memory protection, security etc.
 - What if one program crashes

Third-Generation Computer Systems (1964-1975)

Problem with batching: one-job-at-a-time



Solution: Multiprogramming

- Job pools: have several programs ready to execute
- Keep several programs in memory



New issues:

- Job scheduling,
- Memory management
- Protection

I/O Does not Involve CPU?? Why and How?

- At early stages of computing, CPU was in fact in charge of I/O operations. Let us consider an example
- A 1GHz CPU trying to print a page in the printer (an I/O device), which can print 1 page/sec
- If the CPU had to print 1 page at a time, its rate of operation would reduce from 1B instructions/sec to 1 instruction/sec
 - A 1B-fold slowdown
- To read a byte from disk, which requires >1ms, the CPU would go to 1000 instructions/s
 - A 1M times slowdown
- Clearly, there is room for improvement

Direct Memory Access (DMA)

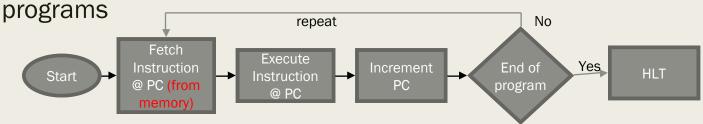
- Instead of involving CPU in every step of I/O, another hardware component called the <u>DMA controller</u> drives the device
 - The CPU only starts the transfer by sending a high-level instruction to DMA controller
 - The DMA controller does the I/O by copying data to/from the given location
 - The device controller speaks to the DMA controller along the way, and does not bother the CPU
 - At the end, the DMA controller notifies CPU that the transfer is done
- This saves a lot of CPU cycles
- This is called **Direct Memory Access (DMA)**
- This is the primary enabler for Multiprogramming

Time Sharing (mid 1960s on)

- Remote interactive access to computer: "Computing as Utility"
- OS interleaves execution of multiple user programs with time quantum
 - CTSS (1961): time quantum 0.2 sec
- User returns to own the machine
- New aspects and issues:
 - On-line file systems
 - resource protection
 - virtual memory
 - sophisticated process scheduling
- Advent of systematic techniques for designing and analyzing OSs.

Multiprogramming(MP)/Time Sharing (TS) – Implementation?

So far, we only have this, which works perfect for sequential



- How do extend this to support either MP or TS
 - Clearly, we need a mechanism to "kick" a program out at some point from the CPU
 - And "bring it back" later into the CPU
- In summary, we need "Asynchrony" in programs

When to Kick-out a Program

Program

Generated by some condition that occurs as a result of an instruction execution, such as arithmetic overflow, division by zero, attempt to execute an illegal machine instruction, and reference outside a user's allowed memory space.

For time sharing

Timer

Generated by a timer within the processor. This allows the operating system to perform certain functions on a regular basis.

For multiprogramming

I/O

Generated by an I/O controller, to signal normal completion of an operation or to signal a variety of error conditions.

Hardware failure Generated by a failure, such as power failure or memory parity error.

- Conceptually, what we need is this:
- Here come "Interrupts"

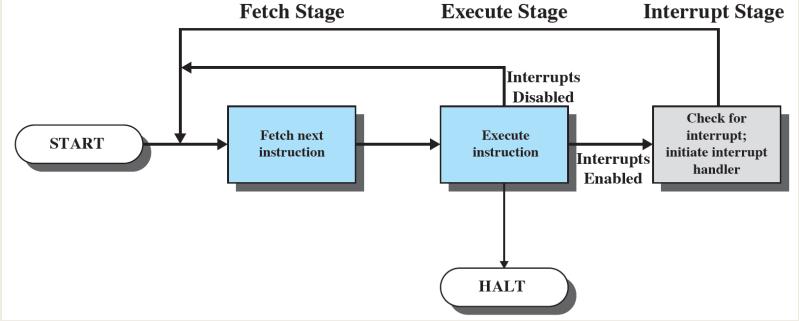
User Program

Interrupt Handler ioccurs here i+1 M

Source: W. Stallings book

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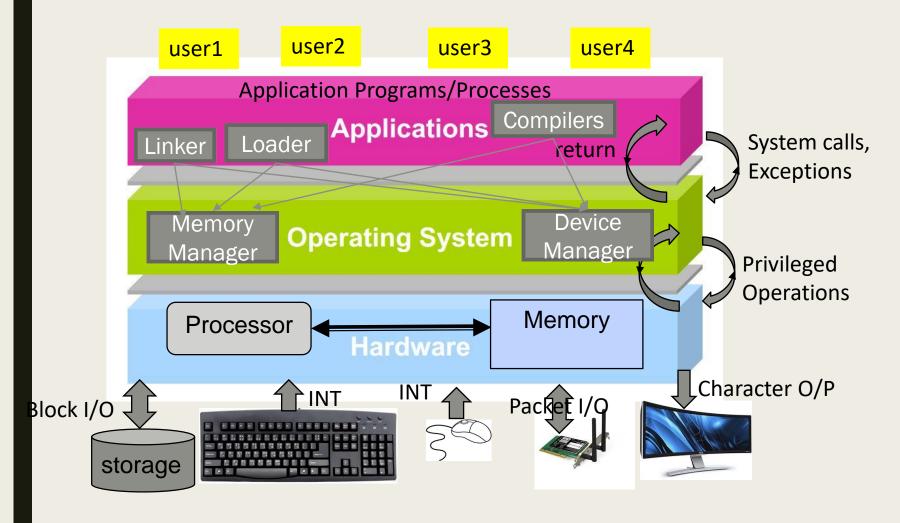
The Computing Model Changes



Computing Model to Support Interrupts

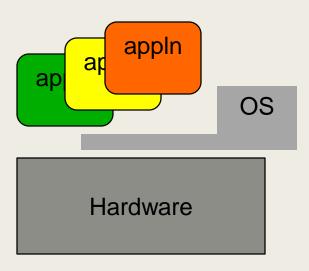
Source: W. Stallings book

A Modern Computer System



What is an operating system?

- Special layer of software that provides application software access to hardware resources
 - Convenient abstraction of complex hardware devices
 - Protected access to shared resources
 - Security and authentication
 - Communication amongst logical entities



Credits

- Slides are based on Dr. Aakash Tyagi and Dr. Riccardo Bettati's lectures
- Programming Assignments are mostly taken from Dr. Bettati
- Other sources are credited along the way