



1a	<pre> Algorithm f(n) k ← 0 i ← 1 while (i &lt; n)     j ← 1     while (j ≤ i)         k ← k + 1         j ← j + 1     i ← 2 * i return k </pre>	Derive the return value of $f(n)$ of the above algorithm. Use asymptotic notations; tighter O-notation or $\Theta$ -notation.	06
1b	<pre> Algorithm f(n) if (n ≤ 1) return 1 k ← f(n-1) k ← k + f(n-1) return k + 1 </pre>	Derive the number of times the function $f$ would be invoked for $f(n)$ . Use tighter O-notation or $\Theta$ -notation.	06
1c	<p>Consider a stack with a capacity to hold 'z' number of elements at the maximum and the following four operations.</p> <p><b>PUSH(S, x)</b> // Pushes element x onto the stack S if the stack is not full, in constant time</p> <p><b>POP(S)</b> // Pops top element from the stack S if the stack is not empty, in constant time</p> <p><b>MULTIPOP(S, k)</b> // Pops at most k elements from the stack while not <b>STACK-EMPTY(S)</b> and <math>k &gt; 0</math> do     <b>POP(S)</b>, <math>k \leftarrow k - 1</math></p> <p><b>Backup(S)</b> // Backs up all the elements in the stack, which is not more than z number of elements in linear time. Backup(S) operation is not invoked by the user, but is automatically invoked after every z number of other operations.</p> <p>Derive a constant amortized costs for each operation using the accounting method of amortized analysis. Explanation of each amortized cost is needed and avoid using other methods to find the amortized costs.</p>		08

