

RISC and CISC Instruction Sets

- One of the most important characteristics that distinguish different computers is the nature of their instructions
- Two approaches in the design of instruction sets for modern computers
 - Reduced Instruction Set Computers(RISC)
 - Complex Instruction Set Computer(CISC)

RISC

- The restriction that each instruction must fit into a **single word** that reduces the complexity and the **number of different types of instructions** that may be included in the instruction set of a computer.
- Such computers are called **Reduced Instruction Set Computers (RISC)**
- **Example-CDC 6600 (1964), ARM**

RISC Features

- Simple addressing modes
- All instructions fitting in a single word
- Fewer instructions in the instruction set, as a consequence of simple addressing modes
- Arithmetic and logic operations that can be performed only on operands in processor registers

RISC Features

- **Load/store architecture** that does not allow direct transfers from one memory location to another; such transfers must take place via a processor register
- Simple instructions that are conducive to **fast execution** by the processing unit using techniques such as pipelining
- Programs that tend to be **larger in size**, because of more, but simpler instructions are needed to perform complex tasks

RISC Instructions

- Memory operands are accessed only using Load and Store instructions
- Example: To execute $C=A+B$, we use three registers R2,R3,R4 ,the four instructions required are:

Load R2, A

Load R3, B

Add R4, R2, R3

Store R4, C

RISC Instruction

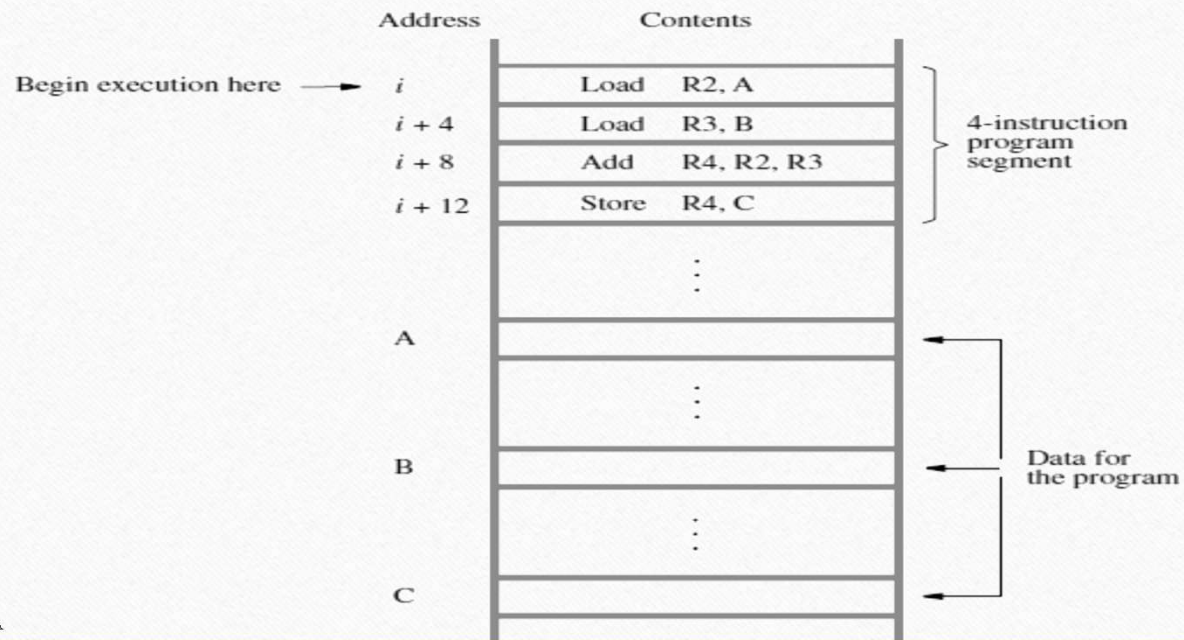


Figure: Instruction Execution and Straight-line Sequencing

RISC-Branching

- Consider the task of adding a list of n numbers
- Two Approaches
- 1) The addresses of the memory locations containing the n numbers are symbolically given as NUM1, NUM2, . . . , NUM n , and separate Load and Add instructions are used to add each number to the contents of register R2
- After all the numbers have been added, the result is placed in memory location SUM.

i	Load	R2, NUM1
$i + 4$	Load	R3, NUM2
$i + 8$	Add	R2, R2, R3
$i + 12$	Load	R3, NUM3
$i + 16$	Add	R2, R2, R3
		⋮
$i + 8n - 12$	Load	R3, NUM n
$i + 8n - 8$	Add	R2, R2, R3
$i + 8n - 4$	Store	R2, SUM
		⋮
SUM		
NUM1		
NUM2		
		⋮
NUM n		

Figure: Adding list
of numbers
without branch

RISC-Branching

- 2) It is possible to implement a **program loop** in which the instructions read the next number in the list and add it to the current sum.
- To **add all numbers**, the loop has to be executed as many times as there are numbers in the list

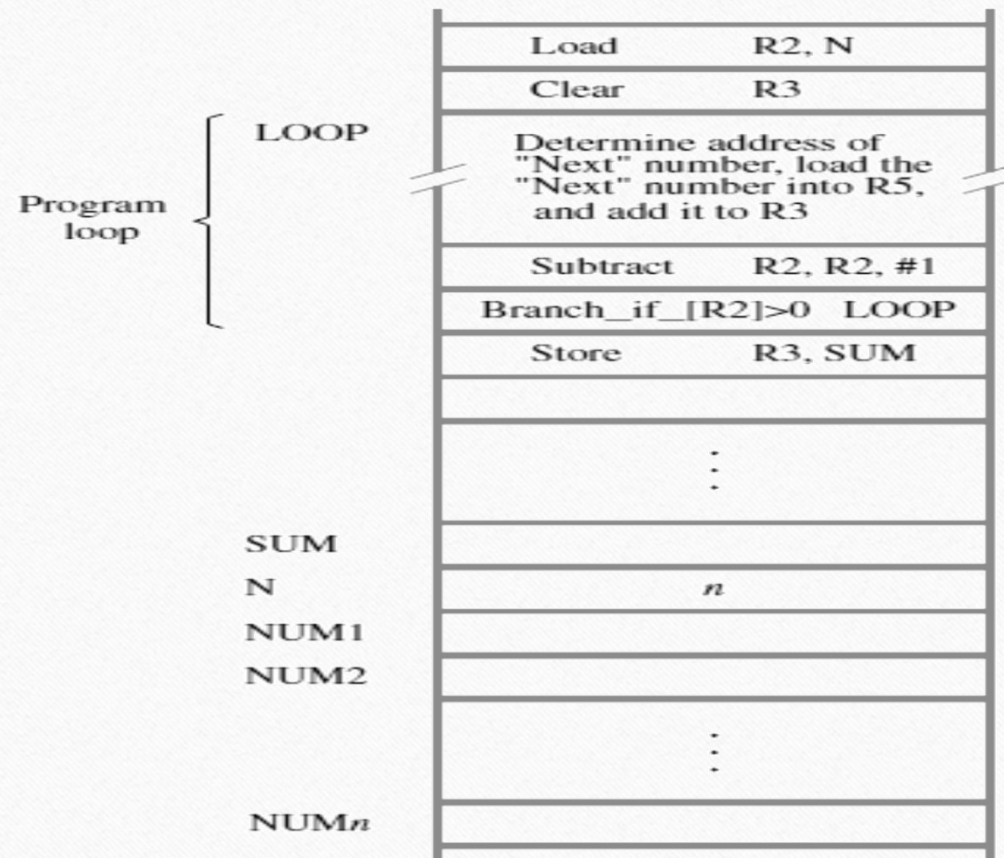


Figure: Adding list of numbers with branch

RISC –Addressing Mode

Name	Assembler syntax	Addressing function
Immediate	#Value	Operand = Value
Register	R_i	$EA = R_i$
Absolute	LOC	$EA = LOC$
Register indirect	(R_i)	$EA = [R_i]$
Index	$X(R_i)$	$EA = [R_i] + X$
Base with index	(R_i, R_j)	$EA = [R_i] + [R_j]$

EA = effective address

Value = a signed number

X = index value

	Load	R2, N	Load the size of the list.
	Clear	R3	Initialize sum to 0.
	Move	R4, #NUM1	Get address of the first number.
LOOP:	Load	R5, (R4)	Get the next number.
	Add	R3, R3, R5	Add this number to sum.
	Add	R4, R4, #4	Increment the pointer to the list.
	Subtract	R2, R2, #1	Decrement the counter.
	Branch_if_[R2]>0	LOOP	Branch back if not finished.
	Store	R3, SUM	Store the final sum.

Figure: Use of Indirect address in adding list of N numbers

CISC(Complex Instruction Set Computer)

- Traditional approach
- Main features :
- A large number of instructions—typically from 100 to 250 instructions.
- Some instructions that perform specialized tasks and are used infrequently
- A large variety of addressing modes—typically from 5 to 20 different modes(R-R,R-M,M-M,Indexed,Indirect)
- Variable-length instruction formats

CISC(Complex Instruction Set Computer)

- Ease of mapping high-level language statement to machine code is possible in CISC, but instruction decoding and control unit design are much more complex
- In the variable length instruction, pipeline implementation will also be quite complex
- Example
 - IBM 360/370 (1960-70)
 - VAX-11/780 (1970-80)
 - Intel x86/Pentium (1985-Present)

CISC(Complex Instruction Set Computer)

- CISC instruction sets are not constrained to the load/store architecture, in which arithmetic and logic operations can be performed only on operands that are in processor registers
- Instructions do not necessarily have to fit into a single word. Some instructions may occupy a single word, but others may span multiple words
- Instructions in modern CISC processors typically do not use a three-address format

CISC(Complex Instruction Set Computer)

- Most arithmetic and logic instructions use the two-address format:

Operation destination, source

- An Add instruction of this type is Add B, A
which performs the operation $B \leftarrow [A] + [B]$ on memory operands

CISC(Complex Instruction Set Computer)

- Consider again the task of adding two numbers $C = A + B$ where all three operands may be in memory locations
- Obviously, this cannot be done with a single two-address instruction. The task can be performed by using another two-address instruction that copies the contents of one memory location into another
- Such an instruction is Move C, B which performs the operation $C \leftarrow [B]$, leaving the contents of location B unchanged

CISC(Complex Instruction Set Computer)

- The operation $C \leftarrow [A] + [B]$ can now be performed by the two-instruction sequence
Move C, B
Add C, A
- Observe that by using this sequence of instructions the contents of neither A nor B locations are overwritten.

CISC(Complex Instruction Set Computer)

- In some CISC processors one operand may be in the memory but the other must be in a register.
- In this case, the instruction sequence for the required task would be

Move Ri, A

Add Ri, B

Move C, Ri

CISC-Addressing Modes

- Most CISC processors have all of the five basic addressing modes —Immediate, Register, Absolute, Indirect, and Index
- Three additional addressing modes are often found in CISC processors
- *Autoincrement and Autodecrement Modes* - There are two modes that are particularly convenient for accessing data items in successive locations in the memory and for implementation of stacks.
- *Relative mode*—The effective address is determined by the Index mode using the program counter in place of the general-purpose register R_i .

CISC-Addressing Modes

- Autoincrement Mode:
- The effective address of the operand is the contents of a register specified in the instruction
- After accessing the operand, the contents of this register are automatically incremented to point to the next operand in memory
- Autoincrement mode is written as $(Ri)+$

CISC-Addressing Modes

- We denote the Autoincrement mode by putting the specified register in parentheses, to show that the contents of the register are used as the effective address, followed by a plus sign to indicate that these contents are to be incremented after the operand is accessed
- Computers that have the Autoincrement mode automatically increment the contents of the register by a value that corresponds to the size of the accessed operand
- Thus, the increment is 1 for byte-sized operands, 2 for 16-bit operands, and 4 for 32-bit operands.

CISC-Addressing Modes

- **Autodecrement Mode:** The contents of a register specified in the instruction are first automatically decremented and are then used as the effective address of the operand.
- We denote the Autodecrement mode by putting the specified register in parentheses, preceded by a minus sign to indicate that the contents of the register are to be decremented before being used as the effective address
- Thus, we write $-(R_i)$
- In this mode, operands are accessed in descending address order

CISC-Addressing Modes

- Example: Main reason for autoincrement and autodecrement mode is to implement a stack structure.
- Instead of needing two instructions to push a new item on the stack
 Subtract SP, #4
 Move (SP), NEWITEM we can use just one instruction `Move -(SP), NEWITEM`
- Similarly, instead of needing two instructions to pop an item from the stack
 - `Move ITEM, (SP)`
 - `Add SP, #4`, we can use just have `Move ITEM, (SP)+`

CISC-Addressing Modes

- **Relative Mode:**
- Some computers have a version of this mode in which the program counter, **PC**, is used instead of a general-purpose register
- Then, **X(PC)** can be used to address a memory location that is X bytes away from the location presently pointed to by the program counter
- Since the addressed location is identified *relative* to the program counter, which always identifies the current execution point in a program, the name Relative mode is associated with this type of addressing.

CISC-Example

	Move	R2, N	Load the size of the list.
	Clear	R3	Initialize sum to 0.
	Move	R4, #NUM1	Load address of the first number.
LOOP:	Add	R3, (R4)+	Add the next number to sum.
	Subtract	R2, #1	Decrement the counter.
	Branch>0	LOOP	Loop back if not finished.
	Move	SUM, R3	Store the final sum.