Optimization of a Finite-Volume Method Application

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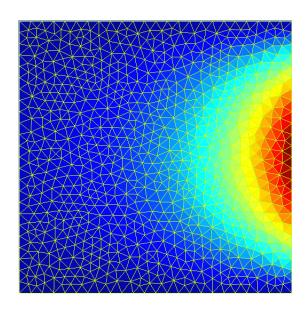
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conv-diff (Recap)

- What? Computes the heat diffusion of a fluid spreading over an area;
 - How? Uses a Finite-Volume method:
 - Why? Represents surface as a mesh, making each cell only dependent of its neighbours;







The main objective is to compute a vector $\overline{\phi}$ such that

$$\overline{\phi} \longrightarrow G(\overline{\phi}) = \begin{pmatrix} 0 \\ 0 \\ \vdots \end{pmatrix}$$
 This is accomplished in three different stages:

- We begin with a candidate vector ϕ
- For each edge, we compute the flux F_{ii} , with i and j being the indexes of the adjacent cells
- **3** For each cell, we compute $\sum |e_{ij}|F_{ij} |c_i|f_i$

Thus:
$$\phi = \begin{pmatrix} \phi_1 \\ \vdots \\ \phi_I \end{pmatrix} \longrightarrow G = \begin{pmatrix} G_1 \\ \vdots \\ G_I \end{pmatrix}$$



makeFlux Compute the contribution from each edge; makeResidual Compute the ϕ vector, adding the flux for each cell from each contribution;

Original implementation

Arrays-of-Pointers;

makeFlux

For all edges:

- Read adjacent cell data;
- 2 Compute edge velocity;
- Ompute flux through edge;

makeResidual

For all edges:

- Subtract flux from right cell;
- Add flux to left cell;

Test Machines (for most of the project)

	compute-511@search	compute-601@search	compute-101@search	MacBookPro
	AMD Opt 6174	Xeon X5650	Xeon E7520	Intel Ivy-Bridge i7
# processors	2	2	2	1
# cores per processor	12	6	1	4
hyper-threading	-	yes	yes	yes
clock frequency(GHz)	2.2	2.66	3.2	2.3
L1 capacity	128KB	32KB	16KB	64KB
L2 capacity	512KB	256KB	2MB	256KB
L3 capacity	12MB	12MB	-	6MB
RAM capacity	64GB	48GB	2GB	16GB

Table: Test machines

Naive Optimizations

- Removed redundant loads and calculations:
- Changed some variable definitions to const;
- Usage of a recent compiler auto-optimizations(SLP);



Identified Problems

- High number of memory accesses;
- Low operational intensity;
- Deep memory indirection chain;
- Bad data locality;

Achieved Results

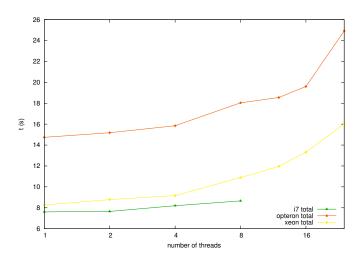


Figure: Total application runtime



Achieved Results

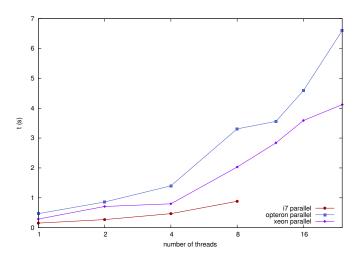


Figure: Parallel section runtime



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