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Algorithm:

We sort the customers by ascending order of their service time, then serve the customer with smallest service time each time.

Proof of Correctness:

The total time can be expressed as

$$T = \sum_{i=1}^n \sum_{j=1}^i t_j$$

which is simply

$$T = \sum_{i=1}^n (n - i + 1)t_i$$

So if the sequence $\{t_1, t_2, t_3, \dots, t_n\}$ is in ascending order, i.e. $t_1 \leq t_2 \leq t_3 \leq \dots \leq t_n$, simply by the **rearrangement inequality** we know such a T is the minimum waiting time because it is the so-called "inverse product sum".

Below is the proof of rearrangement inequality for completeness of this problem.

Denote the two sequence that have already sorted by ascending order by $\{a_i\}$ and $\{b_i\}$. Let $p_i = \sum_{j=1}^i b_j$, $q_i = \sum_{j=1}^i b_{k_j}$. It's easy to see $p_i \leq q_i$ for $i \leq n - 1$ and $p_n = q_n$.

Then

$$\sum_{i=1}^n a_{n-i+1}b_i = \sum_{i=1}^{n-1} a_i(p_{n-i+1} - p_{n-i}) + a_n p_1 \quad (1)$$

$$= \sum_{i=1}^{n-1} (a_{n-i+1} - a_{n-i})p_i + a_1 p_n \quad (2)$$

$$\leq \sum_{i=1}^{n-1} (a_{n-i+1} - a_{n-i})q_i + a_1 q_n \quad (3)$$

$$= \sum_{i=1}^n a_{n-i+1}b_{k_i} \quad (4)$$

Another side of the inequality is not used in our problem, so the proof is omitted here.