CSC 212: Data Structures and Abstractions 06: Dynamic Arrays

Prof. Marco Alvarez

Department of Computer Science and Statistics University of Rhode Island

Fall 2025



C-style arrays

- Contiguous sequence of elements of identical type
 - random access: base_address + index * sizeof(type)



array name: A

array length: n

- Statically allocated arrays
 - ✓ allocated in the stack (fixed-length), size known at compile time
- Dynamic allocated arrays
 - ✓ allocated in the heap (fixed-length), size may be determined at runtime

Dynamic arrays

Example

- Where are each of these variables allocated? (stack vs heap)
- Can the arrays change size during program execution?

```
void sort(int *arr, int size) {
    int i, j, temp;
    // sorting logic here
    // ...
}

int main() {
    int array[100];
    int *ptr;
    // ...
    ptr = new int[100];
    //...
    sort(ptr, 100);
    sort(array, 100);
    //...
    delete[] ptr;
    return 0;
}
```

Dynamic arrays

- Limitations of C-style arrays
 - ✓ size <u>must be known at compile time</u>
 - alternatively, use dynamic memory allocation
 - ✓ once created, array size does not change (inflexible)
- Dynamic arrays
 - ✓ can grow or shrink in size during run-time
 - essential for many applications, for example, a server keeping track of a queue of requests
 - combine the flexibility of dynamic memory allocation with the efficiency of fixed-length arrays
 - e.g. std::vector in C++, ArrayList in Java, List in Python, Array in JavaScript, List in C#, Vec in Rust, etc.

Designing a dynamic array class in C++

```
class DvnamicArrav {
    private:
                                              // pointer to the (internal) array
        int capacity;
                                              // total number of elements that can be stored
        int size;
        DvnamicArrav():
        ~DynamicArray();
                                              // destructor
        void push_back(int val);
                                              // add an element to the end
                                              // remove the last element
        void pop_back();
        const int& operator[](int idx) const; // read-only access at a specific index
        int& operator[](int idx);
                                             // access at a specific index (can modify)
        void insert(int val, int idx);
                                              // insert an element at a specific index
        void erase(int idx);
                                              // remove an element at a specific index
        void resize(int len);
                                              // change the capacity of the array
        int size():
                                              // return the number of elements
        int capacity();
                                              // return the capacity
                                              // check if the array is empty
        void clear();
                                              // remove all elements, maintaining the capacity
        // additional methods can be added here
};
```

A class definition specifies the **data members** and **member functions** of the class. The data members are the attributes of the class, and the member functions are the operations that can be performed on the data members. The class definition is a blueprint for creating objects of the class.

std::vector from C++ STL

```
#include <iostream>
#include <vector>

int main()
{
    // create a vector containing integers
    std::vector<int> v = {8, 4, 5, 9};

    // add two more integers to vector
    v.push_back(6);
    v.push_back(9);

    // overwrite element at position 2
    v[2] = -1;

    // print out the vector
    for (int n : v)
        std::cout << n << ' ';
    std::cout << '\n';
}</pre>
```

Resizing dynamic arrays

- · Grow
 - when the array is full (Size == capacity), allocate a new array with increased capacity, copy elements from old to new array, deallocate old array





- Shrink
 - optional optimization, used when the number of elements is "significantly" less than the capacity, allocate a new array with decreased capacity, copy the elements from old to new array, and deallocate the old array

8

Grow by one

- When array is full, grow capacity to: capacity + 1
 - starting from an empty array, count number of array accesses (reads and writes)
 for appending n elements (ignore cost of allocating/deallocating memory)

element	cost copy	cost append	
1	2 x 0	1	
2	2 x 1	1	
3	2 x 2	1	
4	2 x 3	1	
5			
6			
n-1	2 x (n-2)	1	
n	2 x (n-1)	1	
	read and write	write	

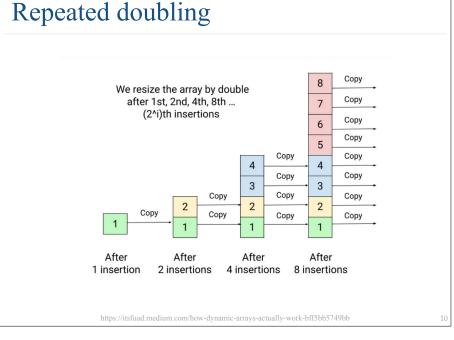
$$T(n) = n + \sum_{i=0}^{n-1} 2i$$

$$= n + 2 \sum_{i=0}^{n-1} i$$

$$= n + 2 \left(\frac{n(n-1)}{2}\right)$$

$$= \Theta(n^2) \xrightarrow{\text{cost of adding n elements}}$$

The amortized cost of inserting an element is $\Theta(n)$, meaning that any sequence of n insertions takes at most $\Theta(n^2)$ time in total.



Grow by factor

- → When array is full, grow capacity to: capacity * factor
 - \cdot called <u>repeated doubling</u> when factor == 2
 - starting from an empty array, count number of array accesses (reads and writes)
 for appending n elements (ignore cost of allocating/deallocating memory)

element	cost copy	cost append	
1	2 x 0	1	
2	2 x 1	1	
3	2 x 2	1	
4	_	1	
5	2 x 4	1	
6	_	1	
7	_	1	
8	_	1	
9	2 x 8	1	
10	_	1	
n-1	_	1	
n	_	1	
	read and write	write	

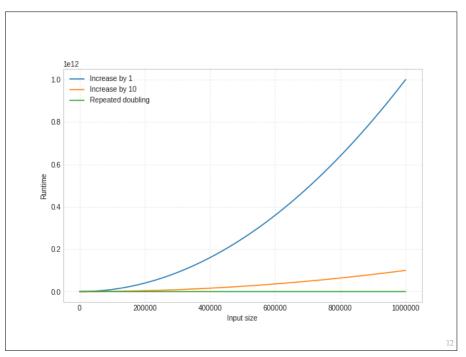
$$T(n) = n + 2 \sum_{i=0}^{\log n - 1} 2^{i}$$

$$= n + 2 \left(2^{\log n} - 1 \right)$$

$$= n + 2n - 2$$

$$= \Theta(n) \xrightarrow{\text{cost of adding } n \text{ elements}}$$

The amortized cost of inserting an element is $\Theta(1)$, meaning that any sequence of n insertions takes at most $\Theta(n)$ time in total.



Shrinking the array

- May half the capacity when array is one-half full
 - worst-case when the array is full and we <u>alternate between adding and</u> removing elements
 - each alternating operation would require resizing the array
- More efficient resizing
 - ✓ <u>half the capacity</u> when the array is <u>one-quarter</u> full
- In practice ...
 - most standard implementations do not automatically shrink capacity
 - avoids performance penalties from frequent resizing
 - instead, they provide explicit operations like shrink_to_fit() that allow the programmer to request size reduction when deemed necessary

Growth factors by language

- . C++ (std::vector)
 - grow by 1.5 or 2.0 times the current capacity (depending on compiler)
 - ✓ shrink when the array is one-quarter full
- Java (ArrayList)
 - ✓ grow by 1.5 of the current capacity
 - ✓ shrink when the array is one-half full
- Python(list)
 - ✓ grow by ~1.125 times the current capacity
 - ✓ shrink when the array is one-quarter full
- Rust (std::vec::Vec)
 - ✓ grow by 2 times the current capacity
 - shrink when the array is one-half full

Information taken from claude.ai (to be confirmed)

bottom line: growth factors range from ~1.12 to ~2 depending on language used

14

Practice

- Complete the following table with rates of growth using Θ notation
 - assume we implement a dynamic array with repeated doubling and no shrinking

Best case	Average case	Worst case
	Best case	Best case Average case

Review references in C++ non-const References int i = 2; int& ri = i; // reference to i ri and i refer to the same object / memory location cout << i <<'\n'; cout << ri <<'\n'; // 2 cout << i <<'\n'; // 5 cout << ri <<'\n': // 5 cout << i <<'\n'; // 88 cout << ri <<'\n'; // 88 · references cannot be "null", i.e., they must always refer to an object a reference must always refer to the same memory location reference type must agree with the type of the referenced object int k = 3: int& ri = i; ri = k; // assigns value of k to i (target of ri) // X COMPILER ERROR: reference must be initialized double& r3 = i; // x COMPILER ERROR: types must agree