

Link -> (Starts at 12:35pm.)

<https://www.designa.in/16043/shaw-coding-question-solution-sde1-september-2023-subarrays>

The logo for DE Shaw & Co features the company name in a blue, serif typeface. Above the text, there is a horizontal line that is broken in the middle by a diagonal line segment sloping upwards from left to right.

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Program Details:->

https://docs.google.com/document/d/1nFt_jEgy1X-UaKU2bu3OnC365ADCL-ls_cpXUMJrvx0/edit

Understanding : We are given an array B of size "N"; Divide the array into 4 continuous parts such that $g = \text{part1} - \text{part2} + \text{part3} - \text{part4}$ is maximized.

Any part is allowed to be empty as well ; but collectively they should cover the whole array

Example:->

$$G = p1 + p3 - (p2 + p4)$$

If you want to make G bigger what should you do ?

According to Tarab's law:-> To maximize G , we should maximize $p1$ and $p3$ and minimize $p2$ and $p4$

-> We have to find two non-intersecting subarrays whose sum is as minimum as possible. The second subarray should always contain the last element

If you find them ; then $p1$ and $p3$ is already fixed so you don't need to check for it

$P1$ should for sure contain the first element of the array or it should be empty

$P4$ should contain the last element of the array or it should be empty

=>[1 2 1 -5]

$$G = p1 + p3 - (p2+p4) = 4 - (-5) = 9.$$

Structure of p4 : [n,n]

[n-1;n]

[n-2;n]

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[1;n]

Or empty = 0

Structure if p2 : [i....j] $i \leq j$; $i \rightarrow 1$ to n ; $j \rightarrow 1$ to n

Subarray :-> Always analyze the condition or subarray at index "i" ;
and fix something to index "i"

Let's take an index "i"; find the p2 ending at index "i" ; then find
p4 ; do these for all "i"; min of all of them will be your answer

Algorithm.-> <https://ideone.com/9pn4SX>

C++ <https://ideone.com/KvyaDm>

Java. <https://ideone.com/tVp8MU>

Py <https://ideone.com/9mHaFD>