## Checking *expression validity* in C++11/14/17

Meeting[0] - C++::London - 05/12/2016

by Vittorio Romeo (@supahvee1234)

```
struct Cat
{
    void meow() const { cout << "meow\n"; }
};

struct Dog
{
    void bark() const { cout << "bark\n"; }
};</pre>
```

```
template <typename T>
void pet(const T& x)
\{
    // Pseudocode:
    if( <`x.meow()` is well-formed> )
        x.meow();
    else if( <`x.bark()` is well-formed> )
        x.bark();
    else
        <compile-time error>
```

```
pet(Cat{}); // "meow"
pet(Dog{}); // "bark"
pet(int{}); // compile-time error
```

```
template <typename T>
void pet(const T& x){ /* ? */ }
```

- C++11: std::void\_t and std::enable\_if.
- C++14: boost::hana::is\_valid and vrm::core::static\_if.
- C++17: if constexpr(...), constexpr lambdas, and std::is\_callable.

$$C + +11$$

std::void\_t and std::enable\_if

Combining std::void\_t with std::enable\_if allows us to detect ill-formed expressions in SFINAE contexts.

"Modern Template Metaprogramming: A Compendium"
 by Walter E. Brown at CppCon 2014

```
template <typename ... >
using void_t = void;
```

( void\_t was standardized in C++17, but can be implemented in C++11.)

```
template <typename, typename = void>
struct has_meow : std::false_type { };

template <typename T>
struct has_meow
<
    T,
    void_t<decltype(std::declval<T>().meow())>
> : std::true_type { };
```

Instantiating has\_meow<T> will attempt to evaluate

void\_t<decltype(std::declval<T>().meow())>

If declval<T>().meow() is well-formed,

```
void_t<decltype(std::declval<T>().meow())>
```

will evaluate to void , and has\_meow 's std::true\_type
specialization will be taken.

If declval<T>().meow() is ill-formed,

```
void_t<decltype(std::declval<T>().meow())>
```

will be ill-formed as well, SFINAE-ing away the std::true\_type specialization. All that's left is the std::false\_type specialization.

```
template <typename, typename = void>  // <<<
struct has_meow : std::false_type { }; // <<</pre>
template <typename T>
struct has_meow<T, void_t<decltype(std::declval<T>().meow())>>
: std::true_type { };
```

How does void\_t work?

```
auto meows = has_meow<Cat>{};
```

...is equivalent to...

```
auto meows = has_meow<Cat, void>{};
```

...because of the default void parameter.

The compiler takes into account partial specializations.

- T is matched to T
- void is matched to

```
void_t<decltype(std::declval<T>().meow())>
```

If void\_t<decltype(std::declval<T>().meow())> doesn't get SFINAE-d away, both the types of the original templates and the specialization match.

The partial specialization is prioritized and chosen.

If the original template's second parameter wasn't defaulted to **void**, it would have been chosen over the partial specialization!

After defining the has\_bark detector class (which is trivial to implement, as well), all that's left to do is use std:: enable\_if to constrain pet.

```
template <typename T>
auto pet(const T& x)
    → typename std::enable_if<has_meow<T>{}>::type
    x.meow();
template <typename T>
auto pet(const T& x)
    → typename std::enable_if<has_bark<T>{}>::type
    x.bark();
```

## C++11 implementation issues:

- A detector class must be defined for every expression to check.
  - The class cannot be defined locally.
- std::enable\_if must be used to constrain multiple versions of the same function.
  - It is not possible to "branch" locally at compile-time.
  - It is necessary to **repeat** the function signature.

$$C + +14$$

boost::hana::is\_valid and vrm::core::static\_if

We can solve the aforementioned issues thanks to



Generic lambdas are "templates in disguise" - they provide a SFINAE-friendly context.

```
auto l = [](auto x){ return x; };
```

...is somewhat equivalent to...

```
struct αμσμησυ2
{
    template <typename T>
    auto operator()(T x) const { return x; }
};
```

Let's take advantage of that...

```
auto has_meow =
    is_valid([](auto& x) → decltype(x.meow()){});

static_assert(has_meow(Cat{}), "");
static_assert(!has_bark(Cat{}), "");
```

- has\_meow can be locally instantiated in any scope.
- Terser & nicer syntax.

How can we implement is\_valid?

```
template <typename TF>
constexpr auto is_valid(TF)
{
   return validity_checker<TF>{};
}
```

```
template <typename TF>
struct validity_checker
{
   template <typename ... Ts>
   constexpr auto operator()(Ts& ...) const
   {
      return is_callable<TF(Ts ...)>{};
   }
};
```

Remember that TF is the type of:

```
[](auto& x) \rightarrow decltype(x.meow()) { }
```

How can we implement is\_callable ?

```
template <typename, typename = void>
struct is_callable : std::false_type { };
template <typename TF, class ... Ts>
struct is_callable
<
    TF(Ts ...),
    void_t<decltype(</pre>
        std::declval<TF>()(std::declval<Ts>() ... )
    )>
: std::true_type { };
```

is\_valid solves the first C++11 annoyance.

• Detectors can be instantiated locally.

Solving the second issue (*local compile-time branching*) is slightly more complicated, but **it can be done**.

I explain how in my CppCon 2016 talk:
"Implementing static control flow in C++14".

```
template <typename T>
auto pet(const T& x)
    auto has_meow = is_valid([](auto& x)
        \rightarrow decltype(x.meow()){};
    auto has_bark = is_valid([](auto& x)
        → decltype(x.bark()){ });
    static_if(has_meow(x))
        .then([&x](auto){ x.meow(); })
        .else_if(has_bark(x))
        .then([&x](auto){ x.bark(); })
        .else_([](auto)
                struct cannot_meow_or_bark;
                cannot_meow_or_bark{};
            })();
```

Is it *nicer* than the C++11 version? **Debatable**.

There are some objective advantages, though:

- Expression validity detector instantiation is local to the function scope.
- There is a **single overload** of pet .
- Compile-time branching is **local to the function scope**.

- boost::hana:is\_valid is a production-ready C++14 implementation of the above is\_valid function.
- You can find my static\_if implementation in vrm::core::static\_if.

## C++14 implementation issues:

- is\_valid has to be assigned to a variable in order to be used in a constant expression.
  - This happens because lambdas are not constexpr.

## Verbosity.

- static\_if makes the code much less readable.
- Having to create a lambda with a decltype(...) trailing return type.

## C + +17

- if constexpr(...)
- constexpr lambdas
- std::is\_callable
- Variadic macro black magic

```
template <typename T>
auto pet(const T& x)
    if constexpr(IS_VALID(_0.meow())(T))
        x.meow();
    else if constexpr(IS_VALID(_0.bark())(T))
        x.bark();
    else
        struct cannot_meow_or_bark;
        cannot_meow_or_bark{};
```

IS\_VALID(\_0.meow())(T) is a variadic macro that:

- Takes an expression built with type placeholders.
- Takes some types.
- Evaluates to true if the expression is valid for the given types.

```
// Can `T` be dereferenced?
IS_VALID(*_0)(T);

// Can `TO` and `T1` be added together?
IS_VALID(_0 + _1)(T0, T1);

// Can `T` be streamed into itself?
IS_VALID(_0 << _0)(T);

// Can a tuple be made out of `TO`, `T1` and `float`?
IS_VALID(std::make_tuple(_0, _1, _2))(T0, T1, float);</pre>
```

IS\_VALID can be used in contexts where only a
constant expression is accepted such as
static\_assert(...) or if constexpr(...).

What is this magic!?

Let's begin by defining some utilities...

```
template <typename T>
struct type_w
{
   using type = T;
};

template <typename T>
constexpr type_w<T> type_c{};
```

type\_c is a *constexpr* variable template that wraps a type into a value.

```
constexpr auto wrapped_int = type_c<int>;
using unwrapped_int =
   typename decltype(wrapped_int)::type;
```

type\_c is useful because it can be passed to *template functions* like a regular value, **retaining the type information**.

```
template <typename TF>
constexpr auto is_valid(TF)
{
   return validity_checker<TF>{};
}
```

Identical to the previous version.

```
template <typename TF>
struct validity_checker
{
    template <typename ... Ts>
    constexpr auto operator()(Ts ... ts)
    {
        return std::is_callable<
            TF(typename decltype(ts)::type ...)
        >{};
    }
};
```

Expects a bunch of type\_c values, then unwraps them into std::is\_callable (standardized in C++17).

The lambda passed as TF is almost identical as well:

```
is_valid([](auto _0) constexpr \rightarrow decltype(_0.meow()){}
```

...but there's a constexpr in there!

## **constexpr** lambdas were standardized in C++17.

```
// Make sure that `int*` can be dereferenced.
static_assert(
   is_valid([](auto _0) constexpr → decltype(*_0){})
   (type_c<int*>)
);
```

This technique is very verbose...

```
template <typename T>
auto pet(const T& x)
    if constexpr(is_valid([](auto _0) constexpr
        \rightarrow decltype(_{0}.meow())(T))
        x.meow();
    else if constexpr(is_valid([](auto _0) constexpr
         \rightarrow decltype(_0.bark())(T))
        x.bark();
    else
        struct cannot_meow_or_bark;
         cannot_meow_or_bark{};
```

That's why we need a macro.

...with some vrm\_pp preprocessor metaprogramming and some pre-generated code...

```
#define IS_VALID( ... ) \
    VRM_PP_CAT( \
        IS_VALID_, VRM_PP_ARGCOUNT(__VA_ARGS__) \
        )(__VA_ARGS__)
```

```
IS_VALID(*_0)(int*)
```

...expands to...

```
is_valid(
    [](auto _0) constexpr → decltype(*_0){})
)(type_c<int*>)
```

...which is equivalent to...

```
std::is_callable<
    decltype(
        [](auto _0) constexpr → decltype(*_0){})
    )(typename decltype(type_c<int*>)::type)
>{}
```

This technique is very useful when combined with if constexpr(...) - it's a barebones in-place concept definition&check.

```
template <typename T0, typename T1>
auto some_generic_function(T0 a, T1 b)
\{
    if constexpr(IS_VALID(foo(_0, _1))(T0, T1))
        return foo(a, b);
    else if constexpr(IS_VALID(_0 + _1)(T0, T1))
        return a + b;
```

```
template <typename TC, typename T>
auto unify_legacy_apis(TC& c, T x)
    if constexpr(IS_VALID(_0.erase(_1))(TC, T))
        return c.erase(x);
    else if constexpr(IS_VALID(_0.remove(_1))(TC, T))
        return c.remove(x);
```

```
template <typename T>
auto poor_man_ufcs(T& x)
    if constexpr(IS_VALID(_0.foo())(T))
        return x.foo();
    else if constexpr(IS_VALID(foo(_0))(T))
        return foo(x);
```

## Small caveat: it does not yet compile.

- clang++ hasn't implemented support for constexpr
   lambdas yet.
- g++ has, but there's a bug I found and reported (as #78131).

IS\_VALID does work properly with g++ trunk in other contexts where a constant expression is required though (e.g. non-template context if constexpr(...) and static\_assert).

## Thanks for your time!

- https://vittorioromeo.info
- vittorio.romeo@outlook.com
- https://github.com/SuperV1234
- @supahvee1234

https://github.com/SuperV1234/cpplondon