

CS 381 Homework 3 – Syntax

Submit a pdf for problems 1 – 4 and a Haskell *.hs file for problem 5.

1. Using the grammar below, show a parse tree and a leftmost derivation for the sentence

$A = B * (C + A)$

.
 $\langle \text{assign} \rangle \rightarrow \langle \text{id} \rangle = \langle \text{expr} \rangle$
 $\langle \text{expr} \rangle \rightarrow \langle \text{expr} \rangle * \langle \text{term} \rangle$
 $\quad \quad \quad | \langle \text{term} \rangle$
 $\langle \text{term} \rangle \rightarrow \langle \text{factor} \rangle + \langle \text{term} \rangle \mid \langle \text{factor} \rangle - \langle \text{term} \rangle$
 $\quad \quad \quad | \langle \text{factor} \rangle$
 $\langle \text{factor} \rangle \rightarrow (\langle \text{expr} \rangle)$
 $\quad \quad \quad | \langle \text{id} \rangle$
 $\langle \text{id} \rangle \rightarrow A \mid B \mid C$

2. Rewrite the following BNF to add the prefix ++ and -- unary operators of Java.

.
 $\langle \text{assign} \rangle \rightarrow \langle \text{id} \rangle = \langle \text{expr} \rangle$
 $\langle \text{expr} \rangle \rightarrow \langle \text{expr} \rangle * \langle \text{term} \rangle$
 $\quad \quad \quad | \langle \text{term} \rangle$
 $\langle \text{term} \rangle \rightarrow \langle \text{factor} \rangle + \langle \text{term} \rangle \mid \langle \text{factor} \rangle - \langle \text{term} \rangle$
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 $\langle \text{factor} \rangle \rightarrow (\langle \text{expr} \rangle)$
 $\quad \quad \quad | \langle \text{id} \rangle$
 $\langle \text{id} \rangle \rightarrow A \mid B \mid C$

3. Show that the following grammar is ambiguous

.
 $\langle \text{compare} \rangle \rightarrow \langle \text{boolexpr} \rangle == \langle \text{boolexpr} \rangle$
 $\langle \text{boolexpr} \rangle \rightarrow \langle \text{boolexpr} \rangle \text{ AND } \langle \text{boolexpr} \rangle$
 $\quad \quad \quad | \langle \text{boolexpr} \rangle \text{ OR } \langle \text{boolexpr} \rangle$
 $\quad \quad \quad | \langle \text{bool} \rangle$
 $\quad \quad \quad | \text{NOT } \langle \text{bool} \rangle$
 $\langle \text{bool} \rangle \rightarrow \langle \text{boolvalue} \rangle \mid \langle \text{boolvar} \rangle$
 $\langle \text{boolvalue} \rangle \rightarrow \text{True} \mid \text{False} \mid 0 \mid 1$
 $\langle \text{boolvar} \rangle \rightarrow \mathbf{u} \mid \mathbf{v} \mid \mathbf{w}$

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4. Write a grammar G for the language L consisting of strings of 0's and 1's that are the binary representation of odd integers greater than 4. For example $11 \notin L$, $101 \in L$, $110 \notin L$. Draw parse trees for the strings 1011 and 1101

5. Below is the EBNF grammar for the animal sentence language

```
<sentence> -> <noun> <verb> [<noun>]
           | <sentence> `and` <sentence>

<noun>      -> <adj> <noun> | <noun> `and` <noun>
           | `cats` | `dogs` | `ducks` | `bunnies`

<verb>      -> `chase` | `cuddle` | `hug` | `scare`
<adj>       -> `silly` | `small` | `old` | `happy`
```

Note: the nonterminals are in < > and the terminals are in ` `.

Using the animal.hs template provided.

- Define the abstract syntax for the animal language as a Haskell data type.
- Provide “pretty printing” functions for the sentences in the language.
- Provide functions to build a sentence.
- Write a function `isNice` to determine if a sentence only contains the verbs `hug` and `cuddle`.
- Write a function to build a sentence that is a conjunction of other sentences.
- Write a function `wordCount` that computes the number of words in a sentence

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1. Using the grammar below, show a parse tree and a leftmost derivation for the sentence

$A = B * (C + A)$

$\langle \text{assign} \rangle \rightarrow \langle \text{id} \rangle = \langle \text{expr} \rangle$

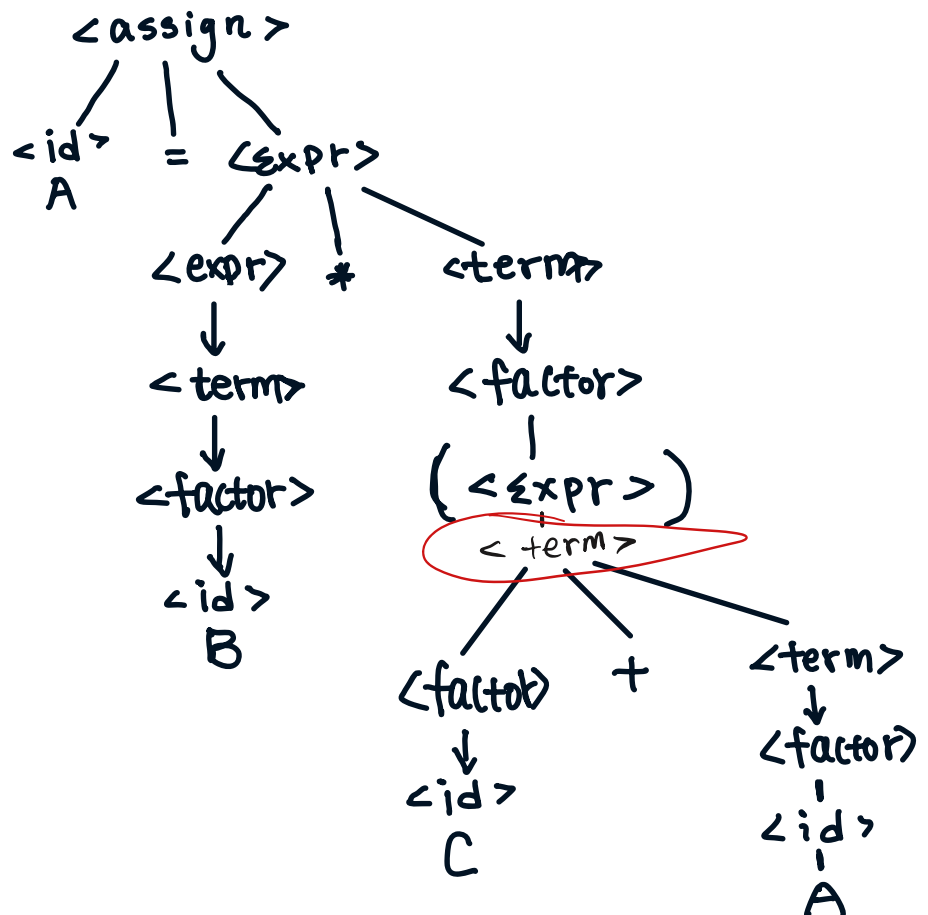
$\langle \text{expr} \rangle \rightarrow \langle \text{expr} \rangle * \langle \text{term} \rangle$
 $\quad \quad \quad | \langle \text{term} \rangle$

$\langle \text{term} \rangle \rightarrow \langle \text{factor} \rangle + \langle \text{term} \rangle \mid \langle \text{factor} \rangle - \langle \text{term} \rangle$
 $\quad \quad \quad | \langle \text{factor} \rangle$

$\langle \text{factor} \rangle \rightarrow (\langle \text{expr} \rangle)$
 $\quad \quad \quad | \langle \text{id} \rangle$

$\langle \text{id} \rangle \rightarrow A \mid B \mid C$

Parse tree:



Leftmost derivation:

$$\text{assign} \Rightarrow \langle \text{id} \rangle = \langle \text{expr} \rangle$$

$$\Rightarrow A = \langle \text{expr} \rangle$$

$$\Rightarrow A = \langle \text{expr} \rangle * \langle \text{term} \rangle$$

$$\Rightarrow A = \langle \text{term} \rangle * \langle \text{factor} \rangle$$

$$\Rightarrow A = \langle \text{factor} \rangle * (\langle \text{expr} \rangle)$$

$A = \langle \text{factor} \rangle * \langle \text{term} \rangle$

$$\Rightarrow A = \langle \text{id} \rangle * (\langle \text{factor} \rangle + \langle \text{term} \rangle)$$

$$\Rightarrow A = B * (\langle \text{id} \rangle + \langle \text{factor} \rangle)$$

$$\Rightarrow A = B * (C + \langle \text{id} \rangle)$$

$$\Rightarrow A = B * (C + A)$$

2. Rewrite the following BNF to add the prefix ++ and -- unary operators of Java.

$\langle \text{assign} \rangle \rightarrow \langle \text{id} \rangle = \langle \text{expr} \rangle$
 $\langle \text{expr} \rangle \rightarrow \langle \text{expr} \rangle * \langle \text{term} \rangle$
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 $\langle \text{factor} \rangle \rightarrow (\langle \text{expr} \rangle)$
 $\quad \quad \quad | \langle \text{id} \rangle$
 $\langle \text{id} \rangle \rightarrow A | B | C$

added the prefix and -- unary operators :

$\langle \text{assign} \rangle \rightarrow \langle \text{id} \rangle = \langle \text{expr} \rangle$
 $\langle \text{expr} \rangle \rightarrow \langle \text{expr} \rangle * \langle \text{term} \rangle | \langle \text{term} \rangle$
 $\langle \text{term} \rangle \rightarrow \langle \text{factor} \rangle + \langle \text{term} \rangle | \langle \text{factor} \rangle - \langle \text{term} \rangle$
 $\quad \quad \quad | \langle \text{factor} \rangle$
 $\langle \text{factor} \rangle \rightarrow (\langle \text{expr} \rangle)$
 $\quad \quad \quad | ++ \langle \text{id} \rangle | -- \langle \text{id} \rangle | \langle \text{id} \rangle$
 $\langle \text{id} \rangle \rightarrow A | B | C$

I add prefix ++ and -- in the front of $\langle \text{id} \rangle$, because according to the given grammar, $\langle \text{id} \rangle$ is terminals which means $\langle \text{id} \rangle$ are variables, according to Java, we can add ++ and -- before a variable.

3. Show that the following grammar is ambiguous

$\langle \text{compare} \rangle \rightarrow \langle \text{boolexpr} \rangle == \langle \text{boolexpr} \rangle$

$\langle \text{boolexpr} \rangle \rightarrow \langle \text{boolexpr} \rangle \text{ AND } \langle \text{boolexpr} \rangle$
| $\langle \text{boolexpr} \rangle \text{ OR } \langle \text{boolexpr} \rangle$
| $\langle \text{bool} \rangle$
| $\text{NOT } \langle \text{bool} \rangle$

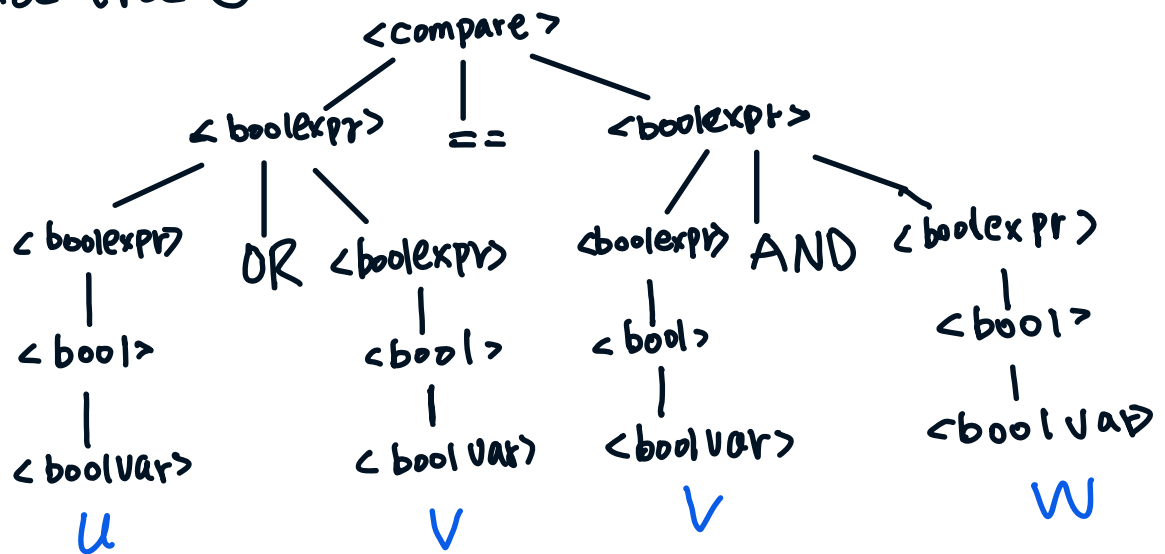
$\langle \text{bool} \rangle \rightarrow \langle \text{boolvalue} \rangle | \langle \text{boolvar} \rangle$

$\langle \text{boolvalue} \rangle \rightarrow \text{True} | \text{False} | 0 | 1$

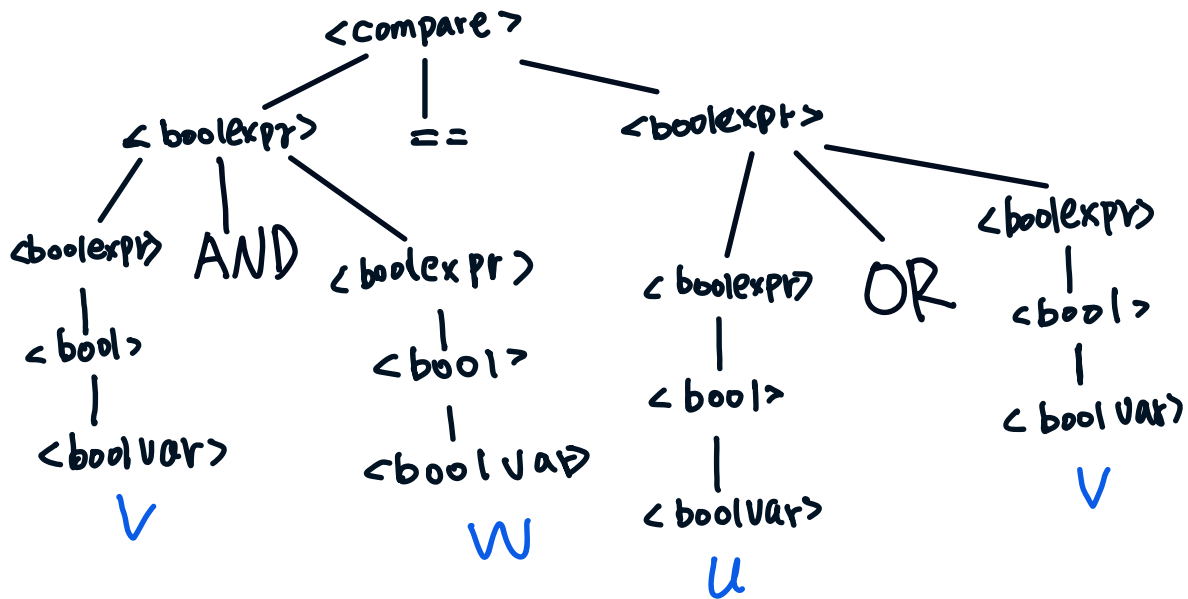
$\langle \text{boolvar} \rangle \rightarrow u | v | w$

Let me generate two parse trees for the sentence : u OR v == v AND w

Parse tree ① :



Parse tree ② :



The first parse tree represents the sentence
U OR V == V AND W and the second
 parse tree represents sentence

V AND W == U OR V. Therefore, the
 grammar is ambiguous, Because the same
 input sentence could be parsed in two
 different ways.

4. Write a grammar G for the language L consisting of strings of 0's and 1's that are the binary representation of odd integers greater than 4. For example $11 \notin L$, $101 \in L$, $110 \notin L$. Draw parse trees for the strings 1011 and 1101

Let $G = (\{V\}, \{0,1\}, P, S)$, where S is the start symbol, $V = \{S, A\}$ is set variable, $\{0,1\}$ is set of terminals, and P is the set of production rules defined as :

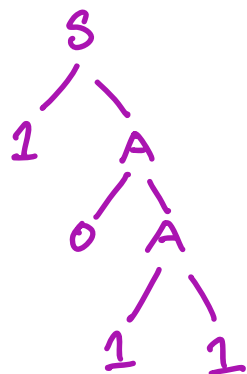
$$S \rightarrow 1A$$

$$A \rightarrow 0A$$

$$A \rightarrow 1$$

This grammar generates binary representation of odd integers greater than 4, where the first character is always 1, followed by zero or more 1's and 0's alternating, ending with a 0.

parse tree for the string '1011'



parse tree for the string '1101'

