

## Lectures

23<sup>rd</sup> (Tuesday)26<sup>th</sup> (Friday)27<sup>th</sup> (Saturday morning)  
(10:00 - 11:30 am)(Dean academic email: shift  
3<sup>rd</sup> Dec Lecture  $\rightarrow$  27<sup>th</sup> Nov)

I. Lecture 16 recording : (Included in Lecture 17 recording)  
(Lecture 17 was revision of Lecture 16)  
plus few topics

III. EXAM: 30<sup>th</sup> November

IV. Syllabus : All topics from Linear transformation onwards  
(including linear transformation)

Problem from previous Lecture:

Q. Define an inner product on  $\mathbb{R}^2$  such that

$$\langle e_1, e_1 \rangle = 2, \quad \langle e_1, e_2 \rangle = -1, \quad \langle e_2, e_2 \rangle = 3$$

$$\langle -, - \rangle : \underset{u}{\mathbb{R}^2} \times \underset{v}{\mathbb{R}^2} \longrightarrow \underset{\text{scalar}}{\mathbb{R}}$$

$$\left\langle (x_1, x_2) \quad (y_1, y_2) \right\rangle \longrightarrow \underbrace{\hspace{2cm}}_{\text{expression}}$$

Exercise

## POLYNOMIALS AND MATRICES

Let  $K$  be a field.

Notation  $K$ , we used  $F$  before

Additional Notes for fields  
(shared before)

By a polynomial over  $K$ , we mean a formal expression

$$f(t) = a_n t^n + \dots + a_0$$

where  $a_i \in K$ ,  $i = 0, 1, \dots, n$

$\hookrightarrow$   $t$  is a "variable".

SUM and PRODUCT:

$$f(t) = a_n t^n + \dots + a_0;$$

$$a_i \in K$$

$$g(t) = b_m t^m + \dots + b_0,$$

$$b_j \in K$$

if  $n > m$ .

re-write  $g(t) = 0 t^n + \dots + b_m t^m + \dots + b_0,$

and then write the sum  $f + g$  as

$$(f+g)(t) = (a_n + b_n) t^n + \dots + (a_0 + b_0)$$

"Polynomial" where  $b_j = 0$  if  $j > m$ .

If  $c \in K$ , then

$$(cf)(t) = ca_n t^n + \dots + ca_0;$$

Hence,  $(cf)$  is also a polynomial.

Thus, polynomials form a vector space over  $K$ .

PRODUCT:

$$(fg)(t) = (a_n b_m) t^{n+m} + \dots + a_0 b_0$$

In general,

$$(fg)(t) = c_{n+m} t^{n+m} + \dots + \underline{c_0}$$

where

$$c_k = \sum_{i=0}^k a_i b_{k-i} = a_0 b_k + a_1 b_{k-1} + \dots + a_k b_0$$

In the space of polynomials, we have product structure too, forming a polynomial algebra structure.

Remark. Let  $f, g$  be <sup>non-zero</sup> polynomials with coefficients in  $K$ . Then

$$\deg(fg) = \deg f + \deg g$$

"Proof is easy".

NOTATION:

$K[t] :=$  set of all polynomials over  $K$ .

Theorem.

(i) Let  $f \in \mathbb{C}[t]$ , with  $\deg f \geq 1$ , then

$f$  has a root in  $\mathbb{C}$  } i.e.  $\exists \alpha \in \mathbb{C}$  s.t.  $f(\alpha) = 0$

How to prove this?

(ii)  $f \in \mathbb{C}[t]$  with  $\deg f = n \geq 1$  and leading coefficient is 1, then there exists complex numbers

$\alpha_1, \alpha_2, \dots, \alpha_n$  such that

$$f(t) = (t - \alpha_1) \cdots (t - \alpha_n).$$

Let  $\alpha_1, \alpha_2, \dots, \alpha_r$  be the distinct roots of the polynomial  $f \in \mathbb{C}[t]$ , then

$$f(t) = (t - \alpha_1)^{m_1} \cdots (t - \alpha_r)^{m_r}$$

with integers  $m_1, \dots, m_r > 0$ , and

$m_i$  is the multiplicity of root  $\alpha_i$ .

Discussion.

Let  $A$  be a square matrix with co-efficients in  $K$ . Let  $f \in K[t]$ , and write

$$f(t) = a_n t^n + \dots + a_0 ; \quad \text{with } a_i \in K.$$

We define

$$f(A) = a_n A^n + \dots + a_0 I.$$

$$f(T) = a_n T^n + \dots + a_0 I$$

Matrix

$$T: V \rightarrow V$$

$A$

Operator

$$I: V \rightarrow V$$

Example. Let  $f(t) = 3t^2 - 2t + 5$ . Let

$$A = \begin{pmatrix} 1 & -1 \\ 2 & 0 \end{pmatrix}.$$

Then,

$$f(A) = 3 \cdot \begin{pmatrix} 1 & -1 \\ 2 & 0 \end{pmatrix}^2 - 2 \begin{pmatrix} 1 & -1 \\ 2 & 0 \end{pmatrix} + 5 \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$$

$$f(A) = \begin{pmatrix} 0 & -1 \\ 2 & -1 \end{pmatrix}$$

Theorem. Let  $f, g \in K[t]$ . Let  $A$  be a square matrix with co-efficients in  $K$ . Then

$$(f+g)(A) = f(A) + g(A)$$

$$(fg)(A) = f(A)g(A).$$

$$\text{If } c \in K, \text{ then } (cf)(A) = cf(A).$$

Proof. Let  $f(t) = a_n t^n + \dots + a_0$

$$g(t) = b_m t^m + \dots + b_0$$

Then

$$(fg)(t) = c_{m+n} t^{m+n} + \dots + c_0$$

$$\text{where } c_k = \sum_{i=0}^k a_i b_{k-i}$$

By definition,

$$(fg)(A) = c_{m+n} A^{m+n} + \dots + c_0 I$$

On the other hand,

$$f(A) = a_n A^n + \dots + a_0 I$$

$$g(A) = b_m A^m + \dots + b_0 I$$

Hence

$$f(A)g(A) = \sum_{i=0}^n \sum_{j=0}^m a_i A^i b_j A^j$$

$$= \sum_{i=0}^n \sum_{j=0}^m a_i b_j A^{i+j}$$

$$= \sum_{k=0}^{m+n} c_k A^k.$$

Thus

$$f(A)g(A) = (fg)(A).$$

The proofs for other cases:  $(f+g)(A) = f(A) + g(A)$

$$\& \quad (cf)(A) = c(f(A))$$

are also easy.



Examples. Let  $f(t) = (t-1)(t+3)$

$$= t^2 + 2t - 3.$$

Then

$$f(A) = A^2 + 2A - 3I$$

$$= (A - I)(A + 3I)$$

$$(A^2 - IA + 3AI - 3I^2)$$

"

$$A^2 + 2A - 3I$$

Examples. Let  $\alpha_1, \alpha_2, \dots, \alpha_n$  be numbers. Let

$$f(t) = (t - \alpha_1) \cdots (t - \alpha_n).$$

Then

$$f(A) = (A - \alpha_1 I) \cdots (A - \alpha_n I).$$

Discussion.

$V$ : vector space over  $K$ .

$A: V \longrightarrow V$  be a linear map

Then  $A^2 = A \circ A = A \cdot A$ .

In general,  $A^n = A \circ \cdots \circ A$  ( $n$  times),  $A^0 = I$  identity

In general,

$$A^{m+n} = A^m A^n \quad \text{for all integers } m, n \geq 0.$$

$$A = \begin{pmatrix} 1 & 2 \\ 3 & 4 \end{pmatrix} \in M_2(\mathbb{R}) = V$$

$$f(A) = 0$$

choose  $f(t) = a_n t^n + \dots + a_0 = 0$   
 $a_n A^n + \dots + a_0 I = 0$

$$\dim_{\mathbb{R}} V = 4$$

Basis of  $V = \{e_{ij} \mid i, j = 1, 2\}$   
 $I \cdot a_0 + \dots + a_n A^n = 0$   
 $\{I, A, A^2, \dots, A^n\} \in M_2(\mathbb{R})$   
 lin. dependent  $n \geq 4$

Theorem. Let  $A$  be an  $n \times n$  matrix in a field  $K$ .

Then there exists a non-zero polynomial

$$f \in K[t] \text{ such that } f(A) = 0.$$

Proof.

$$V = M_n(K)$$

$$\dim_K V = n^2$$

Consider the sequence of vectors in  $V$

$$I, A, A^2, \dots, A^N \quad (\text{where } \underline{N} > n^2)$$

Linearly dependent

$$\Rightarrow a_N A^N + \dots + a_0 I = 0, \text{ set } f(t) = a_N t^n + \dots + a_0$$

choose  $f(t) = a_N t^n + \dots + a_0$

## THEOREM OF HAMILTON - CAYLEY.

Discussion.

$V$  : f.d. vector space over a field  $K$ .

$A : V \longrightarrow V$  be a linear map.

Assume that  $V$  has a basis consisting of eigenvectors of  $A$ , say  $\{v_1, \dots, v_n\}$

Let  $\{\lambda_1, \dots, \lambda_n\}$  be the corresponding eigenvalues.

Then the characteristic polynomial of  $A$  is

$$p(t) = (t - \lambda_1) \cdots (t - \lambda_n)$$

$$p(A) = (A - \lambda_1 I) \cdots (A - \lambda_n I).$$

Observe that

$$p(A)(v_i) = 0$$

$$Av_i = \lambda_i v_i$$

$$\left[ \begin{array}{l} p(A)v_i = (A - \lambda_1 I) \cdots (A - \lambda_n I) \underline{v_i} \\ \quad \quad \quad \neq 0 \\ \quad \quad \quad = 0 \end{array} \right.$$

Since  $Av_i = \lambda_i v_i$

$$p(A) = 0$$

In general, we cannot find basis consisting of eigen vectors.

## [Cayley - Hamilton Theorem]

**Theorem.** Let  $V$  be a f.d. v.s. over  $\mathbb{C}$  with  $\dim V \geq 1$ ,

and let  $A: V \longrightarrow V$  be a linear map. Let

$p$  be its characteristic polynomial. Then

$$p(A) = 0$$

**Proof.**

Note that matrix of  $A: V \longrightarrow V$  is upper triangular.

$$A = \begin{pmatrix} a_{11} & \cdots & a_{1n} \\ 0 & \ddots & \vdots \\ \vdots & & \ddots \\ 0 & & a_{nn} \end{pmatrix}$$

$Av_1 = a_{11}v_1$

Let  $v_1, v_2, \dots, v_n$  be those basis elements such that

$$Av_1 = a_{11}v_1$$

$$Av_2 = a_{12}v_1 + a_{22}v_2$$

$\vdots$

$$Av_n = a_{1n}v_1 + a_{2n}v_2 + \dots + a_{nn}v_n$$

$$\text{Set } V_1 = \{v_1\}$$

$$V_2 = \{v_1, v_2\}$$

$$\vdots$$

$$V_n = \{v_1, v_2, \dots, v_n\}$$

Thus,

$$Av_i = a_{ii}v_i + \text{an element of } V_{i-1}.$$

or

$$(A - a_{ii}I)v_i \in V_{i-1}$$

The characteristic polynomial of  $A$  is

$$p(t) = (t - a_{11}) \cdots (t - a_{nn})$$

hence

$$p(A) = (A - a_{11}I) \cdots (A - a_{nn}I)$$

We shall prove by induction that

$$(A - a_{11}I) \cdots (A - a_{ii}I)v = 0 \quad \text{for all } v \in V_i, \quad i=1, \dots, n.$$

When  $i=n$ , we get our theorem.

$$\text{Let } i=1, \text{ then } (A - a_{11}I)v_1 = Av_1 - a_{11}v_1 = 0 \quad \checkmark$$

Let  $i > 1$ , and assume that statement holds for  $i-1$ .

$$\left[ \begin{array}{l} (A - a_{11} I) \cdots (A - a_{i-1, i-1} I) v = 0 \\ \text{for all } v \in V_i, \quad i = 1, 2, \dots, i-1 \end{array} \right]$$

Let  $v \in V_i$ , then

$$v = v' + c v_i \quad \text{with } v' \in V_{i-1} \text{ \& some scalar } c \in K.$$

Note that  $(A - a_{ii} I) v' \in V_{i-1}$

$$\Downarrow \quad \left[ \because AV_{i-1} \subset V_{i-1} \right]$$

$$a_{ii} v' \in V_{i-1}$$

By induction,

$$(A - a_{11} I) \cdots (A - a_{i-1, i-1} I) (A - a_{ii} I) v' = 0$$

On the other hand,

$$(A - a_{ii} I) c v_i \in V_{i-1}, \quad \text{and hence}$$

by induction,

$$(A - a_{11}I) \cdots (A - a_{i-1,i-1}I) (A - a_{ii}I) v_i = 0.$$

Hence for all  $v$  in  $V_i$ , we have

$$(A - a_{11}I) \cdots (A - a_{ii}I) v = 0$$

$$\Rightarrow P(A)v = 0$$

$$\Rightarrow \boxed{P(A) = 0}$$

Corollary.  $A \in M_n(F)$  and  $p$  be its characteristic polynomial. Then

$$p(A) = \det(tI - A)^{(A)} = 0$$

$\parallel$   
0

Caution:   
Do not substitute directly  $A$  for  $t$

[Cayley-Hamilton Theorem]

Corollary. Let  $V$  be a finite dim. vec. space over the

field  $K$ , and let  $A: V \rightarrow V$  be a linear map. Let

$P$  be the characteristic polynomial of  $A$ . Then

$$P(A) = 0.$$



## Euclidean Algorithm.

**Theorem.** Let  $f, g$  be polynomials over the field  $K$ , assume that  $\deg g \geq 0$ . Then there exist polynomial  $q, r$  in  $K[t]$  such that

$$f(t) = q(t)g(t) + r(t),$$

$g(t) \overline{) f(t)}$   $q(t)$  and  $\deg r < \deg g$ .

The polynomials  $q, r$  are uniquely  
 $r(t)$  determined by these condition.



Corollary. Let  $f \in K[t]$ , a non-zero polynomial.

Let  $\alpha \in K$  be such that  $f(\alpha) = 0$ .

Then there exists a polynomial  $q(t) \in K[t]$  such

that

$$f(t) = (t - \alpha) q(t).$$

Proof.

define

$$f(t) = q(t)(t - \alpha) + r(t)$$

$$\deg r < \deg(t - \alpha)$$

$$\begin{array}{ccc} \uparrow & & \parallel \\ \text{constant,} & & 1 \end{array}$$

$$\hookrightarrow 0 = f(\alpha) = q(\alpha)(\alpha - \alpha) + r(\alpha) = 0$$

$$\Downarrow$$

$$r(\alpha) = 0$$

$$\Rightarrow r = 0.$$

Corollary. Suppose  $f \in K[t]$  has a root in  $K$ , then

every root is in  $K$

there exists elements  $\alpha_1, \dots, \alpha_n \in K$  and  $c \in K$  such

that

$$f(t) = c(t - \alpha_1) \cdots (t - \alpha_n).$$

Corollary. Let  $f$  be a polynomial of degree  $n$  in  $K[t]$ .

There are at most  $n$  roots of  $f$  in  $K$ .

Proof.

Suppose  $f$  has more than  $n$  roots, say  $m$ ,

$\alpha_1, \alpha_2, \dots, \alpha_m$  be distinct roots of  $f$  in  $K$ ,

$$\text{then } f(t) = (t - \alpha_1)(t - \alpha_2) \cdots (t - \alpha_m)g(t)$$

for some polynomial  $g(t)$

$$\left( \text{Contradiction } \deg f > m \right)$$

**Ideal.** By an ideal of  $K[t]$ , we mean a

subset  $I$  of  $K[t]$  satisfying the following conditions:

- $0 \in K[t]$  is in  $I$
- $f, g \in I \Rightarrow f + g \in I$
- $f \in I$  and  $g \in K[t]$ , then  $f \cdot g \in I$

$\langle n \rangle$

$\mathbb{Z}$

Theorem. Let  $I$  be an ideal of  $K[t]$ . Then there exists a polynomial  $g$  which is a generator of  $I$ .

$$I = \langle g \rangle$$

↑  
Single polynomial

Theorem.  $I = \langle f_1, f_2 \rangle \subseteq K[t]$ . Then

$$I = \langle g \rangle, \quad g = \text{gcd}(f_1, f_2)$$

greatest common divisor  
of  $f_1$  and  $f_2$

$\langle 2, 5 \rangle$   
||  
 $\langle \text{gcd}(2, 5) \rangle$

- $K[t]$   
↑  
field
- Unique Factorization Domain  
PID

$$x^2 + 1 \in \mathbb{R}[x] \\ = a(x)b(x)$$

- $f \in K[t]$  is irreducible (over  $K$ ) if its degree  $\geq 1$  and if given a factorization

$$f = h_1 h_2 \text{ with } h_1, h_2 \in K[t], \text{ then}$$

$\deg h_1$  or  $\deg h_2 = 0$  (i.e. one of them is a constant).

M.Sc

$$\mathbb{Z} \quad m, n \quad 0 \quad p \quad a \quad b$$

Lemma. Let  $f \in K[t]$  be irreducible polynomial.

$h_1, h_2 \in K[t]$  and assume  $f$  divides  $h_1 h_2$ .

Then  $f \mid h_1$  or  $f \mid h_2$ .

$$p \mid ab \text{ then } p \mid a \text{ or } p \mid b$$

Corollary.  $f \in K[t]$ ,  $\deg f > 1$ , then

$$f = c \cdot f_1 \cdot f_2 \cdots f_s$$

where  $f_1, f_2, \dots, f_s$  are irreducible polynomials

with leading coefficients 1.

Corollary.

$f \in \mathbb{C}[t]$ ,  $\deg f > 1$ , then

$$f = c \cdot (t - \alpha_1) \cdots (t - \alpha_n).$$

where  $\alpha_i \in \mathbb{C}$  &  $c \in \mathbb{C}$ .

**Theorem.** Let  $f(t) \in K[t]$  and suppose  
that  $f = f_1 \cdot f_2$  and  $\text{GCD}(f_1, f_2) = 1$ .

Let  $T : V \rightarrow V$  be a linear map. Assume that

$$f(T) = 0. \text{ Let}$$

$$W_1 = \text{kernel of } f_1(T)$$

$$W_2 = \text{kernel of } f_2(T).$$

Then  $V = W_1 \oplus W_2.$

**Theorem.**  $V$  : v.s. over  $\mathbb{C}$ ,  $T : V \rightarrow V$

$p(t)$  be a polynomial such that  $p(T) = 0$ ,

Let 
$$p(t) = (t - \alpha_1)^{m_1} \cdots (t - \alpha_r)^{m_r}.$$

Let 
$$W_i = \text{kernel of } (A - \alpha_i I)^{m_i}. \text{ Then}$$

$$V = W_1 \oplus \cdots \oplus W_r$$

## The minimal polynomial.

Recall:  $V$  f.d. v.s. over the field.

$T: V \rightarrow V$  linear operator.

$p(t) = \det(tI - A)$  characteristic equation

By Cayley-Hamilton theorem,

$$p(A) = 0 \quad (\Leftrightarrow p(T) = 0)$$

$$I = \left\{ f \in K[t] \text{ s.t. } f(A) = 0 \right\} \\ (\Leftrightarrow f(T) = 0)$$

collection of all polynomials s.t.  $f(A) = 0$

Algebra

$I$  is an ideal in  $K[t]$

$\uparrow$   
PID

$$I = \langle m \rangle \quad m(T) = 0 \quad \Leftrightarrow \underline{m(A) = 0}$$

unique generator (monic polynomial)

or  
The minimal Polynomial for  $T$

~~The monic polyno~~

The (monic) minimal polynomial is

- monic  $m(t) = \underset{\uparrow}{1} \cdot t^n + a_{n-1} t^{n-1} + \dots + a_0$
- $m(T) = 0$   $m(A) = 0$
- $\deg m \leq \deg f$ , where  $f$  is such that

$$f(T) = 0$$

"minimal" in the sense of degree.

$\left\{ \begin{array}{l} \text{Theorem. The zeros/roots of characteristic polynomial} \\ \text{and the minimal polynomial are same.} \end{array} \right.$