Computational Science on Many-Core Architectures

360.252

Josef Weinbub Karl Rupp



Institute for Microelectronics, TU Wien
http://www.iue.tuwien.ac.at/



Zoom Channel 621 2711 2607 Wednesday, November 29, 2023

Discrete Poisson Equation $-\Delta_h u(x,y) = f_h(x,y)$

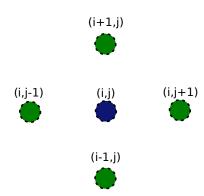
- Rectangular grid
- Assume homogeneous grid spacing h
- Discretize each coordinate direction separately

$$\Delta u(x,y) = u_{xx}(x,y) + u_{yy}(x,y)$$

$$\approx \frac{u(x-h,y) - 2u(x,y) + u(x+h,y)}{h^2} + \frac{u(x,y-h) - 2u(x,y) + u(x,y+h)}{h^2}$$

$$= \frac{u(x-h,y) + u(x+h,y) - 4u(x,y) + u(x,y-h) + u(x,y+h)}{h^2}$$

$$=: \Delta_h u(x,y)$$



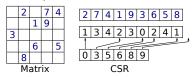
System Matrix Structure

- Interior nodes have 5 nonzero entries
- Boundary nodes have 1 to 4 nonzero entries
- Problem: CSR format requires offset index for each row

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Matrix						CSR									

System Matrix Structure

- Interior nodes have 5 nonzero entries
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Generic GPU matrix assembly skeleton

- 1. Count the nonzero entries for each row
- 2. Deduce offsets for CSR
- 3. Populate column and values arrays

Step 1: Count Nonzeros for an $N \times M$ grid

```
kernel void count nnz(int *row offsets, int N, ...) {
for (int row = blockDim.x * blockIdx.x + threadIdx.x;
        row < N*M;
        row += gridDim.x * blockDim.x)
       int nnz for this node = 1;
       int i = row / N;
       int j = row % N;
       if (i > 0) nnz_for_this_node += 1;
       if (j > 0) nnz_for_this_node += 1;
       if (i < N-1) nnz_for_this_node += 1;</pre>
       if (j < M-1) nnz for this node += 1;
       row offsets[row] = nnz for this node;
```

Parallel Primitives

Prefix Sum

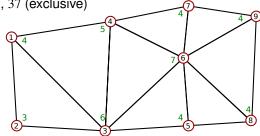
- Inclusive: Determine $y_i = \sum_{k=1}^i x_k$
- Exclusive: Determine $y_i = \sum_{k=1}^{i-1} x_k$, $y_1 = 0$

Example

- x: 4, 3, 6, 5, 4, 7, 4, 4, 4
- y: 4, 7, 13, 18, 22, 29, 33, 37, 41 (inclusive)
- y: 0, 4, 7, 13, 18, 22, 29, 33, 37 (exclusive)

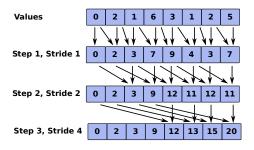
Applications

- Sparse matrix setup
- Graph algorithms



Parallel Primitives

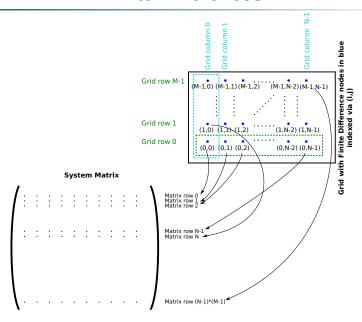
Prefix Sum Implementation



```
for(int stride = 1; stride < blockDim.x; stride *= 2)
{
    __syncthreads();
    shared_m[threadIdx.x] = my_value;
    __syncthreads();
    if (threadIdx.x >= stride)
        my_value += shared_m[threadIdx.x - stride];
}
    __syncthreads();
shared_m[threadIdx.x] = my_value;
```

Step 3: Assembly for an $N \times M$ grid

```
kernel void assembleA(int *row offsets, int N, ...) {
for (int row = blockDim.x * blockIdx.x + threadIdx.x;
          row < N*M;
          row += gridDim.x * blockDim.x) {
     int i = row / N;
     int i = row % N;
     int this row offset = row offsets[row];
     // diagonal entry
     col indices[this row offset] = i * N + j;
     values[this row offset] = 4:
     this row offset += 1;
     if (i > 0) { // bottom neighbor
       col indices[this row offset] = (i-1) * N + j;
       values[this row offset] = -1:
       this row offset += 1;
     if (j > 0) { /* similarly */ }
     if (i < N-1) { /* similarly */ }</pre>
     if (j < M-1) { /* similarly */ }</pre>
```



Other Discretizations

Finite Volumes

- Iterations over vertices: Similar to finite differences
- Beware of advanced schemes (depends)

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Finite Volumes

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Finite Elements

- Iteration over elements (cells)
- BUT: write matrix entries associated with vertices
- Avoid race conditions
- More preparation required before we can address this

