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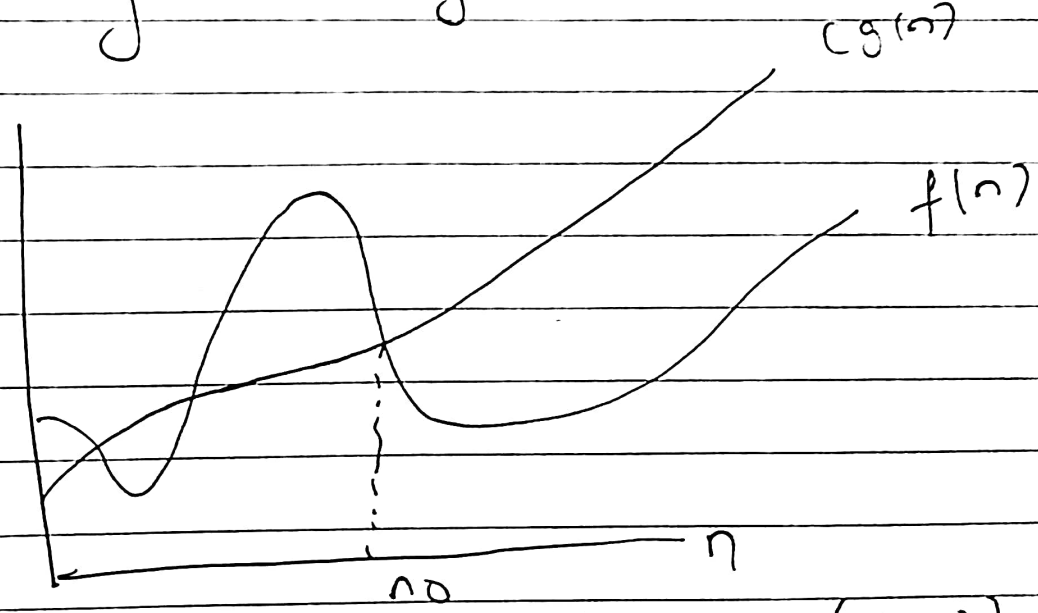
Q1.) Asymptotic Notation :- They are the mathematical notation used to describe the running time of an algorithm when the input tends towards a particular value or a limiting value.

There are mainly three asymptotic notations :-

(i) Big - O - notation.

- provide worst complexity
- provide upper bound of an running time algo.



$f(n) = O(g(n))$

$O(g(n)) = \{ f(n)$: there exist positive constant c & n0 such that $0 \leq f(n) \leq c g(n)$ for all $n \geq n_0 \}$

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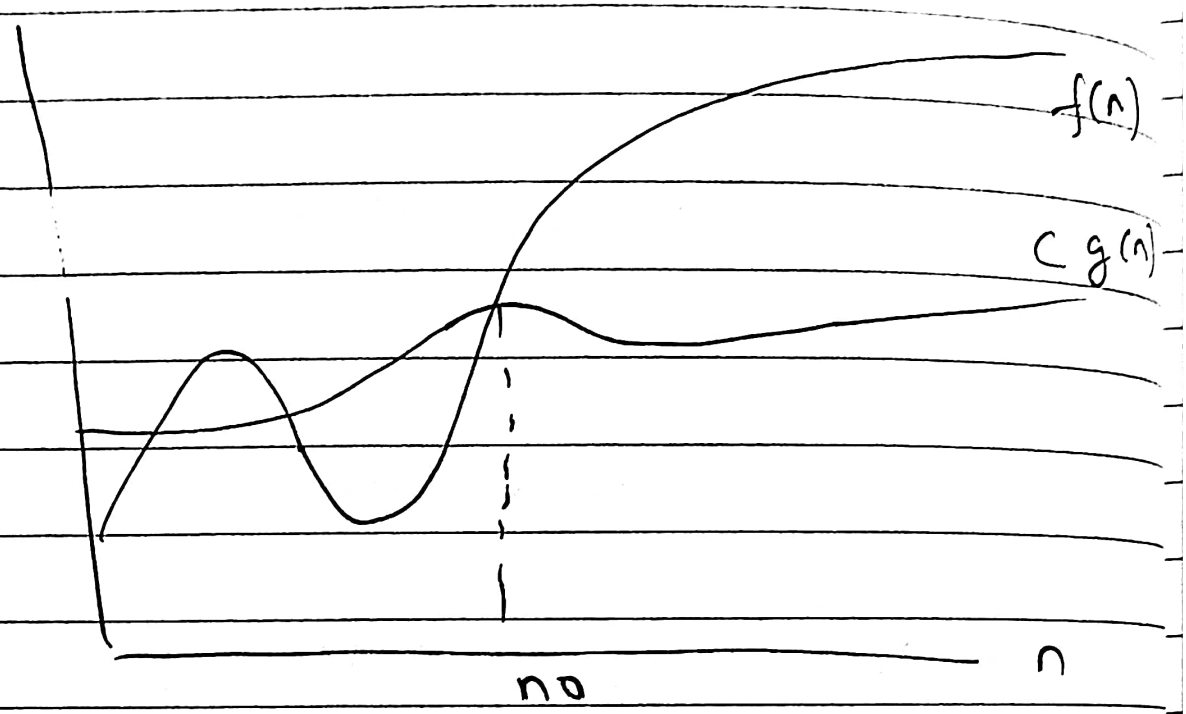
(ii) omega Notation (Ω)

•) provide best case complexity of
•) ref. lower bound running time algo.



$f(n) = \Omega(g(n))$

$\Omega(g(n)) = \{ f(n)$: there exist positive constant c and n_0 such that $0 \leq c\,g(n) \leq f(n)$ for all $n \geq n_0 \}$

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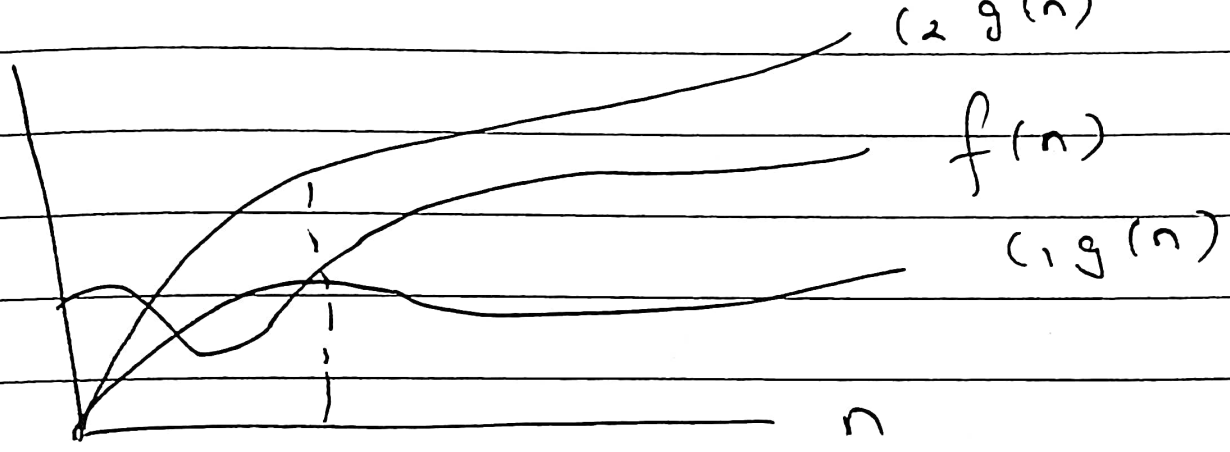
(iii) theta Notation (Θ- notation)

⇒ used for analysing avg. time complexity.



$f(n) = \Theta(g(n))$

$\Theta(g(n)) = \{ f(n)$: there exist positive constant c_1, c_2, no such that $0 \leq c_1 g(n) \leq f(n) \leq c_2 g(n)$ for all $n \geq n_0 \}$

Q2.) for (i = 1 to n)

{

i = i * 2;

}

$i = 1, 2, 2^2 \cdots 2^K$

$2^K \leq n \Rightarrow K = \log_2 n$

$\sum_{i=1}^{K} 1 \Rightarrow 1+1+1 \cdots K \text{ times}$

$T(n) = O(\log n)$

Q3 $T(n) = \begin{cases} 3T(n-1) & , n > 0 \\ 1 & , n \leq 0 \end{cases}$

Using forward substitution:

$T(0) = 1 \quad — O(1)$

$T(1) = 3T(0)$

$T(2) = 3^2 T(0)$

⋮

$T(n) = 3^n \times T(0)$

$T(n) = O(3^n)$

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Q4 $T(n) = \begin{cases} 2T(n-1) - 1, & n > 0 \\ 1, & n \leq 0 \end{cases}$

using forward subs.

$T(1) = 2T(0) - 1$, $T(0) = 1$

$= 1$

$T(2) = 2 \times T(1) - 1 = 1$

$\vdots$

$T(n) = 2 \times T(1) - 1 = 1$

$\therefore T(n) = O(1)$ Ans

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05.)
```
int i=1, s=1;

while (s <= n)
{
    i++ ; s = s+i;
    printf ('#');
}
```

for (i = 1)	for (i = 2)
s = 1+2;	s = 1+2+3

... for (i = k)

$1+2+\cdots+k <= n$

$\frac{k(k+1)}{2} <= n$

$\Rightarrow \frac{(k^2+k)}{2} \leq n$

$\Rightarrow O(k^2) <= n \Rightarrow k <= O(\sqrt{n})$

$\therefore T(n) = O(\sqrt{n})$ Ans

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Q6.)

```
void function (int n) {
    int i, count = 0;

    for (int i = 1; i * i <= n; i++)
        count ++ ⟶ O(1)
}
```

let 'K' be max +ive value such that

$$K^2 \leq n$$

$$\therefore K = \sqrt{n}$$

$$i^2 \leq n$$

$$\therefore \sum_{i=1} 1 \Rightarrow 1+1+ \dots K \text{ times}$$

$$\therefore T(n) = O(\sqrt{n})$$

Q7 =

```
void function (int n){
int i, j, k, count = 0;
for (i = n/2, i<=n, i++)
{
    for (j=1; j<=n; j=j*2)
    {
    for (k=1; k<=n; k=k*2)
        count ++;
    }
}
```

let 'm' be highest value of 2^m such that

$2^m <= n \quad \therefore \quad m = \log_2 n$

$\Rightarrow$ for $i = n/2$ $j = \log n$ $k = \log n$

$i = (n/2+1)$ " "

⋮

$i = n$ " "

$$\therefore \sum_{i=n/2}^{n} j \times k$$

$\Rightarrow n/2 (\log n)^2 \quad \therefore$

$\Rightarrow T(n) = O(n \log^2 n)$

Q8.)

```
func (int n)
{
    if (n==1) return;

    for (i=1 to n)
    {
        for (j=1 to n) {
        printf(...) ⟶ O(1)
        }
    }
    func (n-3);
}
```

for :- for (i=1 to n)

we get j = n times every turn

$\therefore \quad i \times j = n^2$

Now, $T(n) = n^2 + T(n-3)$;

$T(n-3) = (n-3)^2 + T(n-6)$; $(n-3)^2$

$T(n-6) = (n-6)^2 + T(n-9)$; $(n-6)^2$ } K terms

$\vdots$

$T(1) = 1$;

Now subs. each value in $T(n)$

$$T(n) = n^2 + (n-3)^2 + (n-6)^2 \dots + 1$$

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Let

$(n - 3k) = 1$

$\therefore \quad k = (n-1)/3$

$\therefore$ total terms $= k+1$

$T(n) = n^2 + (n-3)^2 + (n-6)^2 \ldots 1$

$T(n) \approx n^2 + n^2 + n^2 \ldots$ (k times +1)

$T(n) \approx kn^2$

$T(n) \approx \frac{(n-1)}{3} \times n^2$

$\therefore T(n) = O(n^3)$

Q9.)

```
function (int n)
{
    for (i=1 to n)
    {
        for (j=1; j<=n; j=j+i)
        printf ("*");
    }
}
```

for :- $i = 1$ $j = 1+2 \cdots (n \text{ terms of } j+i)$

$i = 2$ $j = 1+3+5 \cdots$ ''

$i = 3$ $j = 1+4+7 \cdots$ ''

$\vdots$ $\vdots$

i^{th} term of AP is

$T(m) = a + d \times m$

$T(m) = 1 + d \times m$

$(n-1)/d = m$

$\therefore$ for $i = 1$ j $(n-1)/1$ times

$i = 2$ $(n-1)/2$ times

$i = 3$ $(n-1)/3$ times

$i = n-1$ 1

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We get

$$T(n) = i_1 j_1 + i_2 j_2 \cdots i_{n-1} j_{n-1}$$

$$= \frac{(n-1)}{1} + \frac{(n-2)}{2} + \frac{(n-3)}{3} \cdots +$$

$$= n + n/2 + n/3 \cdots + \frac{n}{n-1} - n \times 1 +$$

$$= n\left[1 + \frac{1}{2} + \frac{1}{3} + \frac{1}{n-1}\right] - n + 1$$

$$= n \times \log n - n + 1$$

Since $\int \frac{1}{x} = \log x$

$$T(n) = O(n \log n)$$

Q10.) We have given

n^k & c^n

at $k \geq 1$ & $c > 1$.

$\Rightarrow$ For values $k \geq 1$, $c > 1$

we have $c^n \geq n^k$

$\therefore$ $n^k = O(c^n)$

$\forall$ $n \geq n_0$, & some constant $k_0 > 0$

$\Rightarrow$ $k_0 c^n \geq n^k$

For $c > 1$ & $n = 1$

we get.

$\Rightarrow$ $k_0 c \geq 1$

$\therefore$ $\boxed{c > 1 \quad \& \quad n_0 = 1}$ Ans