AVL Fees, AVL Sort

I BST Recall

h: height of the bree; longest path from root to leaf.

· Balanced or NOT

Balanced: if $h = \mathcal{H}(\lg n)$

. Keep tree balanced

· Height of made in a tree

2 onyest gath from it down to a leaf.

2 20 63 height of a node

129 50 -1

height of node = height of left child height of right child

JAVL Tree

Adelso - Velski - Landis

. AVL Tree require heights of left & right children of every node to differ by most ± 1 h. The hr

. ALV trees are always bollanced (D(Ign) tree height)
norst case is when right subtree has height one more than
the left for every node.

Define: $N_h = \min \# nodes \ in \ cm \ AVL \ tree \ f \ height \ h$ $N_h = 1 + N_{h-1} + N_{h-2}$ h-1

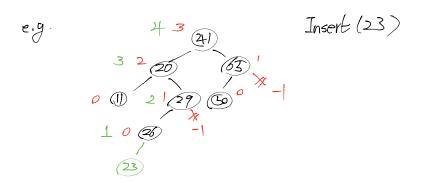
 $N_h > F_n (F_{abinaci}) = \frac{y^h}{\sqrt{5}} (\Psi = 1.618)$ $\downarrow \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \downarrow \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \downarrow \qquad \qquad \downarrow \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad$

 $N_{h} = 1 + N_{h-1} + N_{h-2} > 1 + 2N_{h-2}$ $> 2N_{n-2} = H(2^{h/2})$ $\downarrow > h < 2 |gN$

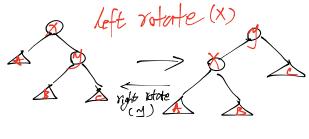
oAVL Insert

O Simple BST insert

3 fix AVL property



· Rotation:

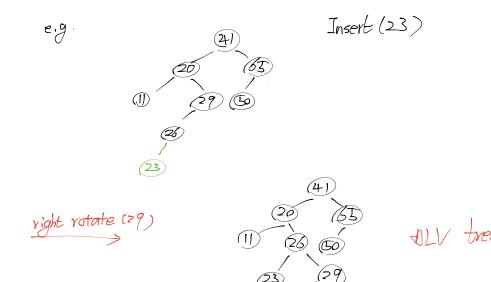


AXBYC -> AXBYC

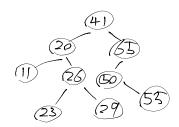
Const. time

(root moves to the left)

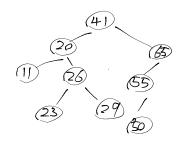
O AVL Insert (Continue)



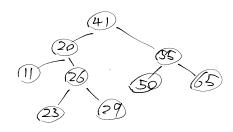
e.g: Insert (55)



left rotate (50)



right votate (65)



@fix All Property

from changed nocle up, suppose X is lowest node volating AUL, assume X. right is heavier — if x's right child is right heavy or balanced do Right rotate of X: RR(X)

