# CS 1010 Discrete Structures Lecture 9: Modular Arithmetic Contd. & Counting Techniques

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January 21, 2021

#### Greatest Common Divisor

- Let  $a, b \in \mathbb{Z}$  s.t. at least one of them is not zero.
- The largest integer d such that  $d \mid a$  and  $d \mid b$  is called the greatest common divisor, gcd(a, b). Does it exist? The set of common divisors of these integers is nonempty and finite.
- One algorithm: find all the positive common divisors and take the largest divisor.
- gcd(18,9): Positive common divisors of 18 and 9 are 1,3,9 so the greatest is 9.
- gcd(4,9): Positive common divisor is only 1. Such integers are said to be relatively prime/co-prime, i.e. gcd(a,b) is 1.

#### Greatest Common Divisor

- The integers  $a_1, a_2, \ldots, a_n$  are pairwise relatively prime or co-prime if  $gcd(a_i, a_j) = 1$  whenever  $1 \le i < j \le n$ .
- Is 10, 17, 21 pairwise co-prime? Yes, since  $\gcd(10, 17) = \gcd(10, 21) = \gcd(17, 21) = 1$ .
- What about 10, 19, 24? No, since gcd(10, 24) = 2 > 1 is not relatively prime.
- Another algo: Find prime factorizations of a and b.

$$a = p_1^{a_1} p_2^{a_2} \cdots p_n^{a_n}, b = p_1^{b_1} p_2^{b_2} \cdots p_n^{b_n},$$

where each  $a_i, b_i \in \mathbb{N}$  and could be zeroes since we are including same primes on both sides.

$$gcd(a,b) = p_1^{min(a_1,b_1)} p_2^{min(a_2,b_2)} \cdots p_n^{min(a_n,b_n)}.$$

#### Least Common Multiple

- Let  $a, b \in \mathbb{N}$ , the lcm(a, b) is the smallest positive integer that is divisible by both a and b.
- Do they always exist? Yes because ab is a common multiple and every nonempty set of positive integer has a least element (WOP).

$$lcm(a,b) = p_1^{max(a_1,b_1)} p_2^{max(a_2,b_2)} \cdots p_n^{max(a_n,b_n)}$$

- Practice question: Let  $a, b \in \mathbb{N}$ . Then  $ab = \gcd(a, b) \cdot lcm(a, b)$ .

#### Euclidean Algorithm

- Finding prime factorizations are time consuming, that is why we explore the Euclidean Algorithm.
- The idea behind the algorithm: Say we have to find gcd(91,287).

$$287 = 91 \cdot 3 + 14$$
 If  $a \mid 91$  and  $a \mid 287$  then  $a \mid (287 - 91 \cdot 3 = 14)$  
$$gcd(287, 91) = gcd(91, 14).$$
 
$$91 = 14 \cdot 6 + 7,$$
 Same argument as above  $gcd(91, 14) = gcd(14, 7).$  
$$14 = 7 \cdot 2.$$
 
$$7 \mid 14 \Rightarrow gcd(14, 7) = 7.$$

### Euclidean Algorithm

- I.e. Use successive divisions to reduce the problem of finding gcd to the same problem with smaller integers, until one of the integers is zero.
- Lemma : Let a = bq + r, where a, b, q, r are integers. Then gcd(a, b) = gcd(b, r).
- Proof Idea: If we can show that the common divisors of a and b are the same as the common divisors of b and r, then we are done because both pairs must have the same greatest common divisor.
- You need to show both directions: If there is a common divisor d that divides a and b then it divides b and r. Also, if there is a common divisor that divides b and r then it divides both a and b as well.

#### Euclidean Algorithm

```
procedure gcd(a, b): positive integers)

x := a

y := b

while y \neq 0

r := x \mod y

x := y

y := r

return x\{gcd(a, b) \text{ is } x\}
```

We analysed the algorithm and saw that the number of divisions required where  $a \ge b$  is  $\mathcal{O}(\log b)$ .

#### Bézout's Theorem

- If  $a, b \in \mathbb{N}$ , then there exist integers s, t such that gcd(a, b) = sa + tb, i.e. an integer linear combination of a and b. The equation is called Bézout's Identity, and s, t are called Bézout coefficients.
- We do not give the proof here but note that we can obtain these coefficients by going backwards through the divisions of the Euclidean algorithm – i.e. a forward and backward pass of the algorithm giving rise to the extended Euclidean algorithm.

#### Bézout's Theorem

- Express gcd(252, 198) = 18 as a l.c. of 252 and 198.

$$252 = 1 \cdot 198 + 54$$

$$198 = 3 \cdot 54 + 36$$

$$54 = 1 \cdot 36 + 18$$

$$36 = 2 \cdot 18$$

- We do the backwards pass:
  - ▶  $18 = 54 1 \cdot 36$
  - ▶ Second division says that  $36 = 198 3 \cdot 54$
  - ► Substituting for 36 we get  $18 = 54 1 \cdot (198 3 \cdot 54) = 4 \cdot 54 1 \cdot 198$
  - First division says that  $54 = 252 1 \cdot 198$ .
  - ▶ Substituting for 54, we get  $18 = 4 \cdot 252 5 \cdot 198$ .

#### Proof of Euclid's Lemma

- We can now give the proof of Euclid's Lemma, in fact we prove something more general.
- If a, b, c are positive integers s.t. gcd(a, b) = 1 and  $a \mid bc$ , then  $a \mid c$ .
- Because gcd(a,b)=1 (which is definitely the case of primes), by Bézout's theorem we know there are integers s,t s.t. sa+tb=1.
- Multiplying with c on both sides, sac + tbc = c.
- a divides LHS and therefore  $a \mid c$ .
- We used Euclid's Lemma to show the uniqueness part of the FTA.

#### Canceling in a Congruence

- Let m be a positive integer and let a, b, c be integers. If  $ac \equiv bc \mod m$  and gcd(c, m) = 1, then  $a \equiv b \mod m$ .
- $m \mid (ac bc) = c(a b)$ .
- Since gcd(c, m) = 1 we have  $m \mid (a b)$ .
- Topics we will not cover include: how to solve congruences, the Chinese Remainder Theorem. Very relevant in cryptography.
- Fermat's Little Theorem : practice question.

# Counting

- Combinatorics study of arrangement of objects, as old as 17th century.
- Counting problems arise in computer science and mathematics including in analysis of algorithms.
- How many leaves in a tree, minimal colorings of a graph, trees with a given set of vertices, five-card hands in a deck of fifty-two, stable marriages given boy's and girl's preferences, probabilities, and so on.

# Counting Principles

#### - THE PRODUCT RULE:

- ► Let a procedure be such that it can be broken down into a sequence of two tasks.
- ► If there are  $n_1$  ways to do the first task and for each of these ways of doing the first task, there are  $n_2$  ways to do the second task,
- ▶ then there are  $n_1 n_2$  ways to do the procedure.
- Example: The chairs of an auditorium are to be labeled with an uppercase letter followed by a positive integer ≤ 100. What

is the largest number of chairs that can be labeled differently?

- ► Two tasks : assigning the seat one of the 26 uppercase English letters,
- ▶ and then assigning to it one of the 100 possible integers.
- ► The product rule tells us there are  $26 \cdot 100 = 2600$  different ways that a chair can be labeled.

#### More Examples

- Extend it to m tasks:A procedure is carried out by performing the tasks  $T_1, T_2, \cdots, T_m$  in sequence. If each task  $T_i, i = 1, 2, \cdots, n$ , can be done in  $n_i$  ways, regardless of how the previous tasks were done, then there are  $n_1 \cdot n_2 \cdots n_m$  ways to carry out the procedure.
- How many different bit strings of length seven are there
  - ► Each of the seven bits can be chosen in two ways, because each bit is either 0 or 1.
  - ▶ Product rule shows there are a total of  $2^7 = 128$  different bit strings of length seven.
- Counting Functions How many functions are there from a set with *m* elements to a set with *n* elements?
  - ▶ A choice of one of the n elements in the codomain for each of the m elements in the domain means by the product rule there are  $n \cdot n \cdots n = n^m$  functions.

#### More Examples

- Counting One-to-One Functions How many one-to-one functions are there from a set with m elements to one with n elements?
  - ▶ When m > n there are no one-to-one functions. Else, suppose the elements in the domain are  $a_1, a_2, ..., a_m$ .
  - ▶ There are n ways to choose the value of the function at  $a_1$ .
  - ▶ But the value at  $a_2$  can be picked in n-1 ways. The value of the function at  $a_k$  can be chosen in n-k+1 ways.
  - ▶ By the product rule, there are  $n(n-1)(n-2)\cdots(n-m+1)$  one-to-one functions.
- Counting Subsets of a Finite Set
  - ▶ One-to-one correspondence between subsets of S and bit strings of length |S|.
  - ▶ By the product rule, there are  $2^{|S|}$  bit strings of length |S|. Hence,  $|\mathcal{P}(S)| = 2^{|S|}$ .

# Counting Principles

- THE SUM RULE: If a task can be done either in one of  $n_1$  ways or in one of  $n_2$  ways, where none of the set of  $n_1$  ways is the same as any of the set of  $n_2$  ways, then there are  $n_1 + n_2$  ways to do the task.
- Either a member of the mathematics faculty or a student who is a mathematics major is chosen as a representative to a university committee.
- How many different choices are there for this representative if there are 37 members of the mathematics faculty and 83 mathematics majors and no one is both a faculty member and a student?
- By the sum rule it follows that there are 37 + 83 = 120 possible ways to pick this representative.

#### Sum Rule

- Extending it to m tasks such that a task can be done in one of  $n_1$  ways, in one of  $n_2$  ways,..., or in one of  $n_m$  ways,where none of the set of  $n_i$  ways of doing the task is the same as any of the set of  $n_j$  ways,for all pairs i and j with  $1 \le i < j \le m$ . Then the number of ways to do the task is  $n_1 + n_2 + \cdots + n_m$ . For example:

```
\begin{array}{l} k := 0 \\ \text{for } i_1 := 1 \text{ to } n_1 \\ k := k + 1 \\ \text{for } i_2 := 1 \text{ to } n_2 \\ k := k + 1 \\ & \cdot \\ & \cdot \\ \text{for } i_m := 1 \text{ to } n_m \\ k := k + 1 \end{array}
```

#### Product Rule

But with nested loops it is the product rule,  $n_1 \cdot n_2 \cdots n_m$ .

```
k := 0

for i_1 := 1 to n_1

for i_2 := 1 to n_2

.

.

for i_m := 1 to n_m

k := k + 1
```

#### More Complex Counting Rules

- Sum rule can be expressed in terms of pairwise disjoint finite sets,  $A_1, A_2, \ldots, A_m$ 

$$|A_1 \cup A_2 \cup \cdots A_m| = |A_1| + |A_2| + \cdots + |A_m|$$
 when  $A_i \cap A_j = \emptyset \forall i, j$ .

- More complicated when sets have elements in common.
- Principle of inclusion-exclusion/THE SUBTRACTION RULE: If a task can be done in either  $n_1$  ways or  $n_2$  ways, then the number of ways to do the task is  $n_1 + n_2$  minus the number of ways to do the task that are common to the two different ways.
- Counting the number of elements in the union of two sets:  $|A_1 \cup A_2| = |A_1| + |A_2| |A_1 \cap A_2|$

#### Examples

- How many bit strings of length 8 that either start with a 1 bit or end with the two bits 00?
- We can construct a bit string of length 8 that starts with a 1 in  $2^7=128$  ways.
- Ends with the two bits 00 can be done in  $2^6 = 64$  ways
- Some of the ways to construct a bit string starting with a 1 are the same as the ways to construct a bit string that ends with 00.
- There are  $2^5=32$  ways to construct such a string that starts with 1 and ends with 00. (How?)
- So the total number of ways : 128 + 64 32 = 160.

#### More Complex Counting Rules

- THE DIVISION RULE: There are n/d ways to do a task if it can be done using a procedure that can be carried out in n ways, and for every way w, exactly d of the n ways correspond to way w.
- I.e. helps you ignore unimportant differences when you are counting things. You can count distinct objects, and then use the division rule to merge the ones that are same.
- Seating at a Round Table. In how many ways can King Arthur seat *n* knights at his round table?
- Two seatings are considered equivalent if one can be obtained from the other by rotation. For eg: if Arthur has only 4 knights, then there are 6 possibilities.
- We use square brackets in clockwise order, starting at an arbitrary point. For eg: [1234] and [4123] are same.

#### Division Rule

- We show using division rule that King Arthur can seat n knights at his round table in (n-1)! ways.
- Let A be the set of ordering of n knights in a line. Let B be the set of orderings of n knights in a ring.
- Define  $f: A \mapsto B$  by  $f(x_1, x_2, \dots, x_n) = [x_1 x_2 \dots x_n]$ , the clockwise arrangement of  $x_1, x_2, \dots, x_n$ .
- What is  $f^{-1}([x_1x_2...x_n])$ ? =  $\{(x_1, x_2, ..., x_n), (x_n, x_1, ..., x_{n-1}), (x_{n-1}, x_n, ..., x_{n-2}) \cdots, (x_2, x_3, ..., x_1)\}$ .
- There are *n* tuples in that list.
- By division rule |A| = n|B|. |A| = n! and so |B| = (n-1)!

#### Pigeonhole Principle

- Suppose that a flock of 20 pigeons flies into a set of 19 pigeonholes to roost. At least one of these 19 pigeonholes must have at least two pigeons in it.
- Why? If each pigeonhole had at most one pigeon in it, at most 19 pigeons, one per hole, could be accommodated.
- This illustrates the pigeonhole principle: If k is a positive integer and k+1 or more objects are placed into k boxes, then there is at least one box containing two or more of the objects.
  - ► Use a proof by contraposition.
  - ► Suppose that none of the *k* boxes contains more than one object.
  - ▶ Then the total number of objects would be at most k. This is a contradiction, because there are at least k + 1 objects.
- Also called the Dirichlet drawer principle, but he was not the first person to use this in his work.

# Simple Applications

- A function f from a set with k+1 or more elements to a set with k elements is not one-to-one.
- Among any group of 367 people, there must be at least two with the same birthday, because there are only 366 possible birthdays.
- In any group of 27 English words, there must be at least two that begin with the same letter, because there are 26 letters in the English alphabet.

# Not So Simple Applications

- In every set of 1000 integers, there are two integers x and y such that 573 | (x y).
- Looks very hard! Our old friend induction seems to be useless here.
- To apply the Pigeonhole Principle, we must identify two things: pigeons and holes. Also, we must have more pigeons than holes. What about 1000 pigeons and 573 holes? Consider numbering the holes 0, 1, ..., 572 and putting a hole *n* all integers congruent modulo 573.

### Not So Simple Applications

- Let S be a set of 1000 integers and let  $M = \{0, 1, \dots 572\}$  and  $f(n) = n \mod 573$ .
- Since  $|S| \ge |M|$  we have x and y in S s.t. f(x) = f(y). Therefore,  $x \mod 573 = y \mod 573$  and  $573 \mid (x y)$ .
- Nothing special about 1000 and 573, we only need n > m then in every set of n integers there are two integers x and y such that  $m \mid (x y)$ .

#### Not So Simple Applications

- S.T. any given set containing 10 distinct positive numbers
   100, there exists two disjoint subsets sum to the same quantity.
- The numbers all vary between 1 and 99. Therefore the maximum sum of any 10 chosen numbers is  $90 + 91 + 92 + \cdots + 99 = 945$ .
- So the sums vary from 1 and 945.
- The number of different subsets of the 10 numbers is  $2^{10} 1$  (excluding the null set) = 1023.
- We have 1023 pigeons and 945 holes.
- Using the pigeonhole principle, we can argue that two different subsets map to the same sum. If these subsets have a common number or numbers, we can always remove the common numbers to produce disjoint subsets that sum to the same

#### Generalized Pigeon Hole Principle

- If N objects are placed into k boxes, then there is at least one box containing at least  $\lceil N/k \rceil$  objects.
  - ► Proof by contraposition.
  - ▶ Suppose that none of the boxes contains more than  $\lceil N/k \rceil 1$  objects.
  - ► Then the total number of objects is at most,

$$k(\lceil N/k \rceil - 1) < k((N/k + 1) - 1) = N,$$

where the inequality  $\lceil N/k \rceil < (N/k) + 1$  has been used. This is a contradiction since there are a total of N objects.

### Applications

- Typically we ask for the minimum number of objects such that at least r of these objects must be in one of k boxes. When we have N objects, the generalized pigeonhole principle tells us there must be at least r objects in one of the boxes as long as  $\lceil N/k \rceil \geq r$ .
- Among 100 people there are at least  $\lceil 100/12 \rceil = 9$  who were born in the same month.
- What is the minimum number of students required in a discrete mathematics class to be sure that at least six will receive the same grade, if there are five possible grades, A, B, C, D and  $F? \lceil N/5 \rceil = 6$ . The smallest such integer is  $N = 5 \cdot 5 + 1 = 26$ .

# Elegant Applications

- A clever application shows the existence of an increasing or a decreasing subsequence of a certain length in a sequence of distinct integers.
- Suppose that  $a_1, a_2, \ldots, a_N$  is a sequence of real numbers. A subsequence of this sequence is a sequence of the form  $a_{i_1}, a_{i_2}, \ldots, a_{i_m}$ , where  $1 \leq i_1 < i_2 < \cdots < i_m \leq N$ .
- A sequence is called **strictly increasing** if each term is larger than the one that precedes it, and it is called **strictly decreasing** if each term is smaller than the one that precedes it.
- Every sequence of  $n^2+1$  distinct real numbers contains a subsequence of length n+1 that is either strictly increasing or strictly decreasing.

#### Proof

- Let  $a_1, a_2, \ldots, a_{n^2+1}$  be a sequence of  $n^2 + 1$  distinct real numbers.
- Associate an ordered pair with each term of the sequence, i.e.  $(i_k, d_k)$  to the term  $a_k$ , where  $i_k$  is the length of the longest increasing subsequence starting at  $a_k$  and  $d_k$  is the length of the longest decreasing subsequence starting at  $a_k$ .
- Suppose that there are no increasing or decreasing subsequences of length n+1.
- Then  $i_k$  and  $d_k$  are  $\leq n$ , for  $k=1,2,\ldots,n^2+1$ . By product rule, there are  $n^2$  choices for  $(i_k,d_k)$ .

#### Proof

- By the pigeonhole principle, two of these  $n^2 + 1$  ordered pairs are equal. I.e. there exist terms in the sequence  $a_s$  and  $a_t$  with s < t such that  $i_s = i_t$  and  $d_s = d_t$ .
- Since terms in the sequence are distinct, either  $a_s < a_t$  or  $a_s > a_t$ .
- If  $a_s < a_t$  then because  $i_s = i_t$  an increasing subsequence of length  $i_{t+1}$  can be built starting at  $a_s$ , by taking  $a_s$  followed by an increasing sequence of length  $i_t$  beginning at  $a_t$ . A contradiction.
- If  $a_s > a_t$  the same reasoning shows that  $d_s$  must be greater than  $d_t$ , which is a contradiction.

# Ramsey Theory

- The generalized pigeonhole principle can be applied to an important part of combinatorics called Ramsey theory, after the English mathematician F. P. Ramsey.
- Ramsey theory deals with the distribution of subsets of elements of sets.
- Assume that in a group of 6, each pair of individuals consists of two friends or two enemies. Show that there are either three mutual friends or three mutual enemies in the group.
  - ► *A*: one of the 6 people.
  - ▶ Of the 5 other people in the group, there are either 3 or more who are friends of A, or 3 or more who are enemies of A.
  - ► Since from the generalized pigeonhole principle, we have when 5 objects are divided into 2 sets, one of the sets has at least  $\lceil 5/2 \rceil = 3$  elements.
  - ► And from this formulate the proof.

# Ramsey Theory

- The Ramsey number R(m, n), where m and  $n \ge 2$ , denotes the minimum number of people at a party such that there are either m mutual friends or n mutual enemies, assuming that every pair of people at the party are friends or enemies.
- What we saw is that  $R(3,3) \leq 6$ .
- Some useful properties of Ramsey numbers:
  - ightharpoonup R(m,n) = R(n,m).
  - ▶ R(2, n) = n for every positive integer  $n \ge 2$ .
  - ► The exact values of only nine Ramsey numbers R(m, n) with  $3 \le m \le n$  are known including R(4, 4) = 18.
  - ▶ Only bounds are known for many other Ramsey numbers, including R(5,5) which is known to satisfy  $43 \le R(5,5) \le 49$ .