▶IPC through PIPE

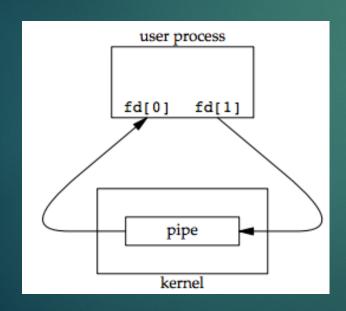
Dr. Manmath. N. Sahoo Dept. Of CSE, NIT Rourkela

## Inter-Process Communication: PIPE

- ▶ PIPE is a means of inter-process communication.
- Each Unix file has a file descriptor, similar to file pointer (FILE \*) in C
- In Unix, all the devices are treated as files.
  - Standard input device is represented as a file with file descriptor '0'
  - Standard output device is represented as a file with file descriptor '1'
- ▶ A PIPE is defined with 02 file descriptors one for read and other for write.

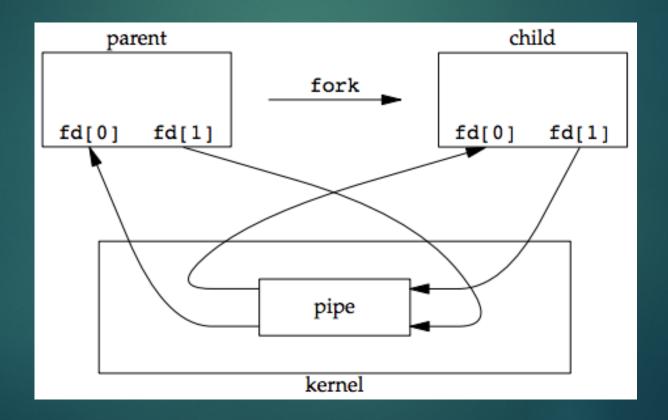
▶ A PIPE can be created using a system call pipe() that takes one integer array of size 2, as argument, e.g.,

```
#include <unistd.h>
int fd[2]; pipe(fd);
/* Returns: 0 if OK, -1 on error */
```

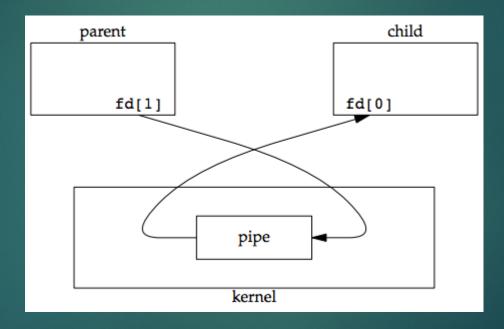


fd[1] is the write file descriptor fd[0] is the read file descriptor.

- ▶ The child inherits all open file descriptors of parent
- However, PIPE supports half-duplex communication



Parent closes its read file descriptor and child closes write file descriptor. // close (filedescriptor)



- ▶ If a process writes to a PIPE without any read fd opened then write returns -1, i.e., write fails.
- ▶ If a process reads from a PIPE without any write end opened, read returns 0 to indicate an end of file after all the data has been read from PIPE.
  - Technically, this end of file is not generated until all the write fds are not closed, and the reader waits until some input is available.
- For writing, at least ONE read fd must be opened.
- For reading, if you want end of pipe (or eof) character then all write fds must be closed.
- Example, ls –l | wc –l
  - While executing wc –l, it won't terminate if it doesn't find eof character.

```
#include<stdio.h> #include<errno.h> #include<unistd.h> #include<string.h>
int main(void){
  int n, fd[2];
  pid t pid;
  char line[20], *msg = "hello world";
  if (pipe(fd) == -1) exit(1);
        printf("PIPE created by parent successfully\n");
 else
  if ((pid = fork()) < 0) exit(1);
  close(fd[0]);
   write(fd[1], msg, strlen(msg)); //write(fd[1], &x, sizeof(int)); --for writing integer x
                               /* child */
 } else {
    close(fd[1]);
    n = read(fd[0], line, 20); //read(fd[0], &y, sizeof(int)); --for reading int
   puts(line)
```

Example: IPC through PIPE

## exec() system call

- int execl(const char \*path, const char \*arg[0], const char \* arg[1], ..., const char \* arg[n], const char \* 0);
- int execlp(const char \*file, const char \*arg[0], const char \* arg[1], ..., const char \* arg[n], const char \* 0);
- int execv(const char \*path, char \* const argv []);
- int execvp(const char \*file, char \* const argv[]);
- path complete path of the executable file
- file name of the executable file
- arg[0] command to be executed
- arg[1] ... arg[n] optional arguments for the command
- char \* 0 argument list is always ended with NULL string.

## exec() system call

- execl("/bin/ls","ls", "-l", NULL);
- execlp("ls", "ls", "-l", NULL);

- char \* const argv[]={"ls", "-l", NULL}
- execv(("/bin/ls",argv);
- execvp(("ls",argv);
- The path of "wc" file is /user/bin/wc
- To know the path of a command(file)
  - which wc
  - whereis ls

# dup(), dup2()

```
close(0);
dup(fd[0]);
```

- duplicates fd[0] as the lowest unused file descriptor.
- this case it is 0

Can be replaced by a single statement

dup2(fd[0],0);

# More on manual page

- ▶ \$man execlp
- ▶ \$man dup
- ▶ \$man dup2