

PROJECT REPORT

MESS MANAGEMENT SYSTEM



Submitted To:-

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PROBLEM STATEMENT

The current mess management system at our University JawaharLal Nehru University is manual,decentralised and inefficient.

CURRENT SITUATION:-

- The students have to sign the Mess Register to get the plates of the Food,but this sometimes leads to duplicity as the manager is unable to identify the legit student and some students eats illegally which lead to inefficiency.
- The data of the Mess is Decentralised so a student of the Hostel1 have to eat in his/her respective hostel only,sometimes it leads to problem as he/she have to go to far and sometimes they miss their classes as well.
- This also sometimes leads to lack of transparency in billing system.

DESIRED FUTURE STATE:-

- If the mess management system will get Online then the student have to enter their respective Details only then he/she will be albe to access the Food.So this will make our Mess Management System more efficient.
- If the Data of the hostels get centralised then the Student will be able to Eat in any food so he/she doesn't have to go to to long distance and can have food in nearby hostel.
- Our project will also lead to transparency in Mess Management System,accurate billing and reduction in wait times for Food.

INTRODUCTION

A database management system (DBMS) refers to the technology for creating and managing databases. Basically DBMS is a software tool to organize (create, retrieve, update and manage) data in a database.

The main aim of a DBMS is to supply a way to store up and retrieve database information that is both convenient and efficient. By data, we mean known facts that can be recorded and that have embedded meaning. Normally people use software such as DBASE IV or V, Microsoft ACCESS, or EXCEL to store data in the form of database.

Database systems are meant to handle large collection of information. Management of data involves both defining structures for storage of information and providing mechanisms that can do the manipulation of those stored information. Moreover, the database system must ensure the safety of the information stored, despite system crashes or attempts at unauthorized access.

This project is aim at computerizing the manual process of Mess Management System. The database is implemented using MySQL.

The project consists of namely four major entities

(i) Hostel, which will include details of hostels allot hostel to students and contains mess details.

(ii) Student, which will have the details of the students, and mess.

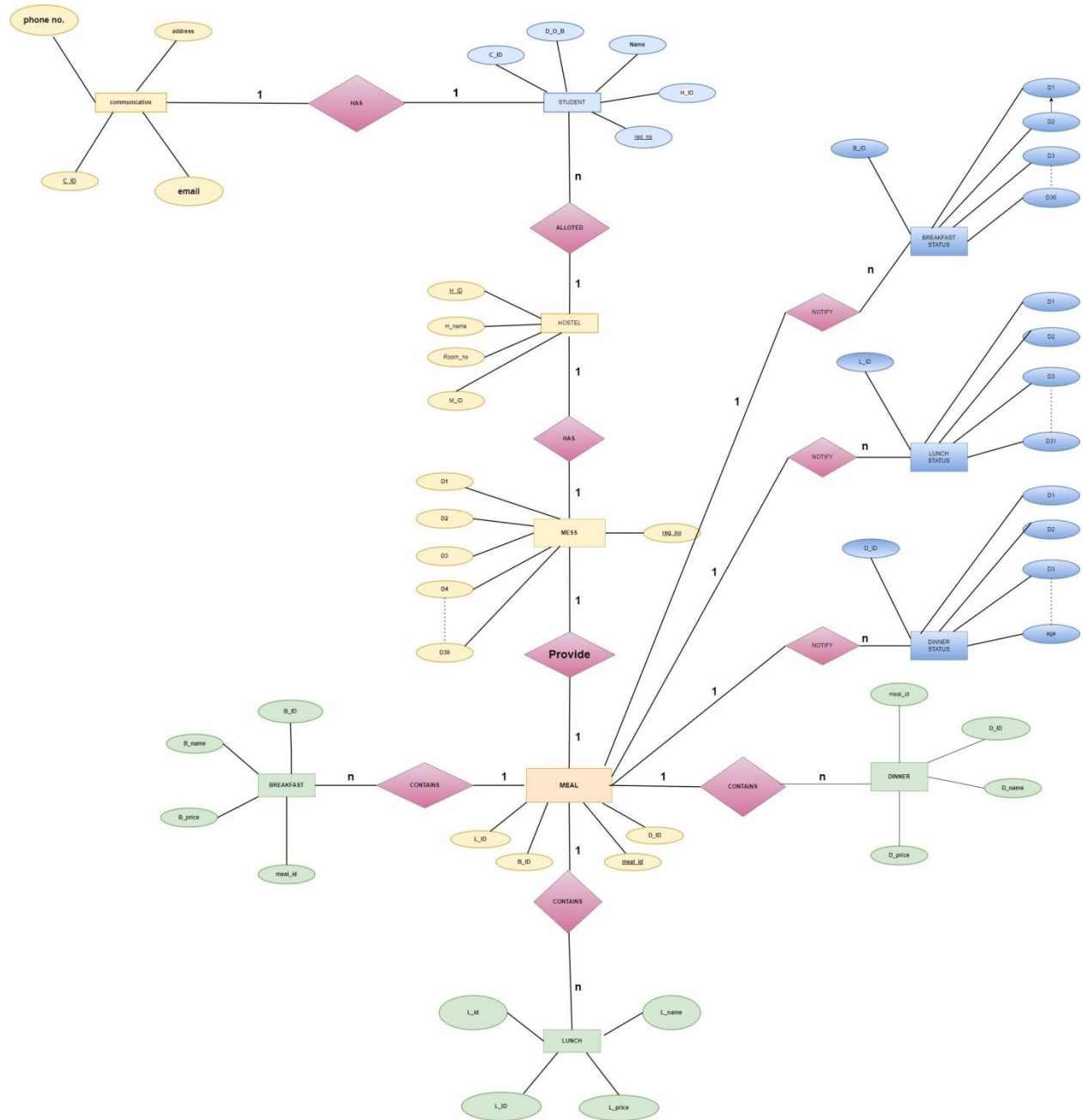
(iii) Mess, it contains the meal details, student's messid, status of student.

(iv)Meal :-This table includes the detail of meal in breakfast,lunch and dinner.

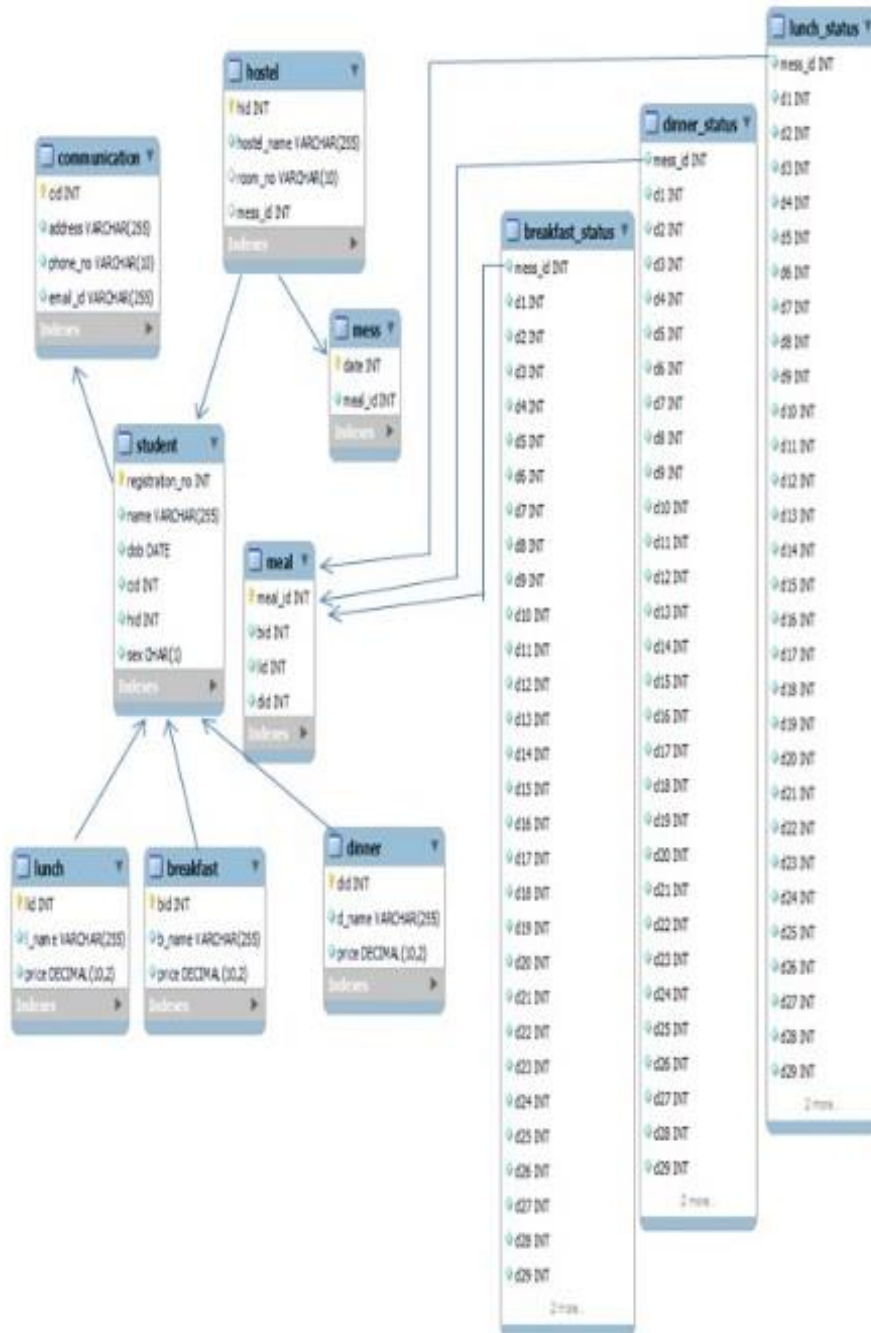
The services of a Mess Management Syste can include:

- Alloting the Hostel to a new Student.**
- Preparing the meal to for the month for respective BreakFast,Lunch and Dinner.**
- Serving the food to student.**
- Checking the status of the student i.e; whether he/she have taken the meal.**
- Preparing the Monthly Bill of the Student.**

ER DIAGRAM



RELATIONAL MODEL



MAPPING

1. HOSTEL-STUDENT(1:N):-A hostel may have many students.
2. HOSTEL-MESS(1:1):-A hostel will have only one mess.
3. HOSTEL-MEAL(1:N):-A hostel provides many Meals.
4. MEAL-BREAKFAST(1:N):-A meal contains many breakfast items.
5. MEAL-LUNCH(1:N):-A meal contains many lunch items.
6. MEAL-BREAKFAST_STATUS(1:N):-A meal contains multiple breakfast_status.
7. MEAL-LUNCH_STATUS(1:N):-A meal contains multiple lunch_status.
8. MEAL-DINNER_STATUS(1:N):-A meal contains multiple dinner_status.
9. STUDENT-COMMUNICATION(1:1):-A student may have related one communication entity

NORMALISATION

Student: To determine the normal forms (NF) for the given set of attributes `reg_id`, `name`, `D_O_B`, `H_ID`, and `C_ID`, we need to analyze the functional dependencies and dependencies among the attributes.

Based on the provided information, let's examine each normal form:

1. First Normal Form (1NF):
 - 1NF requires that each attribute in a relation must have atomic values (no multivalued attributes or repeating groups).
2.
 - From the given set of attributes, it appears that each attribute contains atomic values.
3.
 - Therefore, the relation satisfies 1NF.

2. Second Normal Form (2NF):

- 2NF states that the relation should be in 1NF and there should be no partial dependencies.
- To identify partial dependencies, we need to know the functional dependencies between attributes.
- However, the provided information does not include any explicit functional dependencies.
- Without explicit dependencies, we cannot determine if there are any partial dependencies.
- Assuming there are no partial dependencies, we can consider the relation to satisfy 2NF.

3. Third Normal Form (3NF):

- 3NF states that the relation should be in 2NF, and there should be no transitive dependencies.
- To identify transitive dependencies, we need to know the functional dependencies between attributes.
 - From the given information, we can assume the following functional dependencies:
- `reg_id` \rightarrow `name`
- `reg_id` \rightarrow `D_O_B`
- `reg_id` \rightarrow `H_ID`
- `reg_id` \rightarrow `C_ID`
- There are no transitive dependencies present since each non-key attribute depends solely on the primary key attribute, `reg_id`.
- Therefore, the relation satisfies 3NF.

4. Boyce-Codd Normal Form (BCNF):

- BCNF is a stronger form of normalization that eliminates all non-trivial dependencies in a relation.
- To determine BCNF, we need to identify the candidate keys and non-trivial dependencies.
- From the given information, assuming `reg_id` is the primary key, the candidate key is `{reg_id}`.
- The functional dependencies are:
- `reg_id` \rightarrow `name`
- `reg_id` \rightarrow `D_O_B`
- `reg_id` \rightarrow `H_ID`
- `reg_id` \rightarrow `C_ID`

- The dependencies are non-trivial, and each non-key attribute depends on the candidate key.
- Therefore, the relation satisfies BCNF. In summary, based on the provided information and assumptions, the relation with attributes reg_id, name, D_O_B, H_ID, and C_ID satisfies 1NF, 2NF, 3NF, and BCNF.

Hostel :

1. First Normal Form (1NF):

- From the given functional dependencies, it appears that each attribute contains atomic values.
- Therefore, the functional dependencies satisfy 1NF.

2. Second Normal Form (2NF):

- Assuming H_ID is the candidate key, we have the following functional dependencies:
- $H_ID \rightarrow H_name, M_ID, room_no$
- There are no partial dependencies since each non-key attribute depends on the entire candidate key.
- Therefore, the functional dependencies satisfy 2NF.

3. Third Normal Form (3NF):

- Assuming H_ID is the candidate key, we have the following functional dependencies:
- $H_ID \rightarrow H_name, M_ID, room_no$
- There are no transitive dependencies since each non-key attribute depends directly on the candidate key.
- Therefore, the functional dependencies satisfy 3NF.

4. Boyce-Codd Normal Form (BCNF):

- Assuming H_ID is the candidate key, the functional dependencies are as follows:
- $H_ID \rightarrow H_name, M_ID, room_no$
- There are no non-trivial dependencies since each non-key attribute depends on the candidate key.
- Therefore, the functional dependencies satisfy BCNF. In summary, based on the given functional dependencies $H_ID \rightarrow H_name, M_ID, room_no$, the relation satisfies 1NF, 2NF, 3NF, and BCNF.

Communication:

1. First Normal Form (1NF):

- From the given functional dependencies, it appears that each attribute contains atomic values.
- Therefore, the functional dependencies satisfy 1NF.

2. Second Normal Form (2NF):

- Assuming C_ID is the candidate key; we have the following functional dependencies:
- $C_ID \rightarrow phone_no, email, address$
- There are no partial dependencies since each non-key attribute depends on the entire candidate key.

- Therefore, the functional dependencies satisfy 2NF.
3. Third Normal Form (3NF):
- Assuming C_ID is the candidate key, we have the following functional dependencies:
 - C_ID → phone_no, email, address
 - There are no transitive dependencies since each non-key attribute depends directly on the candidate key.
 - Therefore, the functional dependencies satisfy 3NF.
4. Boyce-Codd Normal Form (BCNF):
- Assuming C_ID is the candidate key, the functional dependencies are as follows:
 - C_ID → phone_no, email, address
 - There are no non-trivial dependencies since each non-key attribute depends on the candidate key.
 - Therefore, the functional dependencies satisfy BCNF. In summary, based on the given functional dependencies C_ID → phone_no, email, address, the relation satisfies 1NF, 2NF, 3NF, and BCNF.

Mess:

1. First Normal Form (1NF):
- From the given functional dependencies, it appears that each attribute contains atomic values.
 - Therefore, the functional dependencies satisfy 1NF.
2. Second Normal Form (2NF):
- Assuming reg_id is the candidate key, we have the following functional dependencies:
 - reg_id → D1, D2, D3, ..., D31
 - There are no partial dependencies since each non-key attribute depends on the entire candidate key.
 - Therefore, the functional dependencies satisfy 2NF.
3. Third Normal Form (3NF):
- Assuming reg_id is the candidate key, we have the following functional dependencies:
 - reg_id → D1, D2, D3, ..., D31
 - There are no transitive dependencies since each non-key attribute depends directly on the candidate key.
 - Therefore, the functional dependencies satisfy 3NF.
4. Boyce-Codd Normal Form (BCNF):
- Assuming reg_id is the candidate key, the functional dependencies are as follows:
 - reg_id → D1, D2, D3, ..., D31
 - There are no non-trivial dependencies since each non-key attribute depends on the candidate key.
 - Therefore, the functional dependencies satisfy BCNF. In summary, based on the given functional dependencies reg_id → D1, D2, D3, ..., D31, the relation satisfies 1NF, 2NF, 3NF, and BCNF.

Meal(Meal_id, BreakfastName, BreakfastPrice, BreakfastStatus, LunchName, LunchPrice, LunchStatus, DinnerName, DinnerPrice, DinnerStatus):-

1. First Normal Form (1NF):

- From the given functional dependencies, it appears that each attribute contains atomic values.
- Therefore, the functional dependencies satisfy 1NF.

2. Second Normal Form (2NF):

- Assuming meal_id is the candidate key, we have the following functional dependencies:
 - meal_id -> **BreakfastName, BreakfastPrice, BreakfastStatus, LunchName, LunchPrice, LunchStatus, DinnerName, DinnerPrice, DinnerStatus**
 - **BreakfastName -> BreakfastPrice, BreakfastStatus**
 - **LunchName -> LunchPrice, LunchStatus**
 - **DinnerName -> DinnerPrice, DinnerStatus**
- There are no partial dependencies since each non-key attribute depends on the entire candidate key.
- Therefore, the functional dependencies satisfy 2NF.

3. Third Normal Form (3NF):

- Assuming meal_id is the candidate key, we have the following functional dependencies:
 - **BreakfastName -> BreakfastPrice, BreakfastStatus**
 - **LunchName -> LunchPrice, LunchStatus**
 - **DinnerName -> DinnerPrice, DinnerStatus**
 - **meal_id -> B_ID, D_ID, L_ID**
- There are transitive dependencies since non-key attribute depends directly on the candidate key.
- Therefore, the functional dependencies satisfy 3NF.

4. Boyce-Codd Normal Form (BCNF):

- Assuming meal_id is the candidate key, the functional dependencies are as follows:
 - meal_id -> B_ID, D_ID, L_ID
- There are no non-trivial dependencies since each non-key attribute depends on the candidate key.
- Therefore, the functional dependencies satisfy BCNF. In summary, based on the given functional dependencies meal_id -> B_ID, D_ID, L_ID, the relation satisfies 1NF, 2NF, 3NF, and BCNF.

Breakfast(BID, Meal_id, BreakfastName, BreakfastPrice, BreakfastStatus):

1. First Normal Form (1NF):

- From the given functional dependencies, it appears that each attribute contains atomic values.
- Therefore, the functional dependencies satisfy 1NF.

2. Second Normal Form (2NF):

- Assuming meal_id is the candidate key, we have the following functional dependencies:
- meal_id -> price, B_ID, B_name, BStatus
- There are no partial dependencies since each non-key attribute depends on the entire candidate key.
- Therefore, the functional dependencies satisfy 2NF.

3. Third Normal Form (3NF):

- Assuming meal_id is the candidate key, we have the following functional dependencies:
- B_ID, -> price, B_name, Breakfaststatus
- There are transitive dependencies since non-key attribute depends directly on the candidate key.
- Therefore, the functional dependencies satisfy 3NF.

BreakFast(B_ID, Meal_d, BreakFastName, BreakFastPrice)

4. Boyce-Codd Normal Form (BCNF):

- Assuming meal_id is the candidate key, the functional dependencies are as follows:
- meal_id -> price, B_ID, B_name
- There are no non-trivial dependencies since each non-key attribute depends on the candidate key.
- Therefore, the functional dependencies satisfy BCNF. In summary, based on the given functional dependencies meal_id -> price, B_ID, B_name, the relation satisfies 1NF, 2NF, 3NF, and BCNF.

Lunch:

1. First Normal Form (1NF):

- From the given functional dependencies, it appears that each attribute contains atomic values.
- Therefore, the functional dependencies satisfy 1NF.

2. Second Normal Form (2NF):

- Assuming meal_id is the candidate key, we have the following functional dependencies:
- meal_id -> L_price, L_ID, L_name
- There are no partial dependencies since each non-key attribute depends on the entire candidate key.
- Therefore, the functional dependencies satisfy 2NF.

3. Third Normal Form (3NF):

- Assuming meal_id is the candidate key, we have the following functional dependencies:
- meal_id -> L_price, L_ID, L_name
- There are no transitive dependencies since each non-key attribute depends directly on the candidate key.
- Therefore, the functional dependencies satisfy 3NF.

4. Boyce-Codd Normal Form (BCNF):

- Assuming meal_id is the candidate key, the functional dependencies are as follows:
- meal_id -> L_price, L_ID, L_name

- There are no non-trivial dependencies since each non-key attribute depends on the candidate key.
- Therefore, the functional dependencies satisfy BCNF. In summary, based on the given functional dependencies $\text{meal_id} \rightarrow \text{L_price}, \text{L_ID}, \text{L_name}$, the relation satisfies 1NF, 2NF, 3NF, and BCNF.

Dinner: To evaluate the functional dependencies $\text{meal_id} \rightarrow \text{D_price}, \text{D_ID}, \text{D_name}$, we need to analyze each dependency and assess the normal forms. Here's the analysis:

1. First Normal Form (1NF):
 - From the given functional dependencies, it appears that each attribute contains atomic values.
 - Therefore, the functional dependencies satisfy 1NF.
2. Second Normal Form (2NF):
 - Assuming meal_id is the candidate key, we have the following functional dependencies:
 - $\text{meal_id} \rightarrow \text{D_price}, \text{D_ID}, \text{D_name}$
 - There are no partial dependencies since each non-key attribute depends on the entire candidate key.
 - Therefore, the functional dependencies satisfy 2NF.
3. Third Normal Form (3NF):
 - Assuming meal_id is the candidate key, we have the following functional dependencies:
 - $\text{meal_id} \rightarrow \text{D_price}, \text{D_ID}, \text{D_name}$
 - There are no transitive dependencies since each non-key attribute depends directly on the candidate key.
 - Therefore, the functional dependencies satisfy 3NF.
4. Boyce-Codd Normal Form (BCNF):
 - Assuming meal_id is the candidate key, the functional dependencies are as follows:
 - $\text{meal_id} \rightarrow \text{D_price}, \text{D_ID}, \text{D_name}$
 - There are no non-trivial dependencies since each non-key attribute depends on the candidate key.
 - Therefore, the functional dependencies satisfy BCNF. In summary, based on the given functional dependencies $\text{meal_id} \rightarrow \text{D_price}, \text{D_ID}, \text{D_name}$, the relation satisfies 1NF, 2NF, 3NF, and BCNF.

Breakfast Status:

1. First Normal Form (1NF):
 - The given functional dependencies do not violate 1NF since each attribute appears to contain atomic values.
 - Therefore, the functional dependencies satisfy 1NF.
2. Second Normal Form (2NF):
 - Assuming B_ID is the candidate key, and $\text{D1}, \text{D2}, \text{D3}, \dots, \text{D31}$ are the attributes, it seems that there is a partial dependency.

- Each attribute D1, D2, D3, ..., D31 depends on B_ID partially since it depends on a specific day (e.g., D1 depends on B_ID for the first day, D2 depends on B_ID for the second day, and so on).
 - Therefore, the functional dependencies do not satisfy 2NF due to partial dependencies.
3. Third Normal Form (3NF):
- Assuming B_ID is the candidate key, and D1, D2, D3, ..., D31 are the attributes, there are no transitive dependencies since each attribute depends directly on the candidate key.
 - Therefore, the functional dependencies satisfy 3NF.
4. Boyce-Codd Normal Form (BCNF):
- Assuming B_ID is the candidate key, and D1, D2, D3, ..., D31 are the attributes, there are no non-trivial dependencies.
 - Each attribute D1, D2, D3, ..., D31 depends entirely on the candidate key B_ID.
 - Therefore, the functional dependencies satisfy BCNF. In summary, based on the given functional dependencies $B_ID \rightarrow D1, D2, D3, \dots, D31$, the relation satisfies 1NF, does not satisfy 2NF, satisfies 3NF, and satisfies BCNF.

Lunch Status :

1. First Normal Form (1NF):
- The given functional dependencies do not violate 1NF since each attribute appears to contain atomic values.
 - Therefore, the functional dependencies satisfy 1NF.
2. Second Normal Form (2NF):
- Assuming L_ID is the candidate key, and D1, D2, D3, ..., D31 are the attributes, it seems that there is a partial dependency.
 - Each attribute D1, D2, D3, ..., D31 depends on L_ID partially since it depends on a specific day (e.g., D1 depends on L_ID for the first day, D2 depends on L_ID for the second day, and so on).
 - Therefore, the functional dependencies do not satisfy 2NF due to partial dependencies.
3. Third Normal Form (3NF):
- Assuming L_ID is the candidate key, and D1, D2, D3, ..., D31 are the attributes, there are no transitive dependencies since each attribute depends directly on the candidate key.
 - Therefore, the functional dependencies satisfy 3NF.
4. Boyce-Codd Normal Form (BCNF):
- Assuming L_ID is the candidate key, and D1, D2, D3, ..., D31 are the attributes, there are no non-trivial dependencies.
 - Each attribute D1, D2, D3, ..., D31 depends entirely on the candidate key L_ID.
 - Therefore, the functional dependencies satisfy BCNF. In summary, based on the given functional dependencies $L_ID \rightarrow D1, D2, D3, \dots, D31$, the relation satisfies 1NF, does not satisfy 2NF, satisfies 3NF, and satisfies BCNF.

Dinner Status:

1. First Normal Form (1NF):
- The given functional dependencies do not violate 1NF since each attribute appears to contain atomic values.

- Therefore, the functional dependencies satisfy 1NF
2. Second Normal Form (2NF):
 - Assuming D_ID is the candidate key, and D1, D2, D3, ..., D31 are the attributes, it seems that there is a partial dependency.
 - Each attribute D1, D2, D3, ..., D31 depends on D_ID partially since it depends on a specific day (e.g., D1 depends on D_ID for the first day, D2 depends on D_ID for the second day, and so on).
 - Therefore, the functional dependencies do not satisfy 2NF due to partial dependencies.
 3. Third Normal Form (3NF):
 - Assuming D_ID is the candidate key, and D1, D2, D3, ..., D31 are the attributes, there are no transitive dependencies since each attribute depends directly on the candidate key.
 - Therefore, the functional dependencies satisfy 3NF.
 4. Boyce-Codd Normal Form (BCNF):
 - Assuming D_ID is the candidate key, and D1, D2, D3, ..., D31 are the attributes, there are no non-trivial dependencies.
 - Each attribute D1, D2, D3, ..., D31 depends entirely on the candidate key D_ID.
 - Therefore, the functional dependencies satisfy BCNF. In summary, based on the given functional dependencies $D_ID \rightarrow D1, D2, D3, \dots, D31$, the relation satisfies 1NF, does not satisfy 2NF, satisfies 3NF, and satisfies BCNF.

RELATION->TABLES

1)Student(registration_no,name,dob,cid,hid,sex)

registration_no	name	dob	cid	hid	sex
202301	Swagatika Hial	2002-02-16	1	1	F
202302	Harshita Yadav	2001-11-26	2	2	F
202303	Mamta Singh	2000-10-02	3	3	F
202304	Shailendra Mohan Sharma	2002-03-02	4	4	M
202305	Priyanshi Agnihotri	2008-05-04	5	5	F

2)Hostel(hid,hostel_name,room_no,mess_id)

hid	hostel_name	room_no	mess_id
1	Koyna	08	1
2	chandrabhaga	09	2
3	chandrabhaga	10	3
4	mahi	26	4
5	koyna	01	5

3)Mess(date,meal_id)

date	meal_id
1	7
2	5
3	4
4	2

5 4
6 3
7 1
8 2
9 2
10 4
11 3
12 5
13 1
14 3
15 5
16 7
17 4
18 6
19 5
20 3
21 5
22 4
23 7
24 6
25 1
26 3
27 4
28 5
29 4
30 5
31 5
+-----+-----+

4)Meal(mead_id,bid,lid,did)

+-----+-----+-----+
meal_id bid lid did
+-----+-----+-----+
1 1 2 1

	2		2		3		1	
	3		2		2		1	
	4		3		1		2	
	5		1		2		3	
	6		3		2		1	
	7		2		1		2	

5)Lunch(lid,l_name,price)

+-----+-----+-----+
lid l_name price
+-----+-----+-----+
1 curry chaawal 35.00
2 dahi baigan 30.00
3 soya bean 35.00)

6)Breakfast(bid,b_name,price)

+-----+-----+-----+
bid b_name price
+-----+-----+-----+
1 chana chura 25.00
2 sprout 25.00
3 chhole bhature 30.00
4 poori sabzi 30.00
5 parantha 20.00

7)Dinner(did,d_name,price)

+-----+-----+-----+
did d_name price
+-----+-----+-----+
1 paneer and rasogulla 50.00
2 aloo katahal 40.00
3 chana and custard 50.00
4 karela 40.00

8)breakfast_status(mess_id,d1,d2,d3.....d31);

9)lunch_status(mess_id,d1,d2,d3.....d31);

10)dinner_status(mess_id,d1,d2,d3.....d31);

SQL SCRIPT

Student table:

```
CREATE TABLE student (  
    registration_no INT NOT NULL AUTO_INCREMENT,  
    name VARCHAR(255) NOT NULL,  
    dob DATE NOT NULL,  
    cid INT NOT NULL,  
    hid INT NOT NULL,  
    sex CHAR(1) NOT NULL,  
    PRIMARY KEY (registration_no)  
);
```

Communication table:

```
CREATE TABLE communication (  
    cid INT NOT NULL,  
    address VARCHAR(255) NOT NULL,  
    phone_no VARCHAR(10) NOT NULL,  
    email_id VARCHAR(255) NOT NULL,  
    PRIMARY KEY (cid)  
);
```

Hostel table:

```
CREATE TABLE hostel (  
    hid INT NOT NULL AUTO_INCREMENT,  
    hostel_name VARCHAR(255) NOT NULL,  
    room_no VARCHAR(10),  
    mess_id INT,  
    PRIMARY KEY (hid)  
);
```

Mess table:

```
CREATE TABLE mess (  
    date INT NOT NULL AUTO_INCREMENT PRIMARY KEY MAXVALUE 31,  
    meal_id INT NOT NULL  
);
```

Breakfast status table:

```
CREATE TABLE breakfast_status (  
    mess_id INT NOT NULL,  
    d1 INT NOT NULL,  
    d2 INT NOT NULL,  
    d3 INT NOT NULL,  
    d4 INT NOT NULL,  
    d5 INT NOT NULL,  
    d6 INT NOT NULL,  
    d7 INT NOT NULL,  
    d8 INT NOT NULL,  
    d9 INT NOT NULL,  
    d10 INT NOT NULL,  
    d11 INT NOT NULL,  
    d12 INT NOT NULL,  
    d13 INT NOT NULL,  
    d14 INT NOT NULL,  
    d15 INT NOT NULL,  
    d16 INT NOT NULL,  
    d17 INT NOT NULL,  
    d18 INT NOT NULL,  
    d19 INT NOT NULL,  
    d20 INT NOT NULL,  
    d21 INT NOT NULL,  
    d22 INT NOT NULL,  
    d23 INT NOT NULL,  
    d24 INT NOT NULL,
```

```
d25 INT NOT NULL,  
d26 INT NOT NULL,  
d27 INT NOT NULL,  
d28 INT NOT NULL,  
d29 INT NOT NULL,  
d30 INT NOT NULL,  
d31 INT NOT NULL  
);
```

Lunch status table:

```
CREATE TABLE lunch_status (  
    mess_id INT NOT NULL,  
    d1 INT NOT NULL,  
    d2 INT NOT NULL,  
    d3 INT NOT NULL,  
    d4 INT NOT NULL,  
    d5 INT NOT NULL,  
    d6 INT NOT NULL,  
    d7 INT NOT NULL,  
    d8 INT NOT NULL,  
    d9 INT NOT NULL,  
    d10 INT NOT NULL,  
    d11 INT NOT NULL,  
    d12 INT NOT NULL,  
    d13 INT NOT NULL,  
    d14 INT NOT NULL,  
    d15 INT NOT NULL,  
    d16 INT NOT NULL,  
    d17 INT NOT NULL,  
    d18 INT NOT NULL,  
    d19 INT NOT NULL,  
    d20 INT NOT NULL,  
    d21 INT NOT NULL,  
    d22 INT NOT NULL,
```

```
d23 INT NOT NULL,  
d24 INT NOT NULL,  
d25 INT NOT NULL,  
d26 INT NOT NULL,  
d27 INT NOT NULL,  
d28 INT NOT NULL,  
d29 INT NOT NULL,  
d30 INT NOT NULL,  
d31 INT NOT NULL  
);
```

Dinner status table:

```
CREATE TABLE dinner_status (  
  mess_id INT NOT NULL,  
  d1 INT NOT NULL,  
  d2 INT NOT NULL,  
  d3 INT NOT NULL,  
  d4 INT NOT NULL,  
  d5 INT NOT NULL,  
  d6 INT NOT NULL,  
  d7 INT NOT NULL,  
  d8 INT NOT NULL,  
  d9 INT NOT NULL,  
  d10 INT NOT NULL,  
  d11 INT NOT NULL,  
  d12 INT NOT NULL,  
  d13 INT NOT NULL,  
  d14 INT NOT NULL,  
  d15 INT NOT NULL,  
  d16 INT NOT NULL,  
  d17 INT NOT NULL,  
  d18 INT NOT NULL,  
  d19 INT NOT NULL,  
  d20 INT NOT NULL,
```

```
    d21 INT NOT NULL,  
    d22 INT NOT NULL,  
    d23 INT NOT NULL,  
    d24 INT NOT NULL,  
    d25 INT NOT NULL,  
    d26 INT NOT NULL,  
    d27 INT NOT NULL,  
    d28 INT NOT NULL,  
    d29 INT NOT NULL,  
    d30 INT NOT NULL,  
    d31 INT NOT NULL  
);
```

Meal Table:

```
CREATE TABLE meal (  
    meal_id INT NOT NULL AUTO_INCREMENT,  
    bid INT NOT NULL,  
    lid INT NOT NULL,  
    did INT NOT NULL,  
    PRIMARY KEY (meal_id)  
);
```

Breakfast Table:

```
CREATE TABLE breakfast (  
    bid INT NOT NULL AUTO_INCREMENT,  
    b_name VARCHAR(255) NOT NULL,  
    price DECIMAL(10,2) NOT NULL,  
    PRIMARY KEY (bid)  
);
```


Lunch Table:

```
CREATE TABLE lunch (  
    lid INT NOT NULL AUTO_INCREMENT,  
    l_name VARCHAR(255) NOT NULL,  
    price DECIMAL(10,2) NOT NULL,  
    PRIMARY KEY (lid)  
);
```

Dinner Table:

```
CREATE TABLE dinner (  
    did INT NOT NULL AUTO_INCREMENT,  
    d_name VARCHAR(255) NOT NULL,  
    price DECIMAL(10,2) NOT NULL,  
    PRIMARY KEY (did)  
);
```

SYSTEM CONFIGURATIONS

Hardware Interface:-

Processor: Intel i5 100Ghz or more

RAM: 4GB or more

Software Interface:-

Operating System: Windows 11

Software : Oracle

CONCLUSION

In this project we have created one database which is easy to access and user friendly.

The database keeps a backup of the mess management data which includes their details. A mess management is a professional app which assists with the mess, meals, accurate bills etc.