# **Coding Practice-7**

# 1.Minimum Path Sum

Given a m x n grid filled with non-negative numbers, find a path from top left to bottom right, which minimizes the sum of all numbers along its path.

Note: You can only move either down or right at any point in time.

## Example 1:

1	3	1
1	5	1
4	2	1

```
Input: grid = [[1,3,1],[1,5,1],[4,2,1]]
Output: 7
Explanation: Because the path 1 \rightarrow 3 \rightarrow 1 \rightarrow 1 \rightarrow 1 minimizes the sum.
```

```
class Solution {  public int minPathSum(int[][] grid) \{ \\ int m = grid.length, n = grid[0].length; \\ for (int j = 1; j < n; j++) \{ \\ grid[0][j] += grid[0][j-1]; \\ \}
```

```
for (int i = 1; i < m; i++) {
       grid[i][0] += grid[i - 1][0];
     }
     for (int i = 1; i < m; i++) {
       for (int j = 1; j < n; j++) {
          grid[i][j] += Math.min(grid[i - 1][j], grid[i][j - 1]);
       }
     }
     return grid[m - 1][n - 1];
  }
}
public class Main {
  public static void main(String[] args) {
     Solution solution = new Solution();
     int[][] grid = {
       \{1, 3, 1\},\
       \{1, 5, 1\},\
       {4, 2, 1}
     };
     int result = solution.minPathSum(grid);
     System.out.println("Minimum Path Sum: " + result);
  }
}
Output: Minimum Path Sum: 7
Time complexity: O(n*m)
```

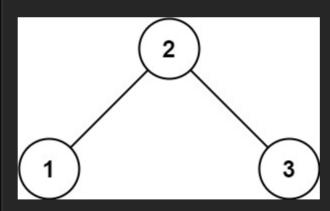
2. Validate binary search tree

Given the root of a binary tree, determine if it is a valid binary search tree (BST).

A valid BST is defined as follows:

- The left subtree of a node contains only nodes with keys less than the node's key.
- The right subtree of a node contains only nodes with keys greater than the node's key.
- Both the left and right subtrees must also be binary search trees.

## Example 1:



Input: root = [2,1,3]
Output: true

```
class Solution {
    public boolean isValidBST(TreeNode root) {
        return validate(root, Long.MIN_VALUE, Long.MAX_VALUE);
    }
    private boolean validate(TreeNode node, long low, long high) {
        if (node == null) return true;
        if (node.val <= low || node.val >= high) return false;
        return validate(node.left, low, node.val) && validate(node.right, node.val, high);
    }
    public static void main(String[] args) {
        TreeNode root = new TreeNode(2, new TreeNode(1), new TreeNode(3));
        Solution sol = new Solution();
    }
}
```

```
System.out.println(sol.isValidBST(root));
  }
}
class TreeNode {
  int val:
  TreeNode left;
  TreeNode right;
  TreeNode() {}
  TreeNode(int val) { this.val = val; }
  TreeNode(int val, TreeNode left, TreeNode right) {
     this.val = val;
     this.left = left:
     this.right = right;
  }
}
Output: true
Time Complexity:O(n)
```

#### 3. Word Ladder

A transformation sequence from word beginWord to word endWord using a dictionary wordList is a sequence of words beginWord -> s<sub>1</sub> -> s<sub>2</sub> -> ... -> s<sub>k</sub> such that:

• Every adjacent pair of words differs by a single letter.

• Every s<sub>i</sub> for 1 <= i <= k is in wordList. Note that beginWord does not need to be in wordList.

• s<sub>k</sub> == endWord

Given two words, beginWord and endWord, and a dictionary wordList, return the number of words in the shortest transformation sequence from beginWord to endWord, or 0 if no such sequence exists.

Example 1:

Input: beginWord = "hit", endWord = "cog", wordList = ["hot", "dot", "dog", "lot", "log", "cog"]

Output: 5

Explanation: One shortest transformation sequence is "hit" ->

"hot" -> "dot" -> "dog" -> cog", which is 5 words long.

```
import java.util.*;
class Solution {
  public int ladderLength(String beginWord, String endWord, List<String> wordList) {
    Set<String> wordSet = new HashSet<>(wordList);
    if (!wordSet.contains(endWord)) return 0;
    Queue<String> queue = new LinkedList<>();
    queue.add(beginWord);
    int level = 1;
    while (!queue.isEmpty()) {
       int size = queue.size();
       for (int i = 0; i < size; i++) {
         String currentWord = queue.poll();
         char[] wordChars = currentWord.toCharArray();
         for (int j = 0; j < wordChars.length; <math>j++) {
            char originalChar = wordChars[j];
            for (char c = 'a'; c \le 'z'; c++) {
              if (c == originalChar) continue;
              wordChars[j] = c;
              String nextWord = new String(wordChars);
              if (nextWord.equals(endWord)) return level + 1;
              if (wordSet.contains(nextWord)) {
                 queue.add(nextWord);
                 wordSet.remove(nextWord);
               }
            wordChars[j] = originalChar;
          }
       level++;
```

```
}
    return 0;
  }
  public static void main(String[] args) {
    Solution solution = new Solution();
    String beginWord = "hit";
    String endWord = "cog";
    List<String> wordList = Arrays.asList("hot", "dot", "dog", "lot", "log", "cog");
    System.out.println(solution.ladderLength(beginWord, endWord, wordList));
  }
}
Output: 5
Time complexity: O(n \times m^2)
4.Word Ladder II
A transformation sequence from word beginword to word endword using a dictionary
wordList is a sequence of words beginword \rightarrow s_1 \rightarrow s_2 \rightarrow \dots \rightarrow s_k such that:

    Every adjacent pair of words differs by a single letter.

    Every s<sub>i</sub> for 1 <= i <= k is in wordList. Note that beginword does not need to be in</li>

  wordList.
• s_k == endWord
Given two words, beginword and endword, and a dictionary wordList, return all the
shortest transformation sequences from beginWord to endWord, or an empty list if no such
sequence exists. Each sequence should be returned as a list of the words [beginWord, s_1, s_2,
..., s<sub>k</sub>].
Example 1:
  Input: beginWord = "hit", endWord = "cog", wordList =
   ["hot","dot","dog","lot","log","cog"]
  Output: [["hit","hot","dot","dog","cog"],
   ["hit","hot","lot","log","cog"]]
  Explanation: There are 2 shortest transformation sequences:
  "hit" -> "hot" -> "dot" -> "dog" -> "cog"
  "hit" -> "hot" -> "lot" -> "log" -> "cog"
```

```
Code:
```

```
import java.util.*;
class Solution {
  String b;
  HashMap<String, Integer> mpp;
  List<List<String>> ans;
  private void dfs(String word, List<String> seq) {
    if (word.equals(b)) {
       List<String> dup = new ArrayList<>(seq);
       Collections.reverse(dup);
       ans.add(dup);
       return;
     }
    int steps = mpp.get(word);
    int sz = word.length();
    for (int i = 0; i < sz; i++) {
       for (char ch = 'a'; ch \leq 'z'; ch++) {
         char[] replacedCharArray = word.toCharArray();
         replacedCharArray[i] = ch;
         String replacedWord = new String(replacedCharArray);
         if (mpp.containsKey(replacedWord) && mpp.get(replacedWord) + 1 == steps) {
            seq.add(replacedWord);
            dfs(replacedWord, seq);
            seq.remove(seq.size() - 1);
          }
       }
  public List<List<String>> findLadders(String beginWord, String endWord, List<String>
wordList) {
```

```
Set<String> st = new HashSet<>();
for (String word : wordList) {
  st.add(word);
}
Queue<String> q = new LinkedList<>();
b = beginWord;
q.add(beginWord);
mpp = new HashMap<>();
mpp.put(beginWord, 1);
st.remove(beginWord);
while (!q.isEmpty()) {
  String word = q.poll();
  int steps = mpp.get(word);
  for (int i = 0; i < word.length(); i++) {
    for (char ch = 'a'; ch <= 'z'; ch++) {
       char[] replacedCharArray = word.toCharArray();
       replacedCharArray[i] = ch;
       String replacedWord = new String(replacedCharArray);
       if (st.contains(replacedWord)) {
         q.add(replacedWord);
         st.remove(replacedWord);
         mpp.put(replacedWord, steps + 1);
       }
     }
  }
ans = new ArrayList<>();
if (mpp.containsKey(endWord)) {
  List<String> seq = new ArrayList<>();
  seq.add(endWord);
```

```
dfs(endWord, seq);
     }
     return ans;
  }
  public static void main(String[] args) {
     Solution solution = new Solution();
     String beginWord = "hit";
     String endWord = "cog";
     List<String> wordList = Arrays.asList("hot", "dot", "dog", "lot", "log", "cog");
     List<List<String>> ladders = solution.findLadders(beginWord, endWord, wordList);
     System.out.println("All shortest transformation sequences:");
     for (List<String> ladder : ladders) {
       System.out.println(ladder);
     }
  }
}
Output:
[hit, hot, dot, dog, cog]
[hit, hot, lot, log, cog]
Time complexity: O(n \cdot m + p \cdot m)
5. Course Schedule
```

```
There are a total of numCourses courses you have to take, labeled from 0 to numCourses
1. You are given an array prerequisites where prerequisites[i] = [a_i, b_i] indicates
that you must take course b<sub>1</sub> first if you want to take course a<sub>1</sub>.

    For example, the pair [0, 1], indicates that to take course 0 you have to first take

  course 1.
Return true if you can finish all courses. Otherwise, return false.
Example 1:
  Input: numCourses = 2, prerequisites = [[1,0]]
  Output: true
  Explanation: There are a total of 2 courses to take.
  To take course 1 you should have finished course 0. So it is
  possible.
Example 2:
  Input: numCourses = 2, prerequisites = [[1,0],[0,1]]
  Output: false
  Explanation: There are a total of 2 courses to take.
  To take course 1 you should have finished course 0, and to take
  course 0 you should also have finished course 1. So it is
  impossible.
```

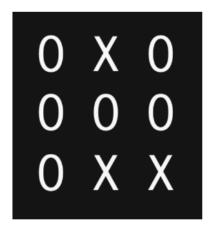
```
class Solution {
   public boolean canFinish(int numCourses, int[][] prerequisites) {
     List<List<Integer>> graph = new ArrayList<>();
     int[] inDegree = new int[numCourses];
     for (int i = 0; i < numCourses; i++) {
        graph.add(new ArrayList<>());
     }
     for (int[] prereq : prerequisites) {
        graph.get(prereq[1]).add(prereq[0]);
        inDegree[prereq[0]]++;
     }
     Queue<Integer> queue = new LinkedList<>();
     for (int i = 0; i < numCourses; i++) {
        if (inDegree[i] == 0) {</pre>
```

```
queue.add(i);
       }
     }
     int count = 0;
     while (!queue.isEmpty()) {
       int course = queue.poll();
       count++;
       for (int neighbor : graph.get(course)) {
          inDegree[neighbor]--;
         if (inDegree[neighbor] == 0) {
            queue.add(neighbor);
          }
       }
     return count == numCourses;
  public static void main(String[] args) {
     Solution sol = new Solution();
     int numCourses = 4;
     int[][] prerequisites = {{1, 0}, {2, 1}, {3, 2}};
     System.out.println(sol.canFinish(numCourses, prerequisites)); // Output: true
  }
}
Output: true
Time complexity: O(n²)
6. Design tic tac toe
```

A Tic-Tac-Toe board of size 3X3 is given after all the moves are played, i.e., all nine spots are filled. Find out if the given board is valid, i.e., is it possible to reach this board position after a set of moves or not.

Note that every arbitrarily filled grid of 9 spaces isn't valid, e.g., a grid filled with 3  $\mathbf{X}$  and 6  $\mathbf{O}$  isn't a valid situation because each player needs to take alternate turns.

Note: The game starts with X



## **Examples:**

```
class Solution {  int[][] \ win = \{\{0,1,2\}, \{3,4,5\}, \{6,7,8\}, \{0,3,6\}, \{1,4,7\}, \{2,5,8\}, \{0,4,8\}, \{2,4,6\}\}
```

```
public int isCWin(char[] board, char c) {
  int count = 0;
  for (int i = 0; i < 8; i++) {
     if (board[win[i][0]] == c \&\& board[win[i][1]] == c \&\& board[win[i][2]] == c) {
       count++;
     }
   }
  return count;
}
public boolean isValid(char[] board) {
  int xCount = 0, oCount = 0;
  for (char c : board) {
     if (c == 'X') {
       xCount++;
     \} else if (c == 'O') {
       oCount++;
     }
  int cx = isCWin(board, 'X');
  int co = isCWin(board, 'O');
  if (xCount != oCount + 1) {
     return false;
  }
  if (co == 1 \&\& cx == 0) {
     return true;
  }
  if (cx == 1 \&\& co == 0) {
     return true;
  }
  if (cx == 0 \&\& co == 0) {
```

```
return true;
}
return false;
}
public static void main(String[] args) {
    Solution solution = new Solution();
    char[] board = {'X', 'X', 'X', 'O', 'O', '', '', '', ''};
    System.out.println(solution.isValid(board));
}
```

**Output: Invalid** 

**Time Complexity: O(1)** 

7. Next permutation

A **permutation** of an array of integers is an arrangement of its members into a sequence or linear order.

• For example, for arr = [1,2,3], the following are all the permutations of arr: [1,2,3], [1,3,2], [2, 1, 3], [2, 3, 1], [3,1,2], [3,2,1].

The **next permutation** of an array of integers is the next lexicographically greater permutation of its integer. More formally, if all the permutations of the array are sorted in one container according to their lexicographical order, then the **next permutation** of that array is the permutation that follows it in the sorted container. If such arrangement is not possible, the array must be rearranged as the lowest possible order (i.e., sorted in ascending order).

- For example, the next permutation of arr = [1,2,3] is [1,3,2].
- Similarly, the next permutation of arr = [2,3,1] is [3,1,2].
- While the next permutation of arr = [3,2,1] is [1,2,3] because [3,2,1] does not have a lexicographical larger rearrangement.

Given an array of integers nums, find the next permutation of nums.

The replacement must be in place and use only constant extra memory.

```
class Solution {
```

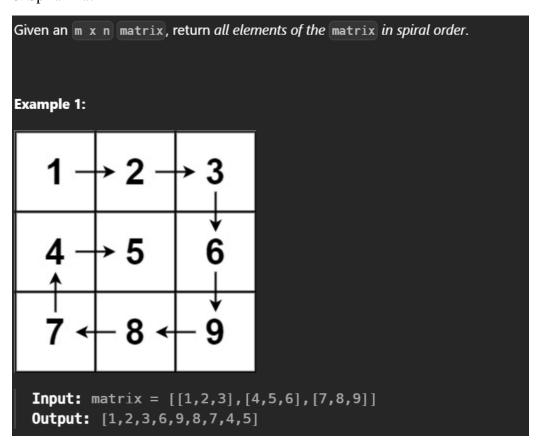
```
public void nextPermutation(int[] nums) {
  int n = nums.length;
  int i = n - 2;
  while (i \ge 0 \&\& nums[i] \ge nums[i + 1]) \{
     i--;
  }
  if (i >= 0) {
     int j = n - 1;
     while (j \ge 0 \&\& nums[j] \le nums[i]) \{
       j--;
     swap(nums, i, j);
  }
  reverse(nums, i + 1, n - 1);
private void swap(int[] nums, int i, int j) {
  int temp = nums[i];
  nums[i] = nums[j];
  nums[j] = temp;
}
private void reverse(int[] nums, int start, int end) {
  while (start < end) {
     swap(nums, start, end);
     start++;
     end--;
  }
}
public static void main(String[] args) {
  Solution solution = new Solution();
  int[] nums = \{1, 2, 3\};
```

```
solution.nextPermutation(nums);
System.out.println(Arrays.toString(nums));
}
```

# Output: [1, 3, 2]

# **Time Complexity:O(n)**

8. Spiral matrix



```
import java.util.ArrayList;
import java.util.List;
class Solution {
  public List<Integer> spiralOrder(int[][] matrix) {
    List<Integer> result = new ArrayList<>();
    if (matrix == null || matrix.length == 0) return result;
    int m = matrix.length;
```

```
int n = matrix[0].length;
  int top = 0, bot = m - 1, l = 0, r = n - 1;
  while (top \le bot && l \le r) {
     for (int i = 1; i \le r; i++) {
        result.add(matrix[top][i]);
     }
     top++;
     for (int i = top; i \le bot; i++) {
        result.add(matrix[i][r]);
     }
     r--;
     if (top <= bot) {
        for (int i = r; i >= l; i--) {
          result.add(matrix[bot][i]);
        }
        bot--;
     if (l <= r) {
        for (int i = bot; i >= top; i--) {
          result.add(matrix[i][l]);
        }
       1++;
     }
  return result;
public static void main(String[] args) {
  Solution solution = new Solution();
  int[][] matrix = \{\{1, 2, 3\}, \{4, 5, 6\}, \{7, 8, 9\}\};
  List<Integer> result = solution.spiralOrder(matrix);
```

}

```
System.out.println(result);
}

Output: [1, 2, 3, 6, 9, 8, 7, 4, 5]

Time Complexity:O(m*n)

9. Longest substring without repeating characters

Given a string s, find the length of the longest substring
```

```
Given a string s, find the length of the longest substring without repeating characters.

Example 1:

Input: s = "abcabcbb"
Output: 3
Explanation: The answer is "abc", with the length of 3.

Example 2:

Input: s = "bbbbb"
Output: 1
Explanation: The answer is "b", with the length of 1.

Example 3:

Input: s = "pwwkew"
Output: 3
Explanation: The answer is "wke", with the length of 3.

Notice that the answer must be a substring, "pwke" is a subsequence and not a substring.
```

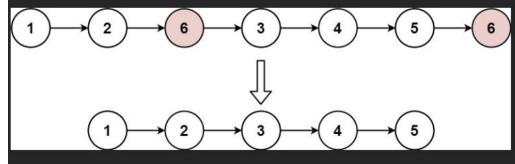
```
import java.util.HashSet;
import java.util.Set;
class Solution {
  public int lengthOfLongestSubstring(String s) {
    Set<Character> set = new HashSet<>();
    int maxLength = 0;
    int left = 0;
    for (int right = 0; right < s.length(); right++) {</pre>
```

```
if (!set.contains(s.charAt(right))) {
          set.add(s.charAt(right));
         maxLength = Math.max(maxLength, right - left + 1);
        } else {
          while (s.charAt(left) != s.charAt(right)) {
            set.remove(s.charAt(left));
            left++;
          }
          set.remove(s.charAt(left));
          left++;
          set.add(s.charAt(right));
       }
     }
     return maxLength;
  public static void main(String[] args) {
     Solution solution = new Solution();
     String s = "abcabcbb";
     System.out.println(solution.lengthOfLongestSubstring(s));\\
  }
}
Output:3
Time Complexity:O(n)
```

10. Remove linked list elements

Given the head of a linked list and an integer val, remove all the nodes of the linked list that has Node.val == val, and return the new head.

# Example 1:



```
Input: head = [1,2,6,3,4,5,6], val = 6
Output: [1,2,3,4,5]
```

## Example 2:

```
Input: head = [], val = 1
Output: []
```

```
class Solution {
  public ListNode removeElements(ListNode head, int val) {
    ListNode temp = null;
    ListNode temp2 = head;

  while(temp2 != null){
    if(temp2.val == val){
        if(temp == null){
        if(head.next == null){
            return null;
        }
        head = head.next;
        temp2 = head;
    }
    else if(temp2.next == null){
```

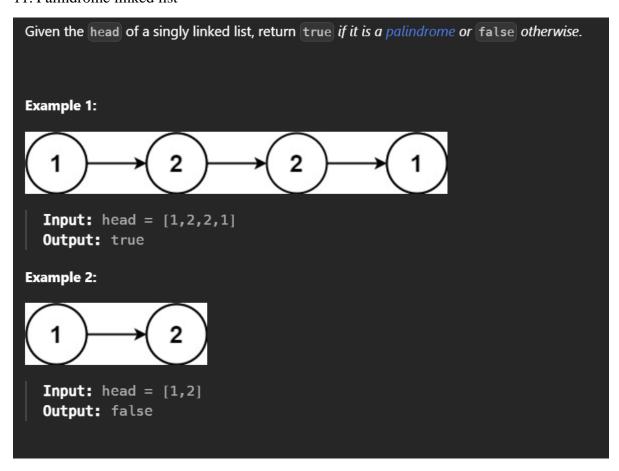
```
temp.next = null;
         break;
       }
       else{
         temp2 = temp2.next;
         temp.next = temp2;
       }
     }
    else{
       temp = temp2;
       temp2 = temp.next;
     }
  }
  return head;
public static void main(String[] args) {
  Solution solution = new Solution();
  ListNode head = new ListNode(1);
  head.next = new ListNode(2);
  head.next.next = new ListNode(6);
  head.next.next.next = new ListNode(3);
  head.next.next.next.next = new ListNode(4);
  head.next.next.next.next.next = new ListNode(5);
  head.next.next.next.next.next.next = new ListNode(6);
  int val = 6;
  ListNode result = solution.removeElements(head, val);
  while (result != null) {
    System.out.print(result.val + " ");
    result = result.next;
```

```
}
}
class ListNode {
  int val;
  ListNode next;
  ListNode(int x) {
    val = x;
    next = null;
  }
}
```

# Output: [1, 2, 3, 4, 5]

# Time Complexity:O(n)

# 11. Palindrome linked list



```
class Solution {
  public boolean isPalindrome(ListNode head) {
```

```
List<Integer> list = new ArrayList<>();
     while (head != null) {
       list.add(head.val);
       head = head.next;
     }
     int l = 0;
     int r = list.size() - 1;
     while (1 < r \&\& list.get(1).equals(list.get(r))) {
       1++;
       r--;
     return 1 \ge r;
  public static void main(String[] args) {
     Solution solution = new Solution();
     ListNode head = new ListNode(1);
     head.next = new ListNode(2);
     head.next.next = new ListNode(2);
     head.next.next.next = new ListNode(1);
     boolean result = solution.isPalindrome(head);
     System.out.println(result);
  }
class ListNode {
  int val;
  ListNode next;
  ListNode() {}
  ListNode(int val) { this.val = val; }
  ListNode(int val, ListNode next) { this.val = val; this.next = next; }
```

}

Outnutstan		
Output:true		
Time Complexity:O(n)		