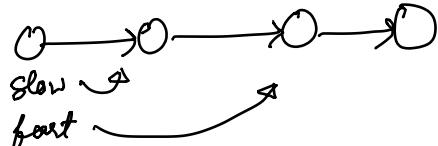
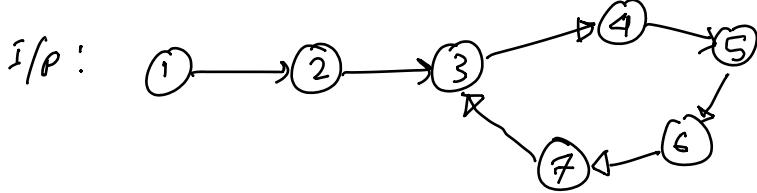


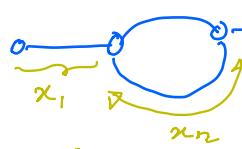
Linked List

→ Find Linked List Cycle



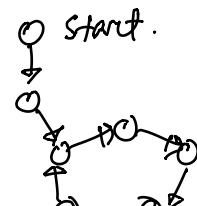
* if there's a cycle they'll meet for sure.

✓ To find the starting position of the cycle.



$$n_1 = x_2 \\ \text{i.e. } l_1 = k - l_2$$

lets assume
a 5 node cycle
and 2 node before
that.



* at his position fast & slow meet.

$$l_1 + l_2 + n_1 k = \text{slow's travel.}$$

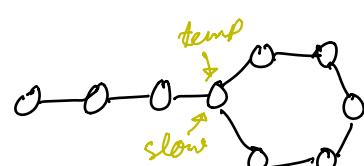
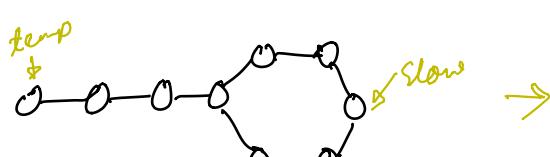
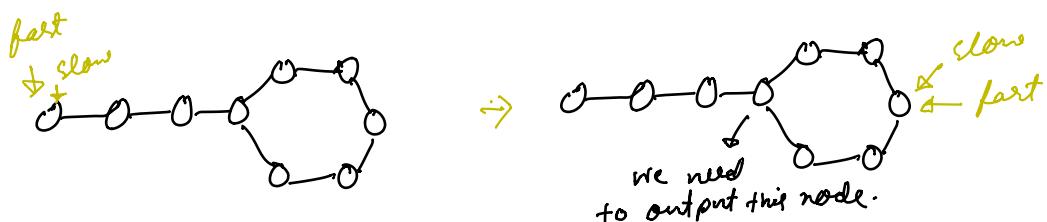
$$l_1 + l_2 + n_2 k = \text{fast's travel.}$$

$$(l_1 + l_2 + n_1 k) 2 = l_1 + l_2 + n_2 k.$$

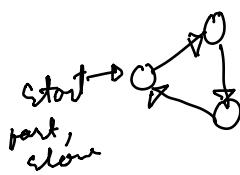
$$l_1 + l_2 + 2n_1 k = n_2 k.$$

$$\text{if } n_1 = 0 \quad l_1 + l_2 = n k.$$

$k \rightarrow$ size of loop
no. of nodes in
loop.

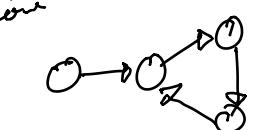


* slow & temp will meet
for sure;
return temp.

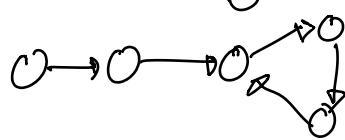


$$k=3 \quad l_1=0 \quad l_2=3$$

slow 3 fast 6



slow 1 + 2 fast $((1+2) \times 2 - 1) \% 3$



slow 2 + k fast $((2+k) \times 2 - 2) \% 3$

$k=1 \quad 3 \quad 1 \quad \checkmark$

$k=2 \quad 4 \quad 6 \% 3 = 0 \quad \times$

$k=3 \quad 5 \quad 10 - 2 \% 3 = 8 \% 3 = 2 \quad \times$

$k=4 \quad 6 \quad 12 - 2 \% 3 = 2 \quad \checkmark$

$k=5 \quad 6 \quad 12 - 2 \% 3 = 1 \quad \checkmark$

$$\begin{array}{ll} \text{alone} & \text{fast} \\ 4+k & ((4+k) \times 2 - 4) \% 3 \\ k=1 & 10 - 4 \% 3 = 0 \quad \times \\ k=2 & 12 - 4 \% 3 = 2 \quad \checkmark \\ k=3 & 14 - 4 \% 3 = 2 \quad \checkmark \\ k=4 & 16 - 4 \% 3 = 2 \quad \checkmark \\ k=5 & 18 - 4 \% 3 = 2 \quad \checkmark \end{array}$$

$$(5+k) \times 2 - 5 \% 3 = 12$$

$k=1$

$$\begin{array}{ll} \text{alone} & \text{fast} \\ 6+k & ((6+k) \times 2 - 6) \% 3 = 12 \\ k=1 & 12 - 6 \% 3 = 1 \quad \checkmark \\ k=2 & 14 - 6 \% 3 = 0 \quad \times \\ k=3 & 16 - 6 \% 3 = 2 \quad \times \\ k=4 & 18 - 6 \% 3 = 1 \quad \checkmark \\ k=5 & \dots \end{array}$$

$$(6+k) \times 2 - 6 \% 3 = 12$$

$$6+k = (6+k) \times 2 \% 3$$

$$k \% 3 = \text{multiple of } 3$$



Remove Nodes from LL :-

i/p: $(5) \rightarrow (3) \rightarrow (13) \rightarrow (3) \rightarrow (8)$ o/p: $(13) \rightarrow (8)$

Approach:-

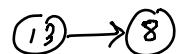
1) using Stack space $O(N)$.

2) Reverse LL $(8) \rightarrow (3) \rightarrow (13) \rightarrow (3) \rightarrow (5) \rightarrow (8) \rightarrow (3)$
then return reverse of ans.

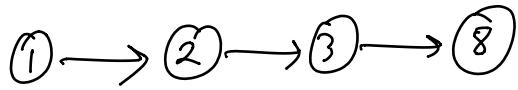
3) Simple Recursion



recursive Remove (head \rightarrow next) will return.



→ Reservoir Sampling :-



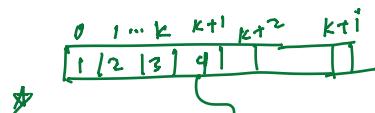
function getRandom() generate a random number from LL.

```

ListNode * Result;
count = 0; temp = head;
while (temp != null) {
    if (rand() % count < 1 / count)
        result = temp->val;
    count++;
    temp = temp->next;
}

```

#Math_algo_numberTheory/Reservoir sampling



* Reservoir sampling: $\frac{1}{k}$ chance of keeping each node. The probability of keeping node i is $\frac{k}{k+i}$. In our question, $k=1$ & i is the count variable.

→ Merge 2 Sorted List

the usual merge ; using recursion -

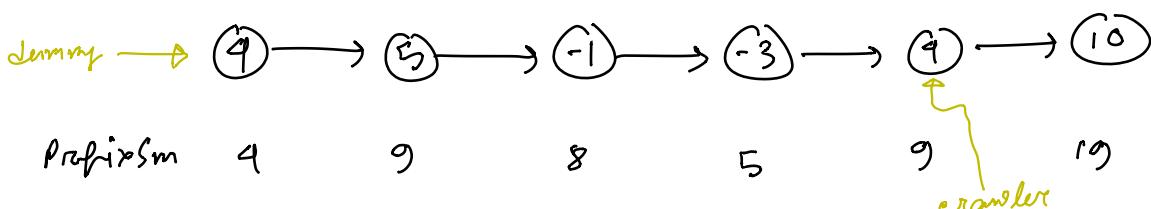
→ Delete Zero Sum

i/p : ④ → ⑤ → ⑥ → ⑦ → ⑧ → ⑨ → ⑩
 $\underbrace{\quad}_{=0}$

o/p : ④ → ⑤ → ⑩

use map to store prefix sum; unordered_map<int, ListNode*
PrefixSum Node>

mp[0] mp[9] mp[9] mp[8] mp[5]

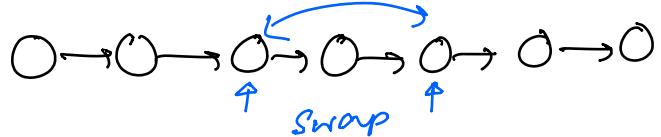


if (mp.find(prefixSum) != mp.end())

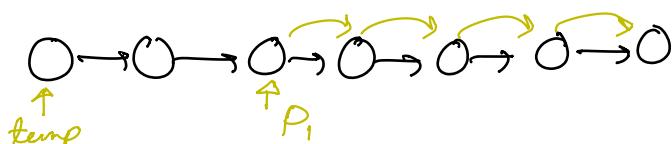
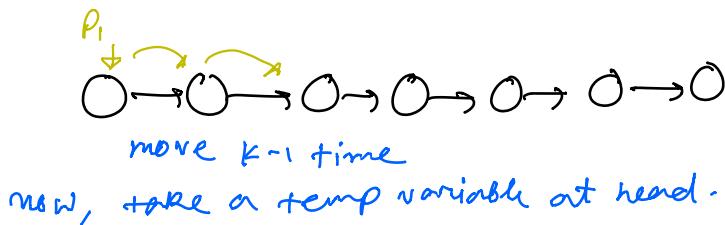
mp[prefixSum] → next = crawler → next.

→ Swapping Kth Node from the beginning & end

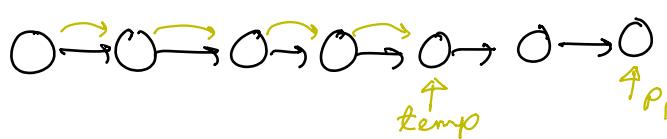
$K = 3$



#Trick



move P_1 until reaches the end,



* when P_1 reaches end
temp reaches K^{th} position from end.

→ Sum Linked List

i/p :- $(\overset{7}{\circlearrowleft}) \rightarrow (\overset{8}{\circlearrowleft}) \rightarrow (\overset{2}{\circlearrowleft}) \rightarrow (\overset{4}{\circlearrowleft})$
 $(\overset{5}{\circlearrowleft}) \rightarrow (\overset{3}{\circlearrowleft}) \rightarrow (\overset{8}{\circlearrowleft})$
o/p : $(\overset{8}{\circlearrowleft}) \rightarrow (\overset{3}{\circlearrowleft}) \rightarrow (\overset{6}{\circlearrowleft}) \rightarrow (\overset{2}{\circlearrowleft})$

The sum $\begin{array}{r} 7 \ 8 \ 2 \ 4 \\ 5 \ 3 \ 8 \\ \hline 8 \ 3 \ 6 \ 2 \end{array}$

✓ Reverse the i/p LL then add & store in another temp LL ($(\overset{2}{\circlearrowleft}) \rightarrow (\overset{6}{\circlearrowleft}) \rightarrow (\overset{3}{\circlearrowleft}) \rightarrow (\overset{8}{\circlearrowleft})$)
then return the reverse of temp.

→ Swap Nodes

#algorithm/recursion i/p: $(\overset{1}{\circlearrowleft}) \rightarrow (\overset{2}{\circlearrowleft}) \rightarrow (\overset{3}{\circlearrowleft}) \rightarrow (\overset{4}{\circlearrowleft})$ o/p: $(\overset{2}{\circlearrowleft}) \rightarrow (\overset{1}{\circlearrowleft}) \rightarrow (\overset{5}{\circlearrowleft}) \rightarrow (\overset{3}{\circlearrowleft})$

✓



(use recursion)

$L_1 \rightarrow \text{next} , L_2 \rightarrow \text{next} = \text{solver}(L_2 \rightarrow \text{next}), L_1$

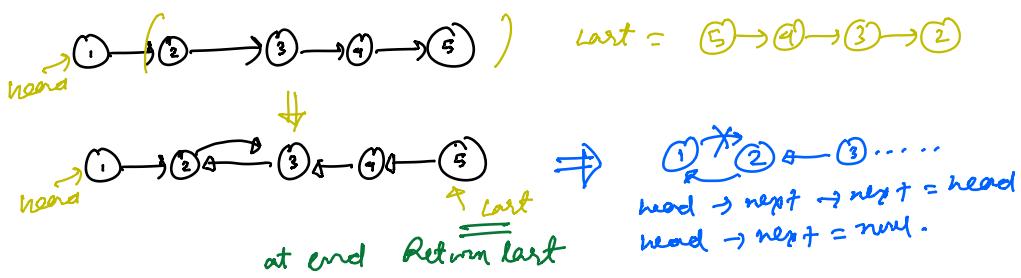
#algorithm/Linked List

→ [Reverse Linked List] :-

```

ListNode* reverse (ListNode* head) {
    if (!head || !head->next) return head;
    ListNode* last = reverse (head->next);
    head->next->next = head;
    head->next = NULL;
    return last;
}

```



→ Reverse Linked List II

using iterative method without recursion.

① → ② → ③ → ④ → ⑤ → ⑥ → ⑦ → ⑧ Reverse from node ③ to ⑥.

i) > ... in each head is also changed, return dummy → next.

ii) assign prev + curr pointer

① → ② → ③ → ④ → ⑤ → ⑥ → ⑦ → ⑧
 prev curr

* Prev points to the starting of reverse can be dummy as well if reversal start from head.

iii) > $\text{ListNode}^* \text{ temp} = \text{prev} \rightarrow \text{next}$

○ $\text{prev} \rightarrow \text{next} = \text{curr} \rightarrow \text{next}$

○ $\text{curr} \rightarrow \text{next} = \text{curr} \rightarrow \text{next} \rightarrow \text{next}$

○ $\text{prev} \rightarrow \text{next} \rightarrow \text{next} = \text{temp}$

