

CSC 212: Data Structures and Abstractions

Binary search trees (part 1)

Prof. Marco Alvarez

Department of Computer Science and Statistics
University of Rhode Island

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Trees

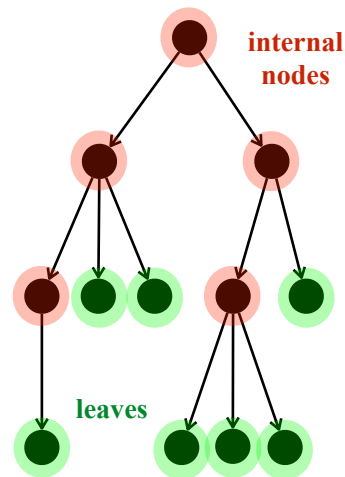
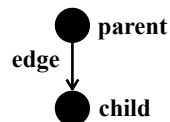
Trees

Definition

- data structure that consists of **nodes** connected by **edges**
 - hierarchical structure, with a single **root** node
 - each node can have zero or more **children**

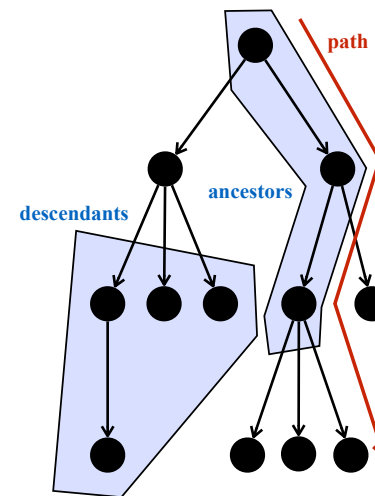
Terminology

- each node is either a **leaf** or an **internal node**
 - leaves are nodes with no children, while internal nodes are nodes with one or more children
- nodes with the same **parent** are **siblings**



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Paths



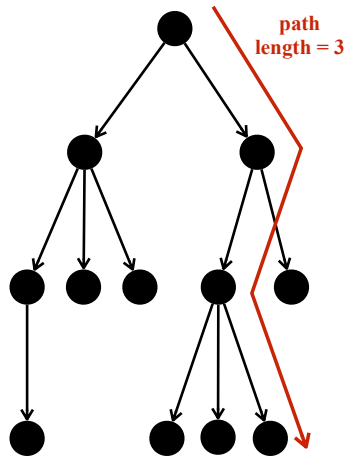
A **path** from node v_0 to v_n is a sequence of nodes v_0, v_1, \dots, v_n , where there is a (directed) edge from one node to the next

The **descendants** of a node v are all nodes reached by a path from node v to the leaf nodes

The **ancestors** of a node v are all nodes found on the path from the root to node v

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Depth and height



The length of a **path** is the number of edges in the path

The **depth** (level) of a node v is the length of the path from the root node to v

The **height** of a node v is the length of the path from v to its deepest descendant

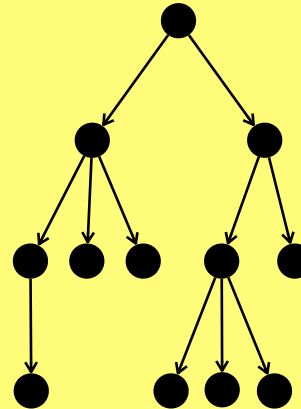
The **depth of the tree** is the depth of deepest node

The **height of the tree** is the height of the root

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Practice

- Label all nodes with height and depth
- ✓ indicate the height and the depth of the tree



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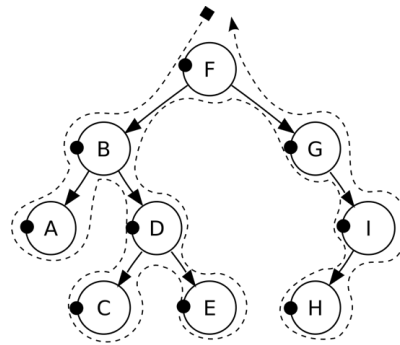
Traversals

Definition

- ✓ a **traversal** is a way of visiting all the nodes in a tree

Types of traversals:

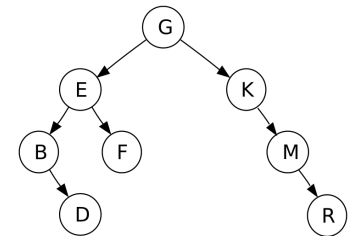
- ✓ **pre-order traversal:** visit the root node first, then recursively visit all subtrees
- ✓ **post-order traversal:** recursively visit all subtrees first, then visit the root node



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Pre-order traversal

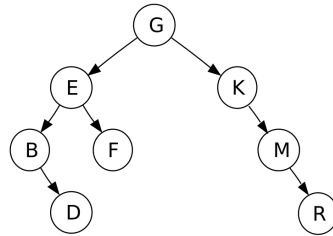
```
algorithm preorder(p) {
    visit(p)
    for each child c of p {
        preorder(c)
    }
}
```



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Post-order traversal

```
algorithm postorder(p) {  
  for each child c of p {  
    postorder(c)  
  }  
  visit(p)  
}
```



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Binary trees

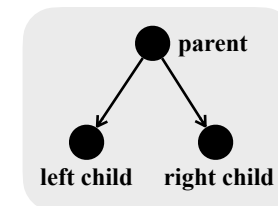
k-ary trees

- **k-ary tree**
 - ✓ every node has between 0 and k children
- **Full k-ary tree**
 - ✓ every node has exactly 0 or k children
- **Complete k-ary tree**
 - ✓ every level is entirely filled
 - ✓ except possibly the deepest, where all nodes are as far left as possible
- **Perfect k-ary tree**
 - ✓ every leaf has the same depth and the tree is full

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Binary trees

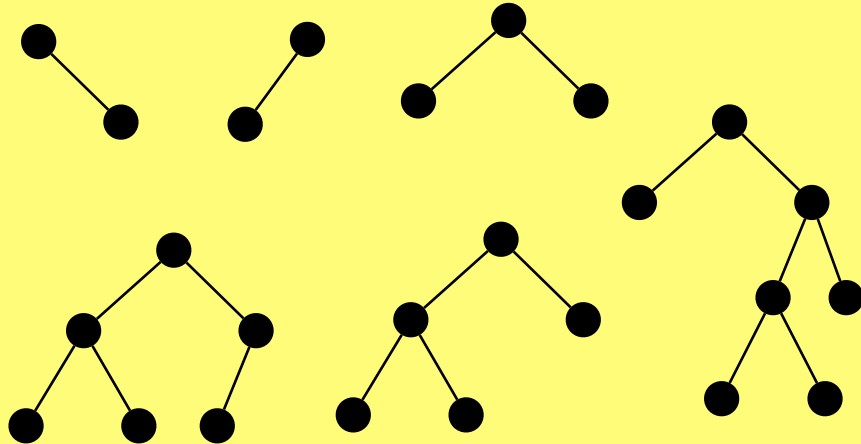
- **Definition**
 - ✓ a special case of a k-ary tree, where $k = 2$



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Practice

- Mark the following binary trees ($k=2$) as full/complete/perfect



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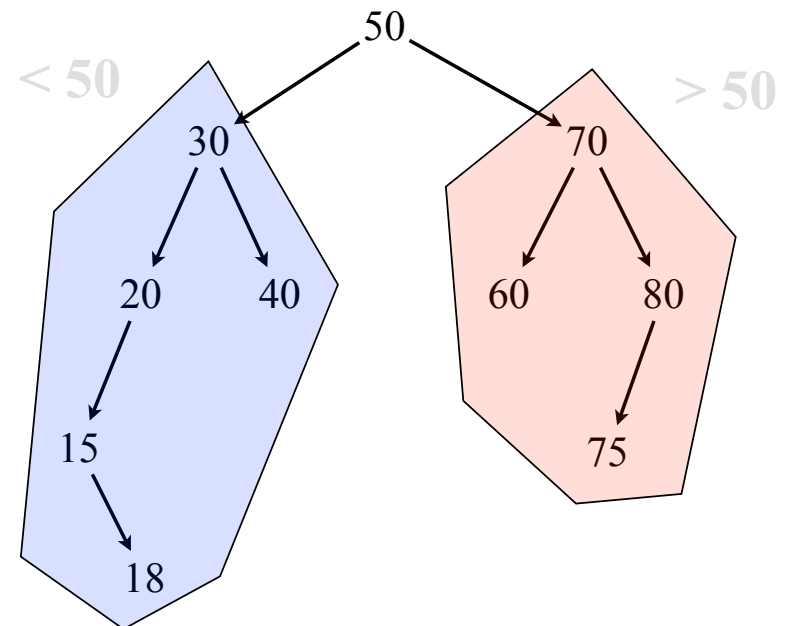
Binary search trees

Binary search tree

- A binary search tree (BST) is a **binary tree**
- A BST has **symmetric order**
 - each node x in a BST has a key denoted by $key(x)$
 - for all nodes y in the left subtree of x , $key(y) < key(x)$ **
 - for all nodes y in the right subtree of x , $key(y) > key(x)$ **

(**) assume that the keys of a BST are pairwise distinct

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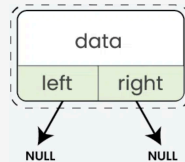


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Representing a node

```
template <typename T>
struct BSTNode {
    T key;
    BSTNode<T> *left, *right;

    BSTNode(const T& value) {
        data = value;
        left = right = nullptr;
    }
};
```



The implementation of a **binary tree node** requires a structure that can accommodate connections to two child nodes

Representing a binary search tree

```
template <typename T>
class BST {
private:
    struct Node {
        T data;
        Node *left, *right;
        Node(const T& value) {
            data = value;
            left = right = nullptr;
        }
    };

    Node *root;
    size_t size;

public:
    BST() : root(nullptr), size(0) {}
    ~BST() { clear(); }
    size_t getSize() const { return size; }
    bool empty() const { return size == 0; }

    void insert(const T& value);
    void remove(const T& value);
    bool contains(const T& value) const;
    void clear();
};
```

