

Chapter 7

Network Flow

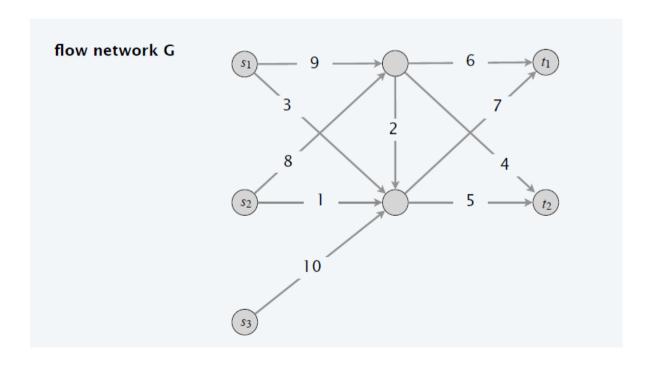


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7.7 Extensions to Maximum-Flow Problem

Multiple sources and sinks

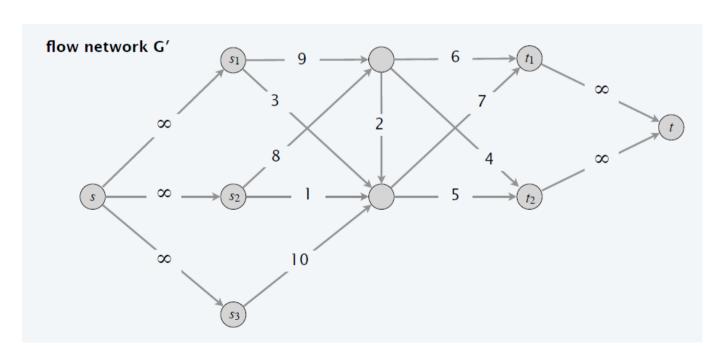
Def. Given a digraph G = (V, E) with edge capacities $c(e) \ge 0$ and multiple source nodes and multiple sink nodes, find max flow that can be sent from the source nodes to the sink nodes.



Multiple sources and sinks: max-flow formulation

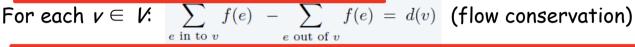
- Add a new source node s and sink node t.
- For each original source node s_i add edge (s, s_i) with capacity ∞ .
- For each original sink node t_i , add edge (t_i, t) with capacity ∞ .

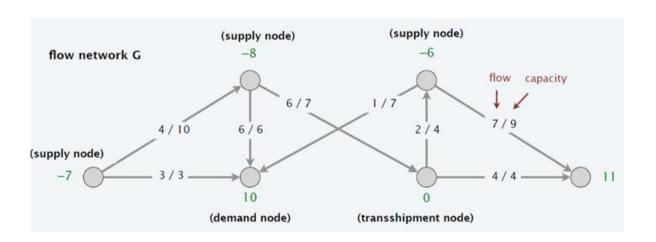
Claim. 1-1 correspondence betweens flows in G and G.



Circulation with supplies and demands

Def. Given a digraph G = (V, E) with edge capacities $c(e) \ge 0$ and node demands d(v), a circulation is a function f(e) that satisfies: For each $e \in E$: $0 \le f(e) \le c(e)$ (capacity)

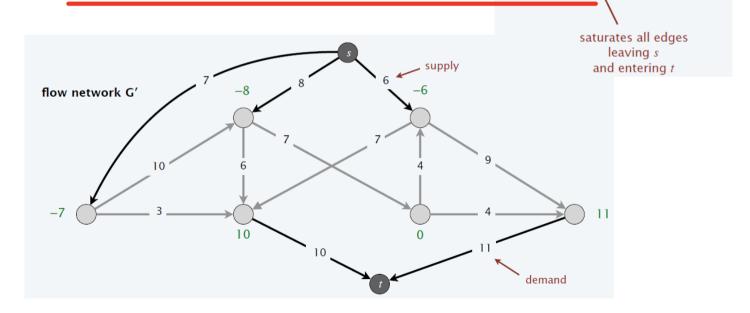




Circulation with supplies and demands: max-flow formulation

- Add new source s and sink t.
- For each ν with $d(\nu) < 0$, add edge (s, ν) with capacity $-d(\nu)$.
- For each v with d(v) > 0, add edge (v, t) with capacity d(v).

Claim. G has circulation iff G' has max flow of value $D = \sum_{v: d(v)>0} d(v) = \sum_{v: d(v)<0} -d(v)$



Circulation with supplies and demands

Integrality theorem. If all capacities and demands are integers, and there exists a circulation, then there exists one that is integer-valued.

Pf. Follows from max-flow formulation + integrality theorem for max flow.

Theorem. Given (V, E, c, d), there does not exist a circulation iff there exists a node partition (A, B) such that $\sum_{v \in B} d(v) > cap(A, B)$.

Pf sketch. Look at min cut in G'.

demand by nodes in B exceeds supply of nodes in B plus max capacity of edges going from A to B