

1 C Memory Management

1. Match the items on the left with the memory segment in which they are stored. Answers may be used more than once, and more than one answer may be required.

1. Static variables **B**
2. Local variables **D**
3. Global variables **B**
4. Constants **A, B, D**
5. Machine Instructions **A**
6. malloc() **C**
7. String Literals **B**

- A. Code
- B. Static
- C. Heap
- D. Stack

4 Explained: With DEFINE, you can replace something in the code. With const declaration, it's stored in the static or stack depending on being declared in a function or not.

2. Write the code necessary to properly allocate memory (on the heap) in the following scenarios

1. An array **arr** of k integers

```
arr = (int *) malloc(sizeof(int) * k);
```

2. A string **str** of length p

```
str = (char *) malloc(sizeof(char) * (p + 1)); // Don't forget the null terminator!
```

3. An $n \times m$ matrix **mat** of integers initialized to zeros

```
mat = (int **) calloc(n, sizeof(int *));
for (int i = 0; i < n; i++) {
    mat[i] = (int *) calloc(m, sizeof(int));
}
```

3. What is wrong with the C code below?

```
int* pi = malloc(314 * sizeof(int));
if(!raspberrypi) pi = malloc(1 * sizeof(int)); // Memory leak if not raspberrypi
return pi;
```

4. Write code to prepend (add to the start) to a linked list, and to free/empty the entire list.

```
struct ll_node { struct ll_node* next; int value; }
```

void prepend(struct ll_node** lst, int val)	void free_ll(struct ll_node** lst)
<pre>struct ll_node* item = (struct ll_node*) malloc(sizeof(struct ll_node)); item->value = val; item->next = *lst; *lst = item;</pre>	<pre>if(*lst) { free_ll(&((*lst)->next)); free(*lst); } *lst = NULL;</pre>

*Note: *lst points to the first element of the list, or is NULL if the list is empty.*

2 Data Structures in C

In this question, we will implement a array-based stack of integers in C. The stack will be represented by the struct below.

```
struct stack {
    int size;          // Number of element in stackArray
    int topIndex;      // Index of the array that is the top of the stack
    int *stackArray;   // Array holding the elements of the stack
}
```

Implement the functions below.

```
// Create a new stack with the given array size
struct stack *init_stack(int size) {
    struct stack *stk = (struct stack *) malloc(sizeof(struct stack *));
    if (!stk) {
        return NULL;
    }
    stk->size = size;
    stk->topIndex = -1;
    stk->stackArray = (int *) malloc(sizeof(int) * size);
    if (!(stk->stackArray)) {
        return NULL;
    }
    return stk;
}

// Add the given element to the stack. Resize by doubling the array size if full.
void push(int x, struct stack* stk) {
    if (stk->topIndex == stk->size - 1) {
        int *stackArray = (int *) realloc(stk->stackArray, stk->size * 2)
        if (!stackArray) {
            return;
        }
        stk->stackArray = stackArray;
        stk->size = stk->size * 2;
    }
    stk->topIndex += 1;
    stk->stackArray[stk->topIndex] = x;
}

// Remove the top element from the stack. Return 0 if empty.
int pop(struct stack *stk) {
    if (stk->topIndex < 0) {
        return 0;
    }
    int x = stk->stackArray[stk->topIndex];
    stk->stackArray[stk->topIndex] = 0;
    stk->topIndex -= 1;
    return x;
}
```

3 RISC-V Intro

1. Assume we have an array in memory that contains `int* arr = {1,2,3,4,5,6,0}`. Let the value of `arr` be a multiple of 4 and stored in register `s0`. What do the snippets of RISC-V code do?

a) `lw t0, 12(s0)`

Sets register `t0` equal to `arr[3]`

b) `slli t1, t0, 2` // Mistakenly shifted by 4 in the original worksheet

`add t2, s0, t1`

`lw t3, 0(t2)`

`addi t3, t3, 1`

`sw t3, 0(t2)`

Increments the array element specified by `t2` (i.e. `arr[t2]`) by 1

c) `lw t0, 0(s0)`

`xori t0, t0, 0xFFF`

`addi t0, t0, 1`

Sets the register `t0` to the two's complement negation of `arr[0]`

2. What are the instructions to branch to `label` on each of the following conditions? The only branch instructions you may use are `beq` and `bne`.

<code>s0 < s1</code>	<code>s0 <= s1</code>	<code>s0 > 1</code>
<code>slt t0, s0, s1</code>	<code>slt t0, s1, s0</code>	<code>sltiu t0, s0, 2</code>
<code>bne t0, 0, label</code>	<code>beq t0, 0, label</code>	<code>beq t0, 0, label</code>

4 Translating between C and RISC-V

Translate between the C and RISC-V code. You may want to use the RISC-V Reference Card for more information on the instruction set and syntax. In all of the C examples, we show you how the different variables map to registers – you don't have to worry about the stack or any memory-related issues. You may assume all registers are initialized to zero.

C	RISC-V
<pre>// s0 -> a, s1 -> b // s2 -> c, s3 -> z int a = 4, b = 5, c = 6, z; z = a + b + c + 10;</pre>	<pre>addi s0, x0, 4 addi s1, x0, 5 addi s2, x0, 6 add s3, s0, s1 add s3, s3, s2 addi s3, s3, 10</pre>
<pre>// s0 -> int * p = intArr; // s1 -> a; *p = 0; int a = 2; p[1] = p[a] = a;</pre>	<pre>sw x0, 0(s0) addi s1, x0, 2 sw s1, 4(s0) slli t0, s1, 2 add t0, t0, s0 sw s1, 0(t0)</pre>
<pre>// s0 -> a, s1 -> b int a = 5, b = 10; if(a + a == b) { a = 0; } else { b = a - 1; }</pre>	<pre>addi s0, x0, 5 addi s1, x0, 10 add t0, s0, s0 bne t0, s1, else xor s0, x0, x0 jal x0, exit else: addi s1, s0, -1 exit:</pre>
<pre>// computes s1 = 2^30 s1 = 1; for(s0=0;s0<30;s++) { s1 *= 2; }</pre>	<pre>addi s0, x0, 0 addi s1, x0, 1 addi t0, x0, 30 loop: beq s0, t0, exit add s1, s1, s1 addi s0, s0, 1 jal x0, loop exit:</pre>
<pre>// s0 -> n, s1 -> sum // assume n > 0 to start int sum; for(sum=0;n>0;sum+=n--);</pre>	<pre>addi s1, s1, 0 loop: beq s0, x0, exit add s1, s1, s0 add s0, s0, -1 jal x0, loop exit:</pre>