,	,	Date	
		Name - Suyash Jan	
\$		Course - BTech CSE	
6	1	Poll No. → 2301010249	i i
6	1		
		Operating System Assignment 01	- Y
	<u> </u>		
	01.		
9	Ano - Eve	n with solvenced hardware, OS resource management (CPU, men	is essential
4	for	resource monagement (CPU, men vides abstraction, multitosking;	L. (O/T, juan
6	mo	wides abstraction, multitasking;	security and
100	lus	sure programs van sun safeli	1 and
_	l'us off	serve programs van sun safeli	q
	J		•
	02.		. \
4	$A_{\text{max}} \rightarrow \mathcal{A}_{\text{max}}$	A real-time, embedded OS is be	st because the
5	ب طلین	ice, needs quick and predicte Le pouver usage, and small	ible response,
4		2 power usage, and small	foolpoint:
	Q 3	<u> </u>	•
	Au -8	A 3 de la constante de la cons	ates doss
· ·		and a train add a superland	May 19+19
		or layered aller better rea	or berlomance
	1		S / P G S IVILLOUS S
	04 -		
	Ans - 9	les Structure impacts sufomance	maintainability.
**		and fault isolation. For en	ample micro
5	k	ernal is secure but slow	rer monolithic
*	3	s faster but less reliable.	
)	J	
	05.	21	
	Aus -> i	y PCB stores registers, PC and sto	te - errord
	-	show misinitialized values	-
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(iii) Content surtch = saving current process state, updating PCB, loading new state. (iii) Mid-execution I/O allocation usually needs non-blocking cell so process continues numing. Part - B Fina of Total content switching time,
save state = 2ms

Load State = 3 ms Schedular overhead = 1 mg Total time - 2+3+1 = 6ms. * Content switching us some overhead (no useful work us done dewing this time.). * Higher dwitching time reduces CPU efficiency as more time is spent switching than encuting process. Aus. Execution time estimate: In Ideal conditions (perfect proallelism, no ownhead). Tuntti = $\frac{1}{n}$ = $\frac{40}{n}$ second.

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Dave.	 	 	 •	•	•	•	•	•	•	•	•

Bingle-thread = 40 sec.
With 2 threads under ideal cond = 40 = 20 sec.

* Multithreading allows overlapping of I/o and computations untilizing multiple cores and reducing waiting time though in practice overhead reduces berfeet gains.

Proces: P1 P2 B Py

But time: 5 3 8 6

a FCFS

Process	Arral	Bust	Completion	Waiting	TAT
j- 16CVS	time (AT)	Jane (Bt)	Time (CT)		i a
P	•	5	5	5-5 = 0	5-0=5
P ₂	0	3	8	8-3=5	8-0=8
P ₃	0	8	16	16-8=8	16-0=16
P.,	0	6	2.2	22-6= 16	22-0=22
٧ -					4

WT = Turn around - Burst (TAT-BT)

Grant Chart

P. P2 P3 P4

0 5 8 16 2

C1(P,) CT(P2) CT(P3)

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					Date.	**********		
Ava	water	ra fin	10. = (0	+5+8+16	1/4 = 7	25 M		
Ava	Avg waiting time = (0+5+8+16)/4 = 7.25 Mg							
	Avg turnaround time = (5+8+16+22)/4 = 12.75 ms							
(b) N	m- me	motive.	STF					
	V. 100	1 1 1 1				i is 🐞		
Process	AT	BT	CT	WT .	TAT			
Pi	Ο.	5	· 8	8-5=3		8=0		
P2	0	3	3	3-3-0	3=	3-0		
p_3	0	8	22	22-8=1		= 22-0		
Pu	6	6	14	14-6=		14-0		
/			, ,					
Grant	Chart,			1	è			
		P2 P	, P4	P ₂		4 · 2 / .		
		3	8 1	4 22				
• ,		G	(,		J)	7		
		CT (P2) Cz(P	(F4)	G(B)			
	Ana	Vaiting	time	= (3+0+14	(+8)/. =	6,25 mg		
	A	twen an	round		14	1		
	1	Twee an	1 =	(8+3+	22+14)/1	y = 11.75 my		
			-9.1	3.		4		
(c)					. (1 -)			
R	mid	Robin	Quande	m = 4ms)				
			1.		3 A . F T			
Process	AT	BT (CT	WT	TAT	1 × 1.		
Pi			16	16-5=11	16-0=16			
Pe	0	5	7	7-3=4	7-0=7	(transport)		
P3	0	8	20	20-8=12	20-0=20			
Py	0	6	22	22-6=16	22-0=23			
		1	,	1/3				
			•					

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brasit Chart	ì
	6
P1 P2 P3 P4 P1 P3 P4	ч
0 4 7 11 15 16 20 22	
\mathcal{L}	
$C_1(P_2)$ $C_7(P_1)$ $C_7(P_4)$	
The first of the f	
Ang waiting time = (11+4+12+16) /4 = 10.75 ms	
Ang turn around = (16+7+20+22)/4 = 16.25 ms	
) 19	
=> SJF gwes lowest writing and tumaround times while RR improves, fairness for interest	
times, while RR improves, fairness for interest	ine
itakes.	
Ou J.	
Aus (i) Cloud migration:	
for a Virtualized cloud, a microkernel is suitable, since it separates core services and provides better security and scalability. Vintu machies (VMs) add isolations by running nultiple. Os instance, on the same hands They also suffort resource shawing, load belowed and living migration.	
suitable since it separates core services and	
provides better security and scalability. Vintu	al
machies (vMs) add aslations by running	
multiple OS instance on the same harde	vore
They also support resource shaving, load balon	<u>ve</u>
and ling mignation.	
0	
(iii) Smout home system:	
Her many 10t devices run together. Th	e
OS uses porority scheduling and TPC do	
Her many 10t devices run together. The OS uses paiority scheduling and IPC do urgents tasks like intrusing detection	get

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CPU mans ed Estaly while lour - Driguity sched	luling
and IPC so rugert takes like int	rusion
detection get CPU immediately, which	low-
browity tasks like lighting wait. Algorith	my such
Os. ODE (Forliest,) Dealling first) of	RMS
(Rote Monotonic Scheduling) can ensure C tasks meets their deadhires without	rut1cal
tasks meets their deadlines without	Masuling -
background process.	
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