## Quiz: DPDA, CFG Normal Forms, Pumping Lemma, Closure Properties

1.	. Which of the following is true about Deterministic Pushdown Automata (DPDA)?	
	A) DPDA can accept all context-free languages	
	B) DPDA can accept some but not all CFLs	
	C) DPDA is equivalent to Turing machine	
	D) DPDA accepts only regular languages	
	er: B	
2	What makes a PDA deterministic?	
_	<b>Answer:</b> It has at most one possible move for each combination of input symbol,	
	stack symbol, and current state.	
3	Which type of languages are accepted by DPDA?	
	Answer: Deterministic context-free languages	
4	DPDA cannot recognize which of the following?	
_	A) Palindromes	
	B) $a^nb^n$	
	C) (a+b)*	
	D) $a^nb^nc^n$	
	Answer: A	
5	What is the major difference between DPDA and NPDA?	
	<b>Answer:</b> DPDA has deterministic transitions, while NPDA can have multiple choices.	
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7 8 9	<ul> <li>Which symbols are considered useless in a CFG?         Answer: Symbols that do not lead to a terminal string or are not reachable from the start symbol. </li> <li>Eliminating useless symbols helps in         Answer: Simplifying the grammar and optimizing parsing. </li> <li>To eliminate unreachable symbols, we must trace from the         Answer: Start symbol </li> <li>If a variable never leads to a terminal string, it is considered         Answer: Non-generating </li> <li>Useless symbols are eliminated in how many passes over the grammar?         Answer: Two </li> </ul>	
1:	<ol> <li>Which of the following is true about ε-productions in CFGs?         Answer: They derive the empty string ε.     </li> <li>What is the first step in removing ε-productions?         Answer: Identify nullable variables.     </li> <li>After removing ε-productions, the resulting CFG is         Answer: ε-free (except possibly for start symbol if ε ∈ L(G))     </li> <li>Which CFG rule shows an ε-production?         A) A → aB         B) B → b         C) C → ε         D) D → d         Answer: C     </li> </ol>	

<ul> <li>16. A unit production is of the form: <ul> <li>A) A → a</li> <li>B) A → B</li> <li>C) A → BC</li> <li>D) A → ε</li> <li>Answer: B</li> </ul> </li> <li>17. Unit productions can be removed by <ul> <li>Answer: Substituting the productions of the variable on the right-hand side.</li> </ul> </li> <li>18. Which step is done after eliminating unit productions? <ul> <li>Answer: Update the production rules accordingly</li> </ul> </li> <li>19. Which of the following is a valid result of removing unit production A → B? <ul> <li>Answer: A inherits all rules of B</li> </ul> </li> <li>20. Removing unit productions helps in converting CFG to <ul> <li>Answer: Chomsky Normal Form (CNF)</li> </ul> </li> </ul>
<ul> <li>21. What is the Chomsky Normal Form (CNF)?     Answer: A CFG where all rules are of the form A → BC or A → a or S → ε</li> <li>22. Can every CFG be converted to CNF?     Answer: Yes</li> <li>23. CNF is used in algorithms.     Answer: Parsing (e.g., CYK algorithm)</li> <li>24. Why convert CFG to CNF?     Answer: It simplifies parsing and analysis.</li> <li>25. Which of the following rules is in CNF?     A) A → aB     B) A → BC     C) A → BCD     D) A → ε     Answer: B</li> </ul>
<ul> <li>26. The pumping lemma for CFLs is used to prove         Answer: That a language is not context-free</li> <li>27. How many parts does the string split into in pumping lemma for CFGs?         Answer: Five: uvwxy</li> <li>28. In the pumping lemma, which parts are pumped?         Answer: v and x</li> <li>29. If a language fails the pumping lemma for CFLs, then it is         Answer: Not context-free</li> <li>30. The pumping lemma holds for all         Answer: Context-free languages</li> </ul>

15. Can every  $\epsilon$ -production be removed without changing the language? Answer: Yes, except when  $\epsilon$  is in the language.