

CSC 212 Data Structures & Algorithms

Fall 2022 | Jonathan Schrader

Merge Sort

Housekeeping

Assignment 3

• Due Friday, 11:59p

Review Project [MEC]

• Due Next Friday, October 28, 11:59pm



Divide & Conquer

Divide the problem into smaller subproblems

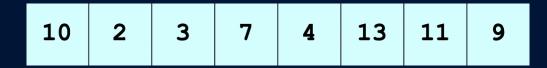
Conquer recursively

• ... each subproblem

Combine Solutions



Example



- ullet sorting with insertion sort is n^2
- we can divide the array into two halves and sort them separately



- each subproblem could be sorted in $pprox rac{n^2}{4}$
- * sorting both halves will require $pprox 2rac{n^2}{4}$ §
- · we need an additional operation to combine both solutions

Time "reduced" from
$$pprox n^2$$
 to $pprox rac{n^2}{2} + n$

Merge Sort

Divide the array into two halves

• just need to calculate the midpoint

Conquer Recursively each half

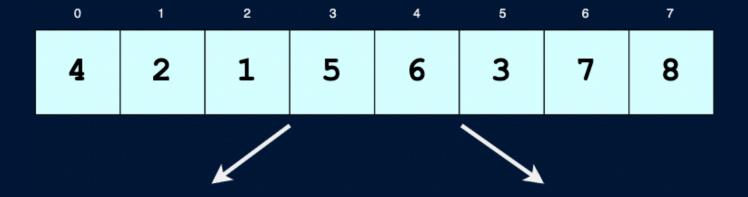
• call Merge Sort on each half (i.e. solve 2 smaller problems)

Merge Solutions

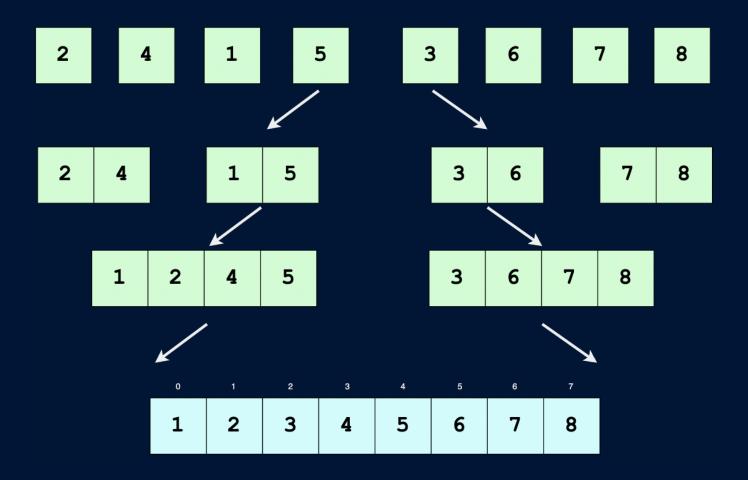
• after both calls are finished, proceed to *merge* the solutions



Divide...



...& Conquer





Merge Sort: pseudocode

```
if (hi <= lo)
    return;
int mid = lo + (hi - lo) / 2;
mergesort(A, lo, mid);
mergesort(A, mid + 1, hi);
merge(A, lo, mid, hi);</pre>
```



Merge Sort

```
void r_mergesort(int *A, int *aux, int lo,int hi) {
    //basecase(single element or empty list)
    if (hi <= lo)
        return;

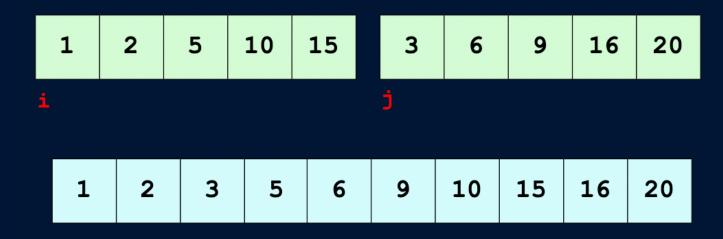
    //divide
    int mid = lo + (hi - lo) / 2;

    //recursively sort halves
    r_mergesort(A, aux, lo, mid);
    r_mergesort(A, aux, mid + 1, hi);

    //merge results
    merge(A, aux, lo, mid, hi);
}

void mergesort(int *A, int n) {
    int *aux = new int[n];
    r_mergesort(A, aux, 0, n - 1);
    delete[] aux;
}</pre>
```

Merging two sorted arrays



A secondary array is necessary

to guarantee a lineartime operation

Merge

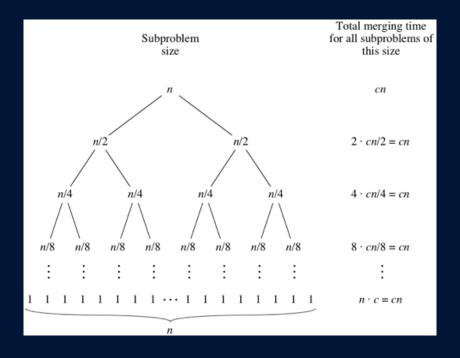
```
void merge (int *A, int *aux, int lo, int mid,int hi) {
    // copy array
    std::memcpy(aux + lo, A + lo, (hi - lo + 1 * sizeof(A)));
    // merge
    int i = lo, j = mid + 1;
    for (int k = lo; k <= hi; k++) {
        if (i > mid)
            A[k]=aux[j++];
        else if (j > hi)
            A[k] = aux[i++];
        else if(aux[j] < aux[i])
            A[k] = aux[j++];
        else
            A[k] = aux[i++];
    }
}</pre>
```

Analysis (recurrence)

```
void r mergesort(int *A, int *aux, int lo,int hi) {
                                                           void merge (int *A, int *aux, int lo, int mid,int hi) {
  //basecase(single element or empty list)
  if (hi <= lo)
                                                             // copy array
                                                             std::memcpy(aux + lo, A + lo, (hi - lo + 1 * sizeof(A)));
    return;
                                                             // merge
  //divide
                                                             int i = lo, j = mid + 1;
 int mid = lo + (hi - lo) / 2;
                                                             for (int k = lo; k <= hi; k++) {
  //recursively sort halves
                                                               if (i > mid)
 r mergesort(A, aux, lo, mid);
                                                                 A[k]=aux[j++];
 r mergesort(A, aux, mid + 1, hi);
                                                               else if (j > hi)
  //merge results
                                                                 A[k] = aux[i++];
 merge(A, aux, lo, mid, hi);
                                                               else if(aux[j] < aux[i])</pre>
                                                                 A[k] = aux[j++];
void mergesort(int *A, int n) {
  int *aux = new int[n];
                                                               else
  r mergesort(A, aux, 0, n - 1);
                                                                 A[k] = aux[i++];
  delete[] aux;
```

	Worst Case	Average Case	Best Case
Time Complexity	$O(n \ log \ n)$	$O(n \ log \ n)$	$O(n \ log \ n)$

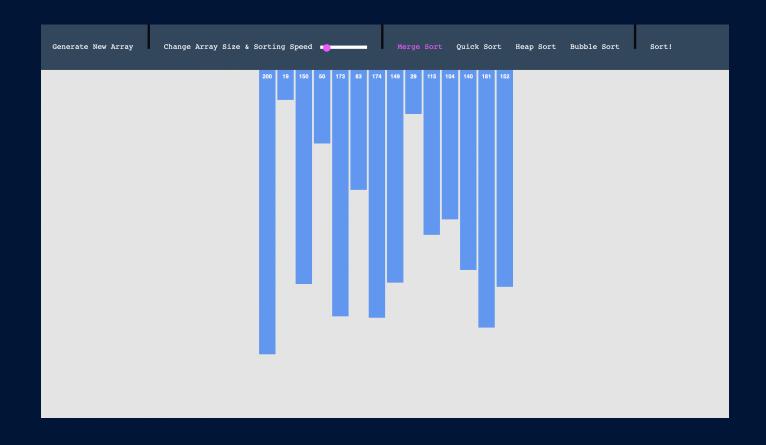
Recursion Tree (trace)



```
void mergesort(int *A, int n) {
  int *aux = new int[n];
  r_mergesort(A, aux, 0, n - 1);
  delete[] aux;
}

void r_mergesort(int *A, int *aux, int lo,int hi) {
  if (hi <= lo) return;
  int mid = lo + (hi - lo) / 2;
  r_mergesort(A, aux, lo, mid);
  r_mergesort(A, aux, mid + 1, hi);
  merge(A, aux, lo, mid, hi);
}</pre>
```

Sorting Visualizer



Comments on Merge Sort

Major disadvantage

- it is *not* in-place
- in-place algorithm exists but it is complex and inefficient

Improvements

- use insertion sort for small arrays
 - avoid overhead on small instances (~10 elements)
- stop if already sorted
 - avoids unnecessary merge
 - works well with partially sorted arrays



In-place Sorting



Example

Think about reversing an array or string

solution 1: use an additional array of equal size

what is the required extra memory?

solution 2: exchange first and last and work recursively on the inner part

- can do it iteratively as well
- what is the required extra memory?



In-place sorting

A sorting algorithm is *in-place* if it uses O(log n) extra memory

Are selection and insertion sorts *in-place*?

```
void selectionSort(int arr[], int n)
{
   int i, j, min_idx;

   // One by one move boundary of
   // unsorted subarray
   for (i = 0; i < n-1; i++) {

        // Find the minimum element in
        // unsorted array
        min_idx = i;
        for (j = i+1; j < n; j++)
        if (arr[j] < arr[min_idx])
            min_idx = j;

        // Swap the found minimum element
        // with the first element
        if (min_idx!=i)
            swap(&arr[min_idx], &arr[i]);
    }
}</pre>
```

```
void insertionSort(int arr[], int n)
{
    int i, key, j;

    for (i = 1; i < n; i++)
    {
        key = arr[i];
        j = i - 1;

        // Move elements of arr[0..i-1],
        // that are greater than key, to one
        // position ahead of their
        // current position
        while (j >= 0 && arr[j] > key)
        {
            arr[j + 1] = arr[j];
            j = j - 1;
        }
        arr[j + 1] = key;
    }
}
```

Stable Sorting

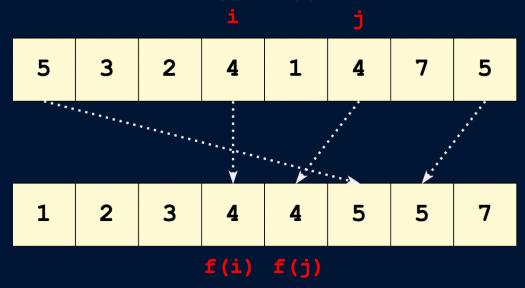


Stability

A sorting algorithm is *stable* if it preserves the order of equal elements

Consider sorting (in ascending order) a list A into a sorted list B. Let f(i) be the index of element A[i] in B. The sorting algorithm is stable if:

for any pair (i,j) such that A[i] = A[j] and i < j, then f(i) < f(j)





Stable?

DL 2273	Detroit	5:30 am	Departed
WN 6240	Chicago - MDW	5:55 am	Departed
AA 489	Philadelphia	6:00 am	Departed
DL 1263	Atlanta	6:00 am	Departed
UA 6208	Washington - IAD	6:00 am	Departed
WN 1138	Baltimore	6:05 am	Departed
AA 5202	Washington - DCA	6:14 am	Departed
B6 475	Orlando	6:15 am	Departed
UA 4894	New York/Newark	6:15 am	Departed
AA 1703	Charlotte	6:17 am	Departed
WN 28	Orlando	6:55 am	Departed
AA 3410	Chicago - ORD	7:02 am	Departed
WN 6235	Tampa	7:05 am	Departed
UA 3615	Chicago - ORD	7:30 am	Departed
AA 1735	Philadelphia	8:02 am	Departed
AA 632	Charlotte	8:07 am	At 9:45 am
WN 6247	Fort Lauderdale	8:30 am	Departed
WN 2640	Washington - DCA	8:45 am	Departed
WN 3420	Chicago - MDW	8:45 am	Departed
AA 4280	Washington - DCA	8:49 am	At 10:20 am
WN 846	Baltimore	9:20 am	Departed
DL 305	Detroit	10:40 am	On time
AA 774	Philadelphia	10:51 am	On time
AA 1981	Charlotte	11:01 am	On time
WN 3020	Baltimore	11:20 am	On time
AA 5524	Washington - DCA	11:46 am	At 2:35 pm
AC 7379	Toronto	11:50 am	On time
AA 5550	Charlotte	11:54 am	On time
DL 5090	Detroit	12:32 pm	On time
WN 6296	Baltimore	12:35 pm	On time
DL 2225	Atlanta	12:48 pm	On time
AA 4424	Washington - DCA	1:38 pm	On time

sort, then sort again

DL 1263	Atlanta	6:00 am	Departed
DL 2225	Atlanta	12:48 pm	On time
WN 1138	Baltimore	6:05 am	Departed
WN 846	Baltimore	9:20 am	Departed
WN 3020	Baltimore	11:20 am	On time
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UA 3615	Chicago - ORD	7:30 am	Departed
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DL 5090	Detroit	12:32 pm	On time
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UA 4894	New York/Newark	6:15 am	Departed
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WN 2640	Washington - DCA	8:45 am	Departed
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Stable: How about now?

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AA 5524	Washington - DCA	11:46 am	At 2:35 pm
AA 4424	Washington - DCA	1:38 pm	On time
UA 6208	Washington - IAD	6:00 am	Departed



Stability

Is selection sort stable?

- long distance swaps
- try: 5238456

Is insertion sort stable?

• equal items never pass each other (depends on correct implementation)



Sorting Algorithms

	Best-Case	Average-Case	Worst-Case	Stable	In-place
Selection Sort	$O(n^2)$	$O(n^2)$	$O(n^2)$	No	Yes
Insertion Sort	O(n)	$O(n^2)$	$O(n^2)$	Yes	Yes
Merge Sort	$O(n \ log \ n)$	$O(n \ log \ n)$	$O(n \ log \ n)$	Yes	No