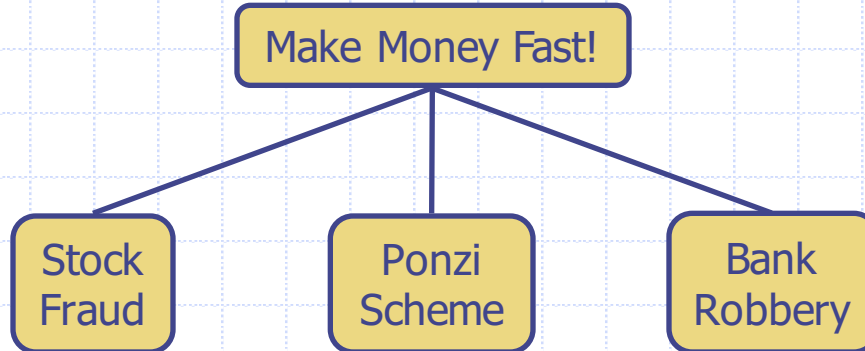
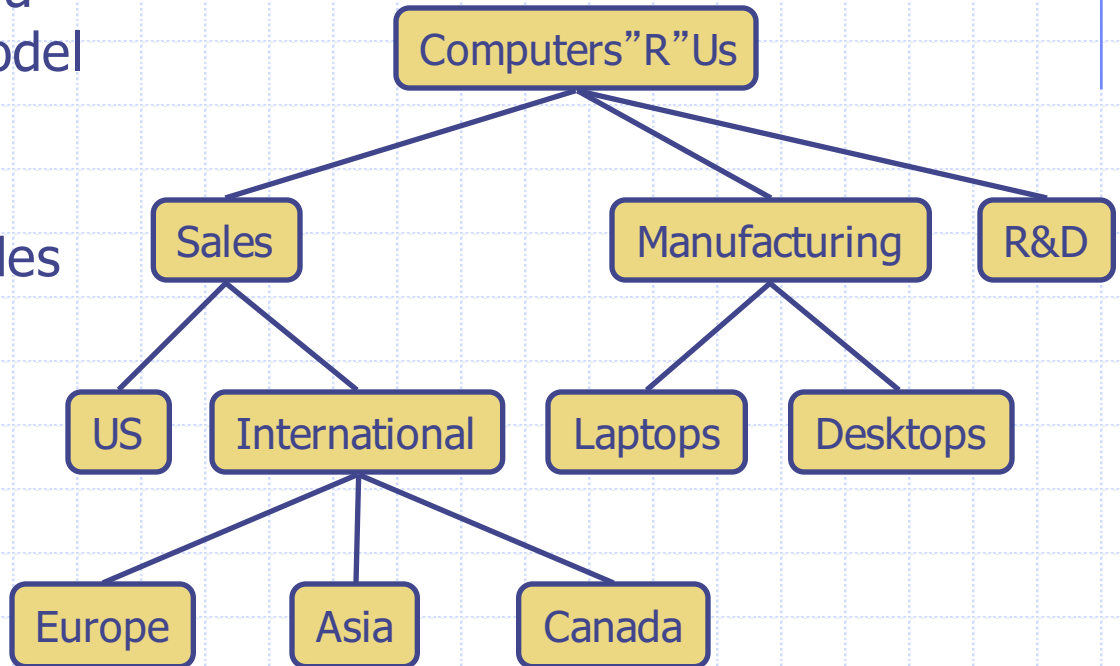


# Trees



# What is a Tree

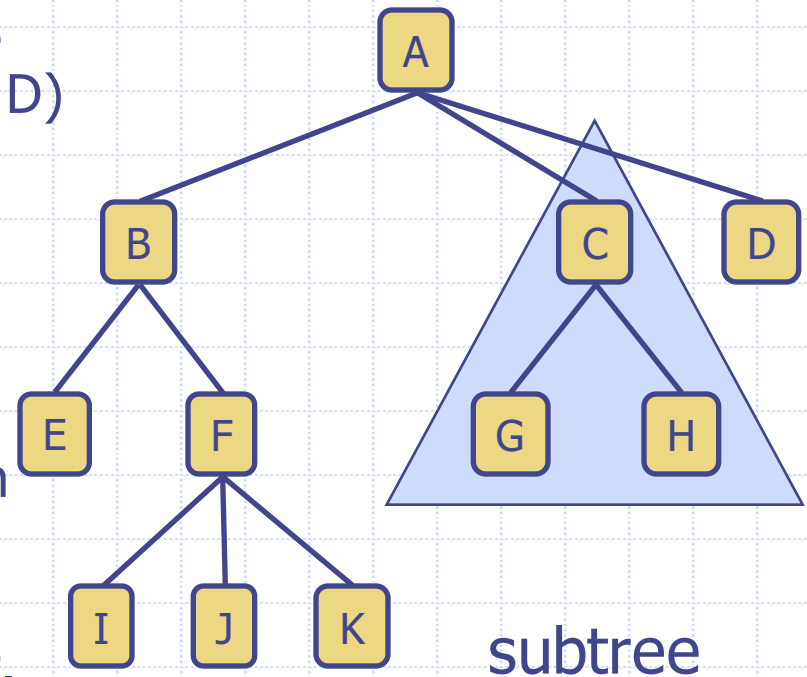
- ❑ In computer science, a tree is an abstract model of a hierarchical structure
- ❑ A tree consists of nodes with a parent-child relation
- ❑ Applications:
  - Organization charts
  - File systems
  - Programming environments



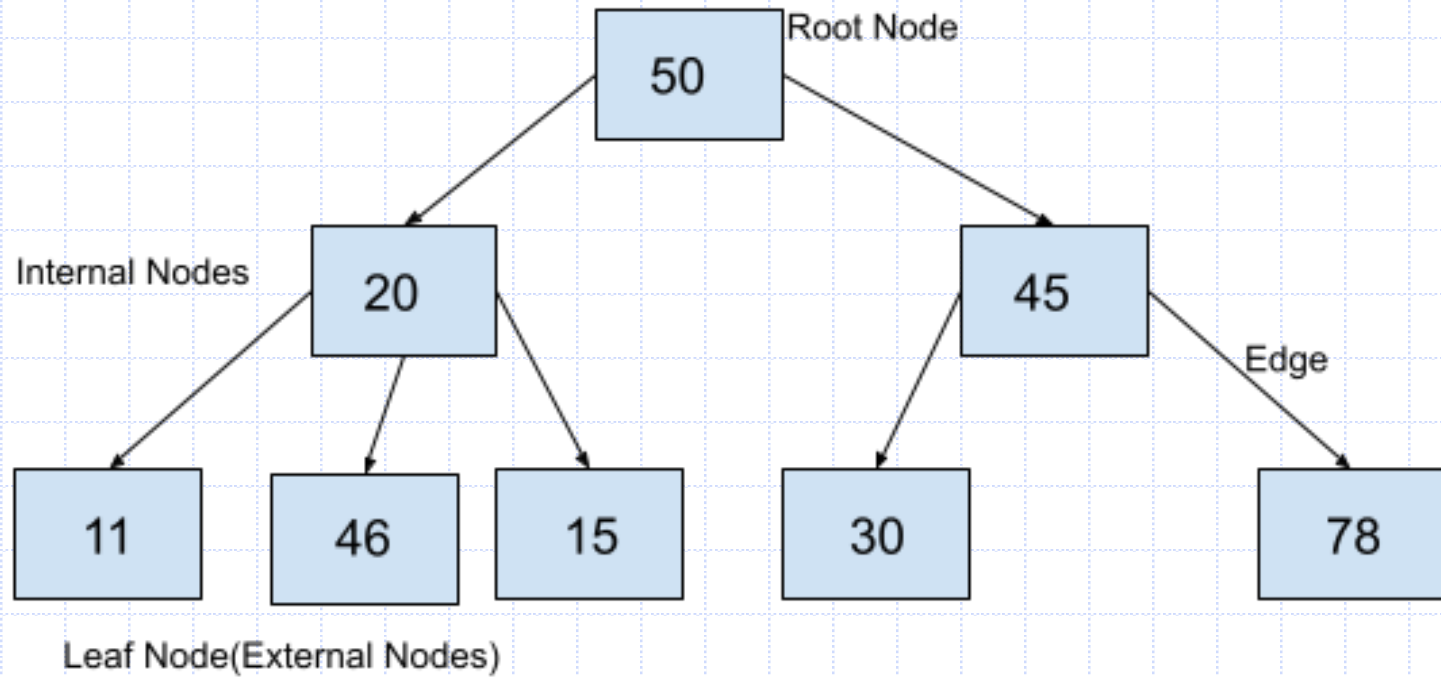
# Tree Terminology

- ❑ Root: node without parent (A)
- ❑ Internal node: node with at least one child (A, B, C, F)
- ❑ External node (a.k.a. leaf): node without children (E, I, J, K, G, H, D)
- ❑ Ancestors of a node: parent, grandparent, grand-grandparent, etc.
- ❑ Depth of a node: number of ancestors
- ❑ Height of a tree: maximum depth of any node (3)
- ❑ Descendant of a node: child, grandchild, grand-grandchild, etc.

- ❑ Subtree: tree consisting of a node and its descendants

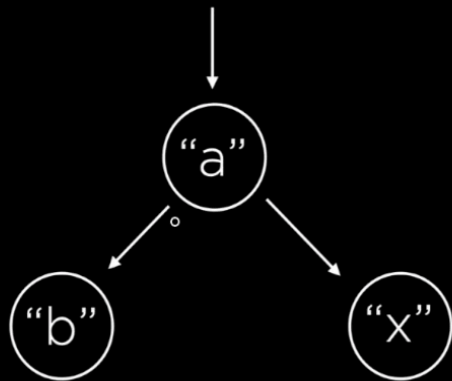


# Tree Terminology

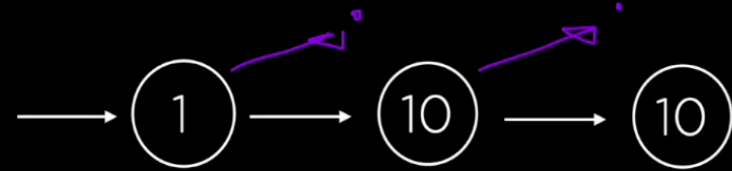


# What is a Tree

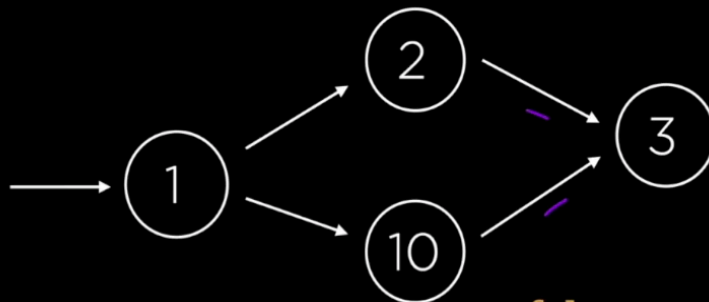
Is this a tree?



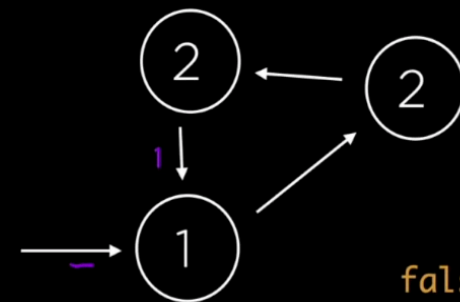
true



true

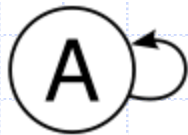
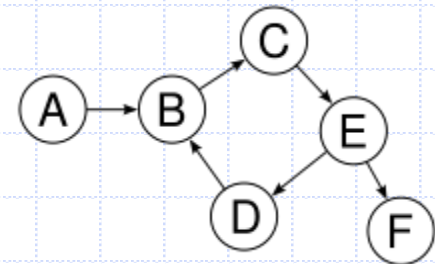
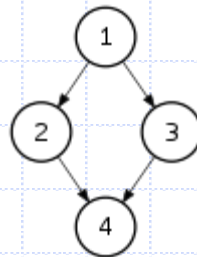
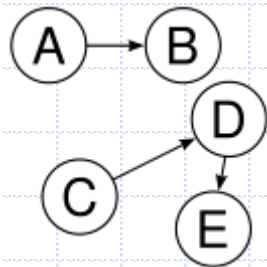


false



false

# What is a Tree



# Tree ADT

- We use positions to abstract nodes
- Generic methods:
  - Integer `len()`
  - Boolean `is_empty()`
  - Iterator `positions()`
  - Iterator `iter()`
- Accessor methods:
  - position `root()`
  - position `parent(p)`
  - Iterator `children(p)`
  - Integer `num_children(p)`
- ◆ Query methods:
  - Boolean `is_leaf(p)`
  - Boolean `is_root(p)`
- ◆ Update method:
  - element `replace (p, o)`
- ◆ Additional update methods may be defined by data structures implementing the Tree ADT

# Abstract Tree Class in Python

```
1 class Tree:
2     """Abstract base class representing a tree structure."""
3
4     #----- nested Position class -----
5     class Position:
6         """An abstraction representing the location of a single element."""
7
8         def element(self):
9             """Return the element stored at this Position."""
10            raise NotImplementedError('must be implemented by subclass')
11
12        def __eq__(self, other):
13            """Return True if other Position represents the same location."""
14            raise NotImplementedError('must be implemented by subclass')
15
16        def __ne__(self, other):
17            """Return True if other does not represent the same location."""
18            return not (self == other)          # opposite of __eq__
19
```

```
20 # ----- abstract methods that concrete subclass must support -----
21 def root(self):
22     """Return Position representing the tree's root (or None if empty)."""
23     raise NotImplementedError('must be implemented by subclass')
24
25 def parent(self, p):
26     """Return Position representing p's parent (or None if p is root)."""
27     raise NotImplementedError('must be implemented by subclass')
28
29 def num_children(self, p):
30     """Return the number of children that Position p has."""
31     raise NotImplementedError('must be implemented by subclass')
32
33 def children(self, p):
34     """Generate an iteration of Positions representing p's children."""
35     raise NotImplementedError('must be implemented by subclass')
36
37 def __len__(self):
38     """Return the total number of elements in the tree."""
39     raise NotImplementedError('must be implemented by subclass')

```

```
40 # ----- concrete methods implemented in this class -----
41 def is_root(self, p):
42     """Return True if Position p represents the root of the tree."""
43     return self.root() == p
44
45 def is_leaf(self, p):
46     """Return True if Position p does not have any children."""
47     return self.num_children(p) == 0
48
49 def is_empty(self):
50     """Return True if the tree is empty."""
51     return len(self) == 0

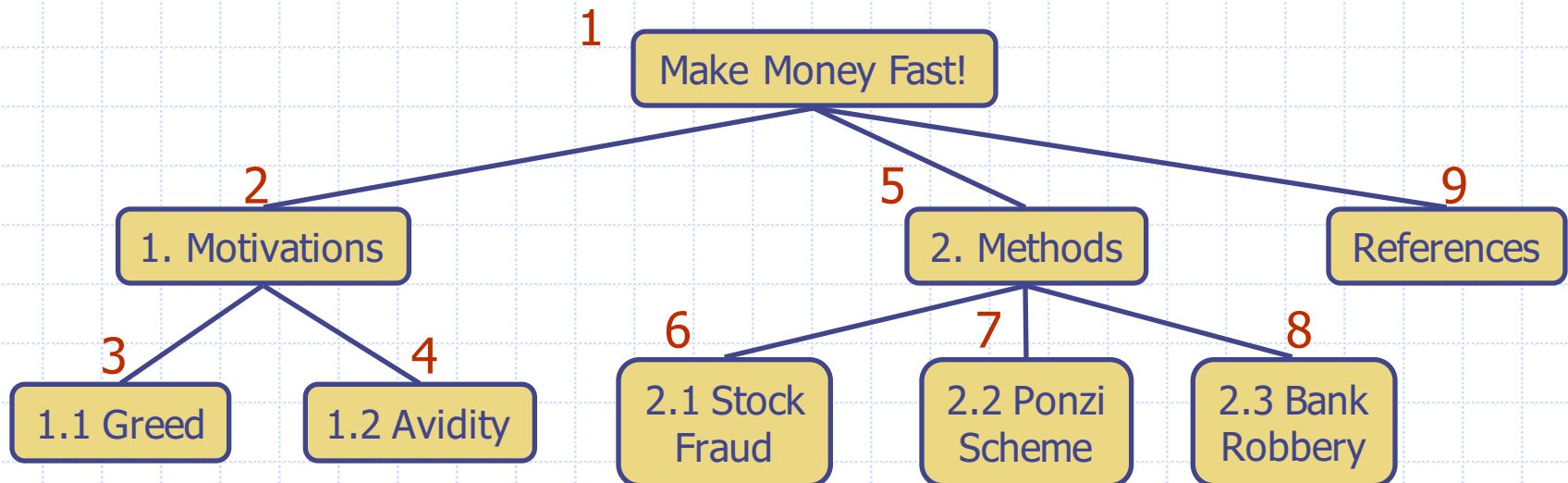
```



# Preorder Traversal

- A traversal visits the nodes of a tree in a systematic manner
- In a preorder traversal, a node is visited before its descendants
- Application: print a structured document

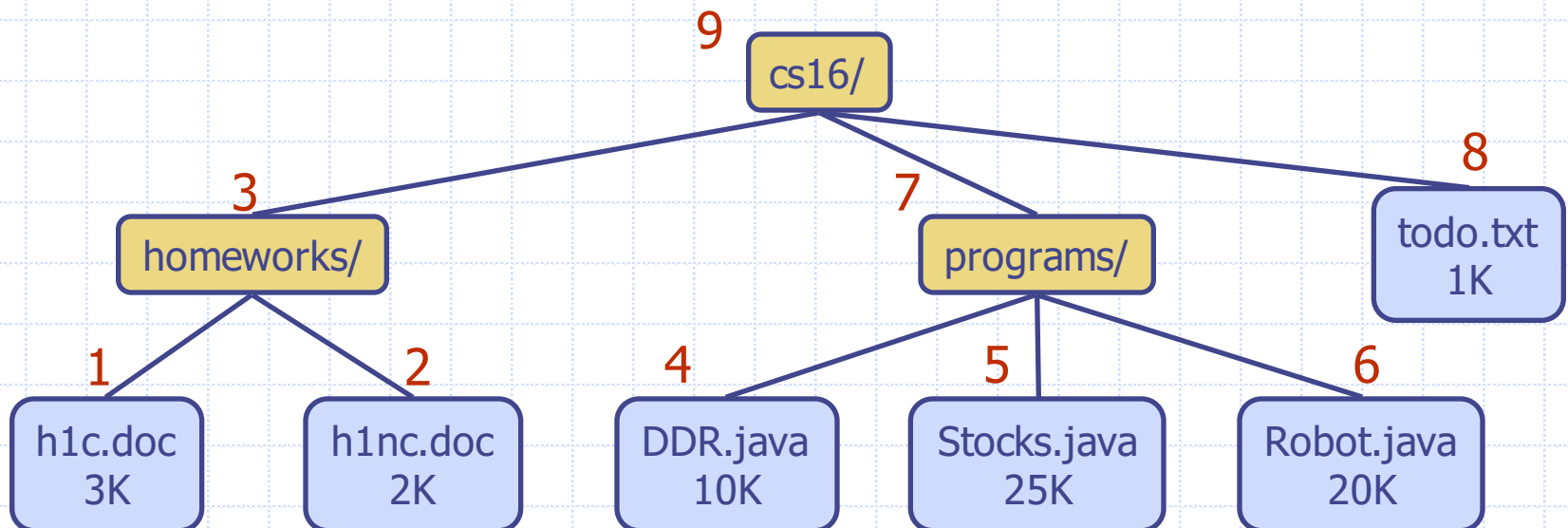
**Algorithm** *preOrder(v)*  
*visit(v)*  
**for each** child *w* of *v*  
*preorder(w)*



# Postorder Traversal

- In a postorder traversal, a node is visited after its descendants
- Application: compute space used by files in a directory and its subdirectories

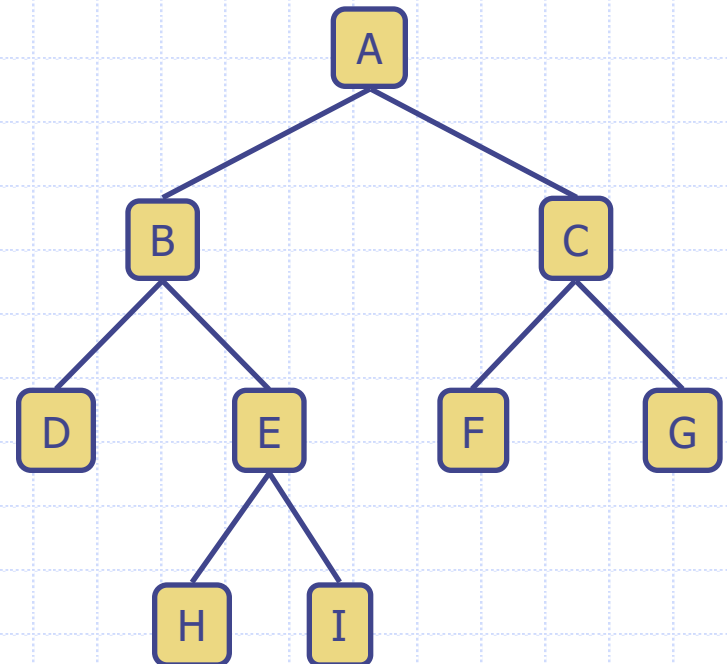
**Algorithm** *postOrder(v)*  
for each child *w* of *v*  
    *postOrder(w)*  
*visit(v)*



# Binary Trees

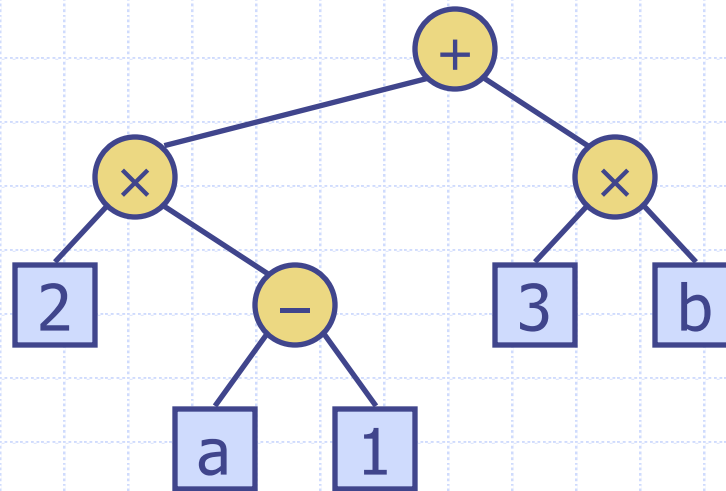
- A binary tree is a tree with the following properties:
  - Each internal node has at most two children (exactly two for **proper** binary trees)
  - The children of a node are an ordered pair
- We call the children of an internal node **left child** and **right child**
- Alternative recursive definition: a binary tree is either
  - a tree consisting of a single node, or
  - a tree whose root has an ordered pair of children, each of which is a binary tree

- Applications:
  - arithmetic expressions
  - decision processes
  - searching



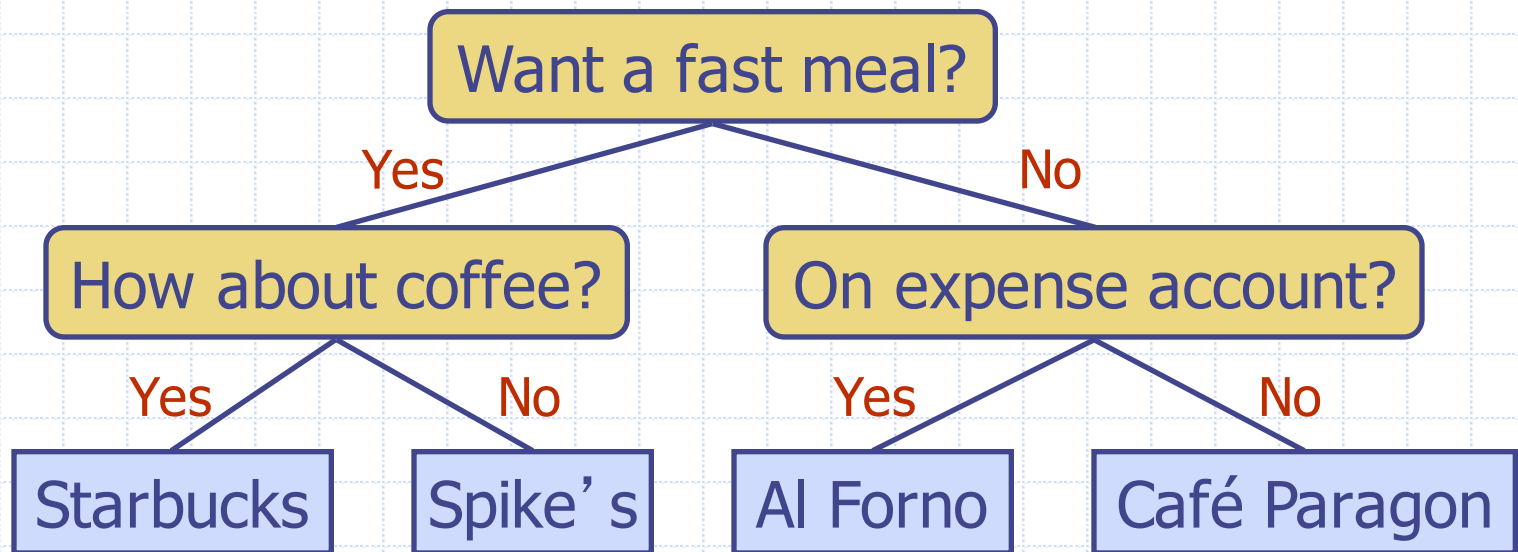
# Arithmetic Expression Tree

- Binary tree associated with an arithmetic expression
  - internal nodes: operators
  - external nodes: operands
- Example: arithmetic expression tree for the expression  $(2 \times (a - 1) + (3 \times b))$



# Decision Tree

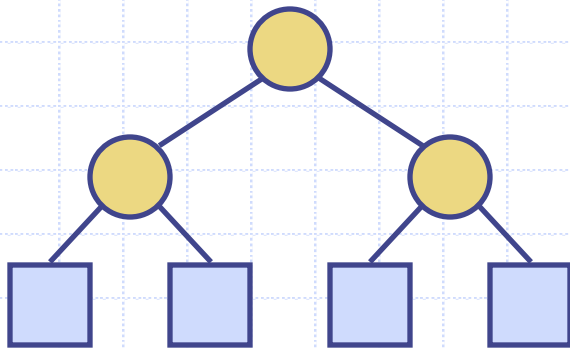
- Binary tree associated with a decision process
  - internal nodes: questions with yes/no answer
  - external nodes: decisions
- Example: dining decision



# Properties of Proper Binary Trees

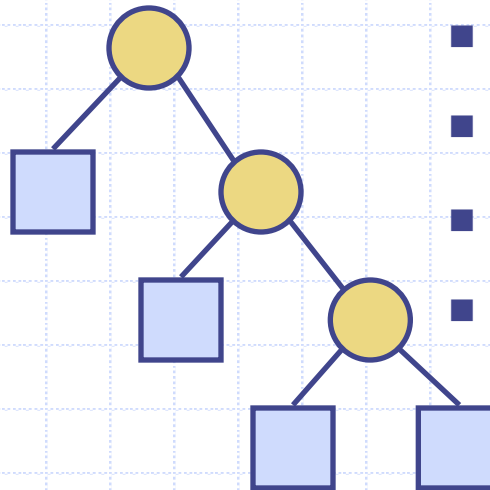
## □ Notation

- $n$  number of nodes
- $e$  number of external nodes
- $i$  number of internal nodes
- $h$  height



## ◆ Properties:

- $e = i + 1$
- $n = 2e - 1$
- $h \leq i$
- $h \leq (n - 1)/2$
- $e \leq 2^h$
- $h \geq \log_2 e$
- $h \geq \log_2 (n + 1) - 1$



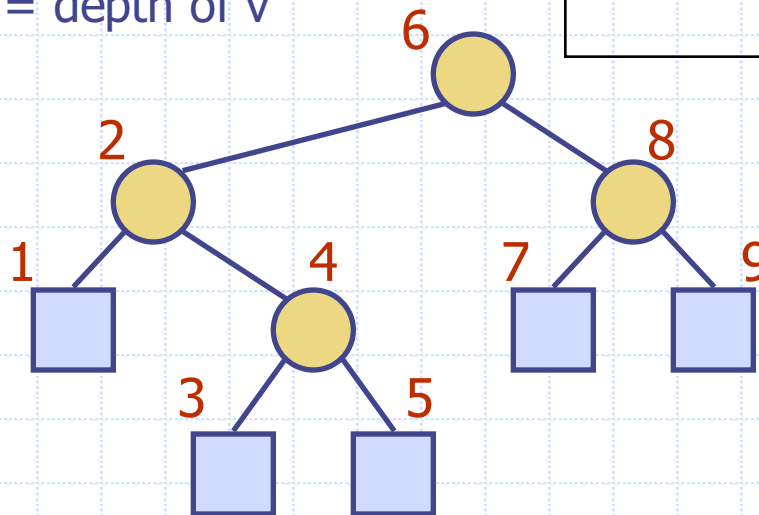
# BinaryTree ADT

- The BinaryTree ADT extends the Tree ADT, i.e., it inherits all the methods of the Tree ADT
- Additional methods:
  - position **left**(p)
  - position **right**(p)
  - position **sibling**(p)
- Update methods may be defined by data structures implementing the BinaryTree ADT

# Inorder Traversal

- In an inorder traversal a node is visited after its left subtree and before its right subtree
- Application: draw a binary tree
  - $x(v)$  = inorder rank of  $v$
  - $y(v)$  = depth of  $v$

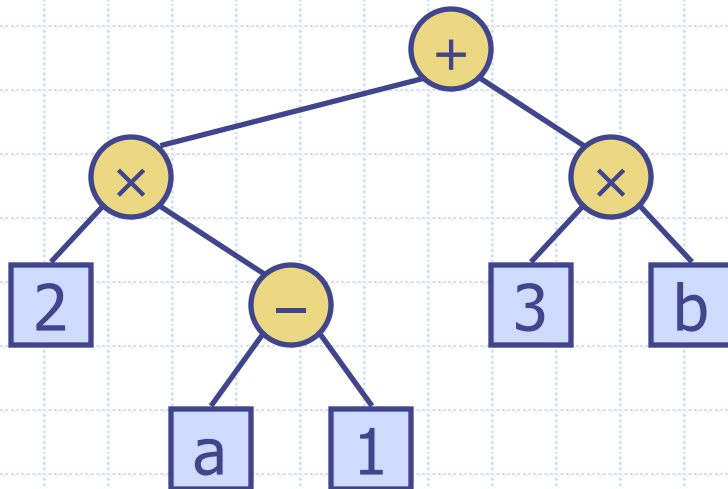
```
Algorithm inOrder( $v$ )  
  if  $v$  has a left child  
    inOrder (left ( $v$ ))  
  visit( $v$ )  
  if  $v$  has a right child  
    inOrder (right ( $v$ ))
```





# Print Arithmetic Expressions

- Specialization of an inorder traversal
  - print operand or operator when visiting node
  - print “(“ before traversing left subtree
  - print “)” after traversing right subtree



**Algorithm** *printExpression(v)*

**if** *v* **has a left child**

*print*“( ’ ”)

*inOrder* (*left(v)*)

*print*(*v.element* ())

**if** *v* **has a right child**

*inOrder* (*right(v)*)

*print*“( ’ ”)

$((2 \times (a - 1)) + (3 \times b))$

# Evaluate Arithmetic Expressions

- Specialization of a postorder traversal
  - recursive method returning the value of a subtree
  - when visiting an internal node, combine the values of the subtrees

**Algorithm** *evalExpr(v)*

**if** *is\_leaf(v)*

**return** *v.element()*

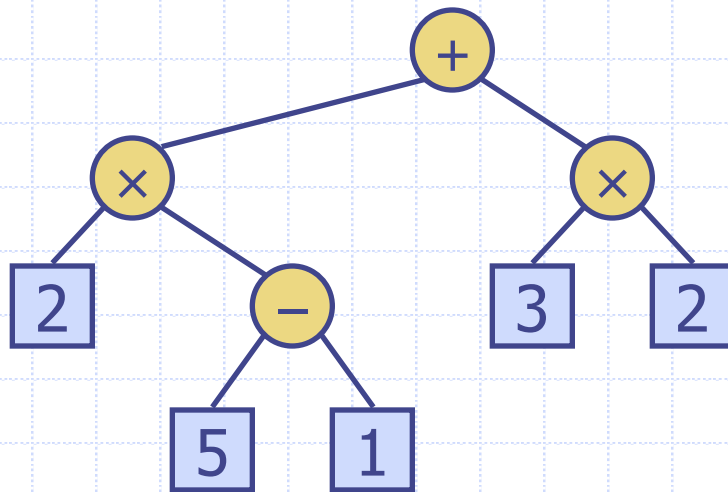
**else**

*x*  $\leftarrow$  *evalExpr(left(v))*

*y*  $\leftarrow$  *evalExpr(right(v))*

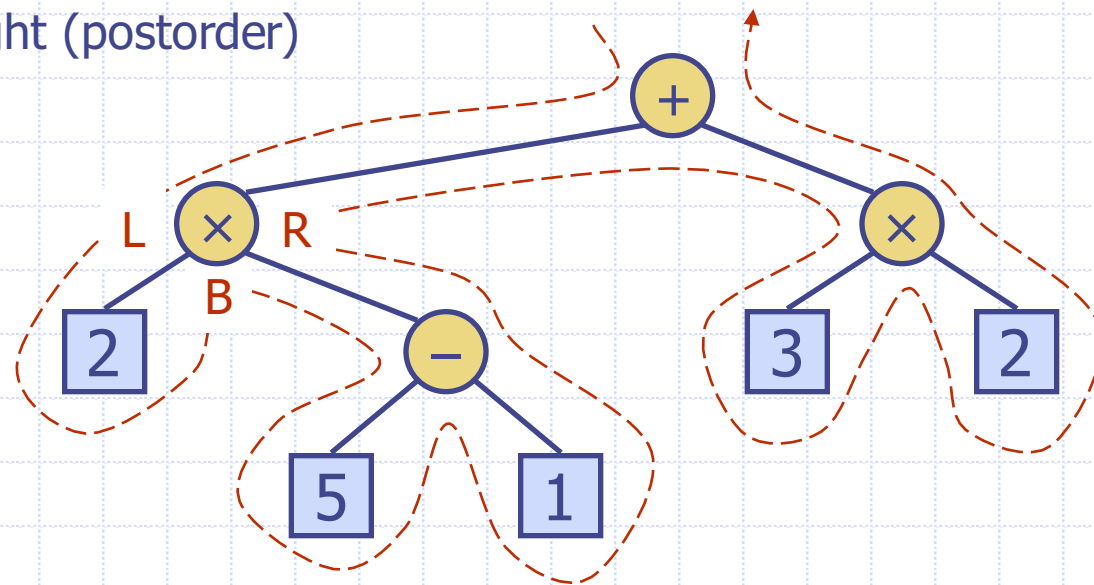
$\diamond$   $\leftarrow$  operator stored at *v*

**return** *x*  $\diamond$  *y*



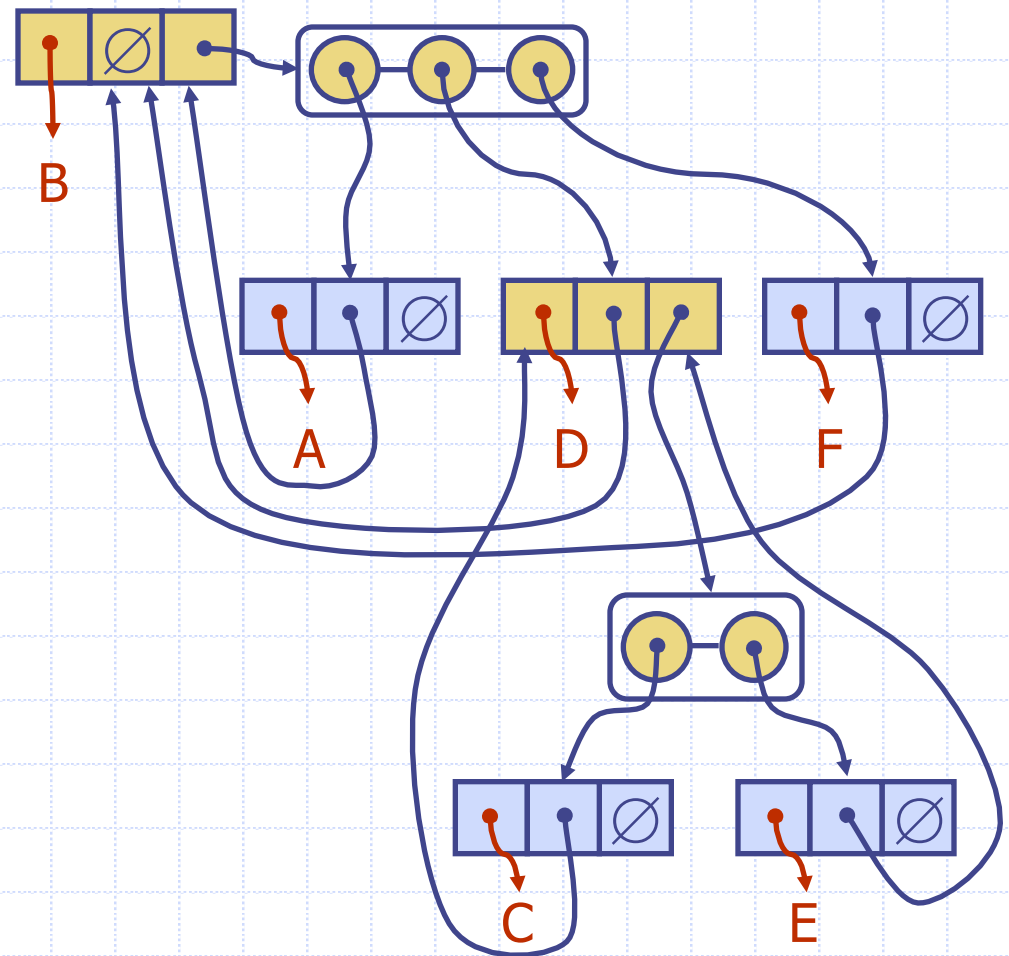
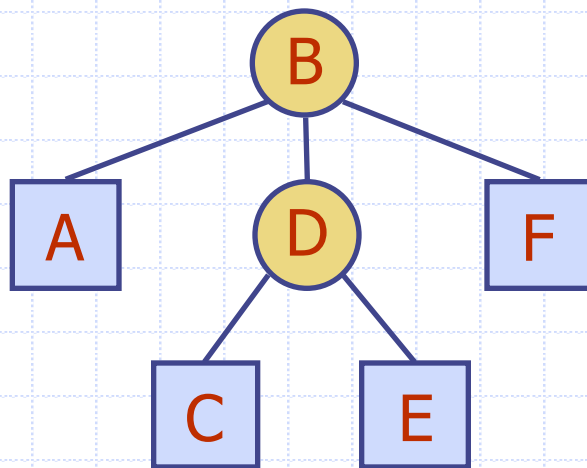
# Euler Tour Traversal

- ❑ Generic traversal of a binary tree
- ❑ Includes a special cases the preorder, postorder and inorder traversals
- ❑ Walk around the tree and visit each node three times:
  - on the left (preorder)
  - from below (inorder)
  - on the right (postorder)



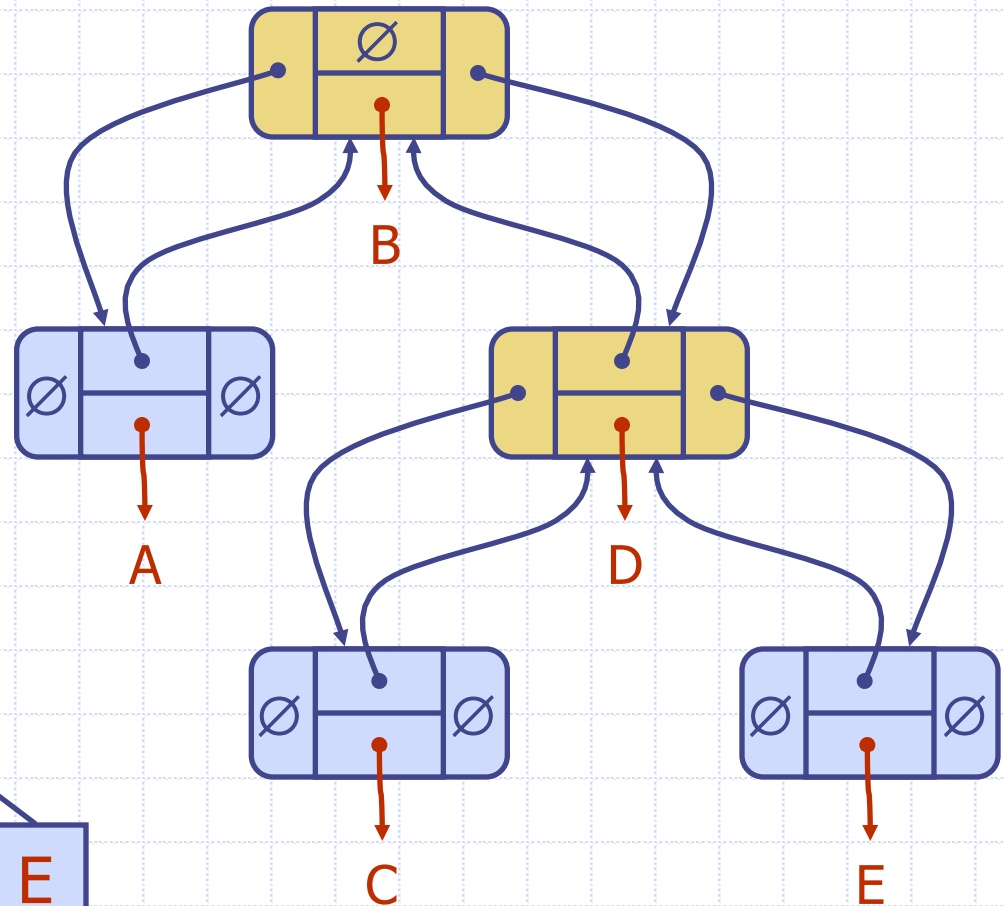
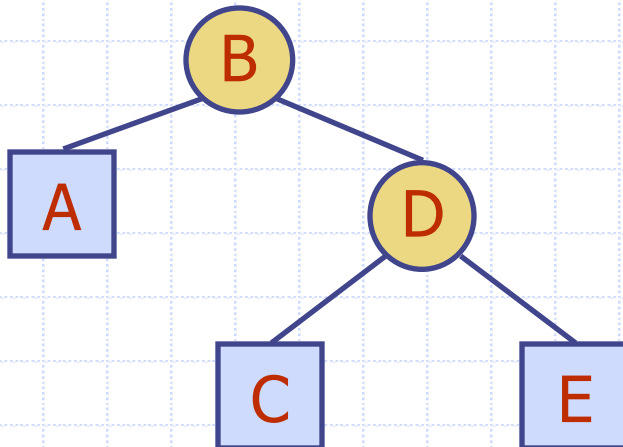
# Linked Structure for Trees

- A node is represented by an object storing
  - Element
  - Parent node
  - Sequence of children nodes
- Node objects implement the Position ADT



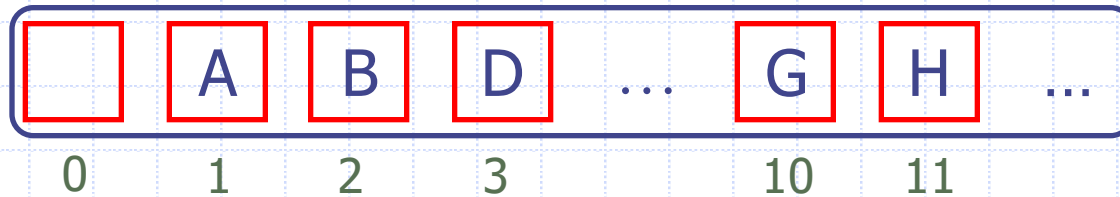
# Linked Structure for Binary Trees

- A node is represented by an object storing
  - Element
  - Parent node
  - Left child node
  - Right child node
- Node objects implement the Position ADT



# Array-Based Representation of Binary Trees

- Nodes are stored in an array A



- Node  $v$  is stored at  $A[\text{rank}(v)]$ 
  - $\text{rank}(\text{root}) = 1$
  - if node is the left child of  $\text{parent}(\text{node})$ ,  
 $\text{rank}(\text{node}) = 2 \cdot \text{rank}(\text{parent}(\text{node}))$
  - if node is the right child of  $\text{parent}(\text{node})$ ,  
 $\text{rank}(\text{node}) = 2 \cdot \text{rank}(\text{parent}(\text{node})) + 1$

