

Formalizing Coq Modules in the MetaCoq project

IMPRS-TRUST Application Talk

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- May-Jul 2022 (after 4th year undergraduate), interned at the Gallinette team in Nantes, France, advised by Nicolas Tabareau.
- Chosen among 3 possible projects: Modules, Parametricity Plugin, Eta-Reduction
- From Aug 2022 onwards: bachelor's thesis in the National University of Singapore: co-advised by Nicolas Tabareau (INRIA), Martin Henz (NUS Computing) and Yue Yang (NUS Mathematics).

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- Coq is a proof assistant based on Type Theory. In particular, the **P**olymorphic, **C**umulative Calculus of (co-)Inductive **C**onstrutions (PCUIC).
- We trust the theory to be correct: PCUIC is proven strongly normalizing, reduction and conversion are shown to be decidable...

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- We trust the theory to be correct: PCUIC is proven strongly normalizing, reduction and conversion are shown to be decidable...
- But when we use Coq, we use the OCaml implementation.
- Is this implementation also reliable? Correct?

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- Bonus: formalize the theory of Coq, PCUIC, using MetaCoq.

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- MetaCoq project in 2018 as a metaprogramming platform for Coq, building on existing systems such as TemplateCoq (2014).
- Bonus: formalize the theory of Coq, PCUIC, using MetaCoq.
- Slogan: the Coq implementation you run every week/month.

What is the theory of Coq?

The type theory behind the Coq proof assistant is PCUIC:

- (Universe) **P**olymorphic
- **C**umulative
- Calculus of (co-)Inductive **C**onstructions (CoC with Inductive Types)
- Calculus of Constructions (Coquand and Huet, 1988)

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- Caveats:
 1. Eta-expansion
 2. Template Polymorphism
 3. SProps (Proof-irrelevant Propositions)
 4. **Modules**

Coq Modules

Example - Definitions

Modules as “collections of definitions”.

```
Inductive nat :=  
  | 0  
  | S : nat -> nat.
```

```
Fixpoint plus (n m: nat) :=  
  match n with  
  | S n' => S (plus n' m)  
  | 0 => m  
  end.
```

Example - Modules

“Packaging” definitions into a Module (Type).

```
(* A magma is a set with a binary (closed) operation. *)  
Module Type Magma.  
    Parameter T: Set.  
    Parameter op: T -> T -> T.  
End Magma.  
  
Module Nat: Magma.  
    Definition T := nat.  
    Definition op := plus.  
End Nat.
```

Example - Aliasing

Modules can be aliased for ease of reference.

```
Module Type M := Magma.
```

```
Module MyNat: M := Nat.
```

Example - Functors

Higher-order modules - Functors.

```
(* A functor transforming a magma into another magma. *)  
Module DoubleMagma (M: Magma): Magma.  
    Definition T := M.T.  
    Definition op x y := M.op (M.op x y) (M.op x y).  
End DoubleMagma.  
  
Module NatWithDoublePlus := DoubleMagma Nat.
```

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- There are many contention/issues regarding the semantics of functors: [OpenIssuesWithModules](#)
- I only work on **plain** (non-parametrized) modules and module types, aliasing.
- 2nd class object; cannot pass around as a term
- Restricted range of operations
 1. Declaration/Definition (including aliasing)
 2. Using modules: access definitions in it using path names
 3. Refining a module: create new module by replacing entries of existing

Representation of Global Environment and Modules

- Before: global environment is a list of name-definition pairs for constants and inductive types.

¹Slightly different from named function application for lambda calculus with constants: $\text{plus } 2 \ 3 \rightarrow_{\delta} 5$

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- After: Modules are tree-like structures: they contain definitions of constants, inductives, modules, and module types.

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Representation of Global Environment and Modules

- Before: global environment is a list of name-definition pairs for constants and inductive types.
- After: Modules are tree-like structures: they contain definitions of constants, inductives, modules, and module types.
- Global environment is just a special case of a Module.
- Items in modules are referred to using pathnames (e.g. $M.a$) instead of defined names (e.g. a). The resolution of path name to its definition extends the δ -reduction¹ in Coq.
- **Goal: Formalize the above in TemplateCoq.**

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Some Properties to Verify

1. Definition: well-definedness of this structure: follow the Typing (Formation) Rules.
2. Definition-Alias: well-definedness of aliasing: no cycles.
3. Path: path points to the correct definition.
4. Path-Alias: aliased path points to the correct definition.
5. Related theorems about environment, typing, translation etc.

Example: TemplateToPCUICCorrectness.v

Lemma trans_lookup_module

```
{cf} {Σ} {cst: kername} m {wfΣ : Typing.wf Σ}:  
  (* If cst refers to a module m in the environment Σ, *)  
  Ast.Env.lookup_env Σ cst = Some (Ast.Env.ModuleDecl m)  
  (* Then one cannot find cst in the translated  
  environment. *)  
  -> lookup_env (trans_global_env Σ) cst = None.
```

Proof.

```
destruct Σ as [univs decls retro].  
intros H.  
cbn -[fold_right].
```

Current Work

1. Formalize modules (data structures, typing rules...)
2. From Coq to TemplateCoq (OCaml plugin)
3. From TemplateCoq To PCUIC (elaborating modules away, correctness...)
4. Verify Properties

Thank You!
