http://courses.csa:il.mit.edu/6.006 Course Werriew

- · Efficient: procedures for solving scale problems
- · Scalability
- · classic data structures
- · Real implementation in Pothan

1 Content

· 8 modules

Algorithmic thinking, peak finding,

Sorting & trees: Event Simulation

Hoshing: Genome Comparisson

Numerics: RSA encryption

Graphs: Rubik's cube

shortest poths: Caltech -> MIT

Jynamic programming: Image compression.

Advanced topics:

1) Peak finder

o One dimentional version

OL.	b	С	d	le	+	9	h	1
1	2	3	4	5	6	7	ક	g

an i are numbers

Position 2 is a peak of and only if b za and b z c
Position 9 is a peak if i z h

Problem: Find a peak if it exists.

Argument: In any will have a peak

o straight forward algorithm

Start from left, go all the way to the night

1, 2, ... n/2...n-1, n

Worst-case Complexity: $\mathfrak{A}(n)$ (on the order of n, a contitues n) downplotic Complexity of this problem is linear

o Pivide and Congher

1,2,... n/2...n-1, n

If $\alpha [n/2] < \alpha [n/2-1]$ then only look at left half $1, 2, \dots, n/2-1$ to look for a peak Else if $\alpha [n/2] < \alpha [n/2] + 1]$ then look at n/2+1, n to look for a peak

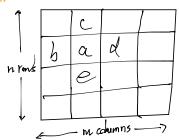
Else a [1/z] is a peak

$$T(n) = T(n/z) + \Theta(11)$$
"Work" algorithm does
on input of size n

Base rase
$$T(1) = \varpi(1)$$

* Can not do better for 17 version

I 27 Version



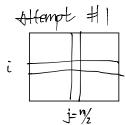
a is a 2) peak if and only if and, and, and, and,

· Greedy ascent algorithm

14	13	12	
75	9		17
	17	19	20

What Complexity (nm), (nz) if m=n

o Tivide and Counque



pick middle column j = m/z, find a 19 pick of (i, j)

use (i) as a start to find a 17-peak on you i

In correct 29 - peak may not exist on you!

Attempt #2

- · pick middle column j = m/z
- · Find global maximinm on column j at (i,j)

 Compare (i,j-1), (i,j+1)

 Pick left cols if (i,j-1) > (i,j), smilerly for right

if (i,j) > (i,j-1), (i',j+1), (i,j) is a 27 peak. Solve the new problem with half the number of columns when have a single column, find the global max \Rightarrow done

$$T(n,m) = T(n, m/z) + H(n)$$

 $T(n,1) = H(n)$

$$T(n,m) = \mathcal{D}(n) + \mathcal{D}(n) + \cdots + \mathcal{D}(n) = \mathcal{D}(n \log_2 m)$$