

***Report on***  
**Mini Compiler for If-Else and While Constructs in Python**

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***Compiler Design***

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# INTRODUCTION

A compiler is a special program that processes statements written in a particular programming language and turns them into machine language or "code" that a computer's processor uses.

A Mini Compiler is Built for Python and handles the If-Elif-Else and the While Constructs using c programming language. Lex and Flex are tools for generating scanners: programs which recognize lexical patterns in text. Flex is a faster version of Lex. Lex/Flex refers to either of the tools. Yacc and Bison are tools for generating parsers: programs which recognize the grammatical structure of programs. Bison is a faster version of Yacc.

The optimizer applies semantics preserving transformations to the annotated parse tree to simplify the structure of the tree and to facilitate the generation of more efficient code. The code generator transforms the simplified annotated parse tree into object code using rules which denote the semantics of the source language. The code generator may be integrated with the parser.

# ARCHITECTURE OF THE LANGUAGE

Python has a very flexible syntax and we have tried to incorporate as much as possible from our experience of using python into the grammar. All lines of code terminate upon seeing a newline character. We have taken care of the following:

1. If-Elif-Else and While constructs
2. Print Statements
3. pass, break and void returns
4. Function definitions and Calls
5. Lists
6. All arithmetic operators and all boolean operators
7. Single Line Comments (#)

Semantically we have checked the following :

- Whether any variable used on the RHS is defined and in the current scope or any Enclosing Scope of the current scope.
- Whether a variable being indexed is List
- Whether all expressions in the If and While Clauses are Boolean expressions

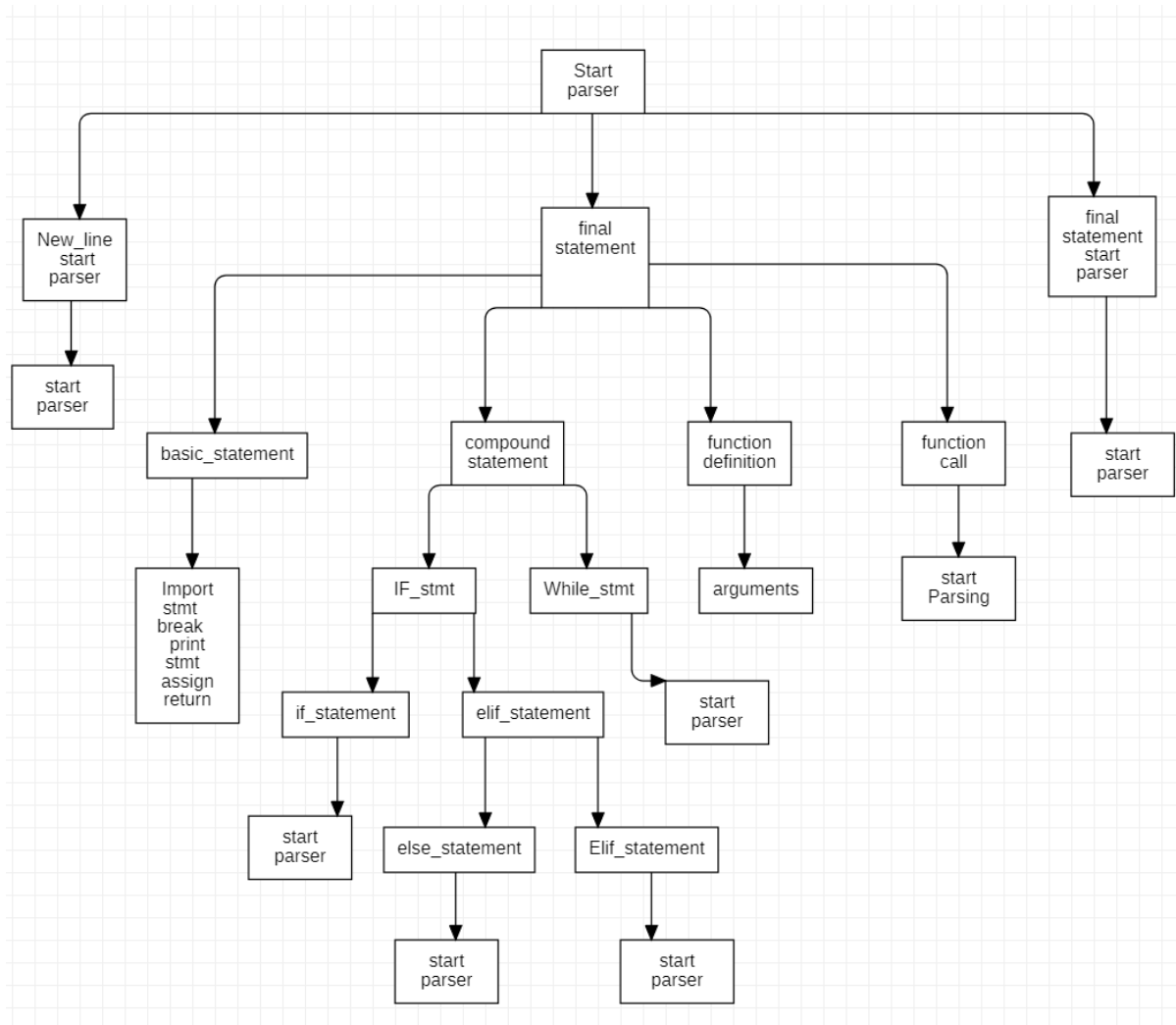
# LITERATURE SURVEY

1. Lex and Yacc Doc by Tom Niemann
2. Official Bison Documentation : <https://www.gnu.org/software/bison/manual>
3. Stackoverflow : <https://stackoverflow.com>
4. Compiler Construction using Flex and Bison - Anthony A. Aaby, Walla Walla College, cs.wwc.edu, aabyan@wwc.edu, Version of February 25, 2004

# THE CONTEXT FREE GRAMMAR

Grammar	Expression	Legend
digit	→ [0-9]	CLN : Colon ( ':' )
constant	→ [digit]+   QUOTE [ _a-zA-Z ]+ QUOTE	NL : New Line ( '\n ' )
variable	→ [ _a-zA-Z ] [ _a-zA-Z0-9 ]*	EQL : Equal ( '=' )
term	→ variable   constant	L_Paren : '('
		R_Paren : ')'
bool_exp	→ True   False   arith_exp ROP arith_exp bool_exp or bool_exp   bool_exp and bool_exp   not bool_exp   L_Paren bool_exp R_Paren	ROP : Relational Operators ('<', '>', '==' ..)
arith_exp	→ term   L_Paren arith_exp R_Paren   arith_exp + arith_exp   arith_exp - arith_exp   arith_exp * arith_exp   arith_exp / arith_exp	ID : Indent
assign_stmt	→ variable EQL term	DD : Dedent
pass_stmt	→ "pass"	QUOTE : "
import_stmt	→ "import" <pkg_name>	STR : Any String
expression_stmt	→ bool_exp   arith_exp	Trip_Quote : "'"
basic_stmt	→ expression_stmt   assign_stmt   pass_stmt	HASH : #
stmt_list	→ basic_stmt NL stmt_list   basic_stmt	
suite	→ stmt_list NL   NL ID stmt + DD	
stmt	→ basic_stmt NL   cmpd_stmt	
cmpd_stmt	→ if_stmt   while_stmt	
while_stmt	→ "while" bool_exp CLN NL ID stmt_list DD	
if_stmt	→ "if" bool_exp CLN suite ("elif" bool_exp CLN suite)* ["else" CLN suite]	
comment	→ HASH STR   Trip_Quote STR multiline_comment	
multiline_comment	→ Trip_Quote   STR NL multiline_comment	

# THE CFG Flow Diagram



## DESIGN STRATEGY

First and foremost, every connection points back to the Symbol Table, as per our Design Strategy. This ensures that the Symbol table is linked to the necessary nodes of the Abstract Syntax tree and the required Quadruples in the Intermediate code. The symbol table also contains all of the Temporaries provided by the compiler.

As you can see in the sample output, the Symbol Table stores 'Records' with four columns, namely

1. Scope: Scope of each variable contained in record
2. Name: Value/Name of each variable contained in record.
3. Type : Type of each variable contained in record
  - i. PackageName
  - ii. Func\_Name
  - iii. Identifier
  - iv. Constant

- v. ListTypeID
  - vi. ICGTempVar
  - vii. ICGTempLabel
4. Declaration : Line of Declaration of each variable contained in record
  5. Last Use Line

The scope is a property of indentation depth, and we have a tuple of the parent's scope and the current scope measured using the indentation depth to make it distinct.

There are two types of nodes in the Abstract Syntax Tree: Leaf nodes and Internal nodes. Depending on the construct it serves, the nodes will have a variable number of children (0-3). Take the If-Else Statement for example.

```

                If
      Condition  CodeBlock  Else
  
```

We take the AST and store it as a matrix of levels in order to represent it. Each stage of the AST has been printed, as seen in the sample output. In addition, next to each Internal node is a number enclosed in brackets that represents the number of children they have in the next step. The identifiers, constants, Lists, and packages leaf nodes in the AST point to a symbol table record.

Recursively stepping through the AST generates the intermediate code. Each line is saved as a quadruple (Operation, Arg1, Arg2, Result) so that code can be quickly optimised in the steps that follow.

In the whole code, we remove dead code, explicitly unused variables. For example, if we have the following lines of code:

```

a = 10
b = 10
c = a + b
  
```

And these 3 variables are not used on any other RHS, then these 3 lines of code are Eliminated during optimization. All the dead variables in the code are removed. We iterate through the Quads to do this.

We introduced a stack and used three tokens to take control of the indentation-based code structure and scoping. The top of the stack always points to the current indentation value; if that value doesn't adjust when scanning the next line, it means we're in the same scope, and we return the token ND, which stands for "No-dent." If the value increases, we are entering a sub-scope, and we return the token 'ID,' which stands for 'Indent,' and if the value reduces, we are returning to one of the enclosing scopes, and we return 'DD,' which stands for 'Dedent.'

Whenever the parser encounters an error, it prints the line number and column number of the error. We also display the following errors:

- Identifier <var> Not declared in scope
- Identifier <var> Not Indexable

All Comments are removed from the code before parsing.

# IMPLEMENTATION DETAILS

The Symbol table uses two Structures,

```
typedef struct record {
    char *type;
    char *name;
    int declLineNo;
    int lastUseLine;
} record;

typedef struct STable {
    int no;
    int noOfElems;
    int scope;
    record *Elements;
    int Parent;
} STable;
```

The “record” arrangement reflects each record in the symbol table. Each symbol table will have a maximum of “MAXRECST” documents, MAXRECST is a macro. One symbol table is represented by the “STable” arrangement. For each scope, a new Symbol table is created. MAXST is a macro that allows a maximum of “MAXST” symbol tables and hence scopes to occur. The Abstract Syntax Tree uses one structure,

```
typedef struct ASTNode {
    int nodeNo;
    /*Operator*/
    char *NType;
    int noOps;
    struct ASTNode** NextLevel;
    /*Identifier or Const*/ record *id;
} node;
```

This ASTNode arrangement looks after all leaf nodes and internal "operator" nodes. Based on the form of the node, the appropriate values are set.

Each node may have a maximum of three children. The AST is printed by storing it in a Matrix of Order "MAXLEVELS" x "MAXCHILDREN" and then printing the matrix Levelwise. This Matrix is a pointer matrix to the AST. The Node's "noOps" element specifies the number of children it has.

The Three-Address Code is represented and stored as Quads that are given by the structure,

```
typedef struct Quad {
    char *R;
    char *A1;
    char *A2;
    char *Op;
    int I;
} Quad;
```

During code optimization, the last element, the integer 'I,' is used. In an array of Quads, all of the Three-Address codes are stored as Quads. A limit of "MAXQUADS" quadruples can be used.

We loop around the code until there is no more code that can be removed for dead code removal. To see if a quadruple has dead code, look to see if the result parameter/element of that quad exists as an argument to all other quads that haven't been deleted. If not, we consider the quad to be dead code and assign the value "-1" to the variable "I".

Scope checking is accomplished by recursively stepping through enclosing scopes and locating the variable's most recent definition. If no definition is found, the error is printed.

If you wish to run the code,

Install flex and bison using the below commands(if not installed),

```
sudo apt-get upgrade
sudo apt-get install bison flex
```

```
lex grammar.l
yacc -dv grammar.y
gcc lex.yy.c y.tab.c -g -ll -o Test.out
./Test.out < InputFile.txt
```

## RESULTS AND SHORTCOMINGS

We now have a mini compiler that parses grammar corresponding to basic Python syntax and produces an efficient intermediate representation as a result. Below are some of the places where our mini compiler fell short:

- Just one very simple optimization is used, and it does not result in a significant reduction in code density.
- There are a few memory leaks in the programme, but the majority of them have been fixed.



# SNAPSHOTS

## 1. TestCase\_1

### 1.1 Input File

```
pb@ubuntu:~/Downloads/Mini-Python-Compiler-in-Lex-and-Yacc-master$ cat input3.txt
#Basic Code

import scipy

x = 2
y = 1

a = 3
b = 4
d = a+b

if(x==1):
    c=1
elif(y==1):
    c=2
else:
    c=1

pb@ubuntu:~/Downloads/Mini-Python-Compiler-in-Lex-and-Yacc-master$
```

### 1.2 Token Sequence

```
-----Token Sequence-----
1 T_NL
2 T_NL
3 T_IMPT T_scipy T_NL
4 T_NL
5 T_x T_EQL T_2 T_NL
6 T_y T_EQL T_1 T_NL
7 T_NL
8 T_a T_EQL T_3 T_NL
9 T_b T_EQL T_4 T_NL
10 T_d T_EQL T_a T_PL T_b T_NL
11 T_NL
12 T_If T_OP T_x T_EQ T_1 T_CP T_Cln T_NL
13 T_ID T_c T_EQL T_1 T_NL
14 T_Elif T_OP T_y T_EQ T_1 T_CP T_Cln T_NL
15 T_ID T_c T_EQL T_2 T_NL
16 T_Else T_Cln T_NL
17 T_ID T_c T_EQL T_1 T_NL
18 T_NL
19 T_EOF

$$ PARSING COMPLETED $$
```

### 1.3 Abstract Syntax Tree

```

-----Abstract Syntax Tree-----
NewLine(2)
import(1)  NewLine(2)
scipy      =(2)  NewLine(2)
  x        2    =(2)  NewLine(2)
    y      1    =(2)  NewLine(2)
      a    3    =(2)  NewLine(2)
        b    4    =(2)  If(3)
          d    +(2)  ==(2)  BeginBlock(2)  Elif(3)
            a    b    x    1    =(2)  EndBlock  ==(2)  BeginBlock(2)  Else(1)
              c    1    y    1    =(2)  EndBlock  BeginBlock(2)
                c    2    =(2)  EndBlock
                  c    1

```

### 1.4 Optimized Quads

```

-----All Quads-----
0      import  scipy  -      -
1      =      2      -      T2
2      =      T2     -      x
3      =      1      -      T5
4      =      T5     -      y
13     =      x      -      T19
14     =      1      -      T20
15     ==     T19     T20    T21
16     If False      T21    -      L0
19     goto      -      -      L1
20     Label     -      -      L0
21     =      y      -      T27
22     =      1      -      T28
23     ==     T27     T28    T29
24     If False      T29    -      L0
27     goto      -      -      L1
28     Label     -      -      L0
31     Label     -      -      L1
32     Label     -      -      L1
-----

```

## 1.5 Intermediate Code

```
import scipy
T2 = 2
x = T2
T5 = 1
y = T5
T8 = 3
a = T8
T11 = 4
b = T11
T14 = a
T15 = b
T16 = T14 + T15
d = T16
T19 = x
T20 = 1
T21 = T19 == T20
If False T21 goto L0
BeginBlockT22 = 1
c = T22
EndBlockgoto L1
L0: T27 = y
T28 = 1
T29 = T27 == T28
If False T29 goto L0
BeginBlockT30 = 2
c = T30
EndBlockgoto L1
L0: BeginBlockT35 = 1
c = T35
EndBlockL1: L1:
```

## 1.6 All Symbol Tables

```
-----All Symbol Tables-----
Scope   Name      Type      Declaration  Last Used Line
(0, 1)  scipy     PackageName  3            3
(0, 1)  2         Constant    5            5
(0, 1)  x         Identifier   5            12
(0, 1)  1         Constant    6            12
(0, 1)  y         Identifier   6            14
(0, 1)  3         Constant    8            8
(0, 1)  a         Identifier   8            10
(0, 1)  4         Constant    9            9
(0, 1)  b         Identifier   9            10
(0, 1)  d         Identifier  10           10
(0, 1)  T2        ICGTempVar  -1           -1
(0, 1)  T5        ICGTempVar  -1           -1
(0, 1)  T8        ICGTempVar  -1           -1
(0, 1)  T11       ICGTempVar  -1           -1
(0, 1)  T14       ICGTempVar  -1           -1
(0, 1)  T15       ICGTempVar  -1           -1
(0, 1)  T16       ICGTempVar  -1           -1
(0, 1)  T19       ICGTempVar  -1           -1
(0, 1)  T20       ICGTempVar  -1           -1
(0, 1)  T21       ICGTempVar  -1           -1
(0, 1)  L0        ICGTempLabel -1           -1
(0, 1)  T22       ICGTempVar  -1           -1
(0, 1)  L1        ICGTempLabel -1           -1
(0, 1)  T27       ICGTempVar  -1           -1
(0, 1)  T28       ICGTempVar  -1           -1
(0, 1)  T29       ICGTempVar  -1           -1
(0, 1)  T30       ICGTempVar  -1           -1
(0, 1)  T35       ICGTempVar  -1           -1
(0, 2)  1         Constant    13           14
(0, 2)  c         Identifier   13           13
(1, 4)  2         Constant    15           15
(1, 4)  c         Identifier   15           15
(1, 8)  1         Constant    17           17
(1, 8)  c         Identifier   17           17
-----
```

## 2. TestCase\_2

### 2.1 Input File

```
pb@ubuntu:~/Downloads/Mini-Python-Compiler-in-Lex-and-Yacc-master$ cat input2.txt
import hWorld
x=10
y=10

#Comment1
listX = []
while(x==10):
    x = 1
#Comment2
if(x==y):
    x=10
else:
    x=10

if(x==y):
    x=10
else:
    y=10
```

### 2.2 Token Sequence

```
-----Token Sequence-----
1 T_IMPT T_hWorld T_NL
2 T_x T_EQL T_10 T_NL
3 T_y T_EQL T_10 T_NL
4 T_NL
5 T_NL
6 T_listX T_EQL T_OB T_CB T_NL
7 T_While T_OP T_x T_EQ T_10 T_CP T_Cln T_NL
8 T_ID T_x T_EQL T_1 T_NL
9 T_NL
10 T_If T_OP T_x T_EQ T_y T_CP T_Cln T_NL
11 T_ID T_x T_EQL T_10 T_NL
12 T_Else T_Cln T_NL
13 T_ID T_x T_EQL T_10 T_NL
14 T_NL
15 T_If T_OP T_x T_EQ T_y T_CP T_Cln T_NL
16 T_ID T_x T_EQL T_10 T_NL
17 T_Else T_Cln T_NL
18 T_ID T_y T_EQL T_10 T_NL
19 T_NL
20 T_EOF

$$ PARSING COMPLETED $$
```

## 2.3 Abstract Syntax Tree

```
-----Abstract Syntax Tree-----
NewLine(2)
import(1) NewLine(2)
hWorld  =(2) NewLine(2)
        x 10  =(2) NewLine(2)
              y 10 listX NewLine(2)
                    While(2) NewLine(2)
                          ==(2) BeginBlock(2) If(3) If(3)
                                x 10  =(2) EndBlock ==(2) BeginBlock(2) Else(1) ==(2) BeginBlock(2) Else(1)
                                      x 1  x y  =(2) EndBlock BeginBlock(2) BeginBlock(2) x y  =(2) EndBlock BeginBlock(2)
                                          10 ==(2) EndBlock x 10  =(2) EndBlock
                                              x 10 y 10
```

## 2.4 Optimized Quads

```
-----All Quads-----
0      import  hWorld  -      -
1      =      10      -      T2
2      =      T2      -      x
3      =      10      -      T5
4      =      T5      -      y
5      =      x      -      T9
6      =      10      -      T10
7      ==     T9      T10    T11
8      Label  -      -      L0
9      If False      T11    -      L1
10     =      1      -      T12
11     =      T12    -      x
12     goto   -      -      L0
13     Label  -      -      L1
14     =      x      -      T18
15     =      y      -      T19
16     ==     T18    T19    T20
17     If False      T20    -      L4
18     =      10      -      T21
19     =      T21    -      x
20     goto   -      -      L5
21     Label  -      -      L4
22     =      10      -      T26
23     =      T26    -      x
24     Label  -      -      L5
25     =      x      -      T33
26     =      y      -      T34
27     ==     T33    T34    T35
28     If False      T35    -      L6
29     =      10      -      T36
30     =      T36    -      x
31     goto   -      -      L7
32     Label  -      -      L6
33     =      10      -      T41
34     =      T41    -      y
35     Label  -      -      L7
-----
```

## 2.5 Intermediate Code

```
import hWorld
T2 = 10
x = T2
T5 = 10
y = T5
T9 = x
T10 = 10
T11 = T9 == T10
L0: If False T11 goto L1
BeginBlockT12 = 1
x = T12
EndBlockgoto L0
L1: T18 = x
T19 = y
T20 = T18 == T19
If False T20 goto L4
BeginBlockT21 = 10
x = T21
EndBlockgoto L5
L4: BeginBlockT26 = 10
x = T26
EndBlockL5: T33 = x
T34 = y
T35 = T33 == T34
If False T35 goto L6
BeginBlockT36 = 10
x = T36
EndBlockgoto L7
L6: BeginBlockT41 = 10
y = T41
EndBlockL7:
```

## 2.6 All Symbol Tables

-----All Symbol Tables-----				
Scope	Name	Type	Declaration	Last Used Line
(0, 1)	hWorld	PackageName	1	1
(0, 1)	10	Constant	2	7
(0, 1)	x	Identifier	2	7
(0, 1)	y	Identifier	3	15
(0, 1)	listX	ListTypeID	6	6
(0, 1)	T2	ICGTempVar	-1	-1
(0, 1)	T5	ICGTempVar	-1	-1
(0, 1)	T9	ICGTempVar	-1	-1
(0, 1)	T10	ICGTempVar	-1	-1
(0, 1)	T11	ICGTempVar	-1	-1
(0, 1)	L0	ICGTempLabel	-1	-1
(0, 1)	L1	ICGTempLabel	-1	-1
(0, 1)	T12	ICGTempVar	-1	-1
(0, 1)	T18	ICGTempVar	-1	-1
(0, 1)	T19	ICGTempVar	-1	-1
(0, 1)	T20	ICGTempVar	-1	-1
(0, 1)	L4	ICGTempLabel	-1	-1
(0, 1)	T21	ICGTempVar	-1	-1
(0, 1)	L5	ICGTempLabel	-1	-1
(0, 1)	T26	ICGTempVar	-1	-1
(0, 1)	T33	ICGTempVar	-1	-1
(0, 1)	T34	ICGTempVar	-1	-1
(0, 1)	T35	ICGTempVar	-1	-1
(0, 1)	L6	ICGTempLabel	-1	-1
(0, 1)	T36	ICGTempVar	-1	-1
(0, 1)	L7	ICGTempLabel	-1	-1
(0, 1)	T41	ICGTempVar	-1	-1
(0, 2)	1	Constant	8	8
(0, 2)	x	Identifier	8	15
(1, 4)	10	Constant	11	11
(1, 4)	x	Identifier	11	11
(1, 8)	10	Constant	13	13
(1, 8)	x	Identifier	13	13
(1, 16)	10	Constant	16	16
(1, 16)	x	Identifier	16	16
(1, 32)	10	Constant	18	18
(1, 32)	y	Identifier	18	18



# CONCLUSION

A Compiler for Python was implemented using lex and yacc tools.. In addition to the constructs specified, the basic python constructs were implemented and function definitions and calls were supported. We have also shown the tokens generated , the abstract syntax tree ,intermediate code and the optimized code separately.

The compiler also reports the basic errors and gives the line number and column number. The Intermediate code was represented by quads , which was later optimized to remove dead code and reduce the usage of temporary variables.

# FURTHER ENHANCEMENTS

The compiler can be further enhanced by adding support for 'For' Loops, classes, and other types of loops. Improved memory management for registers, More effective optimization strategies Semantic analysis for parameter matching, Error recovery can be implemented to enhance the current work.