

# Device Security: Trusted boot and code signing

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# Overview

- **Device security**

- Is code on the device vulnerable to exploits ? (e.g. buffer overflows)
- Is the code authenticated ? (i.e. has not been tampered with)

- **Local data security**

- Is the stored data is accessible to everyone? (e.g. encrypted)
- Is the stored data authenticated?

- **Cloud data security**

- How is data stored in the cloud?
- Who has access to data stored in the cloud?

- **Metadata security**

- What does metadata reveal about stored data?
- Can we tamper the metadata?

# Overview

- Device security
  - Physical device security (e.g. side-channel attacks vulnerabilities)
  - **Firmware/OS security (e.g. bootloader, kernel)**
  - **Application security (e.g. code signing, sandboxing)**

# Introduction

# Security challenges

Protect devices against:

- Malicious applications
- Rootkits

# Malicious applications

Malware distributed using OS specific applications, designed to exploit the operating system vulnerability's.

## Issues:

- High success rate because they (can) masquerade as useful applications
- Are often used as a means to install more dangerous malware: backdoors, rootkits

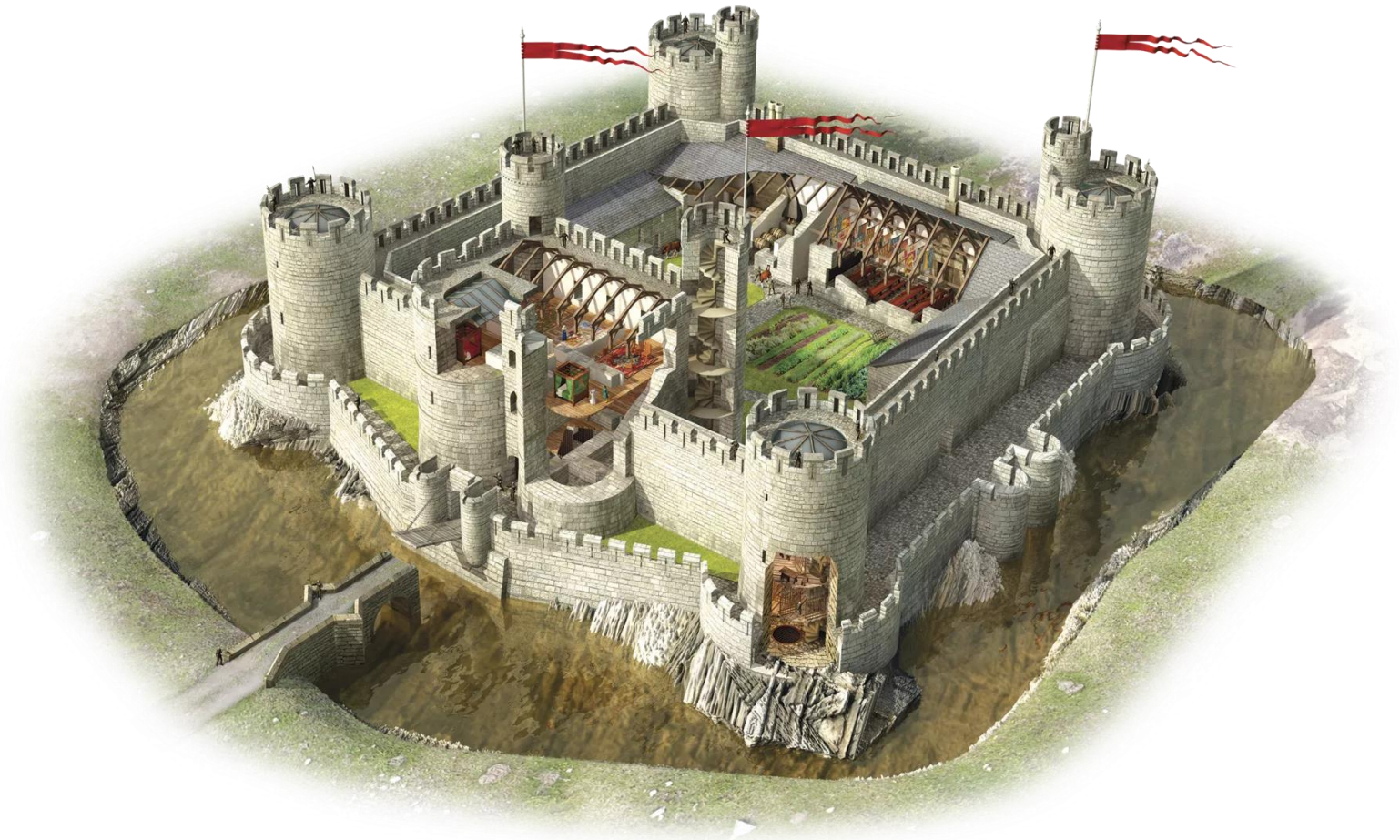
# Rootkits

Code designed to enable access to a computer or areas of its software that would not otherwise be allowed.

## Issues:

- Install themselves with highest privileges
- Prevent detection/removal with anti-malware tools

# Mitigations: defence in depth





# Mitigations

The principle of “defence in depth”:

- Secure the boot process
- Secure the user space
- Secure the distribution channels

# Hash functions and signatures

# Authentication

**Authentication** is the process or action of proving or showing something to be true, genuine, or valid.

Things that can be authenticated:

- In real life: bags, watches, ....
- In CS: **identities**, machines/computers, users, files...

# Authentication

**Authentication** is the process or action of proving or showing something to be true, genuine, or valid.

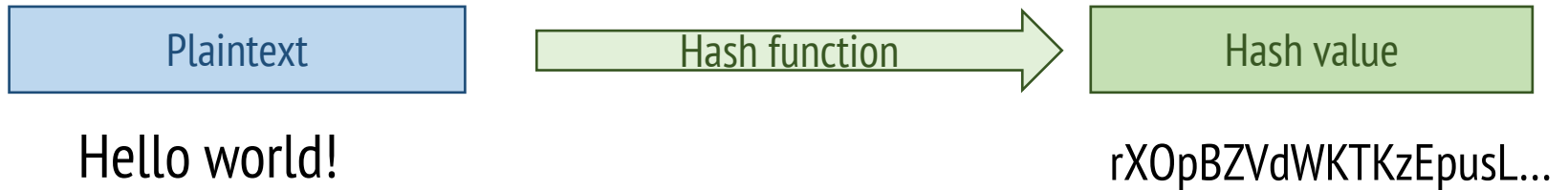
Things that can be **authenticated**:

- In real life: bags, watches, ....
- In CS: **identities**, machines/computers, files,..., also **users** (human beings)

How (in CS):

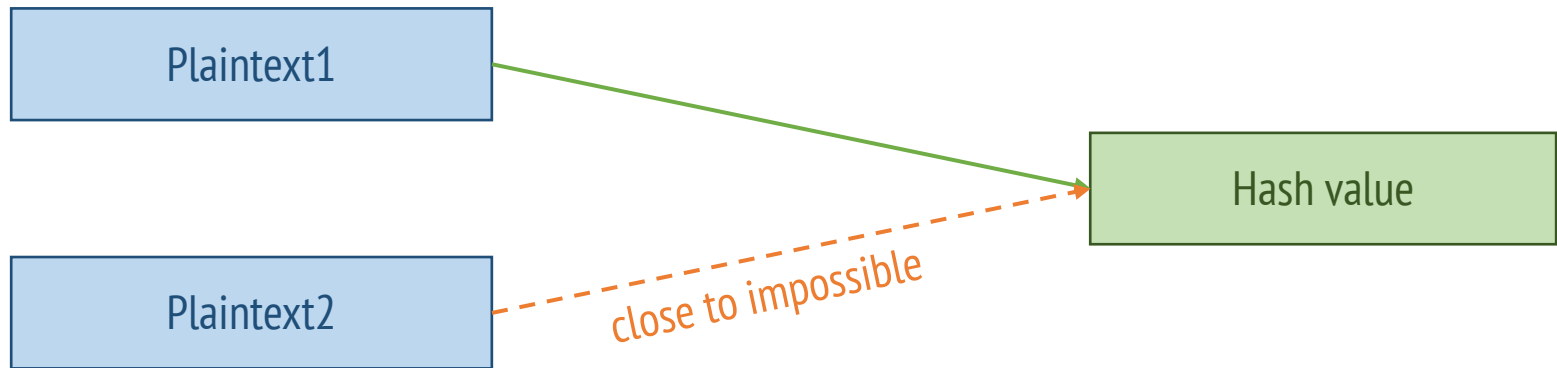
- Hashes – to compute/check **integrity**
- Digital signatures – to **authenticate**

# Hash functions



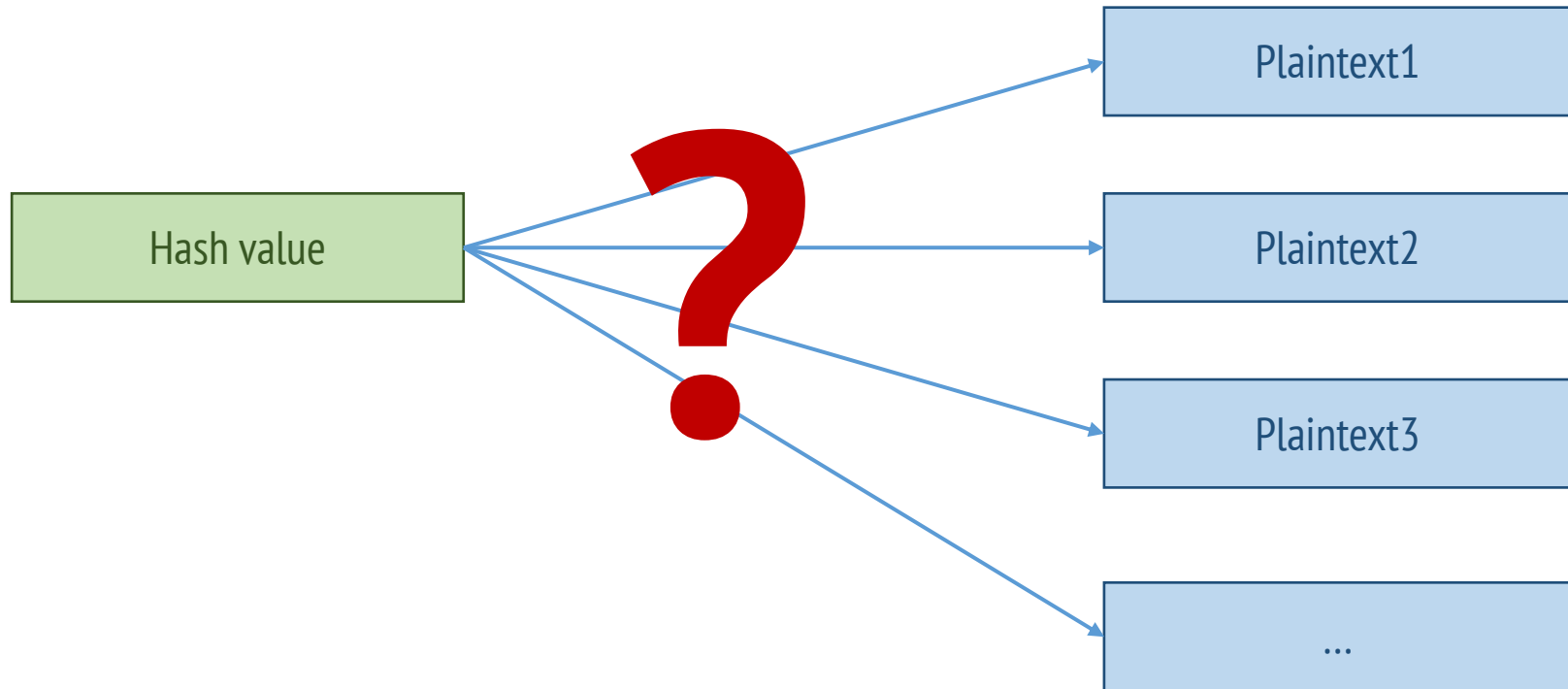
# Hash functions

Unique per plaintext (low collision)



# Hash functions

Difficult to reverse (one way)



# Hash functions

Size of  is independent from size of  .

Size of  is given by the **hash function**.



# Hash functions

Size of  is independent from size of  .

Size of  is given by the **hash function**.

Example:

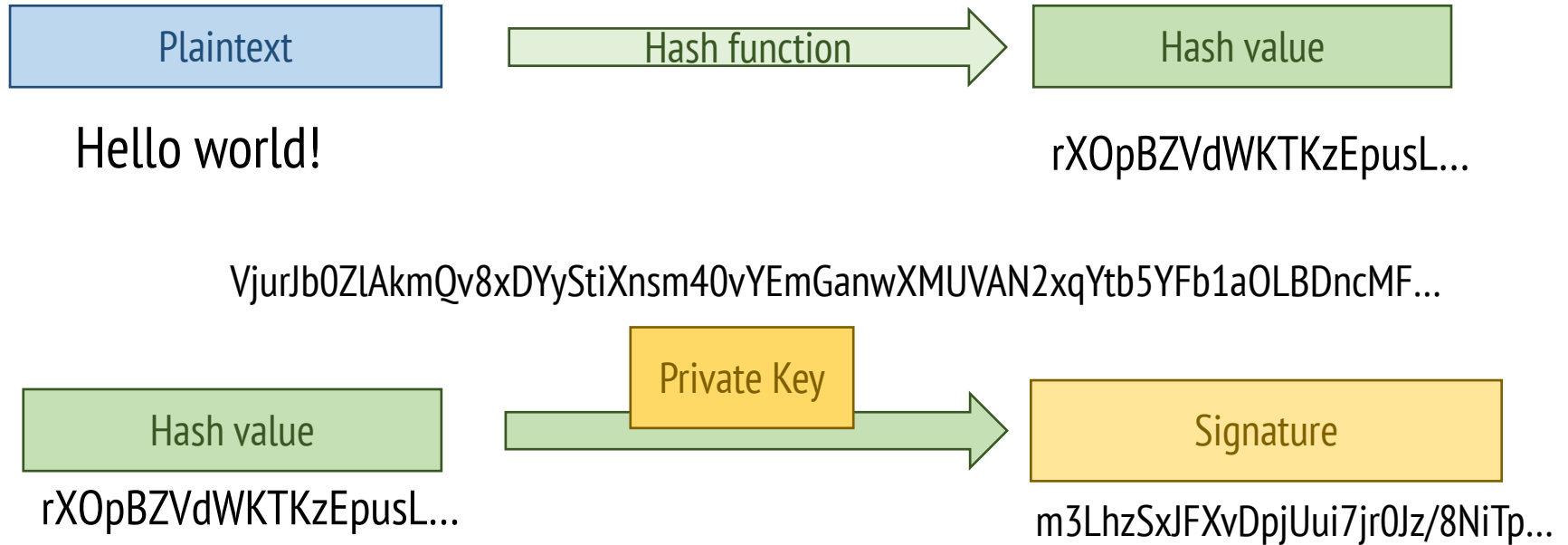
SHA256 = 256 bits

SHA512 = 512 bits

# Signatures



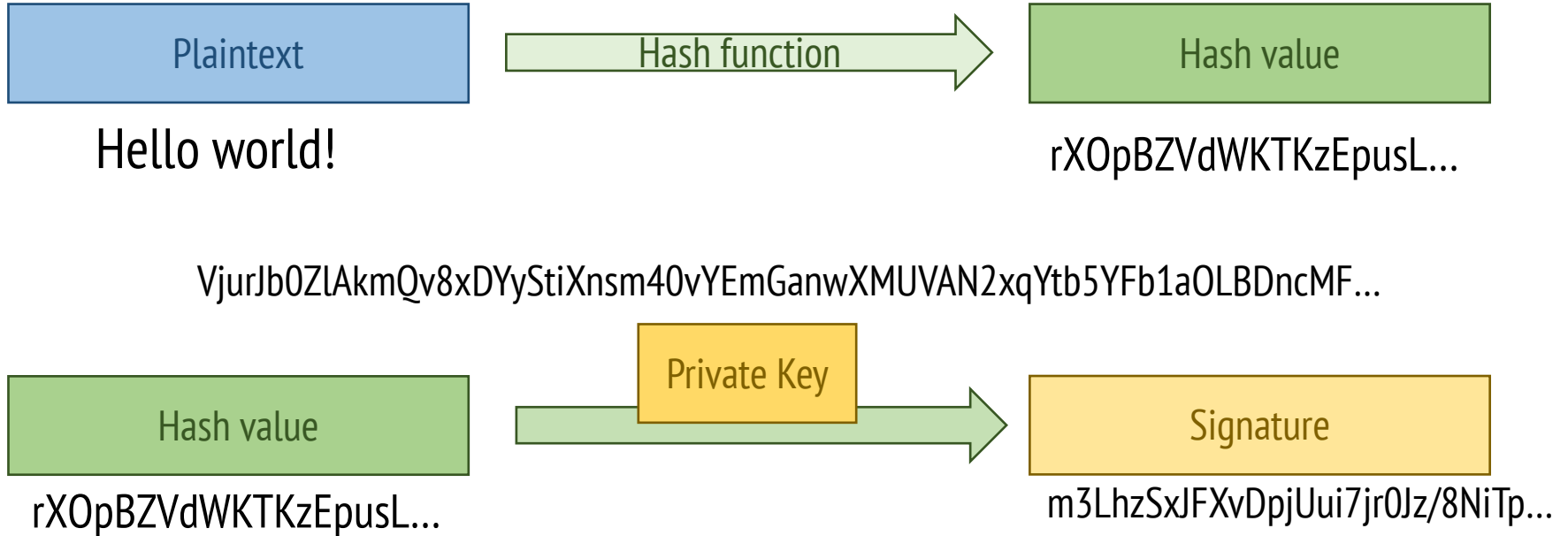
## Signing:



# Signatures



## Signing:



Message:



# Signatures

Public Key

Private Key

## Verification:

Message:

Plaintext

Signature

Plaintext

Hash function

Hash value

rXOpBZVdWKTkzEpusL...

WMWXV1cFZL7B4juLzULK7y2WFFv/9yyRVmDBuy6WbSWYVs...

Signature

Public Key

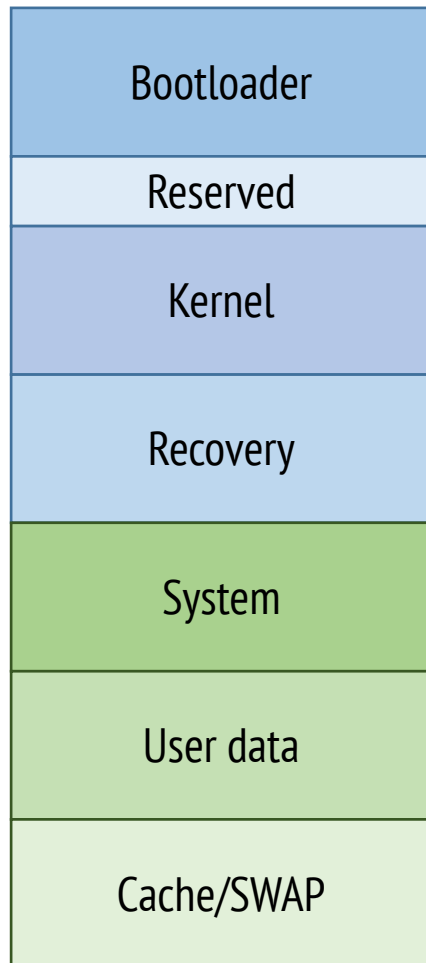
Hash value

m3LhzSxJFXvDpjUui7jr0Jz/8NiTp...

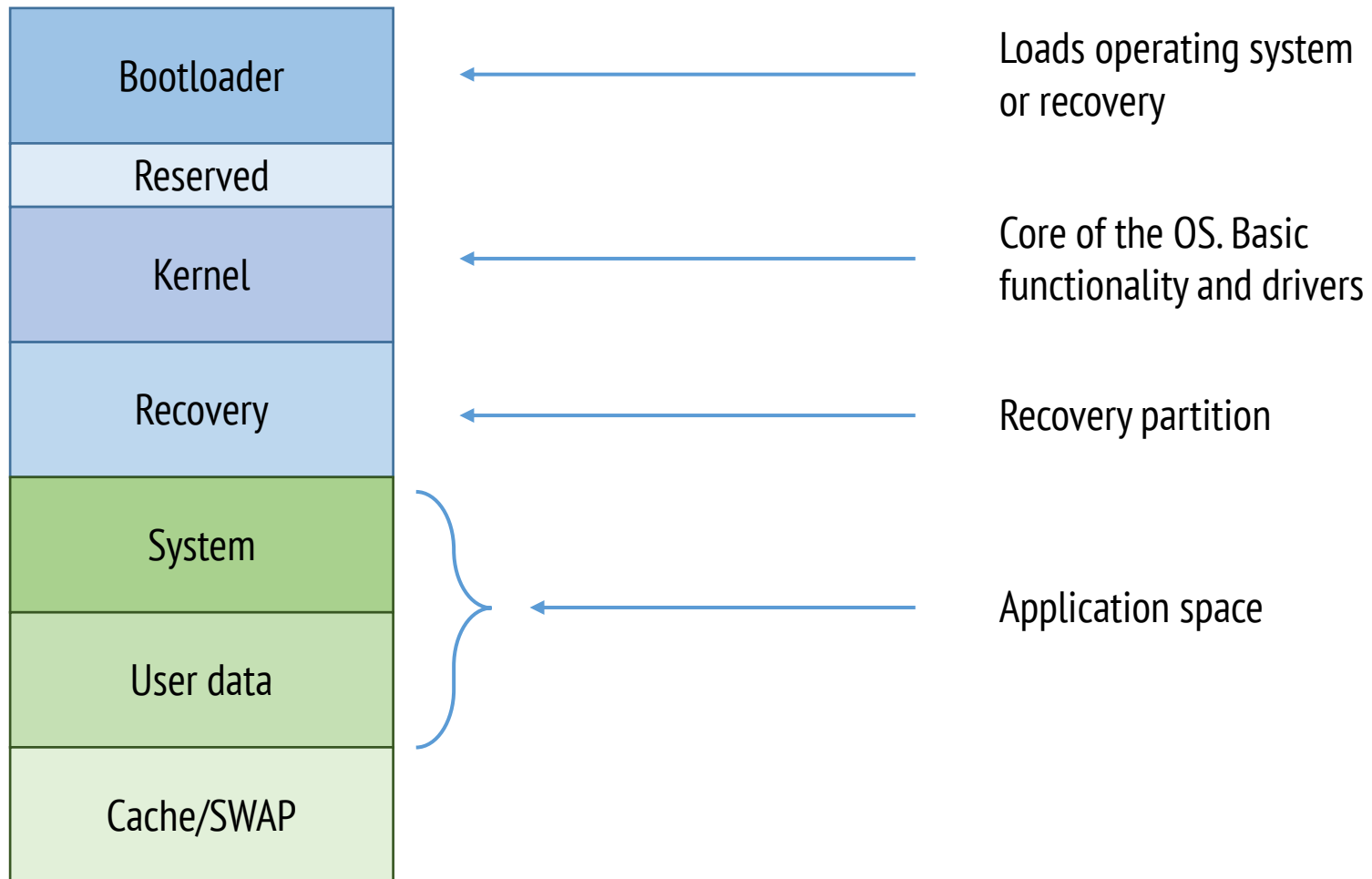
rXOpBZVdWKTkzEpusL...

System architectures

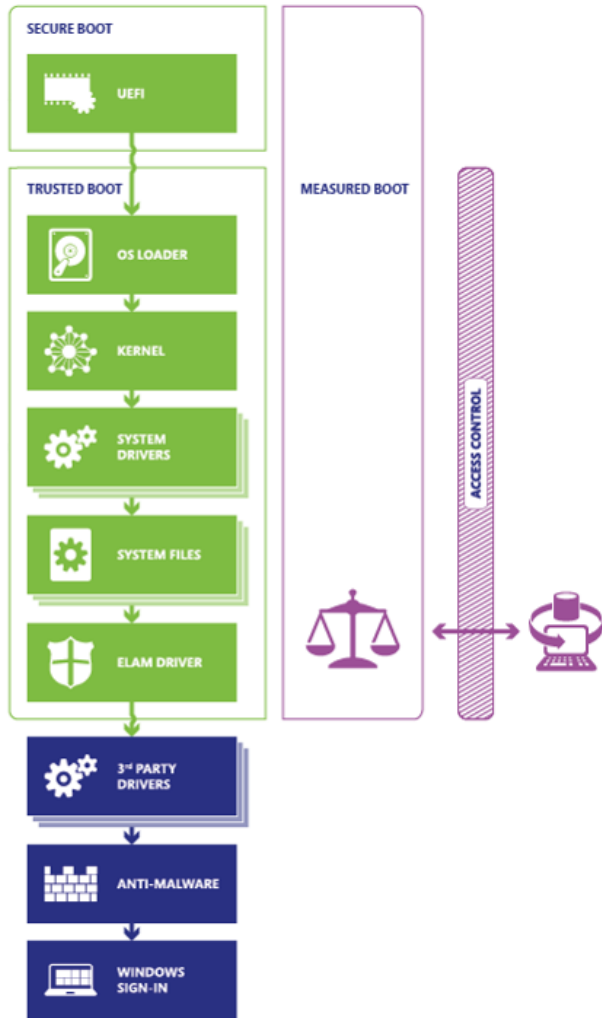
# Generic architecture



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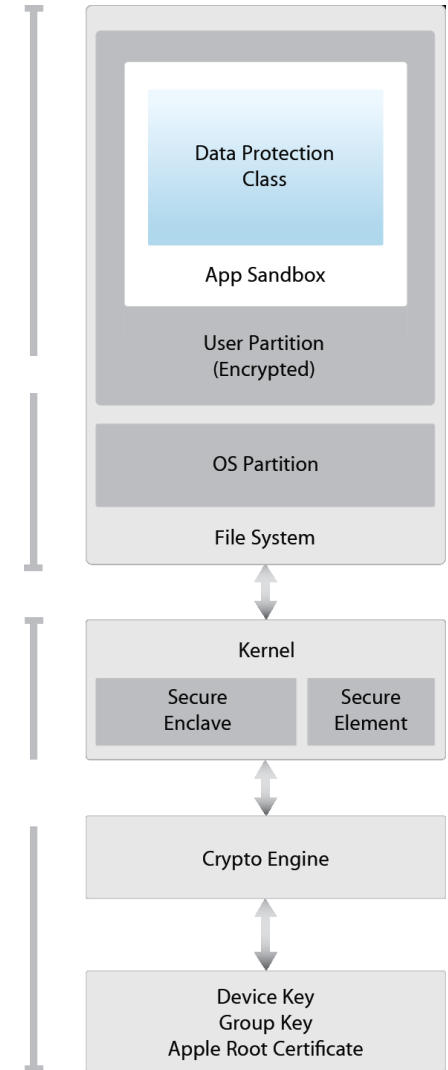
# System architecture



Windows architecture



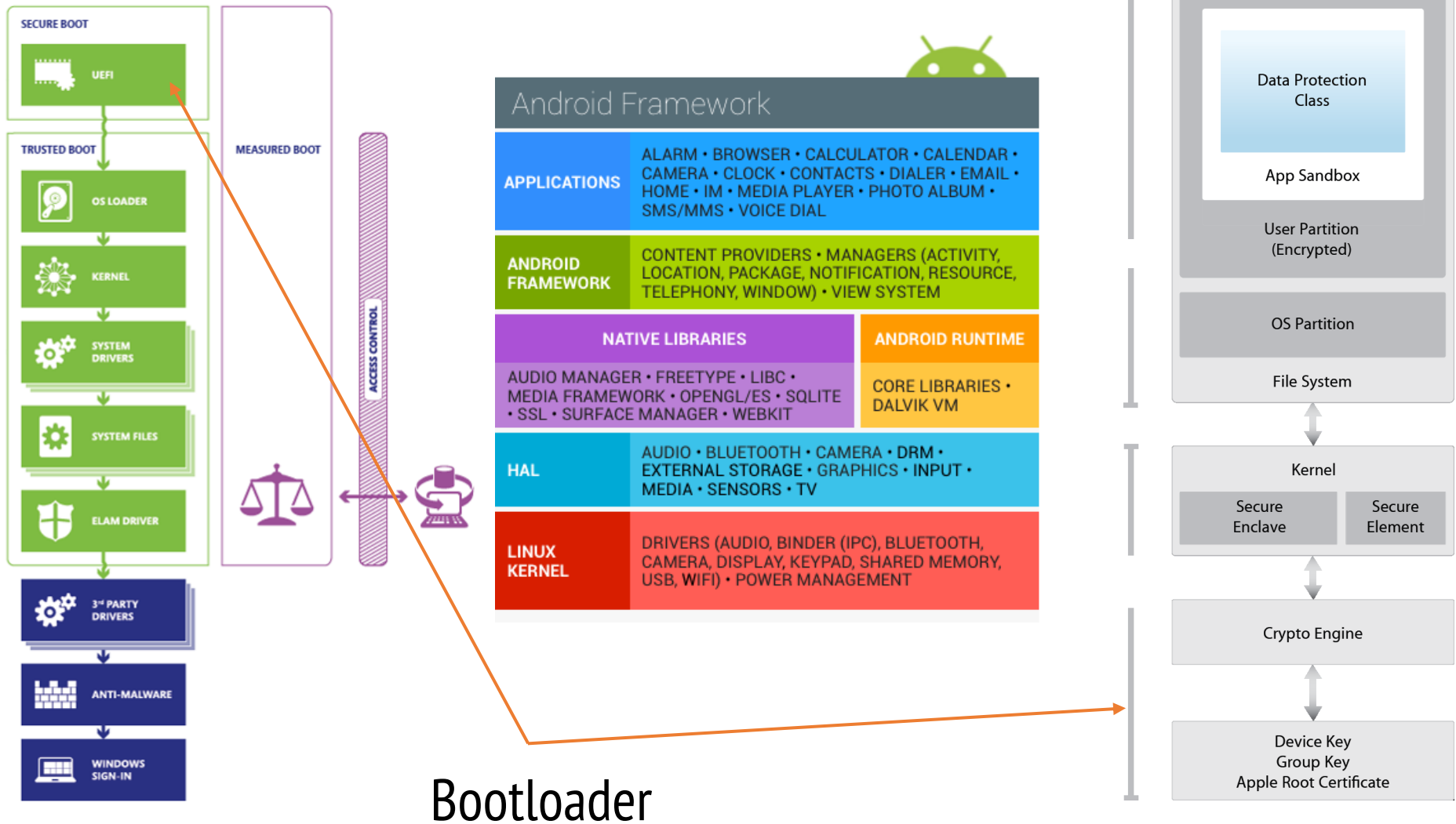
Android architecture



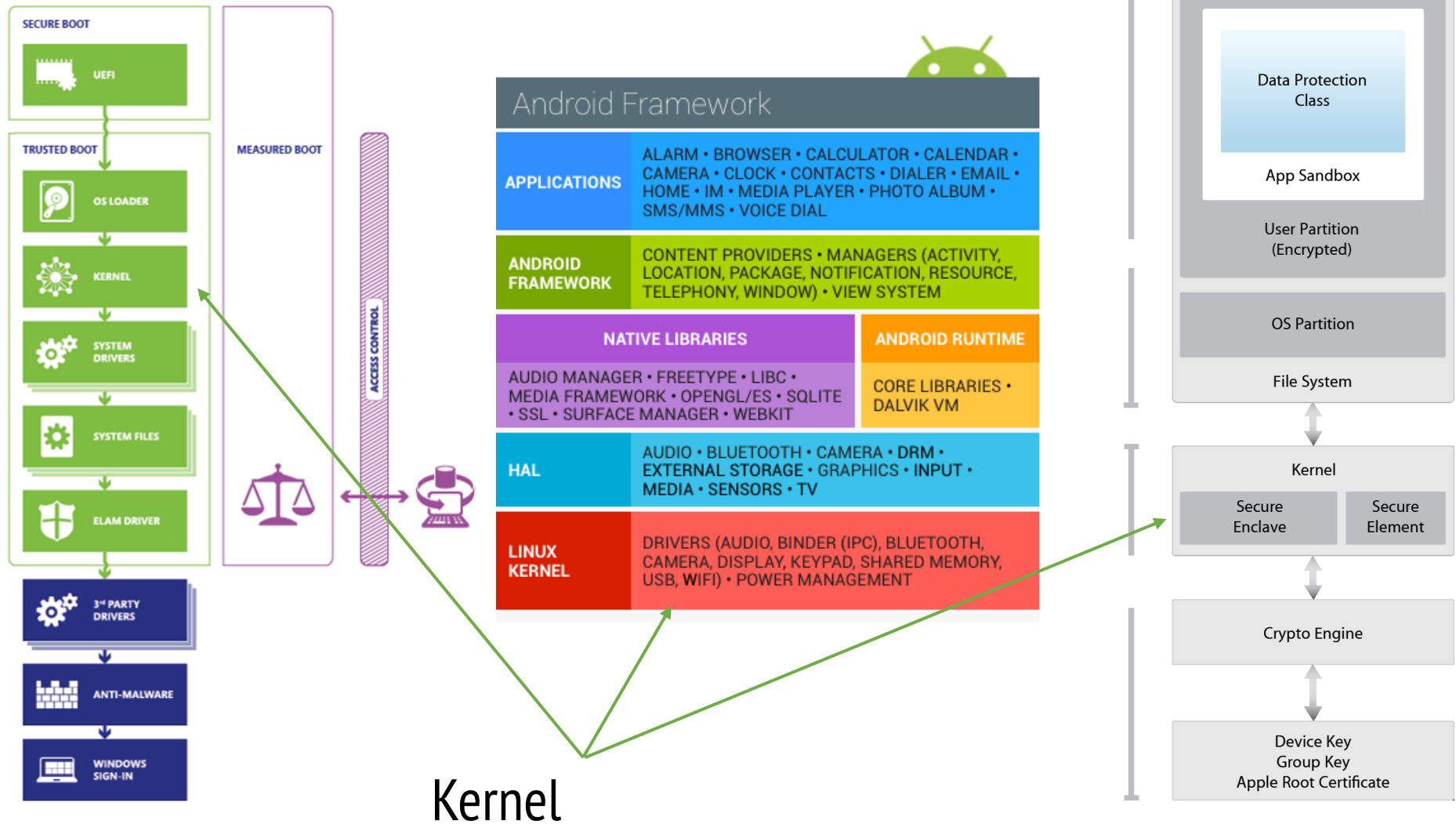
iOS security architecture



# System architecture

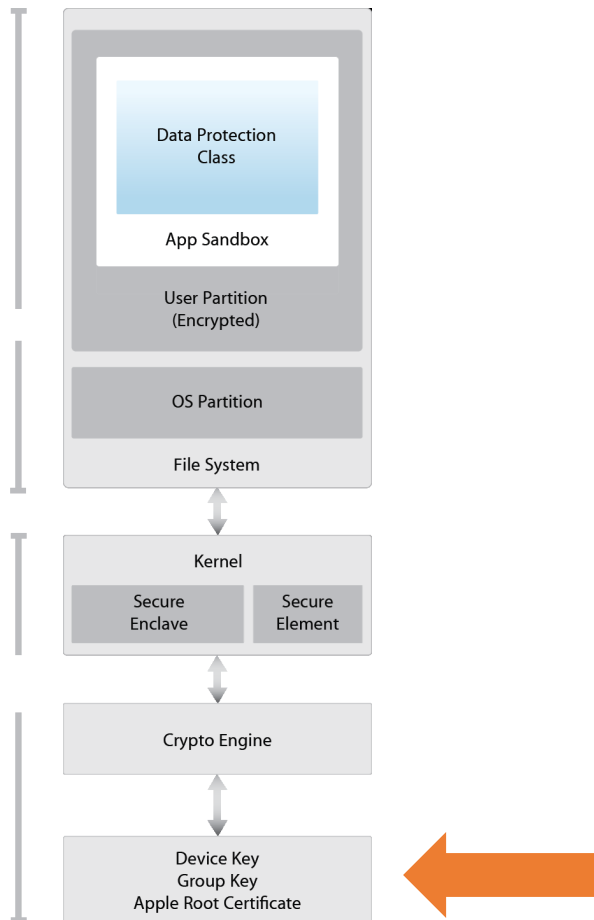


# System architecture



Trusted boot

# Root of trust



1. The bootloader is the guardian of the device state and is responsible for initializing the TEE and binding its root of trust.
2. If **rooting software compromises** the system before the kernel comes up, it will retain that access.
3. The bootloader verifies the integrity of the boot and/or recovery partition before moving execution to the kernel.
4. **Hardware** root of trust is fixed because is laid down during chip fabrication.
5. **Non-hardware** root of trust can be changed because is stored on non-volatile memory (e.g. NAND).

# How to verify?

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Directly hash its contents and compare them to a stored value.

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Some issues:

- Verifying an entire block device can take an extended period
- Will consume much of a device's power

# Dm-verity (android)

## Uses:

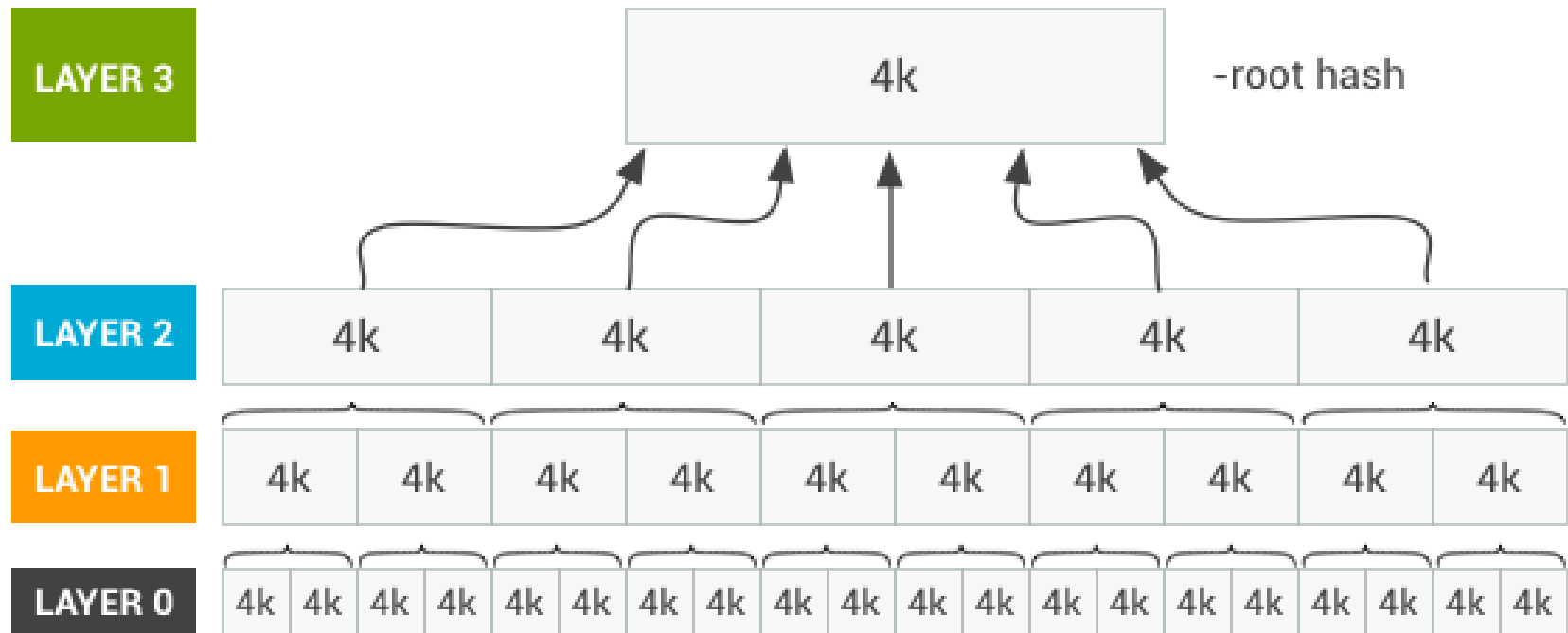
- Block storage
- Cryptographic hash function, e.g. SHA256
- Cryptographic hash tree (i.e. Merkel Tree)



# Benefits of Dm-verity

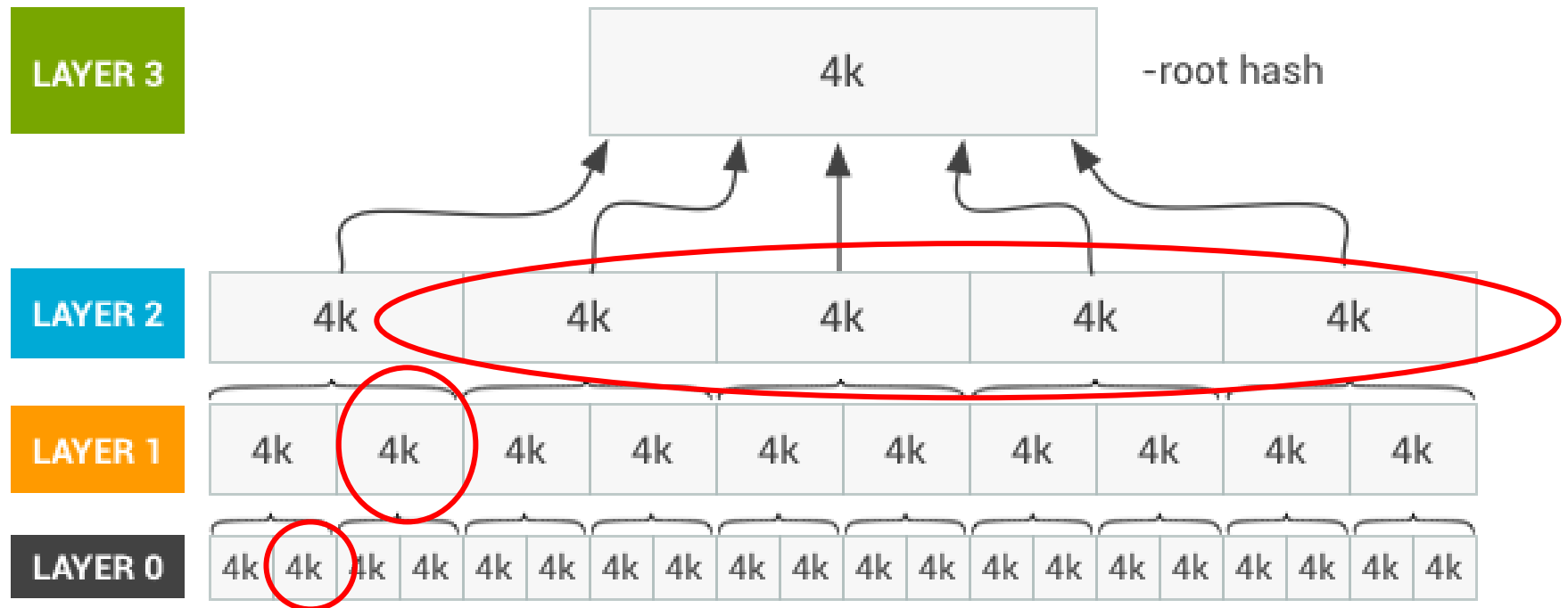
- Verifies blocks individually and only when each one is accessed
- The HASH operation is done when the block is read into memory: the block is hashed in parallel
- The hash is then verified up the tree

# Boot/recovery partition



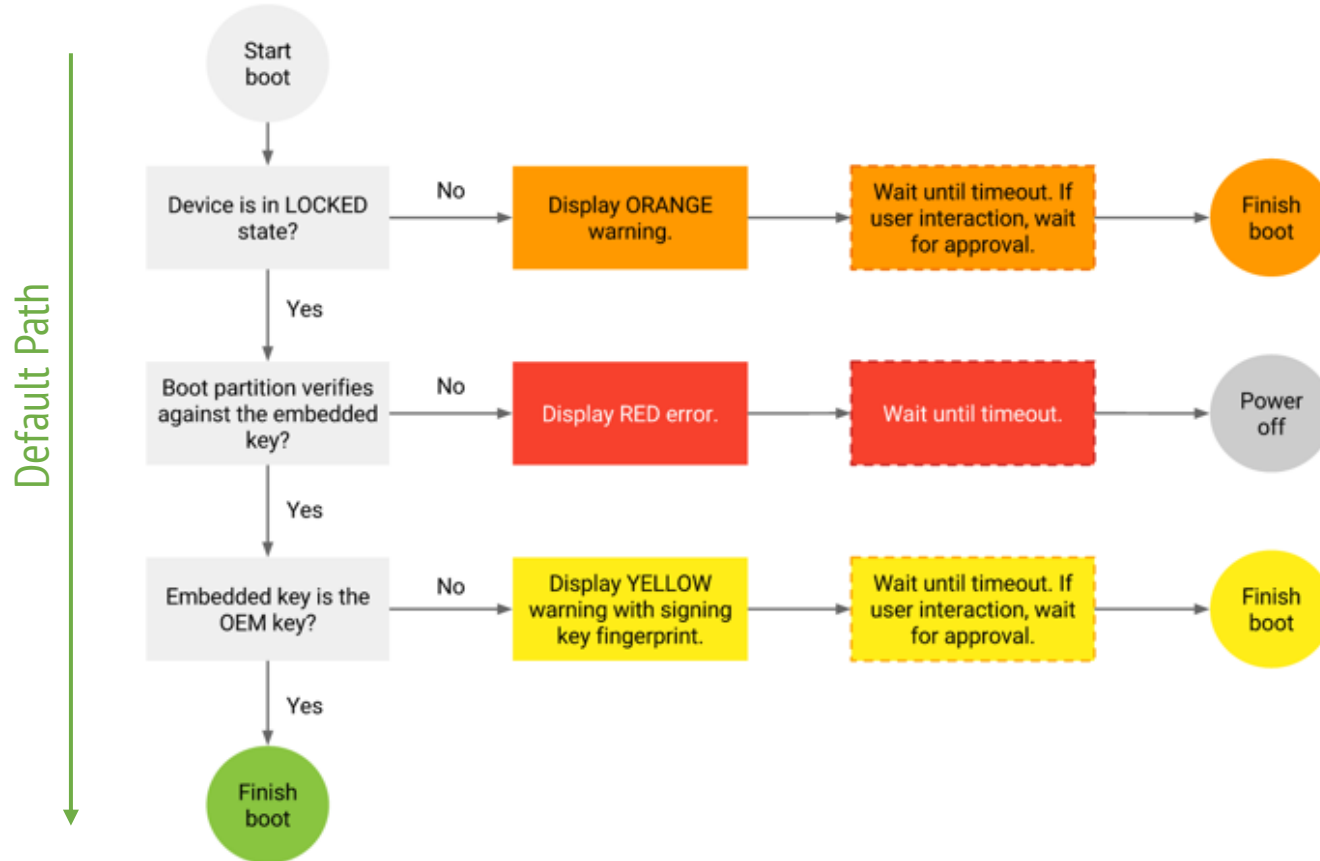
Android partition verification

# Boot/recovery partition

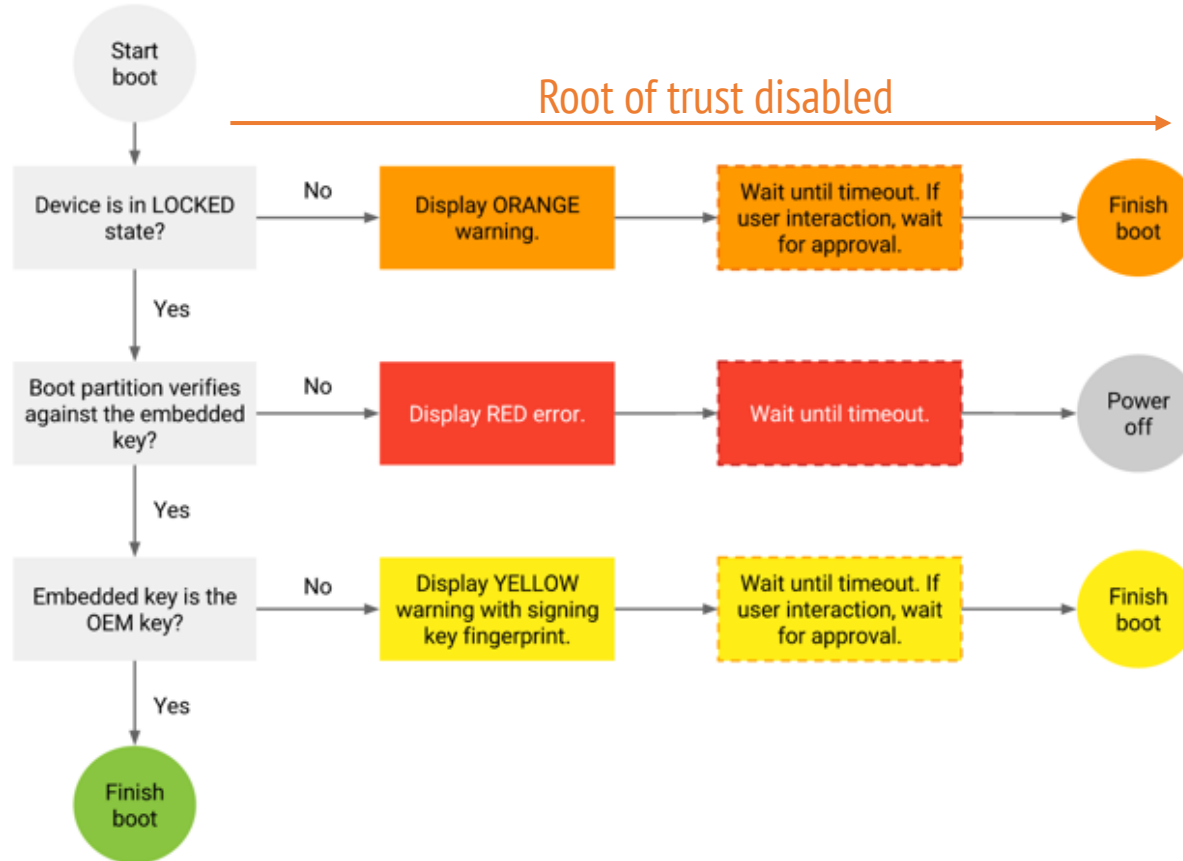


Android partition verification

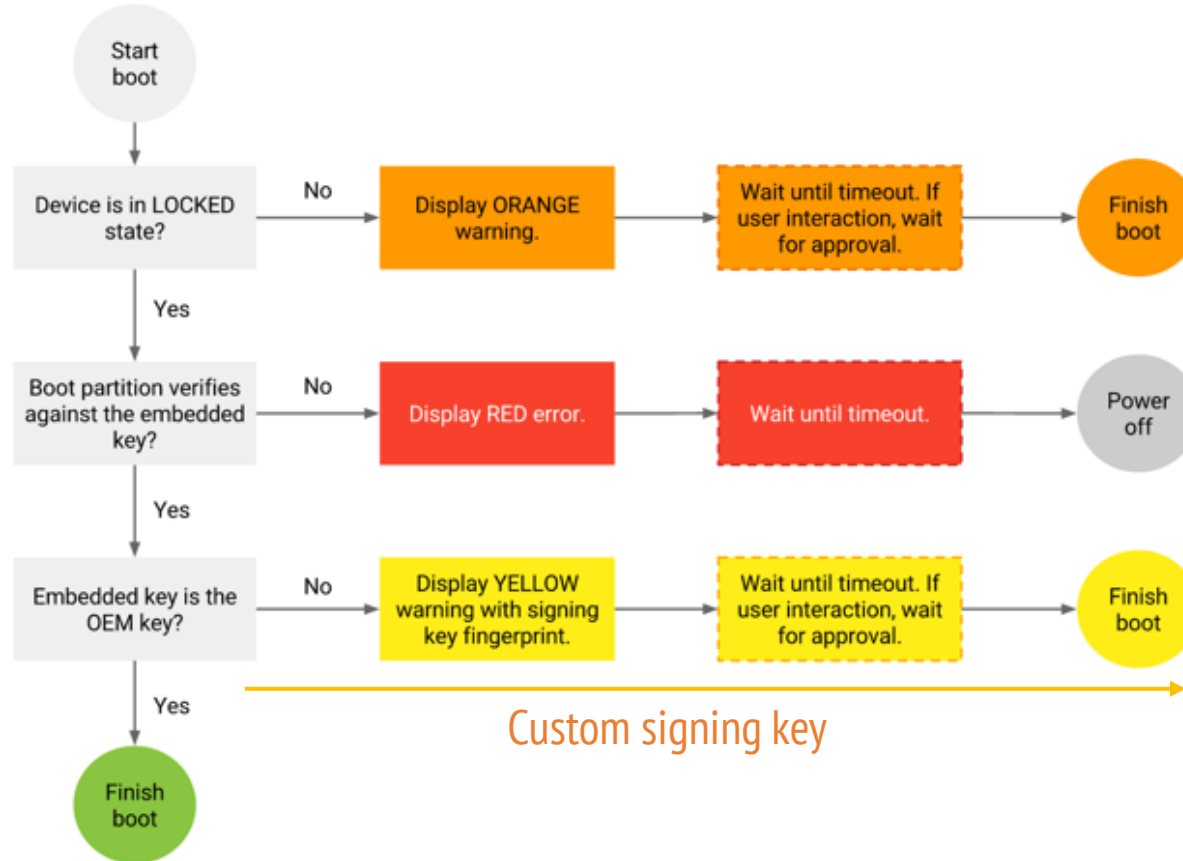
# Boot flow (Android)



# Boot flow (Android)



# Boot flow (Android)



# Main things to remember

- Authentication: every component of the system is authenticated
- Control access: prevent loading the OS without a secure boot loader
- Usable security: simple messages aimed at users
- Fail secure: stop the booting process if anything goes “wrong”
- ...

# What else?

- Kernel secure?
  - Yes
- Can we update OEM keys in case of compromise?
  - No



# What else?

- Kernel secure?
  - Yes
- Can we update OEM keys in case of compromise?
  - No
- What if we can still compromise the system by exploiting some design flaw(s) ?
  - Can we at least detect that and notify the user? (we could before)
  - **Can we prevent sensitive data compromise? (e.g. encryption/signing keys)**

Secure the boot process  
using hardware

# Hardware anchored security

## Trusted Platform Module (TPM)

- a secure cryptoprocessor
- generates cryptographic material (keys and random numbers)
- performs remote attestation
- protects cryptographic material by binding and sealing

## Trusted Execution Environment (TEE)

- secure area of the main processor  
(e.g. TrustZone in ARM, SGX in Intel)
- code and data loaded inside are protected with respect to confidentiality and integrity
- provides isolated execution

# What can I do with a TPM?

- Platform integrity
- Disk encryption
- Password protection
- Digital rights management
- Protection and enforcement of software licenses
- Prevention of cheating in online games

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# Platform integrity

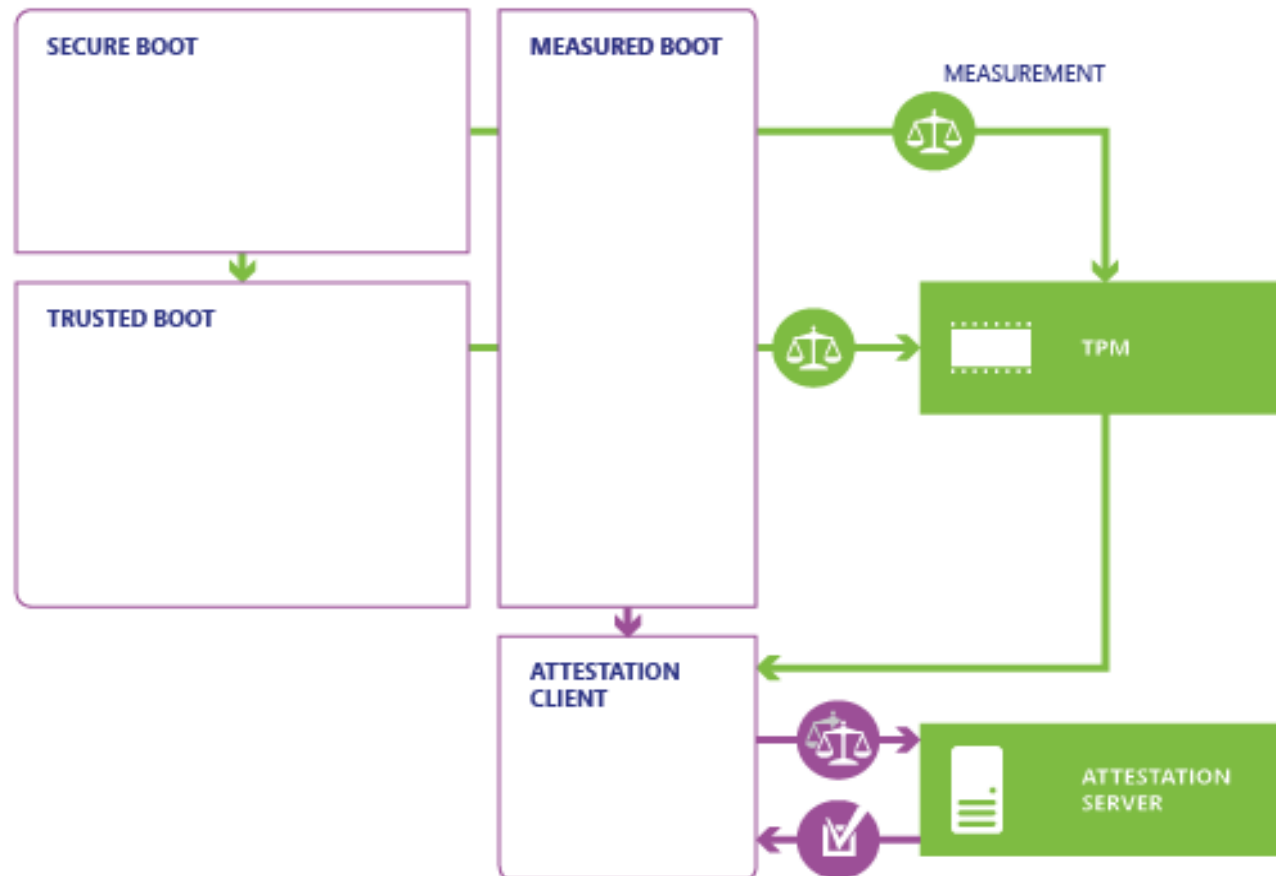
## Requirements

- TPM
- UEFI
- Linux Unified Key Setup (LUKS) or BitLocker Drive Encryption

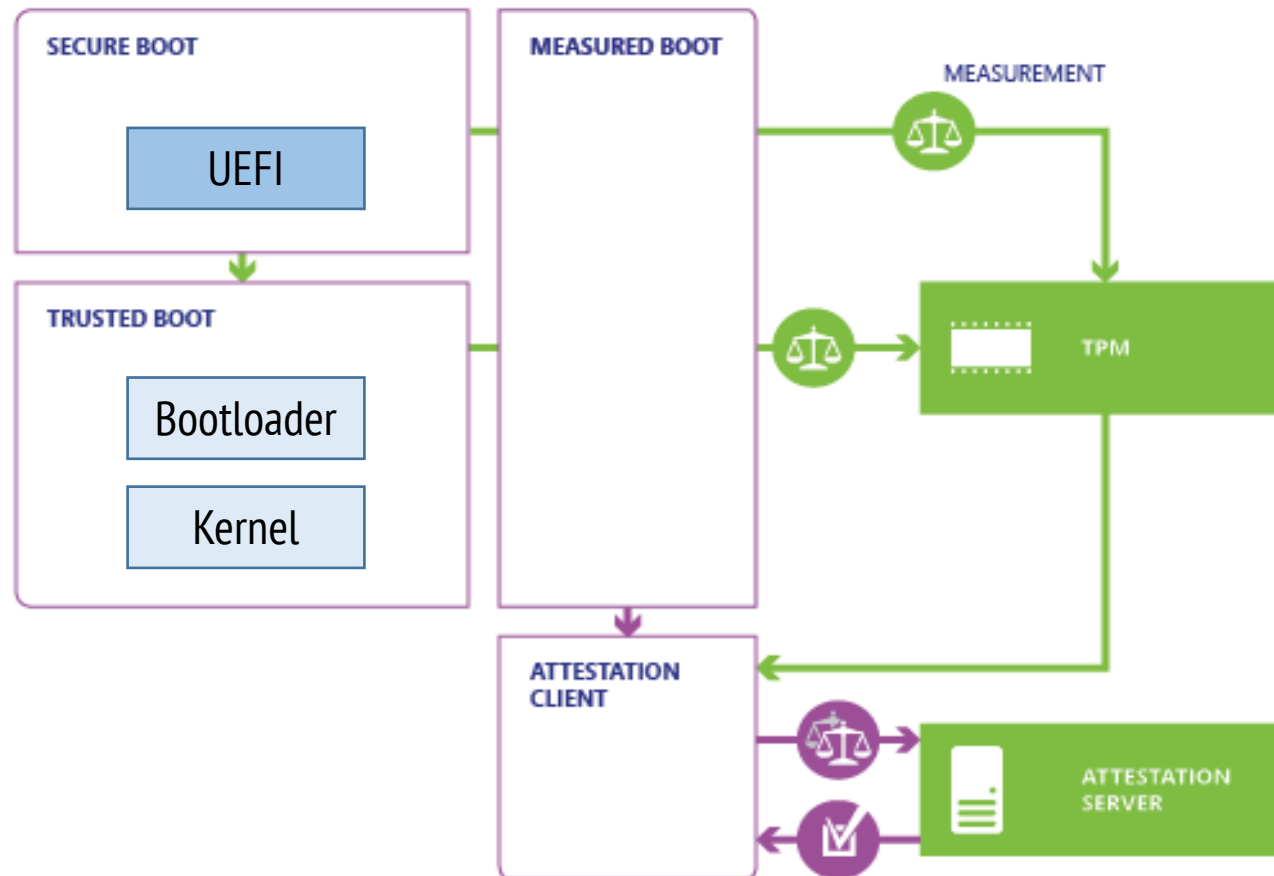
## Enforces

- "root of trust"

# Securing the boot process with a TPM?

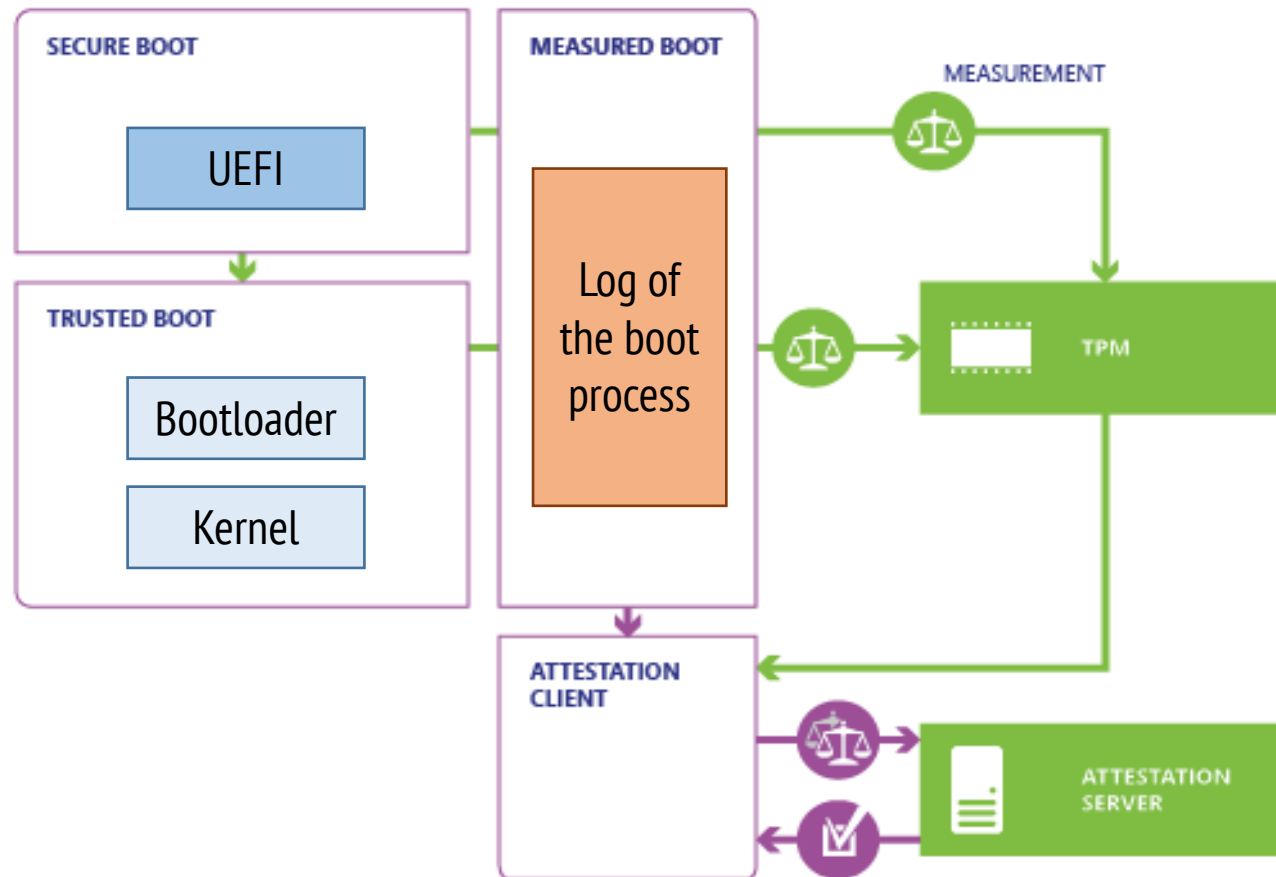


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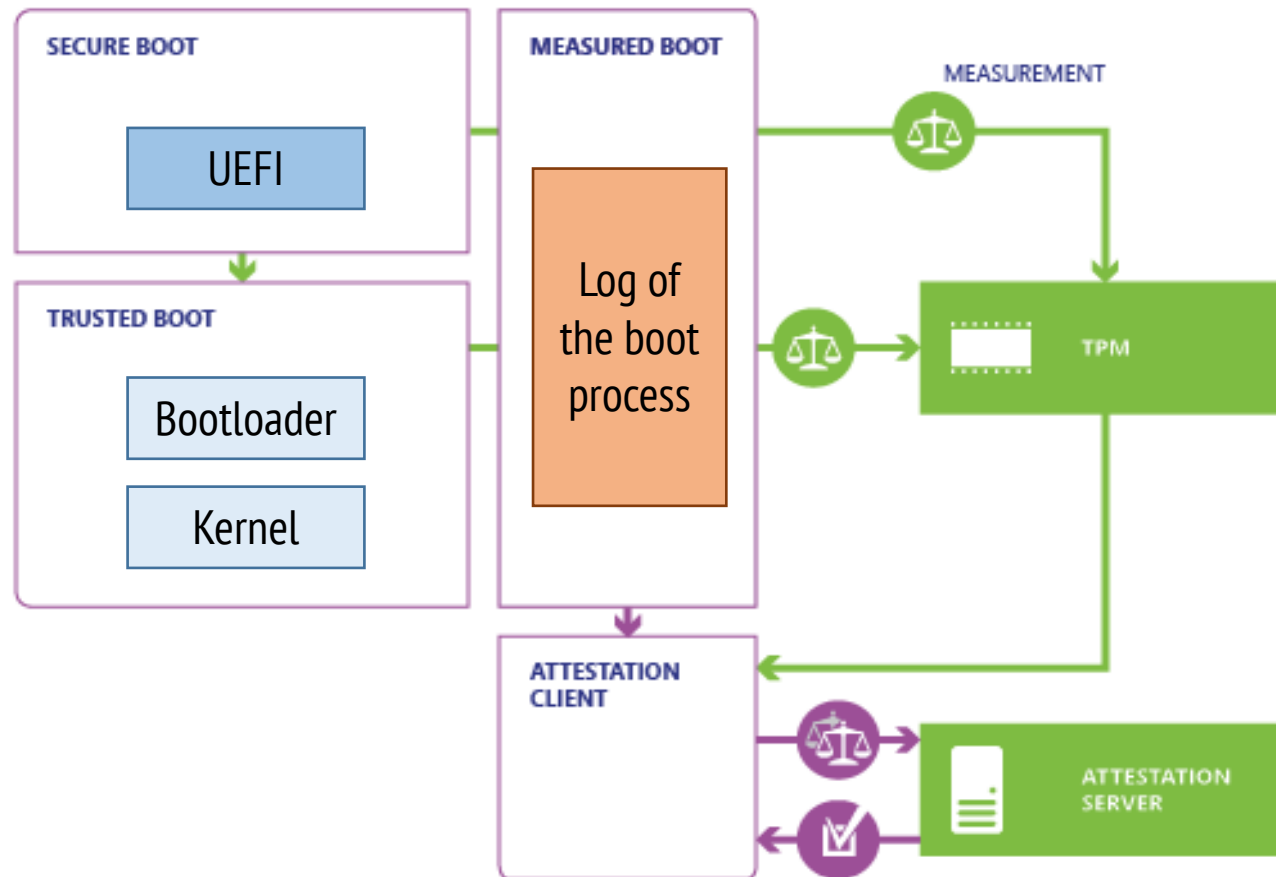




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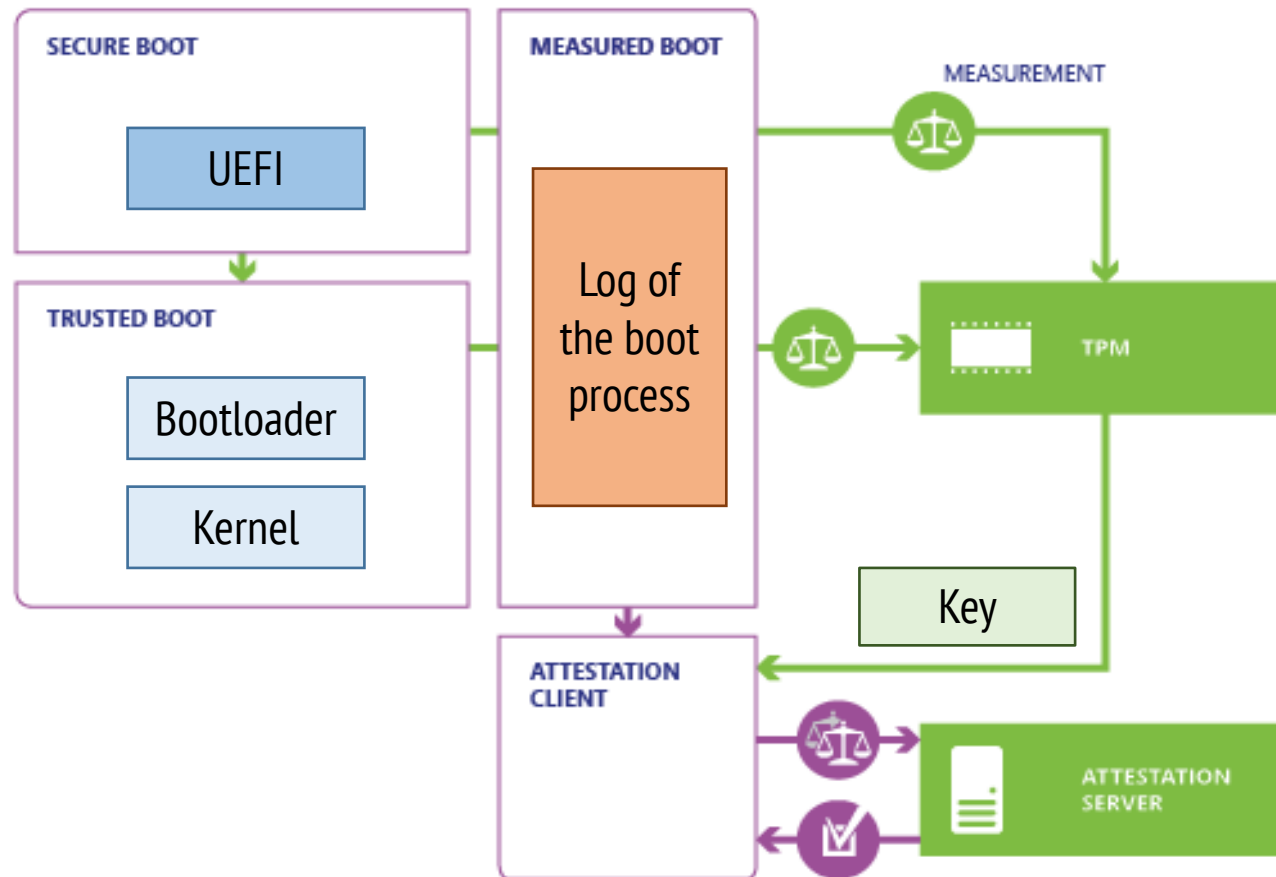
# Securing the boot process with a TPM?



Hashes for:

1. Firmware
2. Bootloader
3. Drivers
4. ...

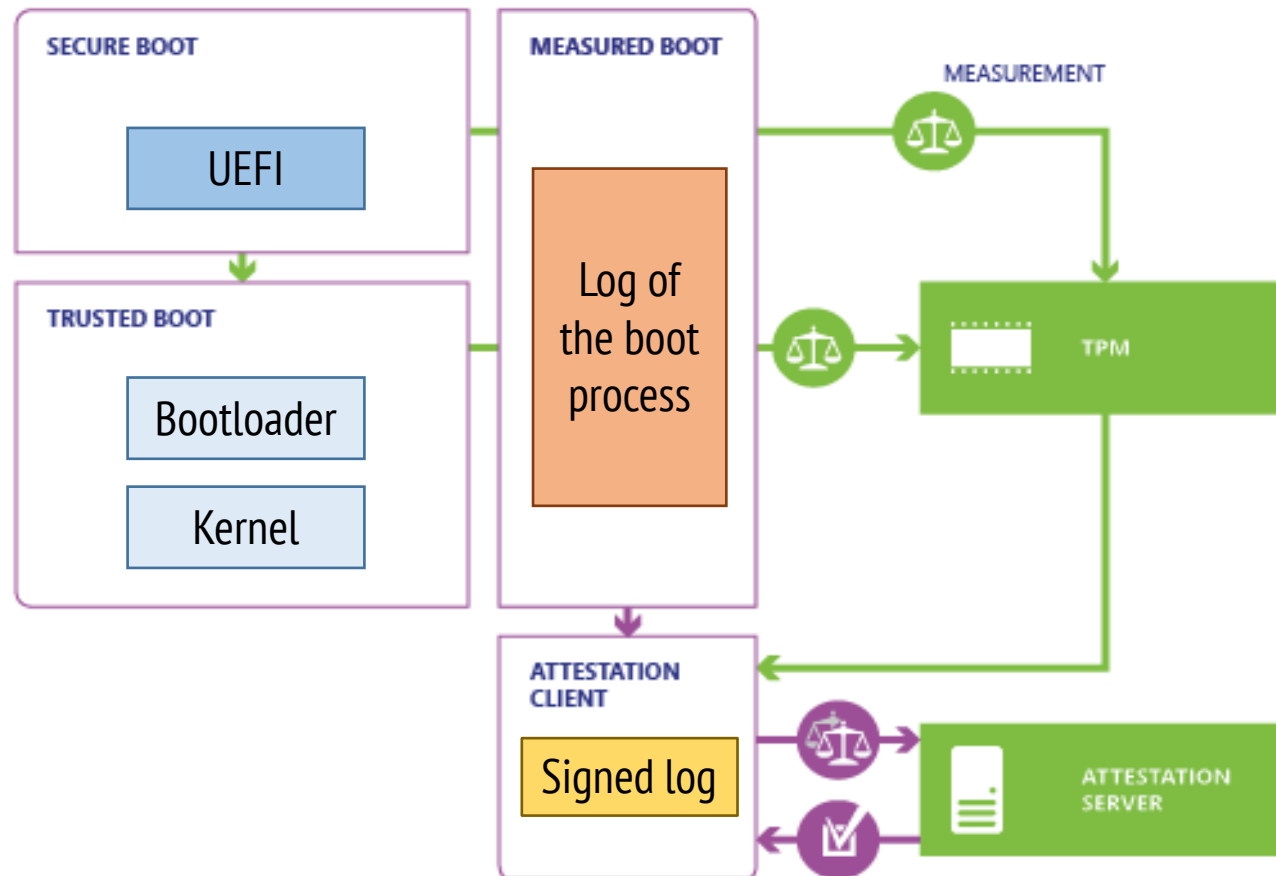
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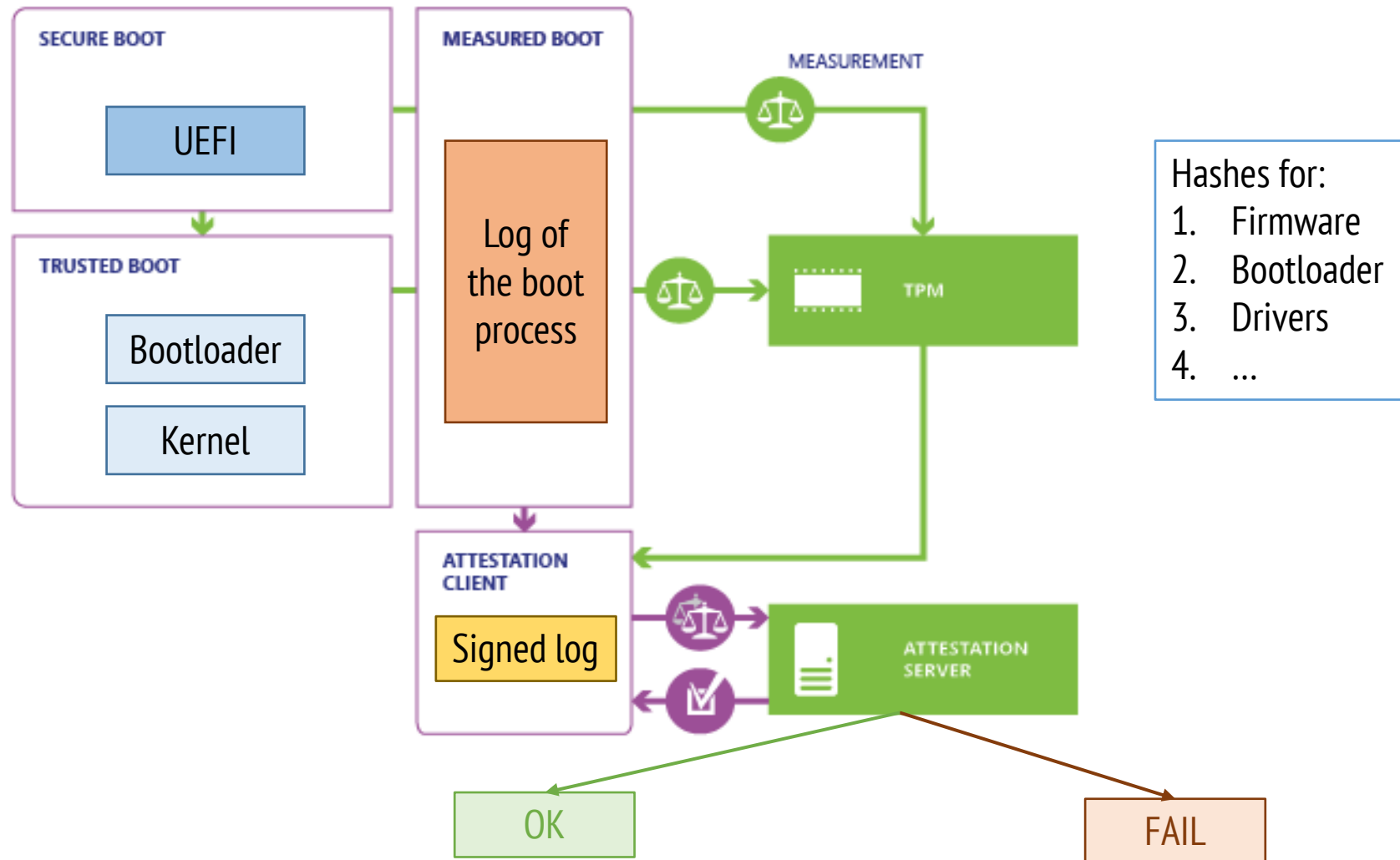
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Hashes for:

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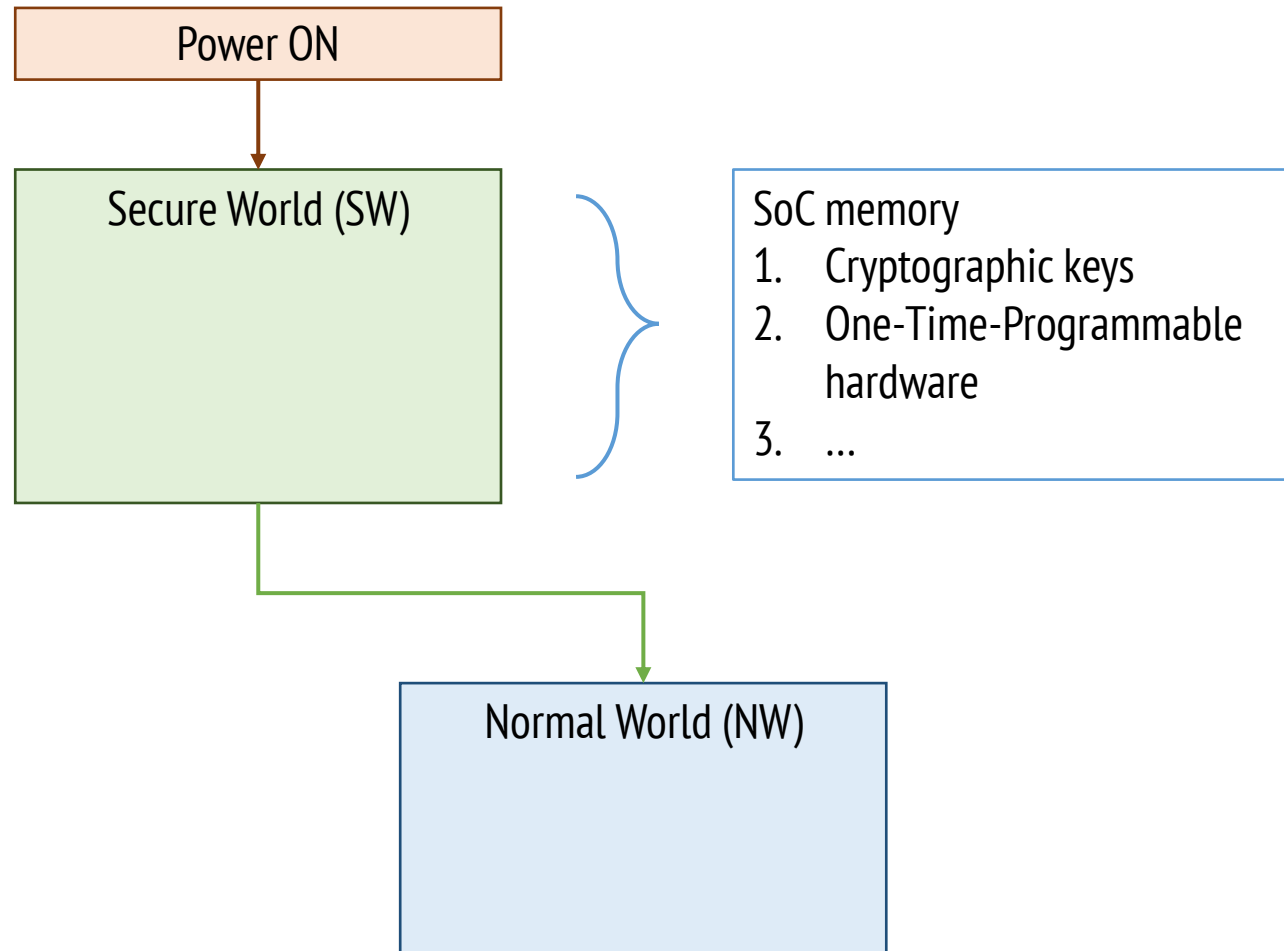
# Securing the boot process with a TPM?



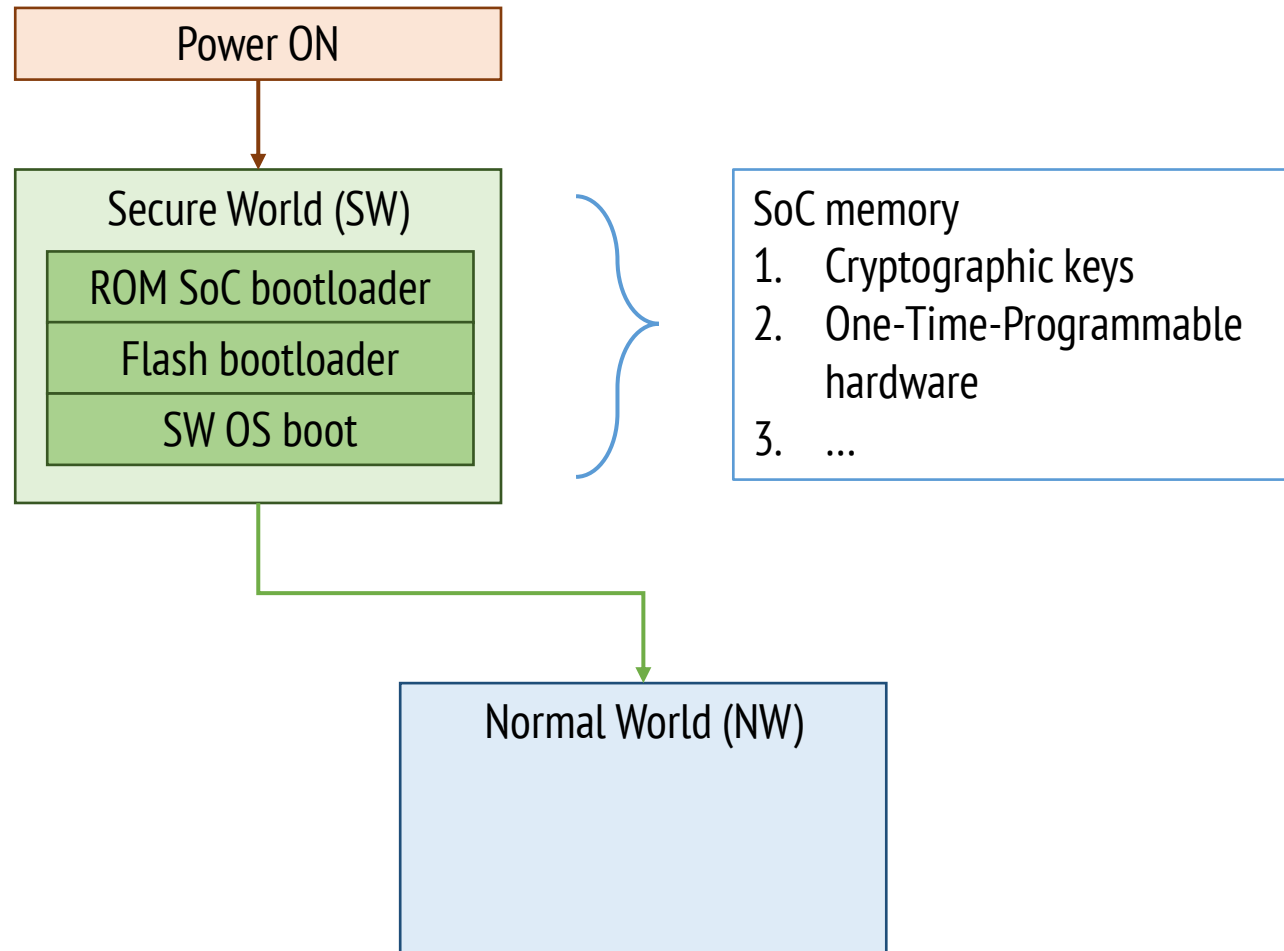
# Things to remember

- Authenticated requests: the TPM authenticates UEFI application and the bootloader
- Fail secure: if things go wrong stop
- Audit and monitor: keep logs of the boot process and verify them against a known and trusted logs
- ...

# Platform integrity with ARM TrustZone

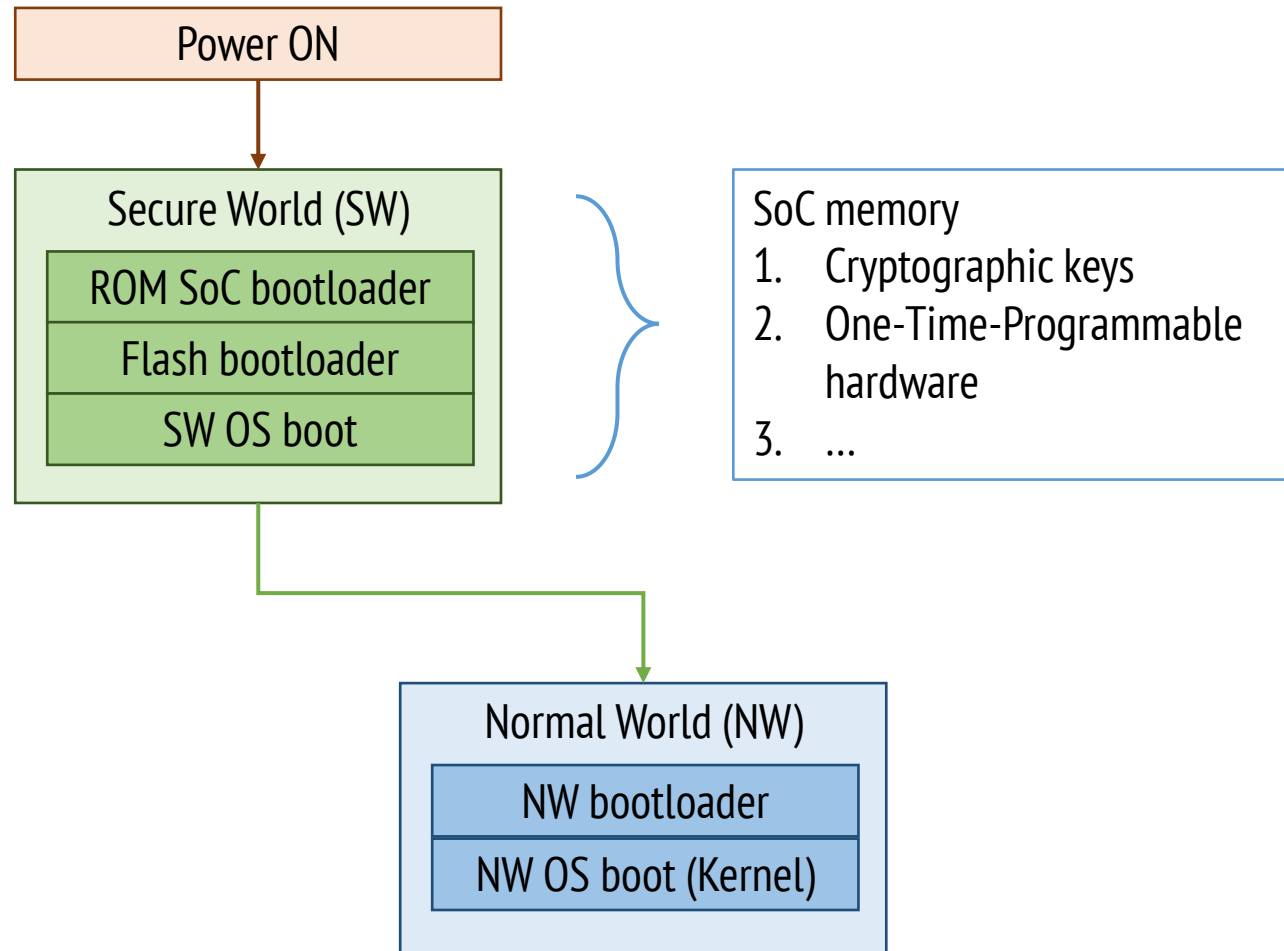


# Platform integrity with ARM TrustZone





# Platform integrity with ARM TrustZone



# Things to remember

- Economise mechanism: keep the system simple i.e. use a secure OS with controlled I/O to load the main kernel
- Fail secure: if things go wrong stop
- ...

# Secure Enclave Processor (SEP)

- Enable sensitive data to be stored securely
- Performs secure services for the rest of the SOC
- Runs its own operating system (SEPOS) which includes: kernel, drivers, services, and applications
- Supports multiple services: TouchID, ApplePay...

# Secure Enclave Processor (SEP)

## Hardware functionality

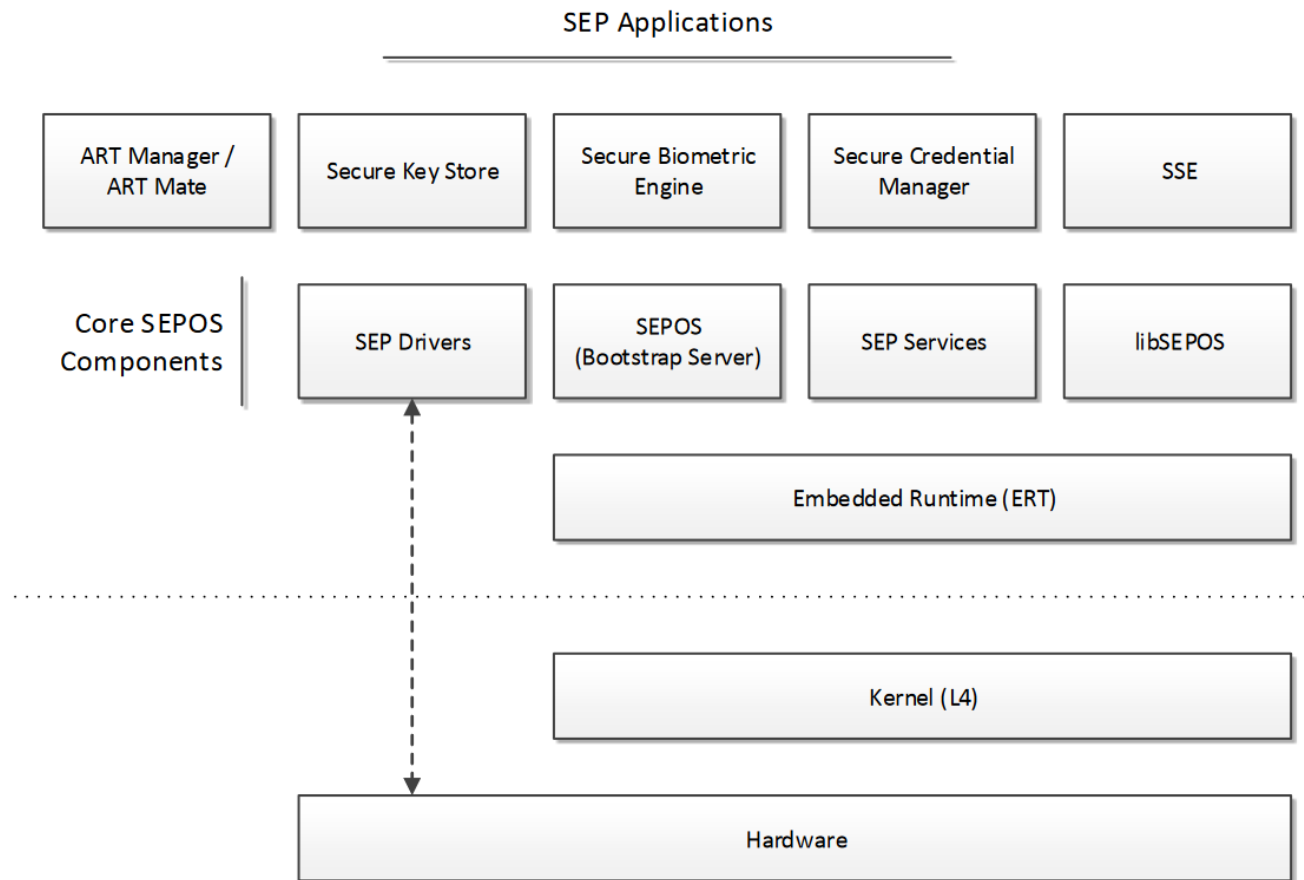
- Crypto engine
- Random Number Generator
- Fuses
- GID/UID
- Dedicated scratch RAM
- Hardware “filter” to prevent application processor (AP) to SEP memory access

# Secure Enclave Processor

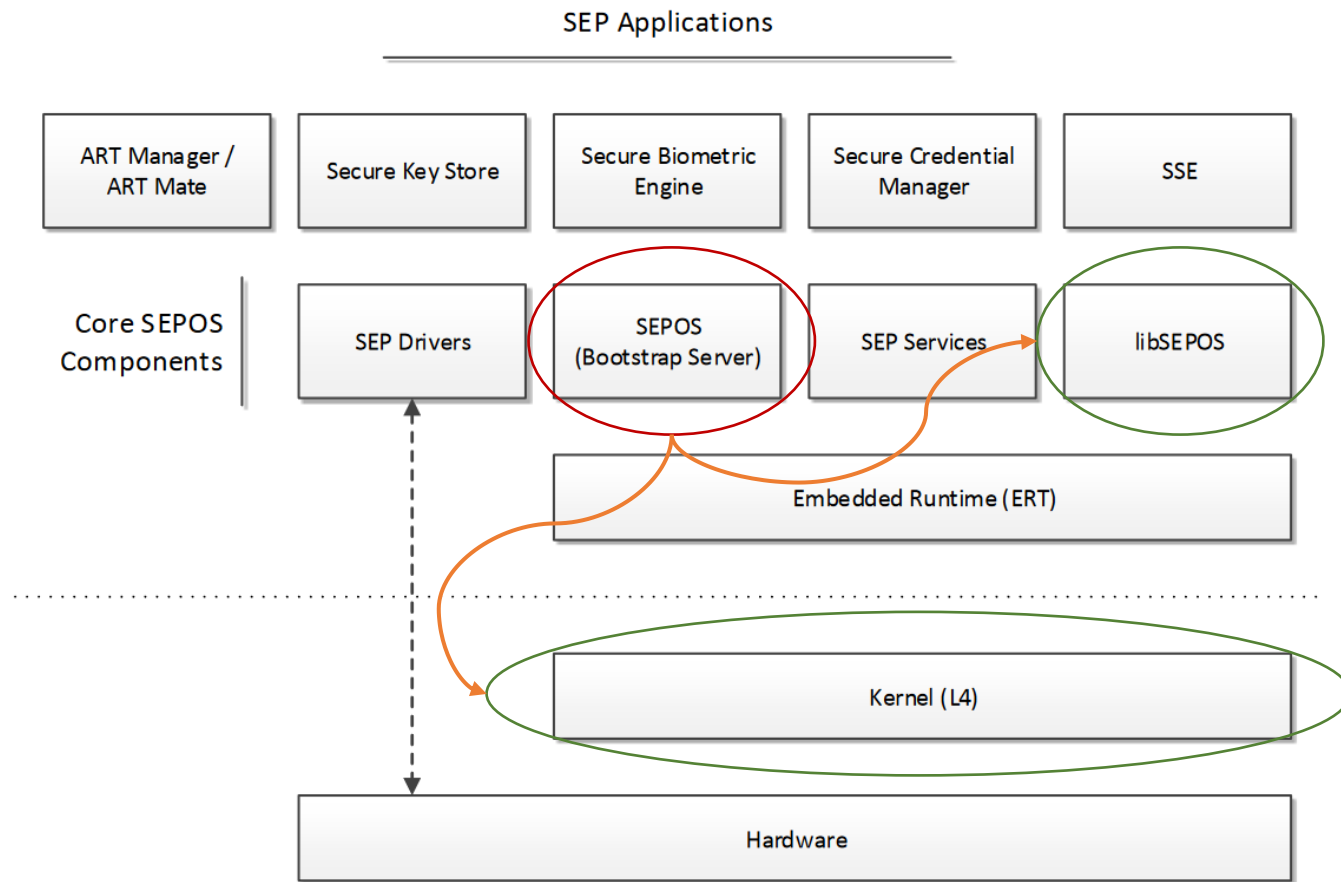
Shared functionality with application processor

- Clock
- RAM
- Power manager
- ...

# SEP Architecture

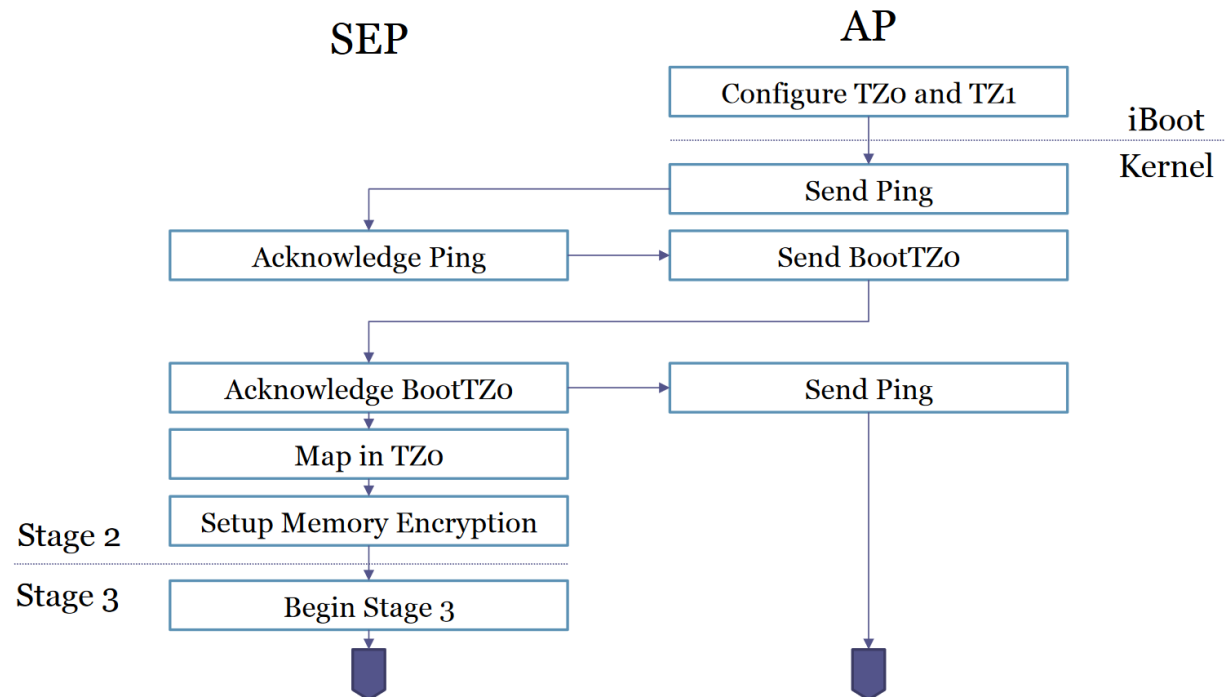


# SEP Architecture



# Secure Enclave Processor (SEP)

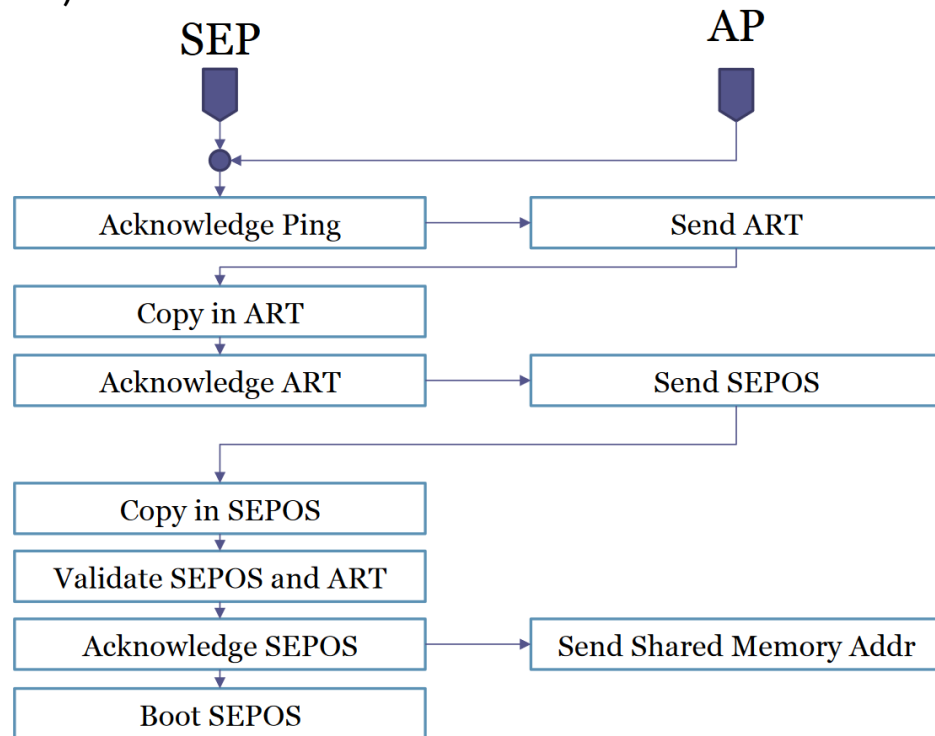
1. Configure Trust Zones 0/1  
(TZ0-SEP, TZ1-AP)
2. Check for SEP
3. Configure memory protection





# Secure Enclave Processor (SEP)

1. Send anti-replay token (ART)
2. Verify ART
3. Verify the SEP operating system (SEPOS)  
i.e. 4096 bytes
4. Establish shared memory location



# SEP communication

- Secure Mailbox allows the AP to communicate with the SEP
- Supported through the SEP Manager API
- Implemented using the SEP device I/O registers

# Things to remember

- Authenticate requests: messages to the SEP are authenticated, apps are authenticated before use
- Fail secure: if anything goes wrong stop execution
- Segregation of duties: all cryptographic operations and accesses to secure hardware are intermediated by the SEP
- Secure the weakest link: protect access to memory

Application security

# Multiple Layers of Defense

Windows Store  
Google Play  
App Store

Unknown  
Sources  
Warning

Install  
Confirmation

Verify Apps  
Consent

Verify Apps  
Warning

Runtime Security  
Checks

Sandbox and  
Permissions



# Main “techniques”

- Code signing
- Runtime security
  - a) Mandatory access control (MAC)
  - b) Sandboxing
  - c) Memory protection (e.g. address space layout randomisation (ASLR), ARM Execute Never (XN))

# Code signing

- Ensure applications have an approved source and haven't been tampered with
- Executable code is signed with store specific certificates
- Prevent applications from loading unsigned code resources and self-modifying code
- Access hardware with OS APIs

# Mandatory access control (MAC)

- Support over all OS
  - Selinux in Linux and Android
  - Mandatory Integrity Control in Windows Vista/7/8/10
  - TrustedBSD variant in Apple IOS and OSX
- MAC is usually enforced over all processes, even processes running with root/superuser privileges.
- MAC usually defaults to denial: anything that is not explicitly allowed is denied.



# Sandboxing

- Restrict applications from using data and resources from other apps.
- Each app has its own random “home directory”
- Run applications as non-privileged user
- Mount “important” partitions as read-only

# Memory protection

- Address space layout randomisation (ASLR)
  - protects memory from corruption
  - protects against attacks that target the stack and/or memory addresses
- Hardware support like ARM's Execute never
  - Supports “permissions” for memory pages
  - Marks memory pages as non-executable

# Conclusion

- Multiple surfaces of attack
- Root of trust
- Layered protection i.e. defence in depth
- Access control
- Audit and monitor
- Make security usable
- Authenticate requests
- Control access
- ...