

David J. Elkind - LIPO

I'm interested in estimating the Lipschitz constant k for functions which approximate as Lipschitz. When the functions are noisy, they can't be Lipschitz, but adding a noise term s can make the functions behave more nicely; including the large penalty 10^6 for s forces it to often be 0.

The function of interest is f and it is only observed via some finite number of points $(x_i, f(x_i))$ for $i = 1, 2, \dots, t$. Define

$$U(y) = \min_{i \in \{1, 2, \dots, t\}} \left[f(x_i) + \sqrt{s_i + (y - x_i)^\top K (y - x_i)} \right] \quad (1)$$

One way to estimate k is to optimize the following program (using a different k_i for each coordinate):

$$\min_{K, s} \quad \|K\|_F^2 + 10^6 \sum_{i=1}^t s_i^2 \quad (2)$$

$$\text{s.t. } U(x_i) \geq f(x_i) \quad i \in \{1, 2, \dots, t\} \quad (3)$$

$$s_i \geq 0 \quad \forall i \in \{1, 2, \dots, t\} \quad (4)$$

$$K_{i,j} \geq 0 \quad \forall i, j \in \{1, 2, \dots, d\}^2 \quad (5)$$

$$K \text{ is a diagonal matrix.} \quad (6)$$

This clearly can be rewritten as

$$\min_{k, s} \quad k^\top k + 10^6 \cdot s^\top s \quad (7)$$

$$\text{s.t. } U(x_i) \geq f(x_i) \quad i \in \{1, 2, \dots, t\} \quad (8)$$

$$s_i \geq 0 \quad \forall i \in \{1, 2, \dots, t\} \quad (9)$$

$$k_i \geq 0 \quad \forall i \in \{1, 2, \dots, d\}^2 \quad (10)$$

But it's not at all clear how to me actually *do* the optimization, since the minimization objective is not obviously a QP. I guess one way to do it would be to roll the two summands together and write

$$\min_{k, s} \quad \begin{bmatrix} k \\ s \end{bmatrix}^\top A \begin{bmatrix} k \\ s \end{bmatrix} \quad (11)$$

$$\text{s.t. } U(x_i) \geq f(x_i) \quad i \in \{1, 2, \dots, t\} \quad (12)$$

$$s_i \geq 0 \quad \forall i \in \{1, 2, \dots, t\} \quad (13)$$

$$k_i \geq 0 \quad \forall i \in \{1, 2, \dots, d\}^2 \quad (14)$$

Just for convenience, we can rewrite U in terms of k :

$$U(y) = \min_{i \in \{1, 2, \dots, t\}} \left[f(x_i) + \sqrt{s_i + \|k^\top (y - x_i)\|_2^2} \right] \quad (15)$$

and observe that all I've done is move stuff around so that A is given by

$$A_{ij} = \begin{cases} 1 & \forall i = j \wedge 1 \leq i \leq d \\ 10^6 & \forall i = j \wedge d < i \leq t + d, \\ 0 & \text{otherwise} \end{cases} \quad (16)$$

i.e. A is a diagonal matrix with 1s corresponding to the elements of k and 10^6 corresponding to the elements of s .

This is “nicer” in a sense, since it’s *obviously* a QP. But what’s less-nice is that U , and therefore its corresponding constraint, is not a linear function, owing to the fact that s and k both appear inside of the square root.

So given that this program is not a QP with linear constraints, what would the appropriate tool be to solve this problem? More specifically, what python library would you recommend using?