

**Team members:**

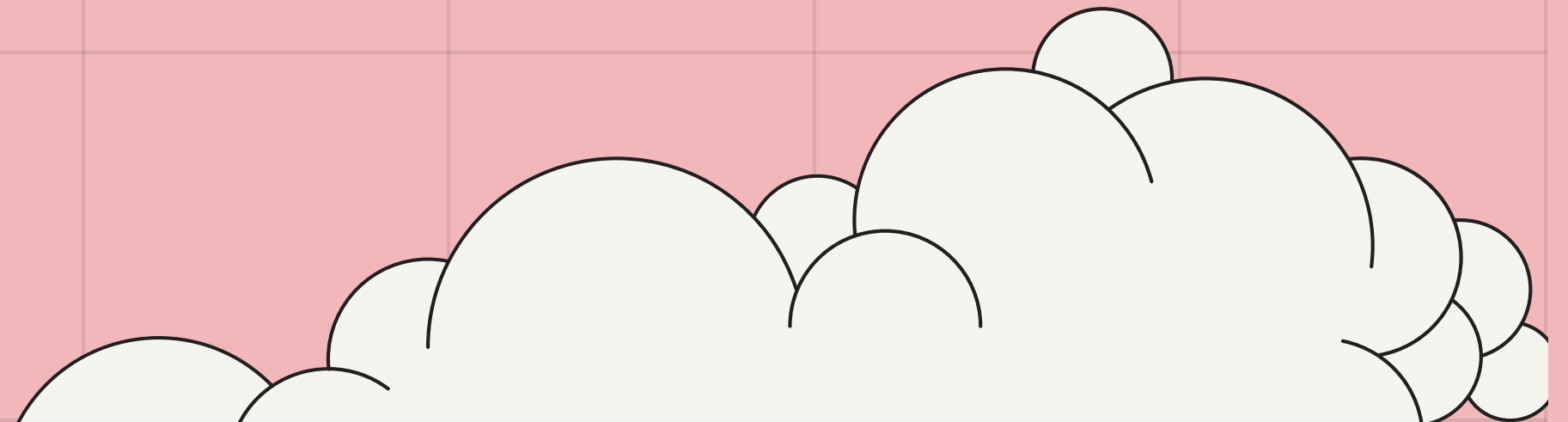
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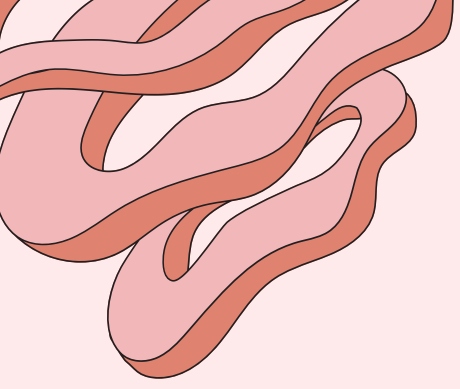
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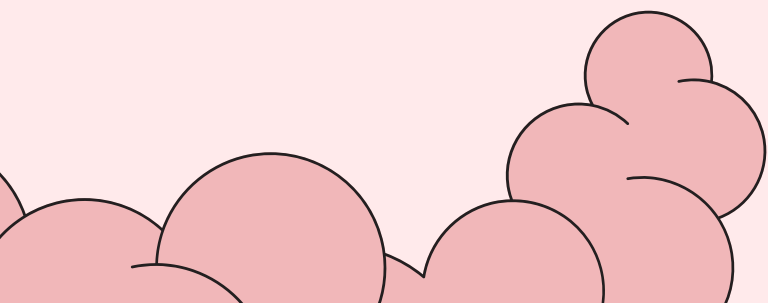


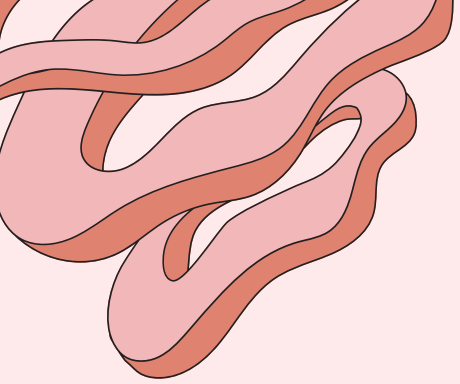
# Real life application of graph

## 1. Cheapest Flights within k stops

### Problem statement:

In the actual world, cities are connected by flights. Either a direct trip or a flight with several stops between cities is an option for travelling between cities. Most people would rather save their money. When travelling by stopover and paying for more than one flight, we can sometimes save money because direct flights are generally more expensive. So, we are going to find the cheapest price to reach from current to destination city with limited stops.

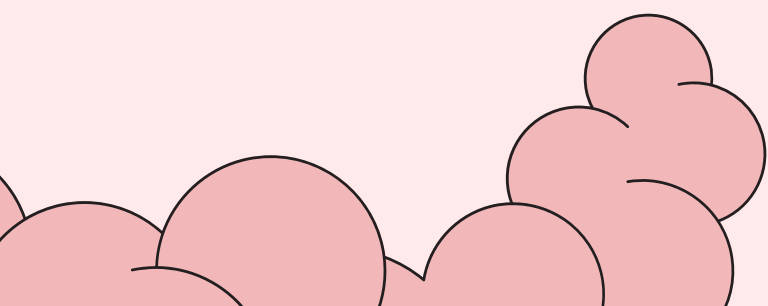




## Given:

We have  $n$  cities connected by some number of flights. We are given an array of flights where  $\text{flights}[i] = [\textit{from}, \textit{to}, \textit{price}]$  indicates that there is a flight from city "*from*" to city "*to*" with cost "*price*". We are representing city by non-negative integers.

Now, we are given three integers *source*, *destination*, and *k*, we will return the cheapest price from *source* to *destination* with at most *k* stops. If there is no such route, return -1.



**Given:**  $n = 4$ , flights =  $[[0,1,1000],[1,2,1000],[2,0,1000],[1,3,6000],[2,3,2000]]$ ,  
source = 0, destination = 3,  $k = 1$

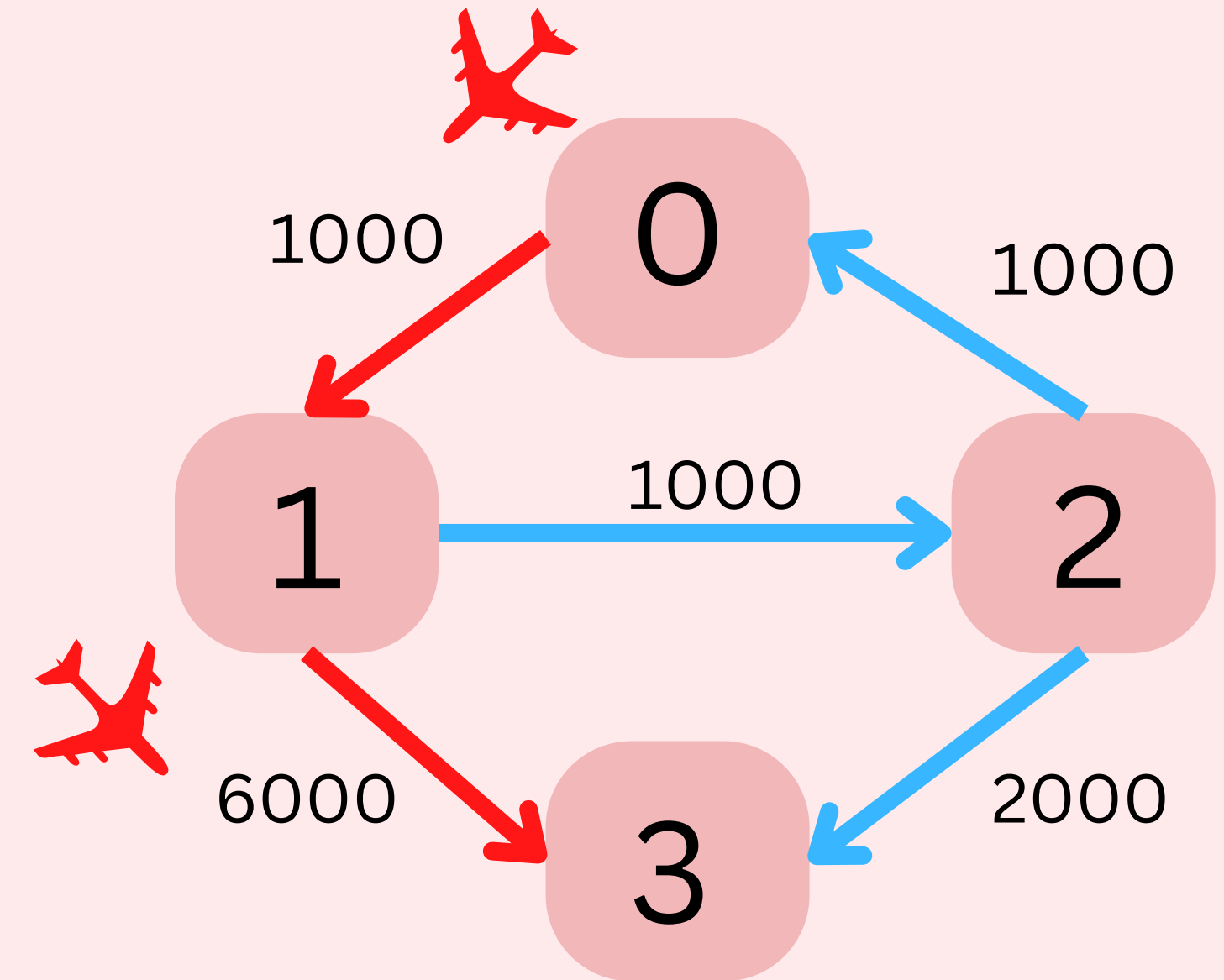
**Answer:** 7000

**Explanation:**

The graph is shown here.

The ideal route from city 0 to city 3 with atmost 1 stop possible is highlighted in red and costs  $1000 + 6000 = 7000$ .

The route through the cities  $[0,1,2,3]$  is less expensive but incorrect because it requires 2 stops.



# Solution:

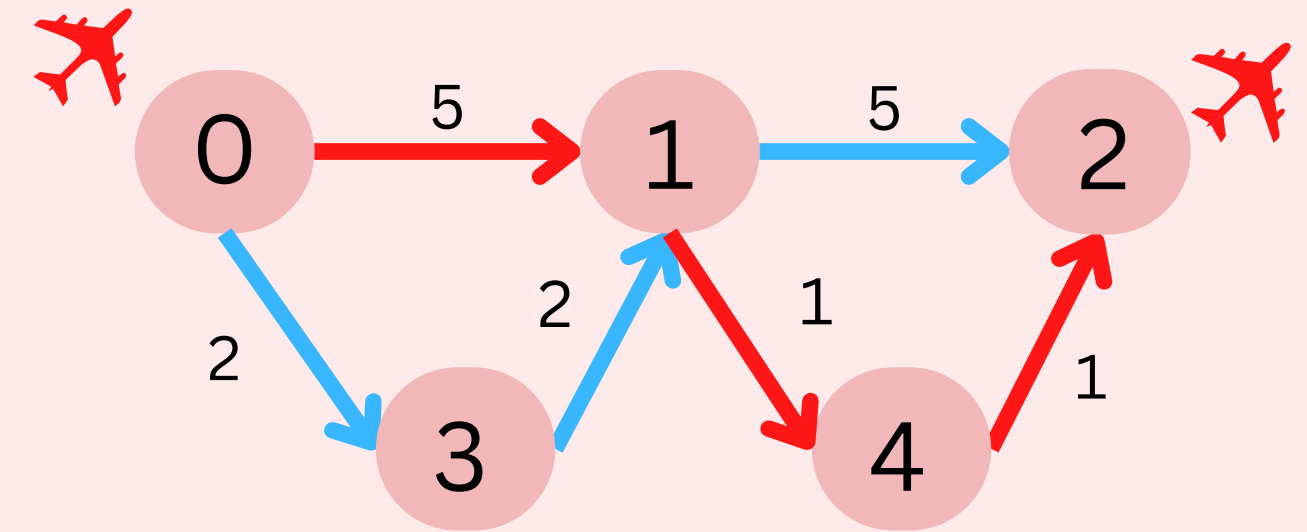
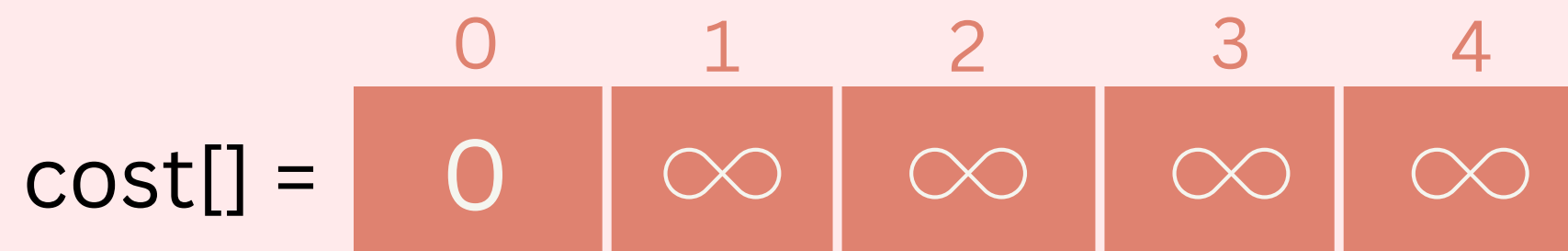
Consider this graph.

source = 0, destination = 2 and k = 2

**Step 1:** Create graph (Adjacency List).

Let's take a Queue which will store set of three values {stops, node, distance} and a cost array which will store cost of ith node. Now, we will initialize cost array with infinity.

**Step 2:** cost[source] = 0. Add {0, 0, 0} to Queue.



0 -> 1, 3  
1 -> 2, 4  
2 ->  
3 -> 1  
4 -> 2

{0, 0, 0}

{stops, node, distance}

**Queue**

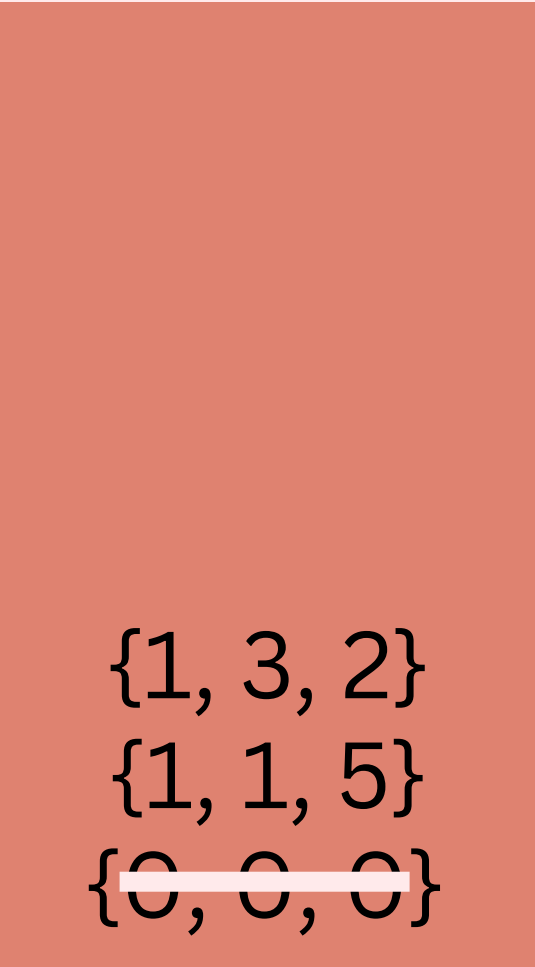
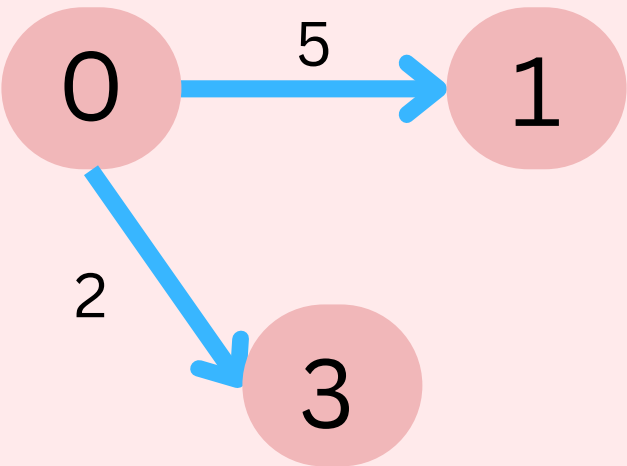
**Step 3:**

We will dequeue {0, 0, 0} initially. Since 0's neighbours are 1 and 3, the following condition will be examined:

*cost of the present node + the cost of travelling to its neighbour < the neighbor's cost.*

If this criterion is met, we will update the cost of the neighbour, which is calculated as follows: cost of the current node + the cost of travelling to its neighbour. Then we will increase stop value (here 0) by 1 and we will add updated {stops, node, cost} (in this case {1, 1, 5} and {1, 3, 2}) into the queue.

	0	1	2	3	4
cost[] =	0	5	$\infty$	2	$\infty$



{stops, node, cost}

**Queue**

## Step 4:

We will dequeue {1, 1, 5}. 1's neighbours are 1 and 4,

since  $\text{cost}[1] + 5 < \text{cost}[2]$  i.e.  $5 + 5 < \infty$

$\Rightarrow \text{cost}[2] = 5 + 5 = 10$

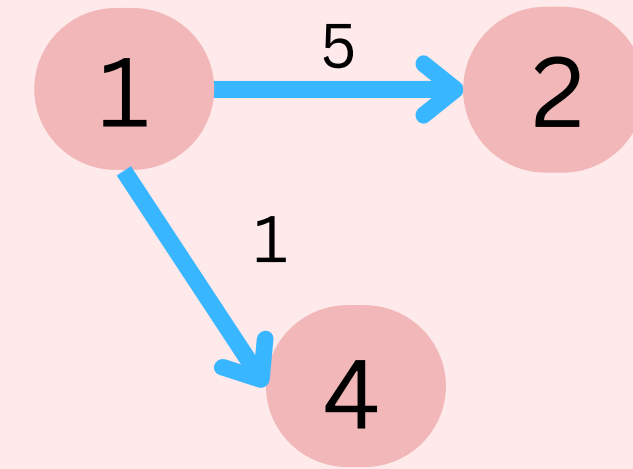
Then we will increase stop value (here 1) by 1 and we will add {2, 2, 10} into the queue.

$\text{cost}[1] + 1 < \text{cost}[4]$  i.e.  $5 + 1 < \infty$

$\Rightarrow \text{cost}[4] = 5 + 1 = 6$

Then we will increase stop value (here 1) by 1 and we will add {2, 4, 6} into the queue.

	0	1	2	3	4
cost[] =	0	5	10	2	6



{2, 4, 6}  
{2, 2, 10}  
{1, 3, 2}  
~~{1, 1, 5}~~  
~~{0, 0, 0}~~

{stops, node, cost}

**Queue**

### Step 5:

We will dequeue {1, 3, 2}. 3's neighbour is 1,  
since  $\text{cost}[3] + 2 < \text{cost}[1]$  i.e.  $2 + 2 < 5$

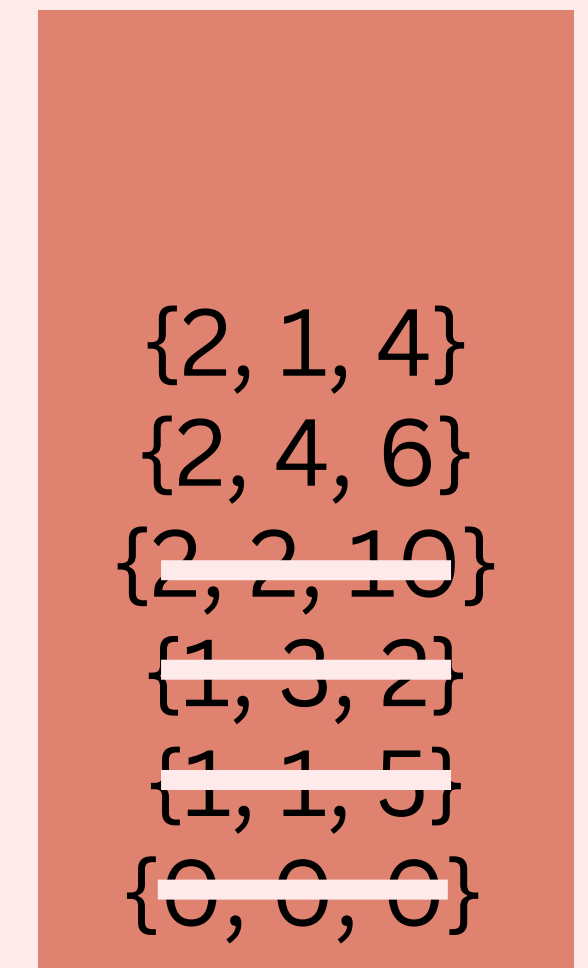
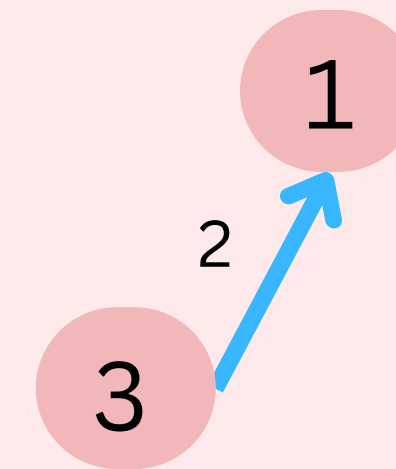
$\Rightarrow \text{cost}[1] = 2 + 2 = 4$

We will increase stop value (here 1) by 1 and  
we will add {2, 1, 4} into the queue.

### Step 6:

Dequeue {2, 2, 10}. 2 doesn't have neighbour.  
So, we don't need to do any changes.

	0	1	2	3	4
cost[] =	0	4	10	2	6



{stops, node, cost}

**Queue**



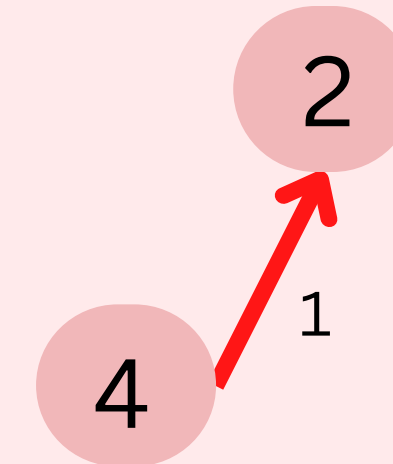
## Step 7:

Dequeue {2, 4, 6}. 4's neighbour is 2,

$\text{cost}[4] + 1 < \text{cost}[2]$  i.e.  $6 + 1 < 10$

$\Rightarrow \text{cost}[2] = 6 + 1 = 7 \Rightarrow \text{stop} = 2 + 1 = 3$

Add {3, 2, 7} into the queue.



## Step 8:

Dequeue {2, 1, 4}. 1's neighbours are 2 and 4.

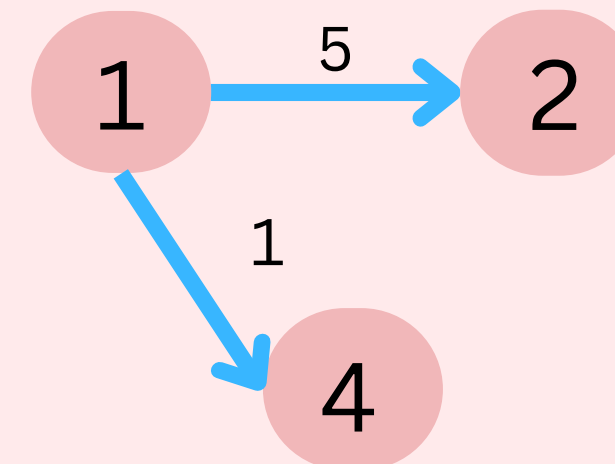
$\text{cost}[1] + 5 \nless \text{cost}[2]$  i.e.  $4 + 5 \nless 7$

So, we will not update  $\text{cost}[2]$ .

$\text{cost}[1] + 1 < \text{cost}[4]$  i.e.  $4 + 1 < 6$ .

$\Rightarrow \text{cost}[4] = 4 + 1 = 5 \Rightarrow \text{stop} = 2 + 1 = 3$ ;

Add {3, 4, 5} into the queue.



	0	1	2	3	4
cost[] =	0	4	7	2	5

{3, 4, 5}
{3, 2, 7}
<del>{2, 1, 4}</del>
<del>{2, 4, 6}</del>
<del>{2, 2, 10}</del>
<del>{1, 3, 2}</del>
<del>{1, 1, 5}</del>
<del>{0, 0, 0}</del>

{stops, node, cost}

**Queue**

## Step 9:

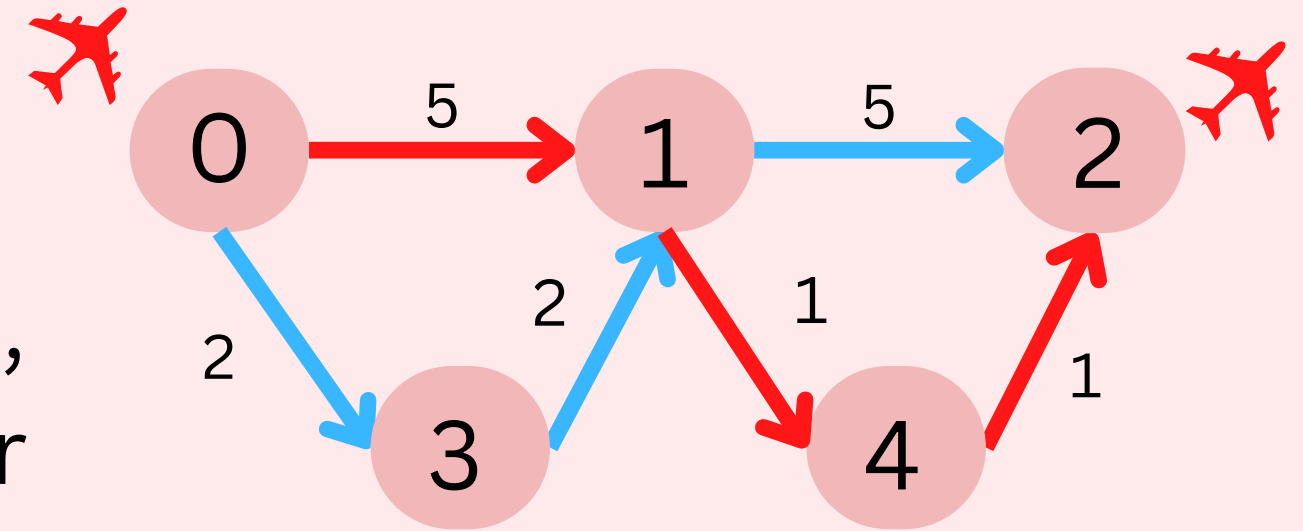
Dequeue {3, 2, 7}.

Now we have reached desired number of stops, even if we go further and found more cheaper alternative we have to increase stop by one, which we can't do since it is already at its maximum.

Thus,  $\text{cost}[2] = 7$  which is our final answer.

**Note:-** If the  $\text{cost}[\text{stop}] = \infty$  then there doesn't exist a path within  $k$  stops from source to destination.

	0	1	2	3	4
cost[] =	0	4	7	2	5



{3, 4, 5}  
~~{3, 2, 7}~~  
~~{2, 1, 4}~~  
~~{2, 4, 6}~~  
~~{2, 2, 10}~~  
~~{1, 3, 2}~~  
~~{1, 1, 5}~~  
~~{0, 0, 0}~~

{stops, node, cost}

**Queue**