

Question No. 20 :-

<https://leetcode.com/problems/valid-parentheses/description/>

Solution Link :-

<https://leetcode.com/problems/valid-parentheses/submissions/1399707177/>

Description :-

Time Complexity :- $O(n)$

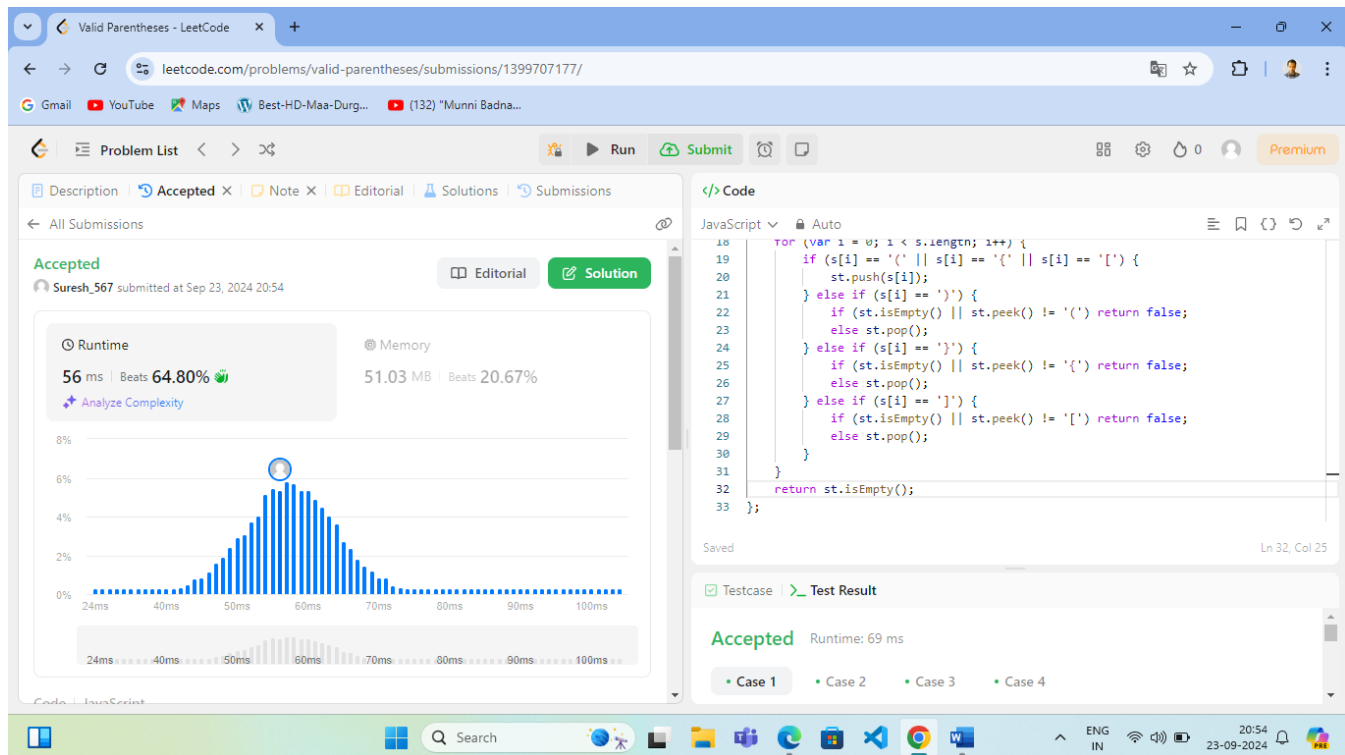
This implementation also iterates over the string s once. The push, pop, peek, and isEmpty methods all operate in constant time, $O(1)$.

Therefore, the time complexity remains **$O(n)$** , where n is the length of the string s .

Space Complexity :- $O(n)$

The worst-case scenario is when the stack holds all the opening brackets, so the space complexity is **$O(n)$** .

Screenshot :-



Question No. 496 :-

<https://leetcode.com/problems/next-greater-element-i/description/>

Solution Link :-

<https://leetcode.com/problems/next-greater-element-i/submissions/1399718722/>

Description :-

Time Complexity :- $O(n+m)$

The outer for loop iterates over all elements of `nums2`, making it $O(n)$, where n is the length of `nums2`.

The inner while loop can only run once for each element (each element is pushed and popped once from the stack). Thus, the while loop in total will also contribute **$O(n)$** .

The second loop iterates over `nums1` and does constant-time lookups in the Map, making it **$O(m)$** , where `m` is the length of `nums1`.

Overall, the time complexity is **$O(n + m)$** , where `n` is the size of `nums2` and `m` is the size of `nums1`.

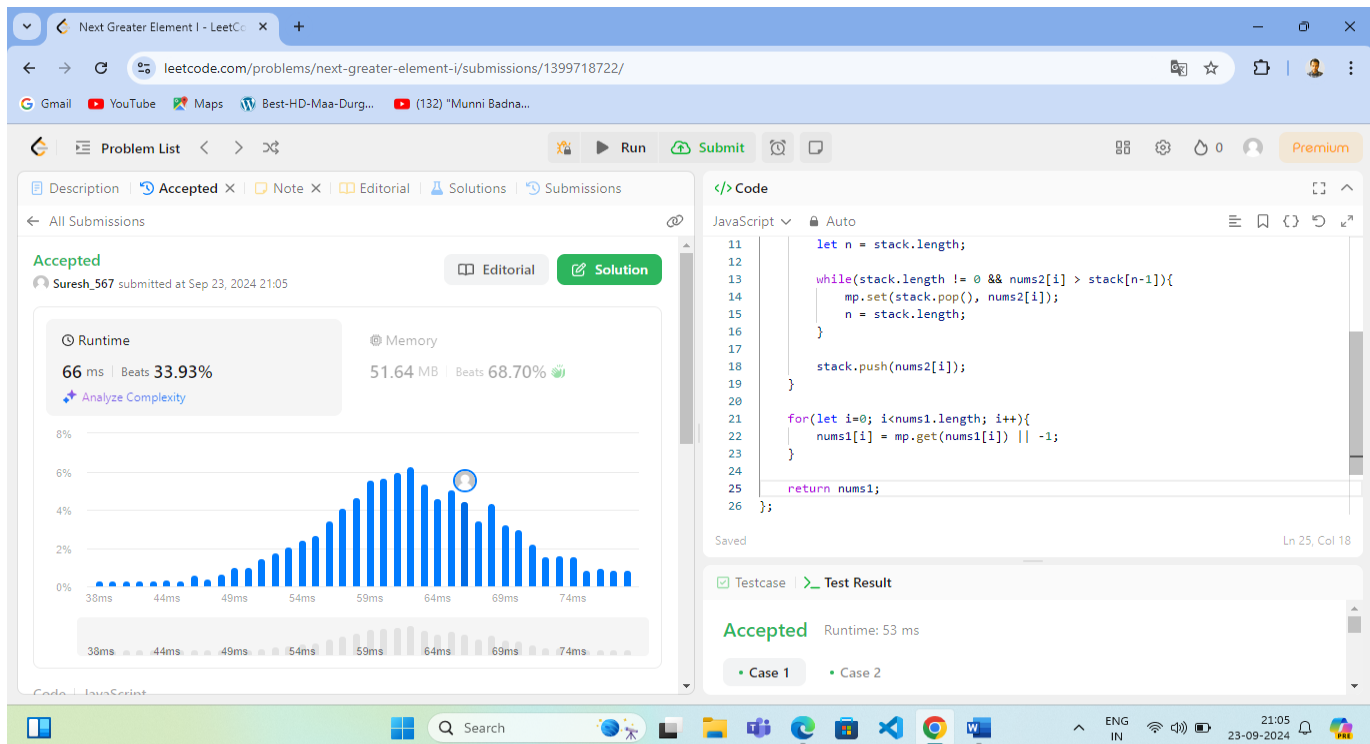
Space Complexity :- **$O(n)$**

The space complexity of the stack is **$O(n)$** in the worst case, as all elements of `nums2` could be pushed onto the stack.

The Map stores a mapping for each element of `nums2`, which also takes **$O(n)$** space.

Therefore, the space complexity is **$O(n)$** .

Screenshot :-



Question No. 1047 :-

<https://leetcode.com/problems/remove-all-adjacent-duplicates-in-string/description/>

Solution Link :-

<https://leetcode.com/problems/remove-all-adjacent-duplicates-in-string/submissions/1399732919/>

Description :-

Time Complexity :- $O(n)$

The loop iterates over each character of the string once, so the time complexity for the loop is $O(n)$, where n is the length of the string.

Stack operations like push and pop are done in constant time **$O(1)$** .

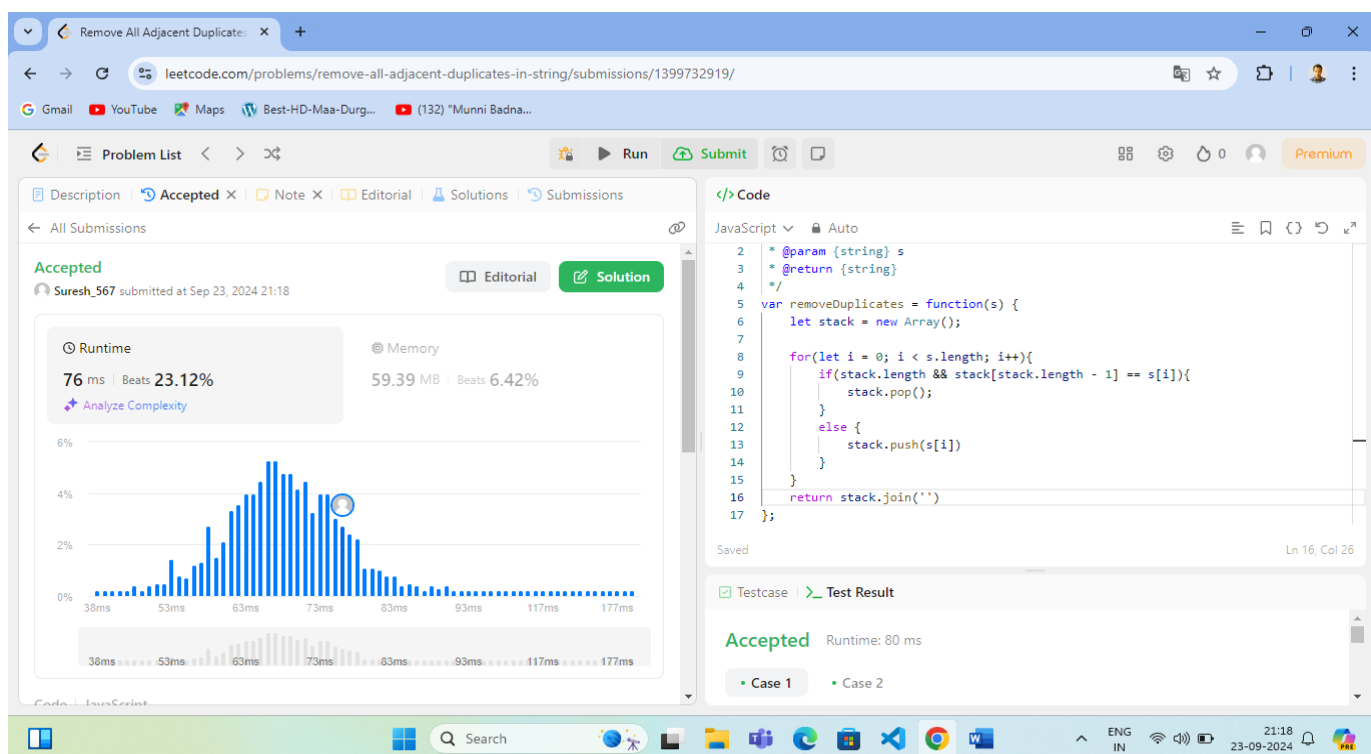
The join operation takes **$O(n)$** to concatenate the string.

Therefore, the total time complexity is **$O(n)$** .

Space Complexity :- $O(n)$

The space complexity is **$O(n)$** in the worst case, where no duplicate characters are found, and all characters are stored in the stack.

Screenshot :-



Question No. 42 :-

<https://leetcode.com/problems/trapping-rain-water/description/>

Solution Link :-

<https://leetcode.com/problems/trapping-rain-water/submissions/1399744726/>

Description :-

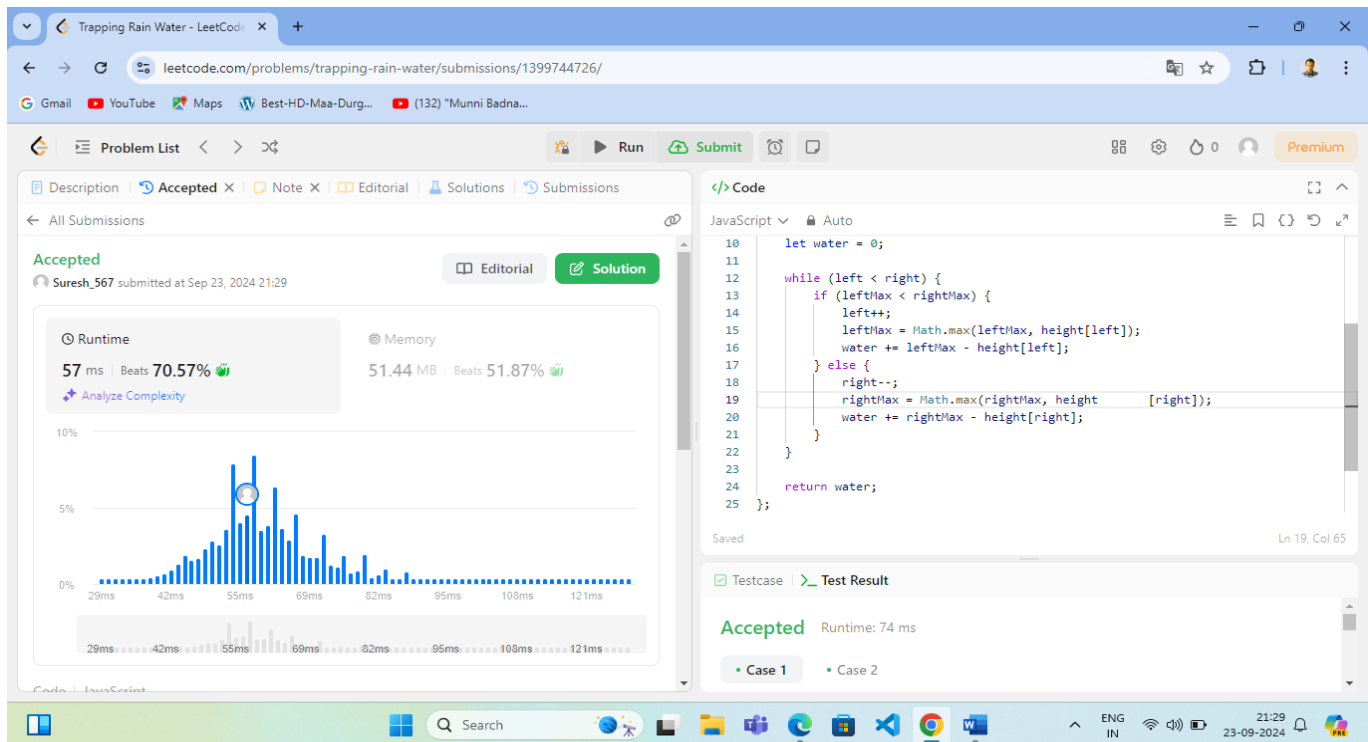
Time Complexity :- $O(n)$

This approach iterates over the array exactly once, moving pointers from both ends of the array toward the center. Each step involves constant time operations, making the time complexity **$O(n)$** , where n is the length of the height array.

Space Complexity :- $O(1)$

This approach uses a constant amount of extra space for the left, right, leftMax, rightMax, and water variables. Therefore, the space complexity is **$O(1)$** .

Screenshot :-



Question No. 84:-

<https://leetcode.com/problems/largest-rectangle-in-histogram/description/>

Solution Link :-

<https://leetcode.com/problems/largest-rectangle-in-histogram/submissions/1399755246/>

Description :-

Time Complexity :- $O(n)$

The algorithm iterates over the heights array once in the main while loop, and during this iteration, each element is pushed onto the stack once and popped from the stack once. This gives an **$O(n)$** time

complexity for the loop, where n is the length of the heights array.

The second while loop processes any remaining bars in the stack. Since each element can only be pushed and popped once, the overall contribution of this part is also **$O(n)$** .

So, the overall time complexity is **$O(n)$** .

Space Complexity :- $O(1)$

The space complexity is determined by the space used by the stack. In the worst case, if the heights are increasing, every element in the heights array will be pushed onto the stack, leading to **$O(n)$** space usage.

So, the overall space complexity is **$O(n)$** .

Screenshot :-

Browser window showing the LeetCode submission page for the problem "Largest Rectangle in Histogram". The page displays the submission status as "Accepted" and the user "Suresh_567" submitted at Sep 23, 2024 21:38.

The submission details show a Runtime of 98 ms (Beats 58.14%) and Memory of 61.02 MB (Beats 89.64%). A bar chart illustrates the performance distribution across various time intervals.

The code editor displays the JavaScript solution for the problem, which uses a stack-based approach to calculate the maximum area of a rectangle in a histogram.

```
17 let area = heights[topOfStack] * width;
18 maxArea = Math.max(maxArea, area);
19 }
20 }
21
22 while (stack.length > 0) {
23   let topOfStack = stack.pop();
24   let width = stack.length === 0 ? index : index - stack[stack.length - 1]
25   - 1;
26   let area = heights[topOfStack] * width;
27   maxArea = Math.max(maxArea, area);
28 }
29 return maxArea;
30 };
```

The test results show the submission is "Accepted" with a Runtime of 49 ms. The test cases are Case 1 and Case 2.

Windows taskbar at the bottom shows the system clock as 21:39 on 23-09-2024.