

# Interpolation & Polynomial Approximation

## Lagrange Interpolating Polynomials II

Numerical Analysis (9th Edition)

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Beamer Presentation Slides

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# Outline

## 1 Interpolating Polynomial Error Bound

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- 2 Example: 2nd Lagrange Interpolating Polynomial Error Bound

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# The Lagrange Polynomial: Theoretical Error Bound

## Theorem

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Suppose  $x_0, x_1, \dots, x_n$  are distinct numbers in the interval  $[a, b]$  and  $f \in C^{n+1}[a, b]$ .

# The Lagrange Polynomial: Theoretical Error Bound

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Suppose  $x_0, x_1, \dots, x_n$  are distinct numbers in the interval  $[a, b]$  and  $f \in C^{n+1}[a, b]$ . Then, for each  $x$  in  $[a, b]$ , a number  $\xi(x)$  (generally unknown) between  $x_0, x_1, \dots, x_n$ , and hence in  $(a, b)$ , exists with

$$f(x) = P(x) + \frac{f^{(n+1)}(\xi(x))}{(n+1)!} (x - x_0)(x - x_1) \cdots (x - x_n)$$



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$$f(x) = P(x) + \frac{f^{(n+1)}(\xi(x))}{(n+1)!} (x - x_0)(x - x_1) \cdots (x - x_n)$$

where  $P(x)$  is the interpolating polynomial given by

$$P(x) = f(x_0)L_{n,0}(x) + \cdots + f(x_n)L_{n,n}(x) = \sum_{k=0}^n f(x_k)L_{n,k}(x)$$

# The Lagrange Polynomial: Theoretical Error Bound

## Error Bound: Proof (1/6)

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Note first that if  $x = x_k$ , for any  $k = 0, 1, \dots, n$ , then  $f(x_k) = P(x_k)$ , and choosing  $\xi(x_k)$  arbitrarily in  $(a, b)$  yields the result:

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If  $x \neq x_k$ , for all  $k = 0, 1, \dots, n$ , define the function  $g$  for  $t$  in  $[a, b]$  by

$$\begin{aligned} g(t) &= f(t) - P(t) - [f(x) - P(x)] \frac{(t-x_0)(t-x_1)\cdots(t-x_n)}{(x-x_0)(x-x_1)\cdots(x-x_n)} \\ &= f(t) - P(t) - [f(x) - P(x)] \prod_{i=0}^n \frac{(t-x_i)}{(x-x_i)} \end{aligned}$$

# The Lagrange Polynomial: Theoretical Error Bound

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## Error Bound: Proof (2/6)

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## Error Bound: Proof (2/6)

Since  $f \in C^{n+1}[a, b]$ , and  $P \in C^\infty[a, b]$ , it follows that  $g \in C^{n+1}[a, b]$ .  
For  $t = x_k$ , we have

$$g(x_k) = f(x_k) - P(x_k) - [f(x) - P(x)] \prod_{i=0}^n \frac{(x_k - x_i)}{(x - x_i)} = 0 - [f(x) - P(x)] \cdot 0 = 0$$

# The Lagrange Polynomial: Theoretical Error Bound

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## Error Bound: Proof (3/6)

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We have seen that  $g(x_k) = 0$ . Furthermore,

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Thus  $g \in C^{n+1}[a, b]$ , and  $g$  is zero at the  $n + 2$  distinct numbers  $x, x_0, x_1, \dots, x_n$ .

# The Lagrange Polynomial: Theoretical Error Bound

## Error Bound: Proof (4/6)

Since  $g \in C^{n+1}[a, b]$ , and  $g$  is zero at the  $n + 2$  distinct numbers  $x, x_0, x_1, \dots, x_n$ , by Generalized Rolle's Theorem [▶ Theorem](#) there exists a number  $\xi$  in  $(a, b)$  for which  $g^{(n+1)}(\xi) = 0$ .

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$$\begin{aligned} 0 &= g^{(n+1)}(\xi) \\ &= f^{(n+1)}(\xi) - P^{(n+1)}(\xi) - [f(x) - P(x)] \frac{d^{n+1}}{dt^{n+1}} \left[ \prod_{i=0}^n \frac{(t - x_i)}{(x - x_i)} \right]_{t=\xi} \end{aligned}$$

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However,  $P(x)$  is a polynomial of degree at most  $n$ , so the  $(n + 1)$ st derivative,  $P^{(n+1)}(x)$ , is identically zero.

# The Lagrange Polynomial: Theoretical Error Bound

## Error Bound: Proof (5/6)

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Also,  $\prod_{i=0}^n \frac{t - x_i}{x - x_i}$  is a polynomial of degree  $(n + 1)$ , so

$$\prod_{i=0}^n \frac{(t - x_i)}{(x - x_i)} = \left[ \frac{1}{\prod_{i=0}^n (x - x_i)} \right] t^{n+1} + (\text{lower-degree terms in } t),$$

# The Lagrange Polynomial: Theoretical Error Bound

## Error Bound: Proof (5/6)

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and

$$\frac{d^{n+1}}{dt^{n+1}} \prod_{i=0}^n \frac{(t - x_i)}{(x - x_i)} = \frac{(n + 1)!}{\prod_{i=0}^n (x - x_i)}$$

# The Lagrange Polynomial: Theoretical Error Bound

## Error Bound: Proof (6/6)



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We therefore have:

$$\begin{aligned} 0 &= f^{(n+1)}(\xi) - P^{(n+1)}(\xi) - [f(x) - P(x)] \frac{d^{n+1}}{dt^{n+1}} \left[ \prod_{i=0}^n \frac{(t - x_i)}{(x - x_i)} \right]_{t=\xi} \\ &= f^{(n+1)}(\xi) - 0 - [f(x) - P(x)] \frac{(n+1)!}{\prod_{i=0}^n (x - x_i)} \end{aligned}$$

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## Error Bound: Proof (6/6)

We therefore have:

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and, upon solving for  $f(x)$ , we get the desired result:

$$f(x) = P(x) + \frac{f^{(n+1)}(\xi)}{(n+1)!} \prod_{i=0}^n (x - x_i)$$

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# Lagrange Interpolating Polynomial Error Bound

## Example: Second Lagrange Polynomial for $f(x) = \frac{1}{x}$

In an earlier example, [▶ Original Example](#) we found the second Lagrange polynomial for  $f(x) = \frac{1}{x}$  on  $[2, 4]$  using the nodes  $x_0 = 2$ ,  $x_1 = 2.75$ , and  $x_2 = 4$ .

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## Note

We will make use of the theoretical result [▶ Theorem](#) written in the form

$$|f(x) - P(x)| \leq \max_{[2,4]} \left| \frac{f^{(n+1)}(\xi)}{(n+1)!} \right| \cdot \max_{[2,4]} \left| \prod_{i=0}^n (x - x_i) \right|$$

with  $n = 2$

# The Lagrange Polynomial: 2nd Degree Error Bound

## Solution (1/3)

Because  $f(x) = x^{-1}$ , we have

$$f'(x) = -\frac{1}{x^2}, \quad f''(x) = \frac{2}{x^3}, \quad \text{and} \quad f'''(x) = -\frac{6}{x^4}$$

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As a consequence, the second Lagrange polynomial has the error form

$$\frac{f'''(\xi(x))}{3!}(x - x_0)(x - x_1)(x - x_2) = -\frac{1}{\xi(x)^4}(x - 2)(x - 2.75)(x - 4)$$

for  $\xi(x)$  in  $(2, 4)$ .



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for  $\xi(x)$  in  $(2, 4)$ . The maximum value of  $\frac{1}{\xi(x)^4}$  on the interval is  $\frac{1}{2^4} = 1/16$ .

# The Lagrange Polynomial: 2nd Degree Error Bound

## Solution (2/3)

We now need to determine the maximum value on  $[2, 4]$  of the absolute value of the polynomial

$$g(x) = (x - 2)(x - 2.75)(x - 4) = x^3 - \frac{35}{4}x^2 + \frac{49}{2}x - 22$$

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Because

$$g'(x) = 3x^2 - \frac{35}{2}x + \frac{49}{2} = \frac{1}{2}(3x - 7)(2x - 7),$$

the critical points occur at

$$x = \frac{7}{3} \text{ with } g\left(\frac{7}{3}\right) = \frac{25}{108} \quad \text{and} \quad x = \frac{7}{2} \text{ with } g\left(\frac{7}{2}\right) = -\frac{9}{16}$$

# The Lagrange Polynomial: 2nd Degree Error Bound

## Solution (3/3)

Hence, the maximum error is

$$\max_{[2,4]} \left| \frac{f'''(\xi(x))}{3!} \right| \cdot \max_{[2,4]} |(x - x_0)(x - x_1)(x - x_2)|$$

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# Use of the Interpolating Polynomial Error Bound

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- Suppose that a table is to be prepared for the function  $f(x) = e^x$ , for  $x$  in  $[0, 1]$ .
- Assume that the number of decimal places to be given per entry is  $d \geq 8$  and that the difference between adjacent  $x$ -values, the step size, is  $h$ .

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- What step size  $h$  will ensure that linear interpolation gives an absolute error of at most  $10^{-6}$  for all  $x$  in  $[0, 1]$ ?

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Let  $x_0, x_1, \dots$  be the numbers at which  $f$  is evaluated,  $x$  be in  $[0, 1]$ , and suppose  $j$  satisfies  $x_j \leq x \leq x_{j+1}$ .

# Use of the Interpolating Polynomial Error Bound

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$$|f(x) - P(x)| = \left| \frac{f^{(2)}(\xi)}{2!} (x - x_j)(x - x_{j+1}) \right| = \frac{|f^{(2)}(\xi)|}{2} |(x - x_j)| |(x - x_{j+1})|$$

# Use of the Interpolating Polynomial Error Bound

## Solution (1/3)

The step size is  $h$ , so  $x_j = jh$ ,  $x_{j+1} = (j+1)h$ , and

$$|f(x) - P(x)| \leq \frac{|f^{(2)}(\xi)|}{2!} |(x - jh)(x - (j+1)h)|.$$

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Hence

$$\begin{aligned} |f(x) - P(x)| &\leq \frac{\max_{\xi \in [0,1]} e^{\xi}}{2} \max_{x_j \leq x \leq x_{j+1}} |(x - jh)(x - (j+1)h)| \\ &\leq \frac{e}{2} \max_{x_j \leq x \leq x_{j+1}} |(x - jh)(x - (j+1)h)|. \end{aligned}$$



# Use of the Interpolating Polynomial Error Bound

## Solution (2/3)

Consider the function  $g(x) = (x - jh)(x - (j + 1)h)$ , for  $jh \leq x \leq (j + 1)h$ .

# Use of the Interpolating Polynomial Error Bound

## Solution (2/3)

Consider the function  $g(x) = (x - jh)(x - (j + 1)h)$ , for  $jh \leq x \leq (j + 1)h$ . Because

$$g'(x) = (x - (j + 1)h) + (x - jh) = 2 \left( x - jh - \frac{h}{2} \right),$$

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the only critical point for  $g$  is at  $x = jh + \frac{h}{2}$ , with

$$g \left( jh + \frac{h}{2} \right) = \left( \frac{h}{2} \right)^2 = \frac{h^2}{4}$$

# Use of the Interpolating Polynomial Error Bound

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$$g \left( jh + \frac{h}{2} \right) = \left( \frac{h}{2} \right)^2 = \frac{h^2}{4}$$

Since  $g(jh) = 0$  and  $g((j + 1)h) = 0$ , the maximum value of  $|g'(x)|$  in  $[jh, (j + 1)h]$  must occur at the critical point.

# Use of the Interpolating Polynomial Error Bound

## Solution (3/3)

This implies that

$$|f(x) - P(x)| \leq \frac{e}{2} \max_{x_j \leq x \leq x_{j+1}} |g(x)| \leq \frac{e}{2} \cdot \frac{h^2}{4} = \frac{eh^2}{8}.$$

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Consequently, to ensure that the the error in linear interpolation is bounded by  $10^{-6}$ , it is sufficient for  $h$  to be chosen so that

$$\frac{eh^2}{8} \leq 10^{-6}. \quad \text{This implies that} \quad h < 1.72 \times 10^{-3}.$$

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Because  $n = \frac{(1-0)}{h}$  must be an integer, a reasonable choice for the step size is  $h = 0.001$ .

Questions?



# Reference Material

# Generalized Rolle's Theorem

Suppose  $f \in C[a, b]$  is  $n$  times differentiable on  $(a, b)$ . If

$$f(x) = 0$$

at the  $n + 1$  distinct numbers  $a \leq x_0 < x_1 < \dots < x_n \leq b$ , then a number  $c$  in  $(x_0, x_n)$ , and hence in  $(a, b)$ , exists with

$$f^{(n)}(c) = 0$$

[◀ Return to Error Bound Theorem](#)

# The Lagrange Polynomial: Theoretical Error Bound

Suppose  $x_0, x_1, \dots, x_n$  are distinct numbers in the interval  $[a, b]$  and  $f \in C^{n+1}[a, b]$ . Then, for each  $x$  in  $[a, b]$ , a number  $\xi(x)$  (generally unknown) between  $x_0, x_1, \dots, x_n$ , and hence in  $(a, b)$ , exists with

$$f(x) = P(x) + \frac{f^{(n+1)}(\xi(x))}{(n+1)!} (x - x_0)(x - x_1) \cdots (x - x_n)$$

where  $P(x)$  is the interpolating polynomial given by

$$P(x) = f(x_0)L_{n,0}(x) + \cdots + f(x_n)L_{n,n}(x) = \sum_{k=0}^n f(x_k)L_{n,k}(x)$$

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# The Lagrange Polynomial: 2nd Degree Polynomial

Example:  $f(x) = \frac{1}{x}$

Use the numbers (called **nodes**)  $x_0 = 2$ ,  $x_1 = 2.75$  and  $x_2 = 4$  to find the second Lagrange interpolating polynomial for  $f(x) = \frac{1}{x}$ .

## Solution (Summary)

$$\begin{aligned}P(x) &= \sum_{k=0}^2 f(x_k) L_k(x) \\&= \frac{1}{3}(x - 2.75)(x - 4) - \frac{64}{165}(x - 2)(x - 4) + \frac{1}{10}(x - 2)(x - 2.75) \\&= \frac{1}{22}x^2 - \frac{35}{88}x + \frac{49}{44}\end{aligned}$$

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