

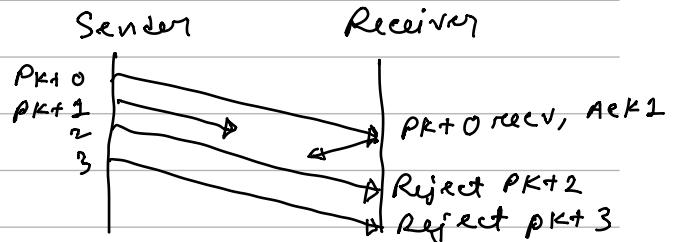
○ Pipelining :

sending data in bulk sequentially.

① Go Back - N (sliding window)

② Selective Repeat

① Go Back N:



② Selective Repeat:

* Individual ACK for

each PKT.

* Out of order

packets

are

sent

in

order

and

ACKs

O

TCP

→ TCP gets continuous data stream → breaks into segments.

Source Port 16 bit	Dest Port 16 bit
Seq. No.	
Ack	
Header Len 4 bit	U A P R S F R C S S T Y E G K H T N N
6 bit	Window Size 16 bit
Check Sum 16 bit	Urgent Data Pointer 16
Options + Data 40 Bytes.	
:	

Byte Streaming

Connection Oriented

Full Duplex

Piggybacking (ACK no., SEQ)

Error Control

Flow Control

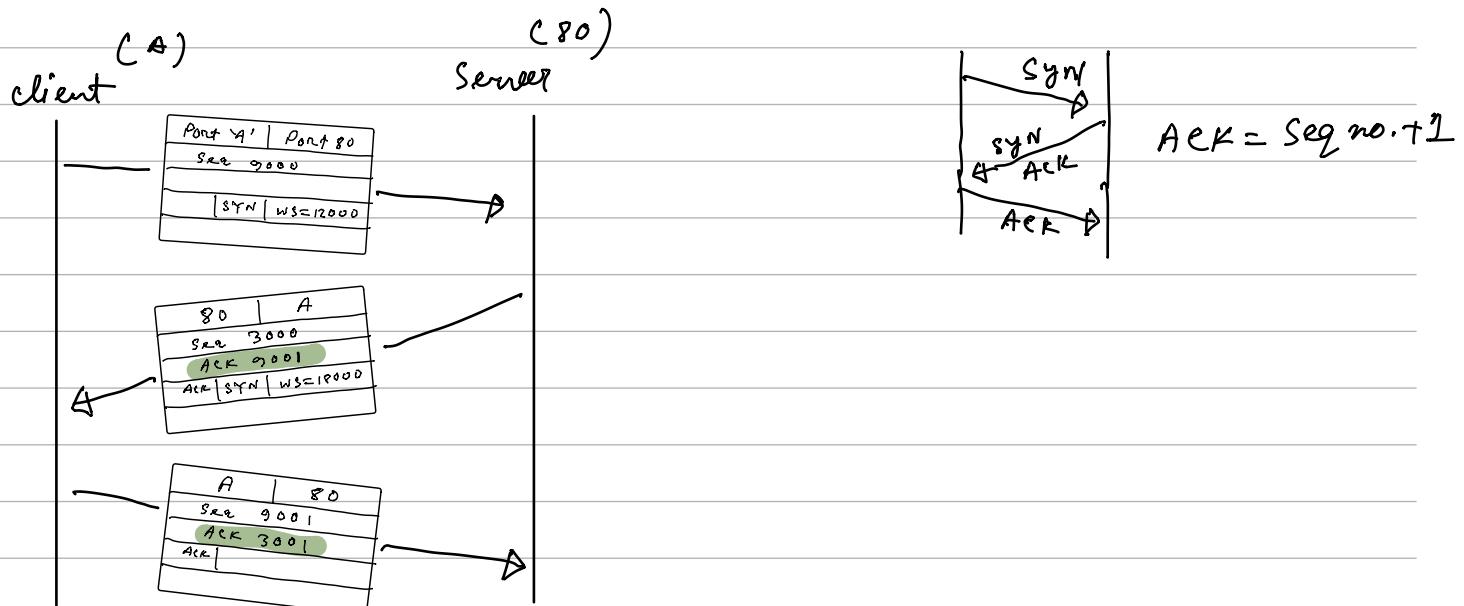
Congestion Control

Ack no. is always next expected Seq. no.

urg : urgent bit ; The Urgent data can be accessed by the urgent data pointer.

PSH: Push ; To send data immediately ; not wait for buffer to full.

□ TCP Connection :



Connection Termination:

i

got FIN=1, Now client deallocate resources

Server no longer can send data



Server deallocates Resources (Half Closed)

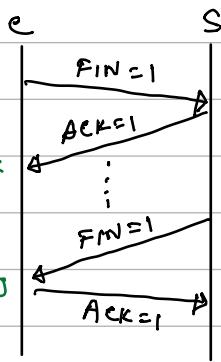
C S client can't send data only ACK.

ii

Server can still send data to client, client did not release resources.

got FIN=1, Now client deallocate resources

Server no longer can send data



Server deallocates Resources (Half Closed)

C S client can't send data only ACK.

O Physical Layer

- Cables & Connectors
 - Physical Topology
 - Hardwares (Repeaters, Hubs)
 - Transmission Mode.
 - Multiplexing → Instead of using multiple channels, use 1 channel & multiplex n signals.
 - Encoding
 - ↳ encode in
 - Analog / Digital
(FM)

Devices in Computer Network :

- Hardware**

 - ① Cables
 - ② Repeaters
 - ③ Hubs
 - ④ Bridges
 - ⑤ Switches
 - ⑥ Routers
 - ⑦ Gateway
 - ⑧ IDS
 - ⑨ Firewall
 - ⑩ Modem (modulator - demodulator)

Security Purposes

Hardware + Software

① Cables:

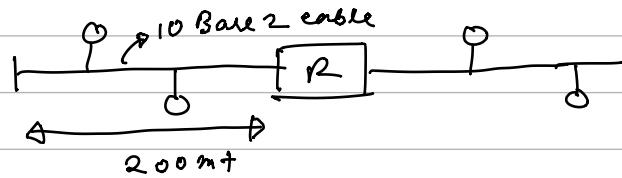
- unshielded Twisted Pair
- Coaxial cable
- Optical Fiber

electrical
10 Mbit/s
Base 2 → attenuation
Base 5
distance = 200mt.

10 Base T
generally 2 Km
light signal.

Attenuation: after 200 meter strength of signal decreases.

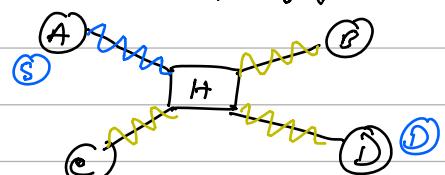
② Repeaters: regenerates signal strength before Attenuation.



- ✓ Repeater is 2 Port device
- ✓ Forwarding Device
- ✓ No filtering.

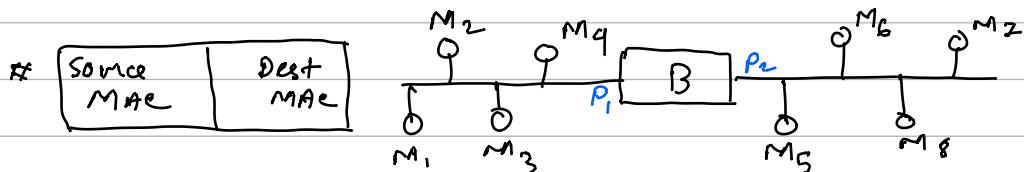
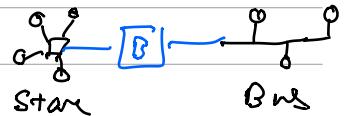
③ Hub: Hub is multipoint Repeater.

used in Star topology



- ✓ Forwarding Device
- ✓ No filtering.

④ Bridge: can connect 2 Different LANs
↳ works in Physical & Link layer*



Types → Static <mac addr, Port no. > Table

→ Dynamic
↓

(Problem: when new device comes constantly update table)

mac	Port
m ₁	P ₁
m ₂	P ₁
m ₃	P ₁
m ₄	P ₁
m ₅	P ₂
m ₆	P ₂
m ₇	P ₂
m ₈	P ₂

Initially Table empty,
when connection happens, that time fill mac, port table.

✓ Forwarding Device

✓ Filters signal

✓ No collision — Has buffer to store incoming data.

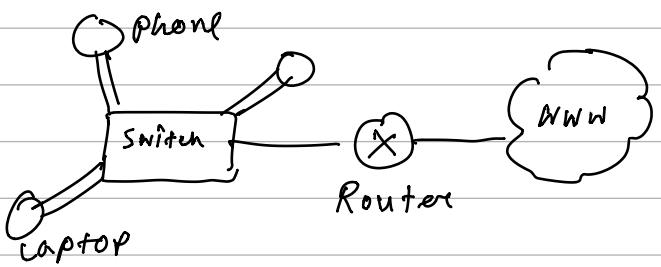
✓ Bridges are Unit Data Protocol

(make spanning Tree to select where to send data) No loop Round Trip count.

⑤ **Switch** : Multipoint Bridge.

↳ Link Layer device

- ✓ Full Duplex links
- ✓ Traffic is minimal
- ✓ Collision domain 0.

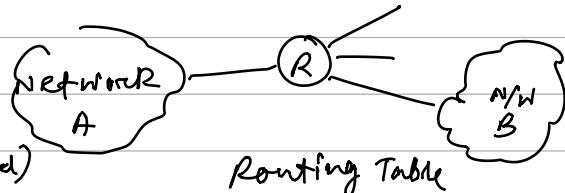


(generally services are connected to switch
not routers)

⑥ **Routers** : Hub, bridge, switch are common for LAN

↳ Routers are the Big Boss holds WAN / Internet
Works on Physical, Link , Network layers.
i.e. Routers work on both MAC & IP addresses.

- ✓ Forwarding
- ✓ Filtering *
- ✓ Routing (can specify forwarded)
- ✓ Flooding (when can't find specific addr in Routing table then send signal to everyone)



- ✓ No collision (Buffer)

* E.g. ARP (Address Resolution Protocol) tell me the Mac Addr corresponding any IP addr.
Router Blocks ARP Requests.

⑦ **Circuit Switching** : Purely physical , Used for Telecommunication.

↳ Dedicated Path & Contiguous Flow (always in Sequence)

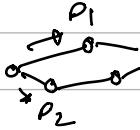
⑧ **Packet Switching** :

↳ Virtual Circuit Switching
Datagram Switching.

- Link &
Network
Layer
- ✓ Store and Forward (Buffer)
 - ✓ Pipelining used

□ Datagram Switching

- Connection less
- Out of Order



- High overhead
- Packet Loss ↑

Virtual Circuit Sw

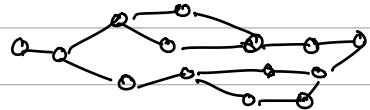
preserves virtual connection
Order maintained

low

Packet loss ↓

○ Message Switching:

Brought the concept
of devision the Data into messages
Store & forward Hop by Hop



↪ Packet switching is just upgrade of Message Switching.

Data Link Layer:

Flow Control : Stop & Wait, Go Back N, Selective Repeat

Error Control : Checksum, CRC (cyclic Redundancy Check)

Access Control: CSMA /
 Collision Detection , Aloha, Token Ring/Bus.

* CRC is integer x $(\text{Data} + 4 \text{Bit}(\text{Data} \cdot 10^x)) \% x = 0 \quad \checkmark$
 $\neq 0 \quad \times$

Medium Access Control (Multiple Access Control) MAC:

(Link Layer)

in mesh topology

all are connected to other

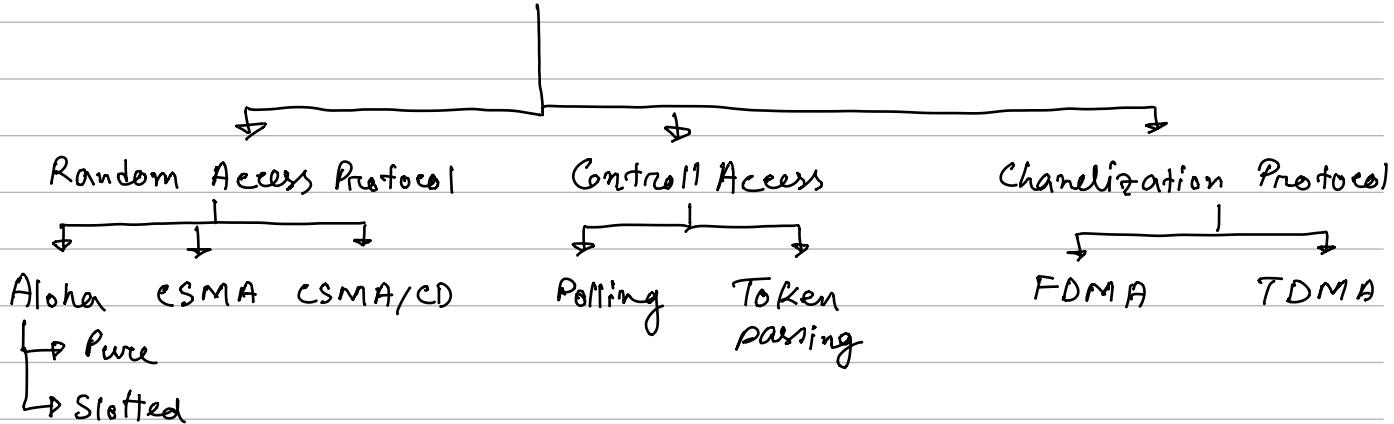


there no

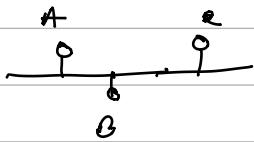
need for MAC

MAC is needed when there's collision or link access

Multiple Access Protocol



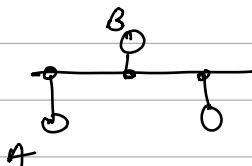
Aloha:



A can send whenever there's need.

There ACK to check whether message was received by C. If not then collision happened, resend.

Carrier-Sense Multiple Access (CSMA)



B before sending; check whether there's already signal on cable.

* Due to propagation delay
False Positive occur



✓ CSMA/CD ; Collision Detection

↳ After sending Data the node can estimate whether there were any collision by calculating how much time it could have taken vs when another incoming signal received.

✓ CSMA/CA ; Collision Avoidance

↳ Used in WLAN.

□ Ethernet Frame Format:

↳ used in LAN

↳ uses CSMA/CD (Mac Protocol)

e.g. 10 Base 2 , 10 Base 5 , 100 Base Tx , 10 Gb Base T.
Thin Thick

○ Topology = Bus

○ 1 mbits/s to 900 Giga bits /s.

↗ Destination Source Address (MAC Addresses)

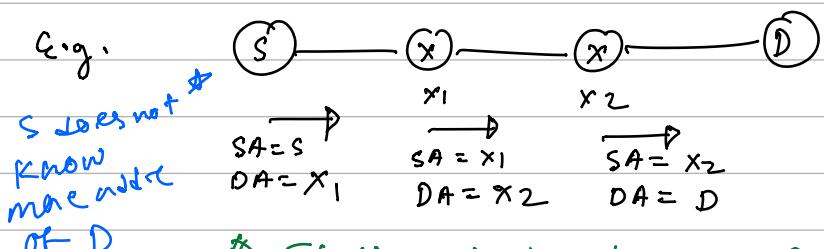
Preamble	SFD	DA	SA	length	Data	CRC
7 B	1 B	6 B	6 B	2 B	96 - 1500 B	4 B

Physical Layer Link Layer

* Preamble : 101010 ... 10 7 B

SFD : 10101011

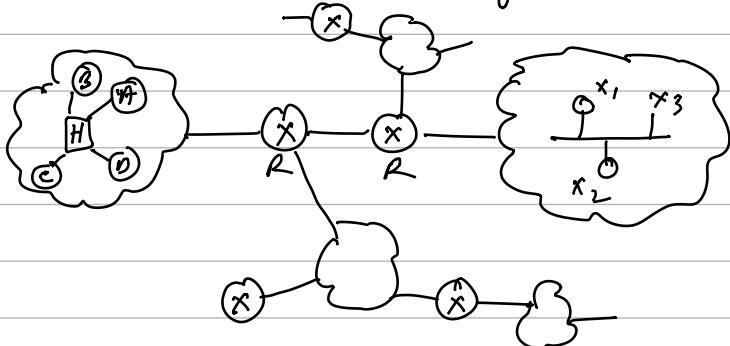
} these 2 are just to make the device alert that data is coming.



* Finding Next node i.e. DA by ARP protocol.

Network Layer :

↳ responsibility Host to Host delivery (source to destination)
does so using logical Address (IP)



IP has info of
① Network that host
Belong
② Host

○ Routing methods : ① RIP (Routing Information Protocol)
major responsibility is shortest path finding
② OSPF

○ Fragmentation :

○ Congestion Control :

in 1960 ARPANET gave the classes.

□ class A : $\underbrace{0}_{\text{Network } 2^7} \dots : \text{Host } 2^{29}$

class B : $\underbrace{10}_{\text{Network } 2^{19}} \dots : \text{Host } 2^{16}$

class C : $\underbrace{110}_{\text{Network } 2^{21}} \dots : \text{Host } 2^8$

class D : $\underbrace{1110}_{\text{Given to Groups, Not given to any Network.}}$

↳ Military purposes.

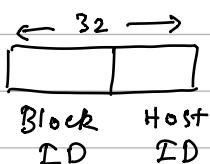
class E : $\underbrace{1111}_{\text{After 1993 classless Addressing become popular.}}$

Classless Addressing :

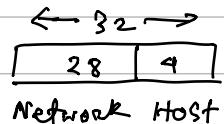
notation

$x.y.z.w/n$

↑ mark (The part that represent Network)



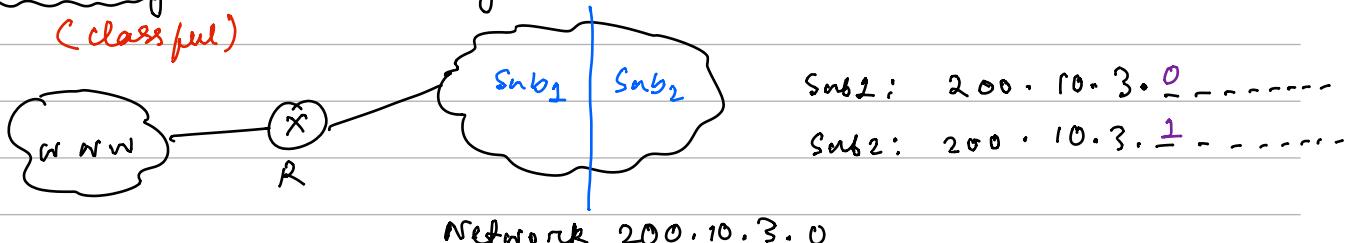
e.g. $200.10.20.9/28 \Rightarrow$



* Unlike Classful Addressing, where mark is predefined for a class, here mark is with the IP.

Subnetting: Divide Big network into small network.

(classful)



For Sub 1

Subnet ID : 200.10.3.0 ✓ * usable addresses 126

Broadcast : 200.10.3.127 ✓

For Sub 2

Subnet ID : 200.10.3.128 ✓ * usable addresses 126

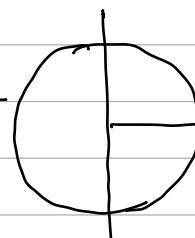
Broadcast : 200.10.3.255 ✓

Variable length Subnetting:

200.10.3.0

ID → 200.10.3.0

B_{nc} → 200.10.3.127



200.10.3.10

200.10.3.128

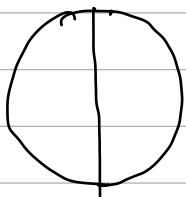
200.10.3.191

200.10.3.11

200.10.3.192

200.10.3.255

0 Subnetting in CIDR (Classless InterDomain Routing) :



$195 \cdot 10 \cdot 20 \cdot 128 / 26$

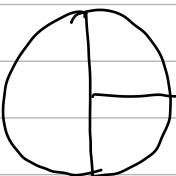
$\boxed{26 \ 16}$

Network ID = $195 \cdot 10 \cdot 20 \cdot 1 \underline{0} \underline{0} \underline{0} \underline{0} \underline{0} \underline{0}$

$195 \cdot 10 \cdot 20 \cdot 10 \underline{0}$ $195 \cdot 10 \cdot 20 \cdot 10 \underline{1}$

0 variable Subnetting in CIDR :

$245 \cdot 248 \cdot 1000 \underline{0} \dots 0$
 $\rightarrow 245 \cdot 248 \cdot 128 \cdot 0 / 21$
 $\rightarrow 245 \cdot 248 \cdot 135 \cdot 255 / 21$



$245 \cdot 248 \cdot 1000 \underline{1} \underline{0} \dots 0$

$\nearrow 245 \cdot 248 \cdot 136 \cdot 0 / 22$
 $\nearrow 245 \cdot 248 \cdot 139 \cdot 255 / 22$

$\rightarrow 245 \cdot 248 \cdot 140 \cdot 0 / 22$
 $\rightarrow 245 \cdot 248 \cdot 143 \cdot 255 / 22$

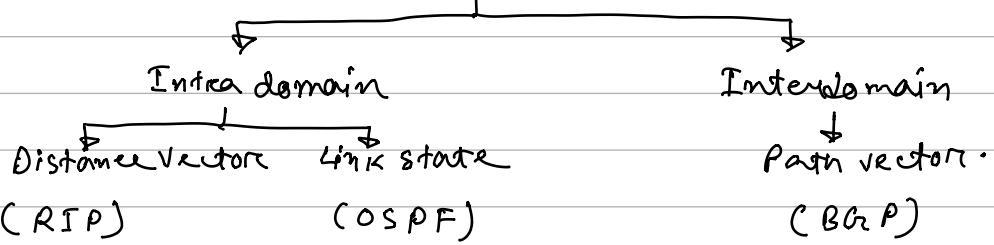
original network $245 \cdot 248 \cdot 128 \cdot 0 / 20$

0 Routing Protocols :

Idea : Sending packet to right path (optimal)

* Routing Table Contains the state of the whole Network.
 ↳ can be static or dynamic.

To generate Dynamic
 Routing table used → Routing Protocols



Routing Information Protocol Open Shortest Path First

- Distance vector: All node have a vector structure
 if any value updated The
 Node share the cost vector with
 its neighbors only.

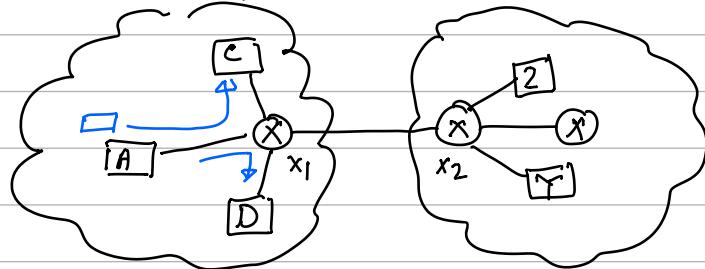
list	Cost	Next
N ₁	0	N ₁
N ₂	1	N ₂
N ₃	∞	-
N ₄	∞	-

- Link State Routing: single source shortest path for all nodes.

- Address Resolution Protocol (ARP): IP → MAC

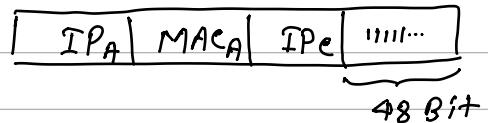
↳ Layer 3 Network Layer Protocol Logical → Physical address.

* ARP Request is always Broadcast Type.
 Reply is unicast Type.



ARP request from A to find mac of C.

A will Broadcast



○ ARP can be done Host → Host Router → Host
 Host → Router Router → Router.

* if A request {IP_A, MAC_A, IP_C, 111... } → {IP_A, MAC_A, IP_C, 111... }
 goes to Router_{x1}
 ↓ goes to Router_{x2}
 {IP_A, MAC_A, IP_C, 111... }

- Network Address Translation (NAT):

Many colleges have 10.8.0.1 IP.

Idea: All hostels have same room no. 001, 002, 101, 102 etc.

Privatize the Network

NAT Translation Table	
Private IP	Public IP
10.8.0.1	101.10.22.33

This is Destination IP (Google.com)

