
CS:1010 DISCRETE STRUCTURES

PRACTICE QUESTIONS LECTURE 2

Instructions

- Try these questions before class. Do not submit!
- (1) Determine the truth value of each of these statements if the domain consists of all integers:
- (a) $\forall n(n + 1 > n)$ True since adding 1 to an integer always makes it greater.
 - (b) $\exists n(2n = 3n)$ True since it is true for $n = 0$.
 - (c) $\exists n(n = -n)$ True since $n = 0$ this is true
 - (d) $\forall n(3n \leq 4n)$ False since if n is a negative number this is not true.
- (2) Let $P(x)$ be the statement $x = x^2$. Domain is \mathbb{Z} , the set of integers. What are the truth values?
- (a) $P(0)$
True
 - (b) $P(1)$
True
 - (c) $P(2)$
False since $2 \neq 2^2 = 4$.
 - (d) $\forall x P(x)$
False since not true for $x = 2$
- (3) Suppose the domain of $P(x)$ is $\{1, 2, 3, 4\}$ then express $\exists x P(x)$ without a quantifier.
 $P(1) \vee P(2) \vee P(3) \vee P(4)$.
- (4) Express each of these statements using logical operators, predicates and quantifiers.
Domain : All propositions. $T(x)$: x is a tautology. $C(x)$: x is a contradiction.
- (a) Some propositions are tautologies.
 $\exists x T(x)$
 - (b) The negation of a contradiction is a tautology.
Implicit is “all/every”. Negation of every contradiction is always a tautology.
That is, if x is a contradiction then $\neg x$ is a tautology.
 $\forall x (C(x) \rightarrow T(\neg x))$.

- (5) What are the truth values of these statements?
- (a) $\exists!xP(x) \rightarrow \exists xP(x)$
True since if there is a unique x satisfying $P(x)$ then there is an x satisfying $P(x)$.
- (b) $\forall xP(x) \rightarrow \exists!xP(x)$
Unless the domain has one item in it, this will not hold.
- (6) Let $S(x)$: x is a student, $F(x)$: x is a faculty member and $A(x, y)$: x has asked y a question. Domain: all people associated with our school. Use quantifiers to express each of these statements.
- (a) Divya has asked Prof. Gupta a question.
 $A(\text{Divya}, \text{Prof. Gupta})$
- (b) Every student has asked Prof. Gupta a question.
 $\forall x(S(x) \rightarrow A(x, \text{Prof. Gupta}))$
- (c) Some student has not asked any faculty member a question.
 $\exists x(S(x) \wedge \forall y(F(y) \rightarrow \neg A(x, y)))$ (there are other ways of writing it too!)
- (7) Express each of the statements using predicated, quantifiers, logical connectives and mathematical operators.
- (a) Every positive real number has exactly two square roots.
Exactly two objects that meet the same condition: we need two existential quantified variables.
 $\forall x > 0 \exists a \exists b (a \neq b \wedge \forall c (c^2 = x \leftrightarrow (c = a \vee c = b)))$.
- (b) A negative real number does not have a square root that is a real number.
 $\forall x ((x < 0) \rightarrow \neg \exists y (x = y^2))$. (Domain is set of all real numbers.)
- (8) Negate the statement such that negation immediately precedes predicates:
 $\forall x \exists y (P(x, y) \rightarrow Q(x, y))$.

$$\begin{aligned} \neg \forall x \exists y (P(x, y) \rightarrow Q(x, y)) &\equiv \exists x \neg \exists y (P(x, y) \rightarrow Q(x, y)) \\ &\equiv \exists x \forall y \neg (P(x, y) \rightarrow Q(x, y)) \\ &\equiv \exists x \forall y (P(x, y) \wedge \neg Q(x, y)). \end{aligned}$$

- (9) Is this argument correct: “ Every computer science student takes discrete mathematics. Natasha is taking discrete mathematics. Therefore, Natasha is a computer science student.”

Invalid. Applying universal instantiation it affirmed the conclusion.

- (10) Explain the rules of inference used in each step. Each of the 93 students in this class own a laptop. Everyone who owns a laptop can use a PDF viewer.

Therefore, Arun, a student in this class can use a PDF viewer.

$C(x)$: x is in this class. $P(x)$: x owns a laptop. $Q(x)$: x can use a PDF viewer.

- (a) $\forall x(C(x) \rightarrow P(x))$ – Hypothesis
- (b) $C(Arun) \rightarrow P(Arun)$ – Universal Instantiation.
- (c) $C(Arun)$ – Hypothesis
- (d) $P(Arun)$ – Modus Ponens
- (e) $\forall x(P(x) \rightarrow Q(x))$ – Hypothesis
- (f) $P(Arun) \rightarrow Q(Arun)$ – Universal instantiation.
- (g) $Q(Arun)$ – Modus Ponens