Lecture 1: introduction to multithreading

concurrency, parallelism, agents, threads, scheduler, threading models, interleaving execution, Amdahl's law, race condition, data race, deadlock, wait-for graph

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https://github.com/Svazars/parallel-programming/blob/main/slides/pdf/l1.pdf

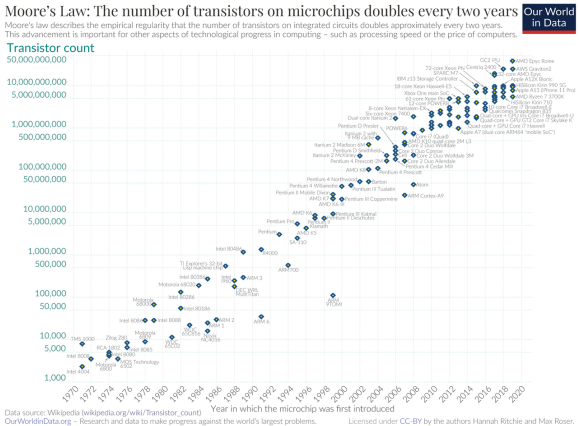
Lecture plan

1. Motivation
2. Concurrency and parallelism
3. Scheduler
4. Threads and processes
5. Multitasking and threading models
6. Concurrent problems
   - Embarrassingly parallel problems and Amdahl's law
   - Race condition
   - Data race
   - Visibility
   - Deadlock
   - Priority inversion
7. Summary

Motivation

Moore's law

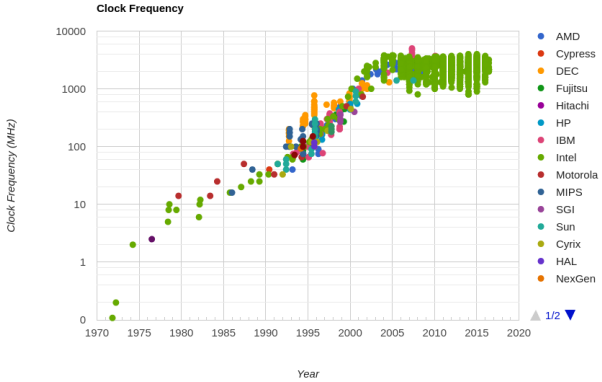
https://en.wikipedia.org/wiki/Moore%27s_law#/media/File:Moore's_Law_Transistor_Count_1970-2020.png



Motivation

Clock frequency

http://cpudb.stanford.edu/visualize/clock_frequency.html



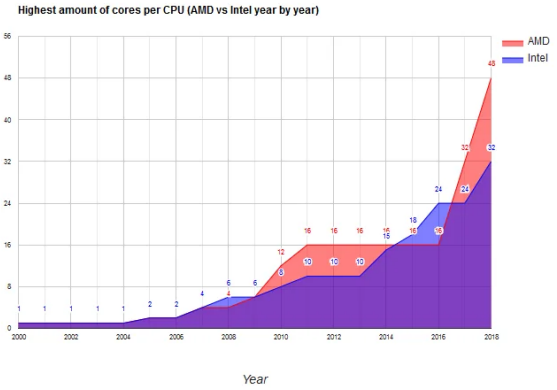
Question time

Question: Where are all these transistors?

Motivation

Core per CPU

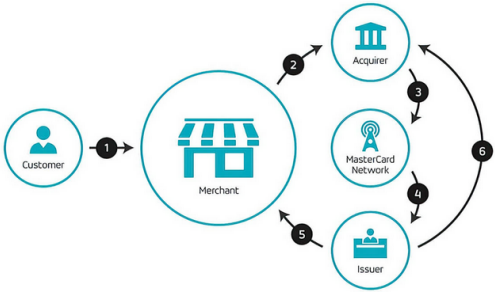
https://www.reddit.com/r/Amd/comments/6cu5ss/highest_amount_of_cores_per_cpu_amd_vs_intel_year



Motivation

Multi-agent systems

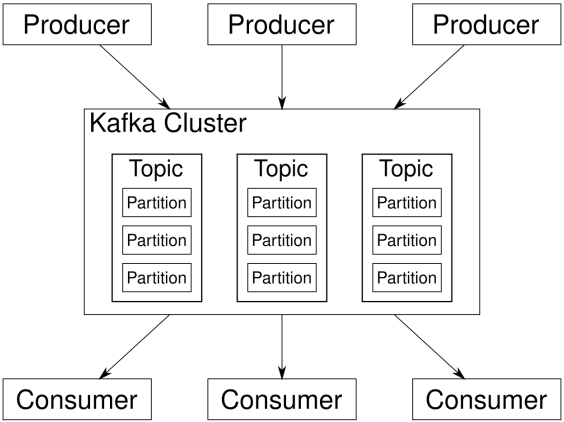
https://sea.mastercard.com/en-region-sea/business/merchants/start-accepting/payment-process.html



Motivation

Multi-agent systems

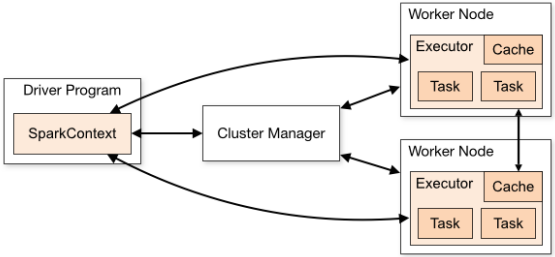
https://en.wikipedia.org/wiki/Apache_Kafka#/media/File:Overview_of_Apache_Kafka.svg



Motivation

Multi-agent systems

https://spark.apache.org/docs/latest/cluster-overview.html



Motivation

Multi-agent systems

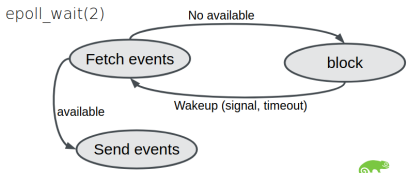
https://events19.linuxfoundation.org/wp-content/uploads/2018/07/dbueso-oss-japan19.pdf

Locking Rules

struct eventpoll
spinlock_t lock
mutex lock
wait_queue_head_t wq
wait_queue_head_t poll_wait
list_head rdlist
rb_root rbr
struct epitem *ovflist
wakeup_source *ws
user_struct *user
file *file
int visited
list_head visited_list_link

Mutex: serialization while transferring events to userspace
copy_to_user might block.
Protect epoll_ctl(2) operations, file exit, etc.

Spinlock: serialization inside IRQ context, cannot sleep.
Protects ready and *overflow* list manipulation.
(Must already hold the ep->mutex)



Common problems

- Many **independent** parts
- Different **speed**
- Need to **communicate** with each other
- Need to **coordinate** execution
- May **fail** spuriously

Agents

This course is about **communicating agents**. We will see

- Alice & Bob
- Processes in operating system
- Threads in operating system
- Separate computers that use networking to pass messages
- ...

Agents will try to:

- Decide who is in charge
- Pass data without corruption
- Coordinate execution of several steps in complicated task
- Find out how many agents take part in the communication
- ...

Different worlds

There are a lot of basic concepts behind the communication of independent agents.
But worlds are *different*:

- timings
  - minutes (conversation)
  - seconds (networking)
  - milliseconds (shared memory)
  - nanoseconds (CPU)
- delivery
  - possible distortion (misunderstanding)
  - possible loss (UDP packet)
  - visibility issue (not-yet-updated memory cell/stale value in CPU cache)
- denial-of-service
  - several people do phone call to the same person
  - network go offline
  - shared memory of single process in consistent (yet weakly ordered)
- ...

Lecture plan

Question time

Question: "concurrency", "parallelism", who could translate both words to Russian?

Concurrency vs parallelism

Situations

Sequential and interruptible = concurrent[1]

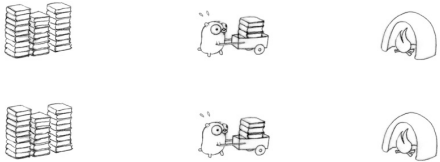
[1] https://go.dev/talks/2012/waza.slide#12

Concurrency vs parallelism

Situations

Sequential and interruptible = concurrent[1]

Independent and simultaneous = parallel[2]



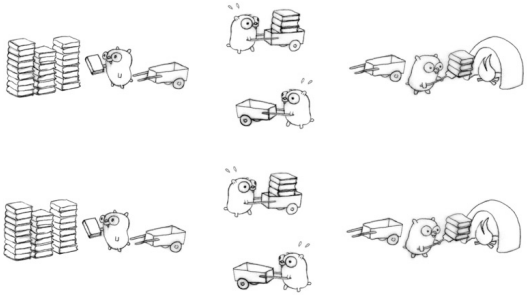
[1] https://go.dev/talks/2012/waza.slide#12

[2] https://go.dev/talks/2012/waza.slide#15

Concurrency vs parallelism

Two sides of the same medal

Real life is mixed[3]



Fun fact: real people in industry/academia may "interpret" parallelism/concurrency in different ways

[3] https://go.dev/talks/2012/waza.slide#22

Concurrency vs parallelism

English

- Parallel
- Concurrent
- Independent
- Simultaneous
- At the same time

Question time

Question: "Parallel", "Concurrent", "Independent", "Simultaneous", "At the same time". Use one Russian word which would be the best translation for all 5 concepts.

Concurrency vs parallelism

English

- Parallel
- Concurrent
- Independent
- Simultaneous
- At the same time

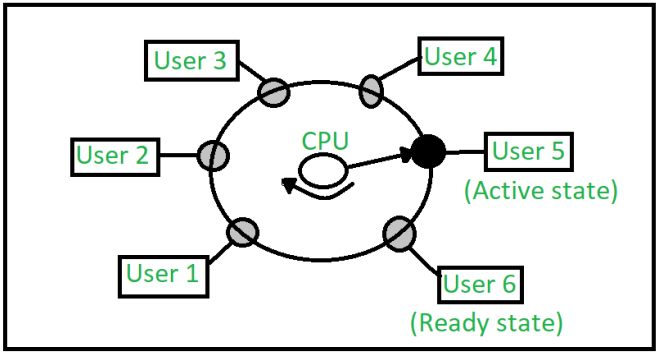
Parallel – *could* execute simultaneously (independent, applicable to different executors)
Concurrent – *could* execute interchangeably (interruptible, applicable to the same executor)
Any real-life task combines parallel and concurrent parts[4].

[4] https://stackoverflow.com/questions/1050222/what-is-the-difference-between-concurrency-and-parallelism

Concurrency vs parallelism

Time-sharing for users

Single CPU, many users[5]

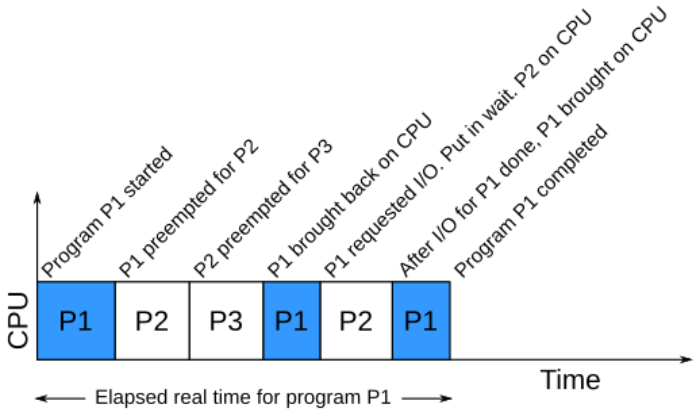


[5] https://www.geeksforgeeks.org/time-sharing-operating-system/

Concurrency vs parallelism

Time-sharing for tasks

Single CPU, pre-emptible tasks[6]

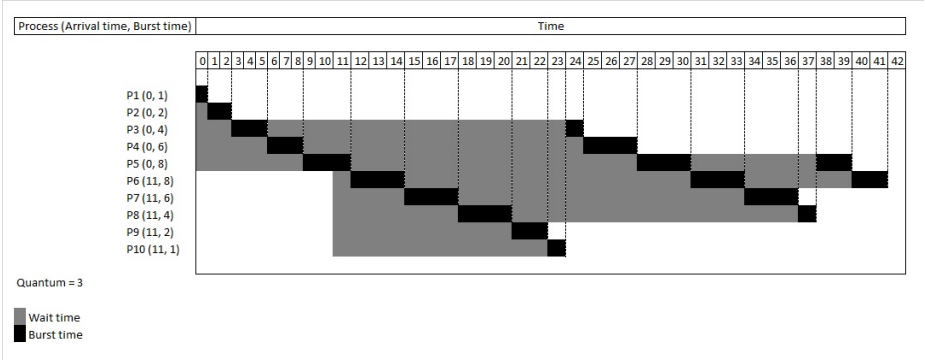


[6] https://en.wikipedia.org/wiki/CPU_time

Concurrency vs parallelism

Time-sharing scheduling strategies

Round-robin scheduling[7]



[7] https://en.wikipedia.org/wiki/Round-robin_scheduling

Concurrency vs parallelism

Conclusion

Concurrency and parallelism are two edges of the same phenomena.
In computer science (and in our course)

- parallelism
- concurrency
- interruptibility
- independence

have more formal meaning than in "usual" communication. Be aware.

Question time

Question: Multitasking, time-sharing, preemption, round-robin. We are going to talk about ...

Lecture plan

Scheduler

Reminders

- Ivan Ugliansky, "Scheduling: from FCFS to EEVDF"

  https://github.com/Svazars/parallel-programming/blob/main/slides/pdf-extra/algo_plus_ose_schedulers.pdf

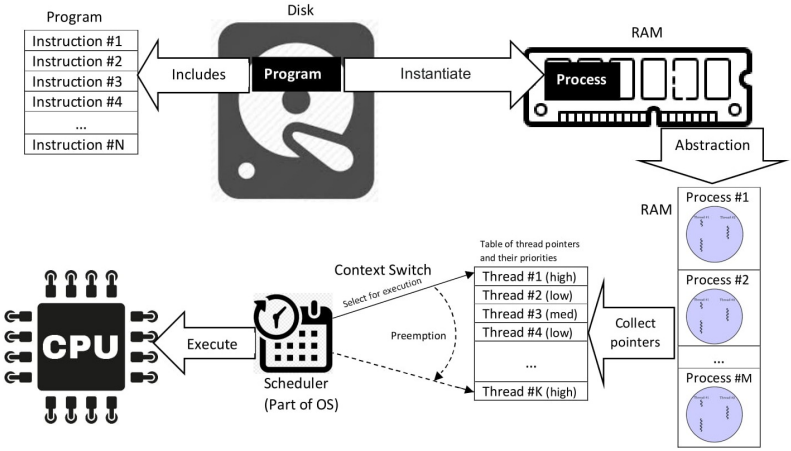
- "Operating Systems: Three Easy Pieces", "Scheduling: Introduction"

  https://pages.cs.wisc.edu/%7Eremzi/OSTEP/cpu-sched.pdf

- https://en.wikipedia.org/wiki/Process_(computing)

Scheduler

Problem overview



Scheduler

Problem overview

Scheduler – magic black box

- "knows and owns" every unit of scheduling (thread, process)
- may pre-empt execution of any unit (save state from CPU to memory)
- may enable execution of any pending unit (load state from memory to CPU)
- implements some scheduling policy (round-robin, shortest remaining time ...)

Scheduler

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Save state / load state == switch CPU context of execution

Question time

Question: What is context switch on modern OS (e.g. Linux) and modern hardware with MMU (e.g. x86_64 intel Skylake)?

Scheduler

Context switch

Do not forget all the important data

- registers (general-purpose, floating-point, other)
- stack pointer
- instruction pointer
- segmentation/page tables
- translation lookaside buffer flush
- ...

Requires `user-space -> kernel space` transition.

Scheduler

Context switch

Do not forget all the important data

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Requires `user-space -> kernel space` transition.
Needs *amortization*.

Scheduler

Context switch

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- ...

Requires `user-space -> kernel space` transition.
Needs *amortization*.
Scheduling quantum.

Scheduler

Scheduling quantum

Quantum may be arbitrary function of `(current task, other tasks, state of OS)`.
Do not forget about

- frequency of context switches
- "price" of context switch
- thread/process priority
- "price" of CPU migration
- occupation of other processors
- ...

It is hard to trust timings[8,9].

[8] `https://shipilev.net/blog/2014/notrusting-nanotime`
[9] `https://pzemtsov.github.io/2017/07/23/the-slow-currenttimemillis.html`

Scheduler

Pre-emptive and cooperative multitasking

When to context switch?

Scheduler

Pre-emptive and cooperative multitasking

When to context switch?

Anytime anywhere:

- flexible
- easy-to-use model
- requires additional care
- has counter-examples
- relies on heuristics

When allowed to:

- fine-grained control
- nontrivial behaviour
- requires additional care
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Scheduler

Pre-emptive and cooperative multitasking

When to context switch?

Anytime anywhere:

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When allowed to:

- fine-grained control
- nontrivial behaviour
- requires additional care
- has counter-examples
- relies on heuristics

In real environments it is always a hybrid solution[10].

[10] The long road to lazy preemption: `https://lwn.net/Articles/994322/`

Scheduler

Goals

- In theory, there exist a *perfect* schedule
- In practice, almost any scheduling problem is NP-complete
- In applications, workload is dynamically generated on-the-fly

Scheduler

Goals

- In theory, there exist a *perfect* schedule
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Realistic goals:

- Max *utilization*, not max *efficiency*
- Reasonable *latency*
- Soft/hard *real-time* constraints
- *Energy efficiency*

Question time

Question: Describe situation when *max utilization* scheduling contradicts *max efficiency* scheduling

Question time

Question: OS scheduler schedules **what**? Scheduler manages which OS **resource**?

Lecture plan

Question time

Question: Difference between "threads" and "processes" in one word?

Threads vs processes

Blast from the past

Process[11]

- instance of computer program (program code and other stuff like .data, .rodata, .bss ...)
- owns resources (RAM, CPU, network ...) allocated by OS
- has logical and physical access permissions (filesystem, user capabilities ...)
- described by state (context)
- managed by OS
- **isolated** from other processes (virtual memory, namespaces ...)

[11] https://en.wikipedia.org/wiki/Process_(computing)

Threads vs processes

Blast from the past

Thread[12]

- part of process (executes some code in some context)
- owns thread-specific resources (stack, TLS ...)
- has specific metadata (priority, TID ...)
- described by state (context)
- managed by OS scheduler
- **shares** memory with other threads

[12] https://en.wikipedia.org/wiki/Thread_(computing)

Threads vs processes

Bird-eye view

Processes and threads are **agents**

- different speed
- could communicate
- need coordination

Question time

Question: Threads need *more* or *less* coordination compared to processes?

Threads vs processes

Bird-eye view

Processes and threads are **agents**

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Threads vs processes

Bird-eye view

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In some sense, threads need *less* coordination (memory is already shared).

Threads vs processes

Bird-eye view

Processes and threads are **agents**

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In some sense, threads need *less* coordination (memory is already shared).
In some sense, threads need *more* coordination (memory is already shared).

Threads vs processes

Examples: process

```
echo "Sequential"
wc 50_000_000.txt
```

```
real    0m 4,900s
user    0m 4,212s
sys     0m 0,240s
```

```
echo "Parallel"
{ ( head -n 25000000 50_000_000.txt | wc) & }
{ ( tail -n 25000000 50_000_000.txt | wc) & }
wait
```

```
real    0m 2,323s
user    0m 4,576s
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Threads vs processes

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```

- Create processes
- Wait completion
- Pass data via pipes, files, advanced IPC

Threads vs processes

Examples: Java threads

```
static long a, b;
static class MyThread extends Thread {
    private int idx; MyThread(int i) { this.idx = i; }
    @Override public void run() {
        for (long i = 0; i < 500_000; i++) {
            if (idx == 1) { a += func(i); }
            else          { b += func(i + 500_000); }
    }}}
public static void main(String... args) throws Exception {
    Thread t1 = new MyThread(1); Thread t2 = new MyThread(2);
    t1.start(); t2.start(); // !!! not t1.run() !!!
    t1.join(); t2.join();
    System.out.println("Result = " + (a + b));
}
```

Your first homework

Your first homework

Read the docs

- https://docs.oracle.com/en/java/javase/11/docs/api/java.base/java/lang/Thread.html
- https://docs.oracle.com/en/java/javase/11/docs/api/java.base/java/lang/Thread.html#start()
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Yes, we use JDK 11 for all homeworks.

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https://github.com/Svazars/parallel-programming/blob/main/hw/block1/1.1/readme.markdown

Homework, mail

Task 1.1 Create 3 Java programs for the following cases, explain results in several sentences.

Thread **A** starts thread **B** and `joins`. Uncaught exception happens in thread **B**.

- What happens in thread **B**?
- What happens in thread **A**?
- What if thread **C** `joins` thread **B** *after* exception happened?
- What if thread **D** `joins` thread **A**?

Question time

Question: Why would somebody use Java for parallel programming course, not C or C++?

Threads vs processes

Conclusion

From our course perspective:

- processes hold all the context and resources
- processes use many different parts of OS
- threads use ultra-fast shared memory for communication
- threads are minimal units eligible for OS scheduler

Almost all lectures will be about

- inter-process multi-threading
- inside high-level programming language
- run by multiprocessor hardware with MMU
- on non-real-time OS with pre-emptive multitasking

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Lecture plan

Threading models

1:1 model

- `N` preemptible tasks
- Every task is allocated on designated computing agent

Threading models

1:1 model

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- Examples:

Threading models

1:1 model

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- Examples:
  - 1 Java thread – 1 OS thread
  - 1 OS thread – 1 CPU

Threading models

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N:1 model

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- Single executor

Threading models

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  - 1 Java thread – 1 OS thread
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N:1 model

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N:M model

- `N` preemptible tasks
- Multiplexed over `M` computing agents. (`1 < M < N`)

Threading models

1:1 model

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- Every task is allocated on designated computing agent
- Examples:
  - 1 Java thread – 1 OS thread
  - 1 OS thread – 1 CPU

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Threading models

1:1 model

- `N` preemptible tasks
- Every task is allocated on designated computing agent
- Examples:
  - 1 Java thread – 1 OS thread
  - 1 OS thread – 1 CPU

N:1 model

- `N` preemptible tasks
- Single executor

N:M model

- `N` preemptible tasks
- Multiplexed over `M` computing agents. (`1 < M < N`)
- Examples:
  - thread pools (Lecture 3)
  - green threads, coroutines (Lecture 12)

N:1 threading model

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- `N` preemptible tasks
- Single executor

N:1 threading model

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Single-core OS

- `N` running processes
- Single CPU

N:1 threading model

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Single-core OS

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Classic JVM run on single-core hardware

- `N` running Java threads
- Each Java thread mapped to corresponding OS thread (`1:1`)
- OS threads run on single CPU (`N:1`)

N:1 threading model

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Classic JVM run on single-core hardware

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- Each Java thread mapped to corresponding OS thread (`1:1`)
- OS threads run on single CPU (`N:1`)

Useful **model** for discussing possible behaviours of concurrent programs

Threading model: not enough

N:1 threading model

- `N` preemptible tasks
- Single executor

Threading model: not enough

N:1 threading model

- `N` preemptible tasks
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Describes allocation of computing resources.

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What about interaction between tasks?
Need additional information

- guarantees to memory update visibility?
- expectations about scheduling policy?
- cross-thread interactions ordering?

Interleaving model

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- At any place (any machine instruction)
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This is **simplified** model of concurrent execution. We will use it up to Lecture 6.

Question time

Question: What is the difference between `N:1` threading model and describing concurrent system using `interleavings`?

Interleaving model: visualization

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static int x, y;
```

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A.1->...->A.5. a_x = 0, a_y = 0;

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Interleaving model: single execution trace

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Interleaving model: possible execution results

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    int a_y = y;
    x++;
    y++;
    log(a_x, a_y);
}}
```

```
class B extends Thread {
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```
a_x = 0, a_y = 0, b_x = 0, b_y = 0
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          ...
```

Interleaving model: key concepts

Concurrent execution trace:

- All operations inside thread ordered sequentially
- All operations across all threads are totally ordered on timeline
- Single deterministic result of execution

Non-determinism:

- Concurrent program may have **many** execution traces.
- Concurrent program may have **many** results of execution.

Question time

Question: single-threaded program that have many execution traces?

Question time

Question: concurrent program with many execution traces but only one result of execution?

Question time

Question: concurrent program with infinite number of execution traces?

Interleaving model: homework

Homework, mail

Task 1.2. For given multithreaded program and execution result provide either

- *concurrent execution trace (`A.1->A.2->B.1 ...`) or*
- *proof of impossibility*

https://github.com/Svazars/parallel-programming/blob/main/hw/block1/1.2/readme.markdown

Homework, mail

Task 1.3. For given multithreaded program:

- *provide all possible execution results*
- *provide only one execution trace per result*
- *prove exhaustiveness of your answer*

https://github.com/Svazars/parallel-programming/blob/main/hw/block1/1.3/readme.markdown

Our main enemy through the course

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Solution:

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- Go to semantic level (Lecture 10)
- ... and beyond (Lectures 11+)

Concurrency: tips and tricks

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- If concurrent tests pass on your local machine

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- If interleaving model proves correctness of your program
  - still run stress-tests.

Lecture plan

1. Motivation
2. Concurrency and parallelism
3. Scheduler
4. Threads and processes
5. Multitasking and threading models
6. Concurrent problems
   - Embarrassingly parallel problems and Amdahl's law
   - Race condition
   - Data race
   - Visibility
   - Deadlock
   - Priority inversion
7. Summary

Concurrent building blocks you know so far

1. Create new thread and `start` it
2. Modify fields concurrently
3. `Join` thread

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- Solve "easy" parallel problems
- Encounter logic errors (race conditions)
- Encounter concurrent memory access effects (data races)

Using **3**:

- Encounter visibility bugs (stale value)
- Encounter no-progress blocking (deadlock)
- Observe slow-progress issues (priority inversion)

Lecture plan

Embarrassingly parallel problems

Introduction

Synonyms:

- Embarrassingly parallel
- Embarrassingly parallelizable
- Perfectly parallel
- Delightfully parallel
- Pleasingly parallel

Key properties:

- Granularity
- Minimal coordination

Question time

Question: Any ideas for embarrassingly parallel problems?

Embarrassingly parallel problems

Examples

- function applied to collection of items (*map*)
  - ray tracing
  - word count in independent files
  - proof-of-work cryptocurrency
  - bioinformatics search (BLAST etc)
  - Discrete Fourier transform
- associative operation applied to large array (order-independent *reduce*)
  - sum of elements in integer array
  - word count for all books in e-library
- ...

Amdahl's law

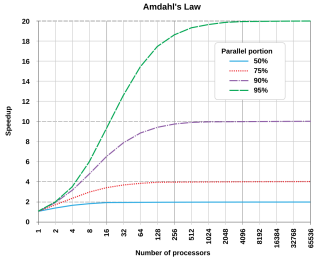
Key ideas

Any problem consists of *serial* part and *parallel* part.

Speed-up of parallel part depends on computational resources. Serial part is "fixed work".

- **Diminishing Returns**: beyond certain point, adding more processors doesn't significantly increase speedup
- **Limited Speedup**: problem cannot be solved faster than serial part

https://en.wikipedia.org/wiki/Amdahl%27s_law



Amdahl's law

Examples

- Programmer enhances a part of the code that represents 10% of the total execution time. This part starts to work 10 000 times faster. What is total speedup of a program?
- Problem could be split into two parts: A and B. A is serial, B is parallel. On single CPU, A takes 50 minutes, B takes 250 minutes. How many CPUs do you need to solve the problem in 100 minutes (achieve 3x speed-up)?
- Prepare an example where 10x speed-up could be achieved by using 100 CPUs only.

Amdahl's law

Examples

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- Prepare an example where 10x speed-up could be achieved by using 100 CPUs only.

https://github.com/Svazars/parallel-programming/blob/main/hw/block1/1.4/readme.markdown

Homework, mail

Task 1.4 Provide answers and explanation for all three problems

Embarrassingly parallel problems

Piece of advice

Such problems usually quite interesting and practically-oriented.
They do **not** require non-trivial coordination so we will not cover them in this course.
In many domains, most of "performance issues" are (almost) embarrassingly parallel.
There are plenty of

- easy-to-use
- straightforward
- safe
- stable
- performant

solutions to such problems. Just use them and get speed-up predicted by Amdahl's law.
If you go real concurrency, it will be **harder**.

Lecture plan

Race condition

Example

```
public static void main(String... args) throws Exception {
    Thread t1 = new Thread(() -> { System.out.println("Thread A"); });
    Thread t2 = new Thread(() -> { System.out.println("Thread B"); });
    t1.start(); t2.start();
    t1.join(); t2.join();
}
```

Race condition

Example

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public static void main(String... args) throws Exception {
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    t1.start(); t2.start();
    t1.join(); t2.join();
}
```

Execution 1: `Thread A Thread B`

Execution 2: `Thread B Thread A`

Race condition

Key idea

- Every single operation in program is OK (properly synchronized, atomic)
- Program as a whole suffers from non-deterministic behaviour
- The only reason to this is different speed of execution

Race condition

Key idea

- Every single operation in program is OK (properly synchronized, atomic)
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https://github.com/Svazars/parallel-programming/blob/main/hw/block1/1.5/readme.markdown

Homework, mail

Task 1.5 Open `https://deadlockempire.github.io`, pass all levels up to "Confused counter", inclusive.

Race condition

Conclusion

Race condition

- at least 2 threads involved
- behaviour of system depends on sequence of events (e.g. timings)
- leads to non-deterministic results

Race condition

Conclusion

Race condition

- at least 2 threads involved
- behaviour of system depends on sequence of events (e.g. timings)
- leads to non-deterministic results

Really **suspicious**

- complicates analysis of algorithm (correctness, performance, resource utilization)
- may lead to rarely reproducible bugs
- may lead to random denial-of-service events

Race condition

Conclusion

Race condition

- at least 2 threads involved
- behaviour of system depends on sequence of events (e.g. timings)
- leads to non-deterministic results

Really **suspicious**

Race condition

Conclusion

Race condition

- at least 2 threads involved
- behaviour of system depends on sequence of events (e.g. timings)
- leads to non-deterministic results

Really **suspicious**
Not **necessarily** a problem

- Server counts number of incoming requests.
- Search engine crawls page graph using several workers. Many pages visited several times, deduplication makes result consistent.
- Torrent client downloads several chunks of a file simultaneously.

Race condition

Conclusion

Race condition

- at least 2 threads involved
- behaviour of system depends on sequence of events (e.g. timings)
- leads to non-deterministic results

Really **suspicious**
Not **necessarily** a problem

Race condition

Conclusion

Race condition

- at least 2 threads involved
- behaviour of system depends on sequence of events (e.g. timings)
- leads to non-deterministic results

Really **suspicious**
Not **necessarily** a problem
Main headache of our course:

- **Race condition may be a bug or intentional behaviour depending on programmer's intention**

In general, cannot be detected by static code analysis/dynamic instrumentation tool.

Question time

Question: Some tools detect race conditions in concurrent software. What is "false positive" detection of race condition?

Lecture plan

Data race

Example 1

```
static volatile int x = 0;
public static void main(String... args) throws Exception {
    Thread t = new Thread(() -> { x++; });
    t.start();
    System.out.println(x);
    t.join();
}
```

Data race

Example 1

```
static volatile int x = 0;
public static void main(String... args) throws Exception {
    Thread t = new Thread(() -> { x++; });
    t.start();
    System.out.println(x);
    t.join();
}
```

Execution 1: `x = 0`

Execution 2: `x = 1`

Data race

Example 2

```
static volatile int x = 0;
public static void main(String... args) throws Exception {
    Thread t = new Thread(() -> { x++; });
    t.start();
    x++;
    System.out.println(x);
    t.join();
}
```

Data race

Example 2

```
static volatile int x = 0;
public static void main(String... args) throws Exception {
    Thread t = new Thread(() -> { x++; });
    t.start();
    x++;
    System.out.println(x);
    t.join();
}
```

Execution 1: `x = 1`
Execution 2: `x = 2`

Data race

Example 2.5

```
static volatile int x = 0;
public static void main(String... args) throws Exception {
    Thread t = new Thread(() -> { x++; });
    t.start();
    x++;
    t.join();
    System.out.println(x);
}
```

Data race

Example 2.5

```
static volatile int x = 0;
public static void main(String... args) throws Exception {
    Thread t = new Thread(() -> { x++; });
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    x++;
    t.join();
    System.out.println(x);
}
```

Execution 1: `x = 2`

Data race

Example 2.5

```
static volatile int x = 0;
public static void main(String... args) throws Exception {
    Thread t = new Thread(() -> { x++; });
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    x++;
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Execution 1: `x = 2`
Execution 2: `x = 1`

Data race

Example 2.5

```
static volatile int x = 0;
public static void main(String... args) throws Exception {
    Thread t = new Thread(() -> { x++; });
    t.start();
    x++;
    t.join();
    System.out.println(x);
}
```

Execution 1: `x = 2`

Execution 2: `x = 1`

`x++ <=> int tmp = x; tmp = tmp + 1; x = tmp`

Data race

Key idea

Actually, formal definition exists and is language-specific. More on this in next lectures.
Data race:

- at least 2 threads
- at least one of threads is writing data
- operating on the same *memory location* (Lecture 9)
- *simultaneously* (Lecture 6)
- without *proper synchronization* (Lecture 10)

Any data race is race condition. Not every race condition is data race.

Data race

Compound data example

```
static volatile ArrayList<String> list = new ArrayList<>();
public static void main(String... args) throws Exception {
    Thread t1 = new Thread(() -> { list.add("x"); });
    Thread t2 = new Thread(() -> { list.add("y"); });
    t1.start(); t2.start();
    t1.join(); t2.join();
    for (String s : list) {
      System.out.println(s);
    }
}
```

Data race

Compound data example

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static volatile ArrayList<String> list = new ArrayList<>();
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```

What is the size of `list`? **1** or **2**?

Data race

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    for (String s : list) {
      System.out.println(s);
    }
}
```

What is the size of `list`? **1** or **2**?

Data race may lead to **inconsistent** state even if every thread executed only "valid" operations.

Data race

Conclusion

Data race is a specific variation of race condition

- concurrent and not properly synchronized
- write access
- to shared data

Always **problematic**

- inconsistent shared state
- unexpected/impossible computation results
- forbidden by language implementation

Could be formalized and avoided

- Strict coding policy (e.g. enforced by linter/compiler)
- Static analysis of source code
- Static analysis of execution traces
- Dynamic instrumenting race finders

Question time

Question: Some tools detect data races in concurrent software. Why do we still have systems with data races?

Lecture plan

Visibility

Introduction

It is one of the most subtle notions in concurrency.

Visibility

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It is one of the most subtle notions in concurrency.
For now, we should expect that

- compiler could aggressively reorder/eliminate field operations
- threads may see stale values, not the "newest" values written by other threads

unless there is "synchronization point" between two threads

Visibility

Introduction

It is one of the most subtle notions in concurrency.
For now, we should expect that

- compiler could aggressively reorder/eliminate field operations
- threads may see stale values, not the "newest" values written by other threads

unless there is "synchronization point" between two threads

- `Thread.start`
- `Thread.join`

Visibility

Incorrect: message passing via plain fields

```
static boolean jobDone = false;
public static void main(String... args) throws Exception {
    Thread t = new Thread(() -> { jobDone = true; });
    t.start();
    while (jobDone == false) { /* loop */ }
    t.join();
}
```

Visibility

Incorrect: message passing via plain fields

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public static void main(String... args) throws Exception {
    Thread t = new Thread(() -> { jobDone = true; });
    t.start();
    while (jobDone == false) { /* loop */ }
    t.join();
}
```

May hang because:

- compiler optimization: `if (!jobDone) while(true);`
- stale value, `jobDone` remains `false` in main thread forever

Lecture plan

Blocking methods

Wait-for graph

Blocking method: execution of current thread is *suspended* until some condition is met.

```
Thread A = new Thread(() -> {
    Thread B = new Thread(() -> { Thread.sleep(10_000); });
    Thread C = new Thread(() -> { B.join(); Thread.sleep(10_000); });
    B.start(); C.start(); B.join(); C.join()
}); A.start(); A.join();
```

Blocking methods

Wait-for graph

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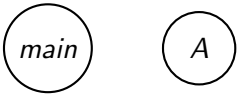

Blocking methods

Wait-for graph

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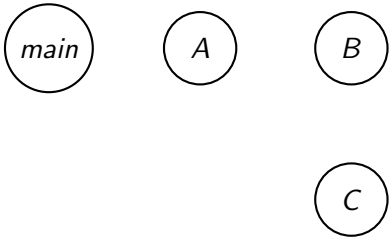

Blocking methods

Wait-for graph

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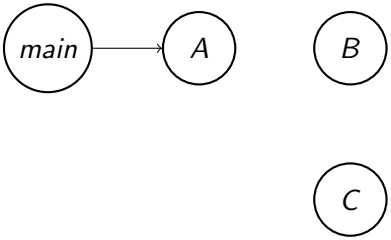

Blocking methods

Wait-for graph

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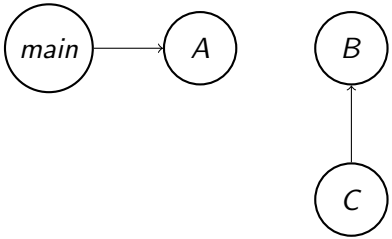

Blocking methods

Wait-for graph

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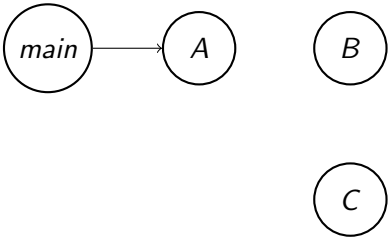

Blocking methods

Deadlock

Deadlock – cycle in wait-for graph.

Blocking methods

Deadlock

Deadlock – cycle in wait-for graph.
Trivial single-threaded deadlock:

Blocking methods

Deadlock

Deadlock – cycle in wait-for graph.
Trivial single-threaded deadlock:

```
Thread.currentThread().join();
```

Blocking methods

Deadlock

Deadlock – cycle in wait-for graph.
Trivial single-threaded deadlock:

```
Thread.currentThread().join();
```

Classic two-thread deadlock:

```
static volatile Runnable lambda = null;
public static void main(String... args) throws Exception {
    Thread A = new Thread(() -> { lambda.run(); });
    Thread B = new Thread(() -> { A.join();     });
    lambda = () -> { B.join(); };
    A.start(); B.start();
    A.join(); B.join();
}
```

Blocking methods

Deadlock visualisation

```
java Sample &
PID="$!" && sleep 1
jstack -l $PID
```

Blocking methods

Deadlock visualisation

```
java Sample &
PID="$!" && sleep 1
jstack -l $PID
```

Output:

```
"main" #1 prio=5 os_prio=0 cpu=38,96ms elapsed=1,25s tid=0x00007fdcc0017000 nid
   java.lang.Thread.State: WAITING (on object monitor)
    at java.lang.Object.wait(java.base@14.0.1/Native Method)
    - waiting on <0x0000000095148930> (a java.lang.Thread)
    ...
    at Sample.main(Sample.java:10)
    ...
```

Blocking methods

Deadlock detection

- If API has at least one blocking method – extreme care required to avoid deadlocks
- If you do not know in which thread your code will be executed – extreme care required to avoid deadlocks
- If you can not control execution (inheritance, virtual methods, function pointers) – extreme care required to avoid deadlocks
- Design of any concurrent system should start with preparing tools for debugging and monitoring (observability API)

Blocking methods

Deadlock detection

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https://github.com/Svazars/parallel-programming/blob/main/hw/block1/1.6/readme.markdown

Homework, mail

Task 1.6 Run `jstack` on all deadlock samples from this lecture.

Question time

Question: Object-oriented languages love to "hide" implementation details using abstractions (inheritance, virtual methods, function pointers, interfaces). How would I know that user of my library/module will not call blocking method and introduce deadlock in my system?

Lecture plan

Blocking methods

Priority inversion

```
public static void main(String... args) {
  Thread A = new Thread(() -> { do1(); B = startThread(); do2(); B.join(); });
  A.start(); doCriticalWork(); A.join();
}
```

`main` waitsFor `A`, `A` waitsFor `B`.
Progress of `main` determined by progress of `B`.
"Soft" version of deadlock (depends-on graph).

- Some Operation Systems **do** have special support to remedy priority inversion problems
- Many Virtual Machines, concurrent libraries, multi-threaded frameworks **do not**

Blocking methods

Priority inversion

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public static void main(String... args) {
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Progress of `main` determined by progress of `B`.
"Soft" version of deadlock (depends-on graph).

- Some Operation Systems **do** have special support to remedy priority inversion problems
- Many Virtual Machines, concurrent libraries, multi-threaded frameworks **do not**

If you launch rovers on Mars, be ready to debug them[13].

[13] https://www.cs.cornell.edu/courses/cs614/1999sp/papers/pathfinder.html

Lecture plan

1. Motivation
2. Concurrency and parallelism
3. Scheduler
4. Threads and processes
5. Multitasking and threading models
6. Concurrent problems
   - Embarrassingly parallel problems and Amdahl's law
   - Race condition
   - Data race
   - Visibility
   - Deadlock
   - Priority inversion
7. Summary

Summary

- Agents with different speed could communicate with each other via shared memory
- Parallel and concurrent execution could look similar and could happen in the same system
  - Concurrency: tasks may be interrupted in pre-defined places or arbitrary
  - Parallel: independent tasks may be executed simultaneously on different cores
- Threads are parts of OS processes and are minimal unit of work for scheduler
- There are different threading models: `1:1`, `N:1`, `N:M`
- Interleaving `N:1` is useful yet simplified model of concurrent execution
- The perfect scenario: problem could be divided into many independent parts with almost zero coordination overheads
- Amdahl's law fundamentally limits possible speed-up of tasks with sequential parts
- Whenever you coordinate threads, you encounter non-determinism, race conditions, data races, visibility issues, deadlocks, priority inversion

Summary: homework

`Lecture 1, Tasks 1.*:` https://github.com/Svazars/parallel-programming/blob/main/hw/block1