

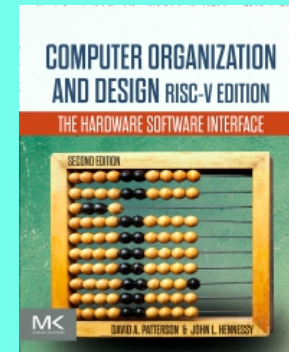
CSC 411

Computer Organization (Spring 2022)
Lecture 6: RISC-V Memory Organization

Prof. Marco Alvarez, University of Rhode Island

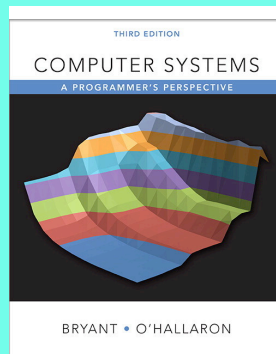
Disclaimer

Some of the following slides are adapted from:
Computer Organization and Design (Patterson and Hennessy)
The Hardware/Software Interface



Disclaimer

Some of the following slides are copied from:
Computer Systems (Bryant and O'Hallaron)
A Programmer's Perspective



Byte-Oriented Memory Organization



■ Programs refer to data by address

- Conceptually, envision it as a very large array of bytes
 - In reality, it's not, but can think of it that way
- An address is like an index into that array
 - and, a pointer variable stores an address

■ Note: system provides private address spaces to each "process"

- Think of a process as a program being executed
- So, a program can clobber its own data, but not that of others

Machine Words

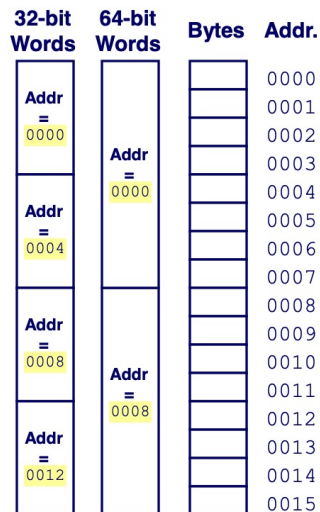
- Any given computer has a “Word Size”
 - Nominal size of integer-valued data
 - and of addresses
 - Until recently, most machines used 32 bits (4 bytes) as word size
 - Limits addresses to 4GB (2^{32} bytes)
 - Increasingly, machines have 64-bit word size
 - Potentially, could have 18 EB (exabytes) of addressable memory
 - That's 18.4×10^{18}
 - Machines still support multiple data formats
 - Fractions or multiples of word size
 - Always integral number of bytes

Example Data Representations

C Data Type	Typical 32-bit	Typical 64-bit	x86-64
char	1	1	1
short	2	2	2
int	4	4	4
long	4	8	8
float	4	4	4
double	8	8	8
pointer	4	8	8

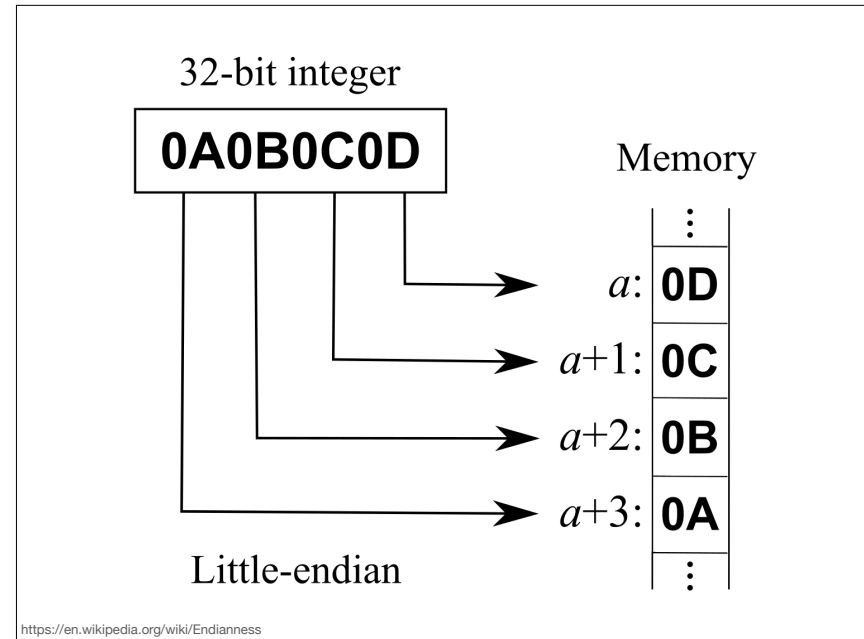
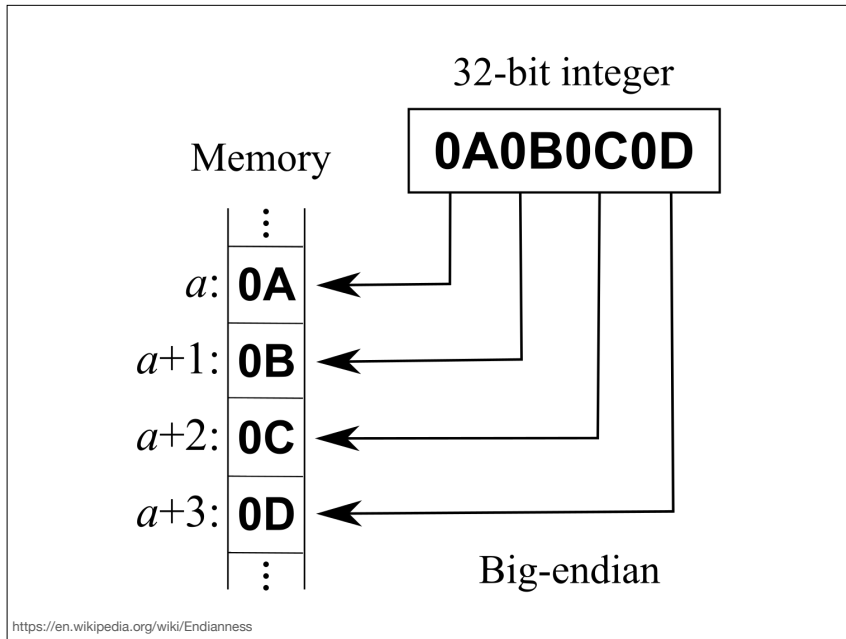
Word-Oriented Memory Organization

- Addresses Specify Byte Locations
 - Address of first byte in word
 - Addresses of successive words differ by 4 (32-bit) or 8 (64-bit)



Byte Ordering

- So, how are the bytes within a multi-byte word ordered in memory?
- Conventions
 - Big Endian: Sun (Oracle SPARC), PPC Mac, *Internet*
 - Least significant byte has highest address
 - Little Endian: *x86*, ARM processors running Android, iOS, and Linux
 - Least significant byte has lowest address



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Byte Ordering Example

- Example**
 - Variable x has 4-byte value of 0x01234567
 - Address given by &x is 0x100

Big Endian

0x100	0x101	0x102	0x103
01	23	45	67

Little Endian

0x100	0x101	0x102	0x103
67	45	23	01

Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition

Carnegie Mellon

Representing Integers

Decimal: 15213

Binary: 0011 1011 0110 1101

Hex: 3 B 6 D

int A = 15213;

long int C = 15213;

int B = -15213;

Two's complement representation

Increasing addresses ↓

Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition

Examining Data Representations

Code to Print Byte Representation of Data

- Casting pointer to unsigned char * allows treatment as a byte array

```
typedef unsigned char *pointer;

void show_bytes(pointer start, size_t len){
    size_t i;
    for (i = 0; i < len; i++){
        printf("%p\t0x%.2x\n", start+i, start[i]);
        printf("\n");
    }
}
```

Printf directives:

%p: Print pointer
%x: Print Hexadecimal

show_bytes Execution Example

```
int a = 15213;
printf("int a = 15213;\n");
show_bytes((pointer) &a, sizeof(int));
```

Result (Linux x86-64):

```
int a = 15213;
0x7fffb7f71dbc 6d
0x7fffb7f71dbd 3b
0x7fffb7f71dbe 00
0x7fffb7f71dbf 00
```

Representing Strings

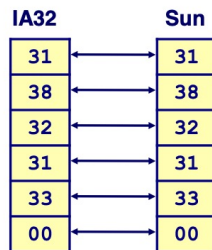
Strings in C

- Represented by array of characters
- Each character encoded in ASCII format
 - Standard 7-bit encoding of character set
 - Character "0" has code 0x30
 - Digit i has code $0x30+i$
- String should be null-terminated
 - Final character = 0

Compatibility

- Byte ordering not an issue

```
char S[6] = "18213";
```



ASCII control characters				ASCII printable characters				Extended ASCII characters			
00	NULL	(Null character)		32	space	64	@	96	.		
01	SOH	(Start of Header)		33	!	65	A	97	a		
02	STX	(Start of Text)		34	"	66	B	98	b		
03	ETX	(End of Text)		35	#	67	C	99	c		
04	EOT	(End of Trans.)		36	\$	68	D	100	d		
05	ENQ	(Enquiry)		37	%	69	E	101	e		
06	ACK	(Acknowledgement)		38	&	70	F	102	f		
07	BEL	(Bell)		39	'	71	G	103	g		
08	BS	(Backspace)		40	(72	H	104	h		
09	HT	(Horizontal Tab)		41)	73	I	105	i		
10	LF	(Line feed)		42	*	74	J	106	j		
11	VT	(Vertical Tab)		43	+	75	K	107	k		
12	FF	(Form feed)		44	,	76	L	108	l		
13	CR	(Carriage return)		45	-	77	M	109	m		
14	SO	(Shift Out)		46	.	78	N	110	n		
15	SI	(Shift In)		47	/	79	O	111	o		
16	DLE	(Data link escape)		48	0	80	P	112	p		
17	DC1	(Device control 1)		49	1	81	Q	113	q		
18	DC2	(Device control 2)		50	2	82	R	114	r		
19	DC3	(Device control 3)		51	3	83	S	115	s		
20	DC4	(Device control 4)		52	4	84	T	116	t		
21	NAK	(Negative acknowl.)		53	5	85	U	117	u		
22	SYN	(Synchronous idle)		54	6	86	V	118	v		
23	ETB	(End of trans. block)		55	7	87	W	119	w		
24	CAN	(Cancel)		56	8	88	X	120	x		
25	EM	(End of medium)		57	9	89	Y	121	y		
26	SUB	(Substitute)		58	:	90	Z	122	z		
27	ESC	(Escape)		59	;	91	[123	{		
28	FS	(File separator)		60	<	92	\	124			
29	GS	(Group separator)		61	=	93]	125	}		
30	RS	(Record separator)		62	>	94	^	126	~		
31	US	(Unit separator)		63	?	95	_				
127	DEL	(Delete)									
128	Ç			160	à			192	À		
129	ù			161	á			193	Á		
130	é			162	â			194	Â		
131	â			163	ã			195	Ã		
132	ä			164	ä			196	Ä		
133	å			165	Å			197	Å		
134	ä			166	ä			198	ä		
135	ç			167	ç			199	Ç		
136	ê			168	ê			200	Ê		
137	ë			169	ë			201	Ë		
138	è			170	è			202	È		
139	í			171	í			203	Í		
140	ì			172	ì			204	Ì		
141	î			173	î			205	Î		
142	Á			174	«			206	»		
143	À			175	»			207	«		
144	É			176				208			
145	Ê			177				209			
146	Ë			178				210			
147				179				211			
148				180				212			
149				181				213			
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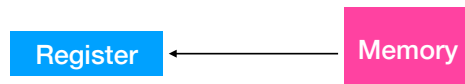
Memory operands

MIPS memory operands

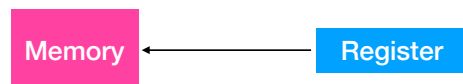
- Main memory used for composite data
 - arrays, structures, dynamic data
- To apply arithmetic operations:
 - **load** values from memory into registers
 - **store** result from register to memory
- Memory is byte addressed
 - each address identifies an 8-bit byte
 - address must be a multiple of 4 (32-bit words)

Fetching and storing data

- Values can be fetched from memory
 - using a **load** instruction



- Values can be stored in memory
 - using a **store** instruction



Load instruction

- Load word from memory

destination register source address

```
lw   $t0, 32($s3)
```

- destination can be any register
- source address uses a constant (offset) added to the register in brackets (base)

Store instruction

- Store word into memory

source register destination address
SW \$t0, 48(\$s3)

- source can be any register
- destination address uses a constant (offset) added to the register in brackets (base)

Example with memory operand

- C code

```
g = h + A[8];
```

- Compiled MIPS code

- index 8 requires offset of 32 (4 bytes)

```
# assume g in $s1, h in $s2  
# assume base address of A in $s3  
lw  $t0, 32($s3)    # load word  
add $s1, $s2, $t0
```

Example with memory operand

- C code

```
A[12] = h + A[8];
```

- Compiled MIPS code

- try yourself

```
# assume h in $s2  
# assume base address of A in $s3  
lw  $t0, 32($s3)    # load word  
add $t0, $s2, $t0  
sw  $t0, 48($s3)    # store word
```

Registers vs memory

- Operating on memory data requires **loads** and **stores**

- consider memory latency and additional instructions to be executed

- **Registers** are a fast read/write memory right on the CPU that can hold values

- Compiler must use registers for variables as much as possible

- only spill to memory for less frequently used variables
- **register optimization** is important!