# OPERATING SYSTEMS LAB REPORT

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# LAB 4

#### **GOAL**

Designing SJF and RR schedulers

- (1) compute various performance measures (turnaround time, waiting time, penalty ratio for each process and system averages and system throughput)
- (2) analyze the behavior of your schedulers. Your output should include results per process and the system's overall performance

# **SHORTEST JOB FIRST (SJF)**

# **Explanation of SJF**

It is a non-preemptive scheduling method which is used in OS to manage CPU time allocation to various processes. SJF operates on the principle that the CPU should be assigned to the process with the smallest anticipated execution duration.

The priority queue (pq\_cpu) is utilized for selecting the process with the shortest CPU burst time during CPU scheduling. The priority queue is implemented using a custom comparator class (Compare) which ensures that processes are ordered first by their CPU burst times and, in case of a tie, by their arrival times.

During the execution, the program continually selects processes from the priority queue for CPU execution.

Processes are added to the priority queue when they become available for execution based on their arrival times. If a process arrives while the CPU is idle, it is immediately added to the priority queue.

As a process runs on the CPU, its remaining CPU burst time is decremented. Upon completion of a CPU burst, processes are checked for any remaining IO burst times. If present, the process is enqueued into a separate queue (q\_io) to simulate IO operations.

If the process still has CPU burst time remaining, it is re-inserted into the priority queue to be selected for execution again.

According to SJF, the loop continues until both the CPU and IO queues are empty.

# **Expected job characteristics**

- > Non-preemptive: It does not allow for interruptions & a process will continue to execute until it is finished or blocked/suspended state which is gone for I/O.
- > Priority based algorithm: Since it uses cpu burst time as priority the process with the shortest burst time will be given highest priority.
- > Optimize Throughput: Since the shortest processes are executed first thus overall waiting time for processes is reduced.
- > Starvation: SJF may cause very long turn-around times or starvation.
- > Job completion: In SJF job completion time must be known earlier, but sometimes it is hard to predict.

# Test process data to bring out the suitability of SJF

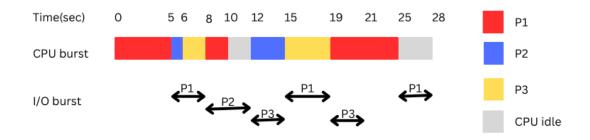
Since the scheduler is non-preemptive it is suitable for prioritizing jobs having short execution time and also they won't have to undergo starvation.

The time sharing will be according to the execution time.

A sample test case is given below: suiatable\_sif.dat

Process	Arrival time	CPU burst	I/O burst	CPU burst	I/O burst	CPU burst	I/O burst	Total CPU burst
P1	0	5	3	2	4	6	3	13
P2	2	1	4	3	-	-	-	4
Р3	5	2	3	4	2	-	-	6

# **Gantt Chart :**



Desired output on running above test case :

```
1 Process P1:
  Turnaround Time: 28
  Waiting Time: 0
  Penalty Ratio: 1.21739
  Completion Time: 28
  8 Process P2:
9 Turnaround Time: 13
10 Waiting Time: 3
  Penalty Ratio: 1.625
  Completion Time: 15
  15 Process P3:
16 Turnaround Time: 16
  Waiting Time: 1
18 Penalty Ratio: 1.45455
19 Completion Time: 21
  Average Turnaround Time: 19
23 Average Waiting Time: 1.33333
  Average Penalty Ratio: 1.43231
25 Throughput: 0.107143
```

#### **Conclusion:**

- SJF is suitable for P2 as it prioritizes the shortest job, leading to efficient CPU utilization and minimizing the turnaround time for P2.
- SJF is suitable for all processes due to short total CPU burst, reducing waiting time.
- SJF is suitable for all processes since it has almost similar penalty ratio for all.

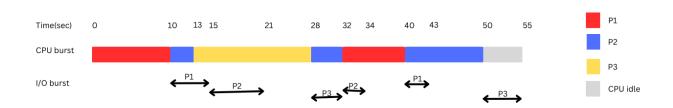
# Test process data to bring out the shortcomings of SJF

To highlight the shortcomings of the SJF code, let's create a test case that involves a mix of short and long processes with varying arrival times and burst times. This will help demonstrate how SJF may not always provide optimal results and may lead to increased waiting times for certain processes and starvation.

Below given a test case: shortcoming\_sjf.dat

Process	Arrival time	CPU burst	I/O burst	CPU burst	I/O burst	Total CPU burst
P1	0	10	5	8	3	18
P2	5	3	6	4	2	7
Р3	8	15	4	10	5	25

#### **Gantt Chart:**



#### Desired output on running above test case :

```
Process P1:
Turnaround Time: 43
Waiting Time: 17
Penalty Ratio: 1.65385
Completion Time: 43
Process P2:
Turnaround Time: 29
Waiting Time: 12
Penalty Ratio: 1.93333
Completion Time: 34
Process P3:
Turnaround Time: 47
Waiting Time: 13
Penalty Ratio: 1.38235
Completion Time: 55
Average Turnaround Time: 39.6667
Average Waiting Time: 14
Average Penalty Ratio: 1.65651
Throughput: 0.0545455
```

#### **Conclusion:**

- P3, arrives early and prevents shorter processes from getting CPU time correctly. In this case, P2 has to wait until P3 completes its execution, causing a delay.
- Since P3 is executing at the end it has remaining I/O thus causing the CPU to be idle.
- SJF may not always be suitable in scenarios where there is a mix of short(P2) and

long (P1,P3) processes, as it tends to favor the shortest job without considering the overall waiting time of processes.

# The analysis of the performance of SJF when run on the test cases

#### Process1.dat :

```
1 Process P1:
2 Turnaround Time: 1097
 Waiting Time: 676
4 Penalty Ratio: 2.6057
5 Completion Time: 1097
  8 Process P2:
9 Turnaround Time: 1436
10 Waiting Time: 1095
11 Penalty Ratio: 4.21114
12 Completion Time: 1438
15 Process P3:
16 Turnaround Time: 1084
17 Waiting Time: 793
18 Penalty Ratio: 3.72509
19 Completion Time: 1087
22 Process P4:
23 Turnaround Time: 573
24 Waiting Time: 372
25 Penalty Ratio: 2.85075
26 Completion Time: 577
```

```
1 Process P5:
2 Turnaround Time: 179
3 Waiting Time: 166
4 Penalty Ratio: 13.7692
  8 Process P6:
9 Turnaround Time: 99
10 Waiting Time: 94
11 Penalty Ratio: 19.8
12 Completion Time: 105
17 Waiting Time: 272
18 Penalty Ratio: 2.32683
  Completion Time: 487
23 Average Waiting Time: 495.429
25 Throughput: 0.00486787
```

#### Process2.dat:

```
Process P1:
3 Waiting Time: 0
  8 Process P2:
9 Turnaround Time: 83
11 Penalty Ratio: 6.38462
12 Completion Time: 84
  17 Waiting Time: 280
 Penalty Ratio: 2.39303
19 Completion Time: 487
22 Process P4:
24 Waiting Time: 156
25 Penalty Ratio: 13
30 Turnaround Time: 813
33 Completion Time: 837
36 Process P6:
38 Waiting Time: 151
39 Penalty Ratio: 12.6154
```

```
. .
4 Penalty Ratio: 2.74194
 Completion Time: 961
8 Process P8:
9 Turnaround Time: 147
 Woiting Time: 134
22 Process P10:
23 Turnaround Time: 142
24 Woiting Time: 129
29 Process Pll:
 Woiting Time: 112
32 Penalty Ratio: 9.61539
38 Woiting Time: 107
46 Penalty Ratio: 7.92308
```

```
. .
1 Process P14:
2 Turnaround Time: 80
3 Waiting Time: 67
4 Penalty Ratio: 6.15385
5 Completion Time: 120
9 Turnaround Time: 95
10 Waiting Time: 82
11 Penalty Ratio: 7.30769
12 Completion Time: 135
15 Process Pl6:
16 Turnaround Time: 75
17 Waiting Time: 62
18 Penalty Ratio: 5.76923
19 Completion Time: 117
22 Process P17:
23 Turnaround Time: 59
24 Waiting Time: 46
26 Penalty Ratio: 4.53846
26 Completion Time: 102
30 Turnaround Time: 54
31 Waiting Time: 41
32 Penalty Ratio: 4.15385
33 Completion Time: 99
36 Average Turnaround Time: 215.5
37 Average Waiting Time: 157.5
38 Average Penalty Ratio: 6.89117
39 Throughput: 0.0187305
```

#### Process3.dat:

```
.
1 Process P1:
2 Turnaround Time: 593
3 Waiting Time: 365
  Penalty Ratio: 2.60088
8 Process P2:
9 Turnaround Time: 1336
10 Waiting Time: 1042
11 Penalty Ratio: 4.54422
12 Completion Time: 1338
15 Process P3:
16 Turnaround Time: 1378
17 Waiting Time: 1074
18 Penalty Ratio: 4.53289
22 Process P4:
23 Turnoround Time: 315
24 Woiting Time: 147
25 Penalty Ratio: 1.875
26 Completion Time: 323
30 Turnaround Time: 586
31 Waiting Time: 417
32 Penalty Ratio: 3.46746
33 Completion Time: 598
36 Process P6:
37 Turnaround Time: 1799
38 Waiting Time: 1417
39 Penalty Ratio: 4.70942
```

```
. .
  Turnaround Time: 973
3 Waiting Time: 737
4 Penalty Ratio: 4.12288
  Completion Time: 1003
   9 Turnoround Time: 51
11 Penalty Ratio: 8.5
12 Completion Time: 86
16 Turnaround Time: 167
17 Woiting Time: 107
18 Penalty Ratio: 2,78333
19 Completion Time: 203
  23 Turnoround Time: 164
24 Waiting Time: 64
25 Penalty Ratio: 1.64
30 Turnground Time: 63
31 Woiting Time: 64
32 Penalty Ratio: 7
36 Process Pl2
37 Turnaround Time: 2083
39 Penalty Ratio: 5,35476
43 Average Turnaround Time: 792,333
44 Average Waiting Time: 596.917
45 Average Penalty Ratio: 4.2609
45 Throughput: 0 00565238
```

#### How to run code?

Open a terminal or command prompt and navigate to the directory where your *SJF.cpp* file is located.

suyas@Z-Sparrow:/mnt/c/Users/suyas/OneDrive/Documents/Minix3/LAB 4\$ g++ SJF.cpp -o sjf

Once the compilation is successful, you can run the compiled program. Our program expects a **command-line argument in the form of a test case file** (e.g., *test\_cases/process1.dat*).

suyas@Z-Sparrow:/mnt/c/Users/suyas/OneDrive/Documents/Minix3/LAB\_4\$ ./sjf test\_cases/process1.dat

The program will print the results to a file named *SJF\_output.txt* which is in the *Results* folder.

# **ROUND ROBIN (RR)**

# **Explanation of SJF**

It is a preemption-based approach that allocates CPU time to processes in a cyclic manner. This algorithm maintains a ready queue, and processes are scheduled for execution based on the specified time slice. If a process's CPU burst is not completed within the time slice, it is moved to the back of the queue to allow other processes to execute.

The code uses a queue to represent the ready queue and another queue for handling I/O bursts. Processes are scheduled in a time-sliced manner until all processes complete their execution.

The algorithm ensures fairness by giving each process a turn to execute, but it may lead to higher turnaround times for processes with longer burst times.

# **Expected job characteristics:**

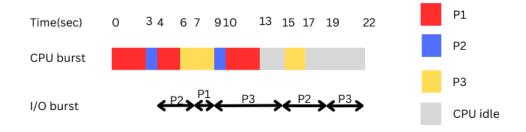
- > Preemptive: A running process can be interrupted before its full time slice is consumed, allowing other processes in the ready queue to get a chance to execute.
- > Starvation-Free: Round Robin is designed to be starvation-free, ensuring that all processes get a fair share of the CPU. No process is left waiting indefinitely, as each process gets a chance to execute within a fixed time SLICE.
- > Fairness: The algorithm provides fairness by ensuring that every process receives an equal share of CPU time.
- > Low throughput: The Round Robin algorithm might have lower throughput, especially when using a large time quantum. This happens because processes may not switch quickly, causing the CPU to not be fully used.
- > Context Switches: Round Robin scheduling involves frequent context switches, as processes are preempted after their time quantum expires.

# Test process data to bring out the suitability of RR

A sample test case is given below with <u>time slice</u> 3 sec: *suitable\_rr.dat* 

Process	Arrival time	CPU burst	I/O burst	CPU burst	I/O burst
P1	0	5	2	3	-
P2	2	1	3	1	4
Р3	4	3	6	2	3

#### **Gantt Chart:**



#### Desired output on running above test case :

```
Process Pl:
  Turnaround Time: 13
  Waiting Time: 2
  Penalty Ratio: 1.3
  Completion Time: 13
  8 Process P2:
9 Turnaround Time: 17
10 Waiting Time: 3
11 Penalty Ratio: 1.88889
12 Completion Time: 19
  15 Process P3:
16 Turnaround Time: 18
17 Waiting Time: 2
18 Penalty Ratio: 1.28571
19 Completion Time: 22
  Average Turnaround Time: 16
23 Average Waiting Time: 2.33333
24 Average Penalty Ratio: 1.49153
25 Throughput: 0.136364
```

#### **Conclusion:**

- Each process gets a time slice to execute its CPU burst, preventing any single process from occupying the CPU.
- This responsiveness is shown in the quick running of P1 and P2 on CPU, which have shorter initial CPU burst times.
- Shorter processes, like P1 and P2, get opportunities to execute within their time slices despite the presence of a longer process (P3).

# Test process data to bring out the shortcomings of RR:

Below is given a test example to show shortcomings of RR named *shortcoming\_rr.dat* with *time slice* 7.

Process	Arrival time	CPU burst	I/O burst	CPU burst	I/O burst
P1	0	10	5	8	4
P2	2	5	3	7	2
Р3	4	12	4	6	6

#### **Conclusion:**

- The small time quantum of 7 ms leads to frequent context switches, which can increase overhead.
- Processes are often interrupted before completing their CPU bursts, leading to longer waiting times.
- P1 experiences interruptions in its CPU burst due to the short time quantum. It gets interrupted after using only 7 sec of its initial 10 ms burst.
- P1 is frequently context-switched with other processes, leading to additional overhead and a fragmented execution pattern.

#### Desired output on running above test case :



# The analysis of the performance of SJF when run on the test cases

Process1.dat: timeslice 3

```
Turnaround Time: 1430
                      Waiting Time: 1009
                      Penalty Ratio: 3.39667
     8 Process P2:
  10 Waiting Time: 1014
 11 Penalty Matio: 3.97361
27 . Here is a real and the 
 38 Waiting Time: 27
44 Turnaround Time: 999
 50 Average Turnaround Time: 874.857
```

#### Process2.dat: timeslice 3

```
1 Process P1:
2 Turnaround Time: 8
3 Waiting Time: 3
4 Penalty Ratio: 1.6
6 Completion Time: 8
  8 Process P2:
9 Turnaround Time: 22
10 Waiting Time: 9
11 Penalty Ratio: 1.69231
12 Completion Time: 23
15 Process P3:
16 Turnaround Time: 711
17 Waiting Time: 510
18 Penalty Ratio: 3.53731
19 Completion Time: 717
22 Process P4:
23 Turnaround Time: 79
24 Waiting Time: 66
25 Penalty Ratio: 6.07692
26 Completion Time: 102
```

```
. .
1 Process P5:
2 Turnaround Time: 889
3 Waiting Time: 598
4 Penalty Ratio: 3.05498
5 Completion Time: 913
8 Process P6:
10 Waiting Time: 79
11 Penalty Ratio: 7.07692
12 Completion Time: 117
15 Process P7:
17 Waiting Time: 592
19 Completion Time: 959
22 Process P8:
23 Turnaround Time: 105
24 Waiting Time: 92
25 Penalty Ratio: 8.07692
29 Process P9:
30 Turnaround Time: 259
31 Waiting Time: 222
32 Penalty Ratio: 7
33 Completion Time: 287
```

```
1 Process P10:
2 Turnaround Time: 118
3 Waiting Time: 105
4 Penalty Ratio: 9.07692
5 Completion Time: 147
6 ****************
8 Process Pll:
9 Turnaround Time: 119
10 Waiting Time: 106
11 Penalty Ratio: 9.15385
12 Completion Time: 150
15 Process P12:
16 Turnaround Time: 123
17 Waiting Time: 110
18 Penalty Ratio: 9.46154
19 Completion Time: 156
22 Process P13:
23 Turnaround Time: 124
24 Waiting Time: 111
25 Penalty Ratio: 9.53846
26 Completion Time: 159
29 Process P14:
30 Turnaround Time: 122
31 Waiting Time: 109
32 Penalty Ratio: 9.38461
33 Completion Time: 162
```

```
. .
1 Process P15:
2 Turnaround Time: 128
3 Waiting Time: 115
4 Penalty Ratio: 9.84615
6 Completion Time: 168
8 Process Pl6:
9 Turnaround Time: 129
10 Waiting Time: 116
11 Penalty Ratio: 9.92308
12 Completion Time: 171
16 Process P17:
16 Turnaround Time: 134
17 Waiting Time: 121
18 Penalty Ratio: 10.3077
19 Completion Time: 177
22 Process P18:
23 Turnaround Time: 138
24 Waiting Time: 125
25 Penalty Ratio: 10.6154
26 Completion Time: 183
29 Average Turnaround Time: 235.167
30 Average Waiting Time: 177.167
31 Average Penalty Ratio: 7.11995
32 Throughput: 0.0187696
```

#### Process3.dat: timeslice 3

```
. .
1 Process P1:
4 Penalty Ratio: 7.01316
5 Completion Time: 1599
9 Turnground Time: 1922
10 Waiting Time: 1628
11 Penalty Ratio: 6.53742
12 Completion Time: 1924
15 Process P3:
16 Turnaround Time: 1888
17 Waiting Time: 1584
18 Penalty Ratio: 6.21053
19 Completion Time: 1893
22 Process P4:
23 Turnoround Time: 1355
24 Waiting Time: 1187
25 Penalty Ratio: 8.06548
26 Completion Time: 1363
29 Process P5:
30 Turnaround Time: 1324
31 Waiting Time: 1119
32 Penalty Ratio: 7.83432
  Completion Time: 1336
3A Process P61
37 Turnaround Time: 2069
38 Waiting Time: 1687
39 Penalty Ratio: 5.41623
```

```
. .
2 Turnaround Time: 1638
  Completion Time: 1668
  8 Process P8:
9 Turnaround Time: 58
10 Woiting Time: 62
11 Penalty Ratio: 9.66667
16 Turnaround Time: 163
18 Penalty Ratio: 2.71667
19 Completion Time: 199
22 Process P10:
24 Woiting Time: 786
29 Process Pll:
30 Turnaround Time: 100
32 Penalty Ratio: 11.1111
36 Process Pl2:
38 Woiting Time: 1664
  Penalty Ratio: 5.25193
43 Average Turnaround Time: 1253.75
44 Average Waiting Time: 1051.08
```

# How to run code?

Open a terminal and navigate to the directory where your *rr.cpp* file is located.

```
suyas@Z-Sparrow:/mnt/c/Users/suyas/OneDrive/Documents/Minix3/LAB_4$ g++ RR.cpp -o rr
```

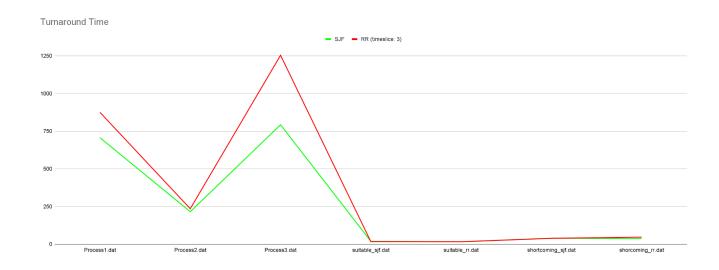
Once the compilation is successful, you can run the compiled program. Our program expects **two command-line arguments in the form of a test case file and followed by timeslice** (e.g., *test\_cases/process1.dat 3*).

```
suyas@Z-Sparrow:/mnt/c/Users/suyas/OneDrive/Documents/Minix3/LAB 4$ ./rr test cases/process1.dat 3
```

The program will print the results to a file named *RR\_output.txt* which is in the *Results* folder.

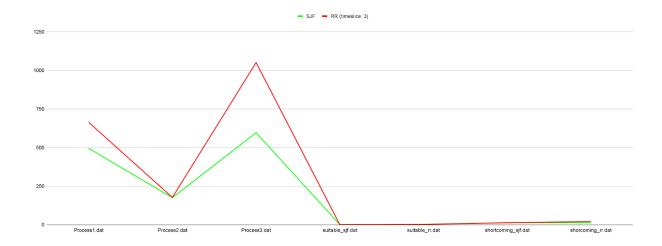
# Graphs capturing the variations in performance :

Turnaround Time SJF V/S RR



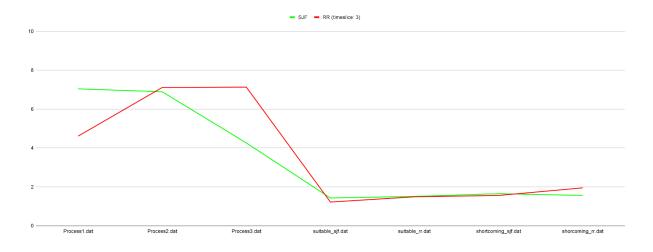
The green line for SJF generally shows lower turnaround times compared to the red line for RR, indicating better performance of SJF.

### Waiting time SJF v/s RR



The SJF algorithm performs best on "Process3.dat", while the RR algorithm performs consistently across all processes. Both algorithms perform equally well on the processes after "suitable\_sjf/rr.dat".

# Penalty ratio SJF v/s RR



RR shows more fluctuation in penalty ratio compared to SJF as processes change. However, both algorithms show a decline in penalty ratio after Process3.dat.

#### Variation of time slice of RR for process1.dat

Turnaround time

# RR Turnaround Time Timeslice=1 Timeslice=3 Timeslice = 5 Timeslice=1 Timeslice=5

Process3.dat

suitable\_rr.dat

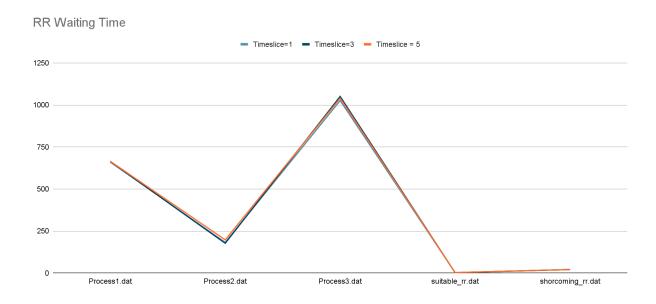
shorcoming\_rr.dat

Process2.dat

# Waiting Time

Process1.dat

500



# Penalty Ratio

# Penalty Ratio

