

1 Optimal Coupling

1.1 Basic Ideas

To reach the lower bound $D_{TV}(\mu, \nu)$, the intuition is that we maximize the terms where $X = Y$. Denote our optimal coupling by ω^* , then $\forall (x, y) \in \Omega^2$ such that $x = y$, the maximum of $\omega^*(x, y)$ can only be $\min\{\mu(x), \nu(y)\}$. As we want to maximize these cases, we can directly set it as $\min\{\mu(x), \nu(y)\}$.

Then we need a formula to define $\omega^*(x, y)$ where $x \neq y$. Here the intuition is that $\omega^*(x, y)$ should be some constant times $\max\{\mu(x) - \nu(x), 0\} \cdot \max\{\nu(y) - \mu(y), 0\}$. So up to here, I construct a coupling as follows:

$$\omega^*(x, y) = \begin{cases} \min\{\mu(x), \nu(y)\}, & x = y \\ C \cdot \max\{\mu(x) - \nu(x), 0\} \cdot \max\{\nu(y) - \mu(y), 0\}, & o.w. \end{cases}$$

Now I will determine C so that ω^* is indeed a valid coupling, and then show $Pr_{(X,Y) \sim \omega^*}(X \neq Y)$ is actually $D_{TV}(\mu, \nu)$.

1.2 Determine C and proof of valid coupling

Define $A = \{x \in \Omega | \mu(x) \geq \nu(x)\}$. It follows that $\bar{A} = \{x \in \Omega | \mu(x) < \nu(x)\}$. We now calculate the marginal distribution of X under ω^* .

$$\forall x \in A, \quad \sum_{y \in \Omega} \omega^*(x, y) = \sum_{y \in \Omega \wedge y=x} \omega^*(x, y) + \sum_{y \in \Omega \wedge y \neq x} \omega^*(x, y) \quad (1)$$

$$= \min\{\mu(x), \nu(x)\} + \sum_{y \in \Omega \wedge y \neq x} \omega^*(x, y) \quad (2)$$

$$= \nu(x) + C \max\{\mu(x) - \nu(x), 0\} \sum_{y \in \Omega \wedge y \neq x} \max\{\nu(y) - \mu(y), 0\} \quad (3)$$

$$= \nu(x) + C(\mu(x) - \nu(x)) \sum_{y \in \bar{A}} (\nu(y) - \mu(y)) \quad (4)$$

Note that $\sum_{y \in \bar{A}} (\nu(y) - \mu(y)) = D_{TV}(\mu, \nu)$ by our definition of \bar{A} , so by setting $C = 1/D_{TV}(\mu, \nu)$ we can get

$$\forall x \in A, \quad \sum_{y \in \Omega} \omega^*(x, y) = \nu(x) + \mu(x) - \nu(x) \quad (5)$$

$$= \mu(x) \quad (6)$$

We use the same C and consider the cases where $x \in \bar{A}$. Similarly

$$\forall x \in \bar{A}, \quad \sum_{y \in \Omega} \omega^*(x, y) = \min\{\mu(x), \nu(x)\} + \sum_{y \in \Omega \wedge y \neq x} \omega^*(x, y) \quad (7)$$

$$= \mu(x) + D_{TV}(\mu, \nu) \max\{\mu(x) - \nu(x), 0\} \sum_y \max\{\nu(y) - \mu(y), 0\} \quad (8)$$

$$= \mu(x) + D_{TV}(\mu, \nu) \cdot 0 \cdot \sum_y \max\{\nu(y) - \mu(y), 0\} \quad (9)$$

$$= \mu(x) \quad (10)$$

Now we have proved such a coupling satisfies that the marginal distribution of X is indeed $\mu(x)$. The same argument works for Y as well. So by setting $C = 1/D_{TV}(\mu, \nu)$ we actually constructed

$$\omega^*(x, y) = \begin{cases} \min\{\mu(x), \nu(y)\}, & x = y \\ \frac{1}{D_{TV}(\mu, \nu)} \cdot \max\{\mu(x) - \nu(x), 0\} \cdot \max\{\nu(y) - \mu(y), 0\}, & o.w. \end{cases}$$

Now we only need to show $Pr_{(X,Y) \sim \omega^*}(X \neq Y)$ is actually $D_{TV}(\mu, \nu)$.

1.3 Proof of optimal coupling

We show by calculating $Pr_{(X,Y) \sim \omega^*}(X = Y)$.

$$Pr_{(X,Y) \sim \omega^*}(X = Y) = \sum_{x \in \Omega} \min\{\mu(x), \nu(x)\} \quad (11)$$

$$= \sum_{x \in A} \nu(x) + \sum_{x \in \bar{A}} \mu(x) \quad (12)$$

$$= \nu(A) + \mu(\bar{A}) \quad (13)$$

$$= \nu(A) + (1 - \mu(A)) \quad (14)$$

$$= 1 - (\mu(A) - \nu(A)) \quad (15)$$

$$= 1 - D_{TV}(\mu, \nu) \quad (16)$$

So $Pr_{(X,Y) \sim \omega^*}(X \neq Y) = D_{TV}(\mu, \nu)$

2 Stochastic Dominance

The idea of proof comes from our teacher's lecture notes in previous year.

Following the idea of that lecture note, I will first prove the proposition about monotone coupling, because the other two questions are actually applications of this proposition.

2.1 Monotone Coupling

Sufficiency:

If a monotone coupling of μ and ν exists, denoted by ω . Then

$$\forall a \in \Omega, Pr_{Y \sim \nu}(Y \geq a) = Pr_{(X,Y) \sim \omega}(Y \geq a) \quad (17)$$

$$= Pr_{(X,Y) \sim \omega}(X \geq Y \wedge Y \geq a) + Pr_{(X,Y) \sim \omega}(X < Y \wedge Y \geq a) \quad (18)$$

$$= Pr_{(X,Y) \sim \omega}(X \geq Y \wedge Y \geq a) \quad (19)$$

$$\leq Pr_{(X,Y) \sim \omega}(X \geq a) \quad (20)$$

$$= Pr_{X \sim \mu}(X \geq a) \quad (21)$$

Hence $Pr_{X \sim \mu}(X \geq a) \geq Pr_{Y \sim \nu}(Y \geq a), \forall a \in \Omega$.

Necessity:

2.2 Binomial Distribution

Sufficiency:

If $p \geq q$, suppose $X \sim \text{Binom}(n, p)$, $Y \sim \text{Binom}(n, q)$. We define such a coupling ω of these two Binomial Distributions where we do n trials and for these n trials we independently pick a real r in $[0, 1]$ uniformly at random and every trial is independent. Then let $X = x$ where x is the number of these trials with $r \leq p$ and let $Y = y$ where y is the number of these trials with $r \leq q$.

By our definition we can see clearly $Pr_{(X,Y) \sim \omega}(X \geq Y) = 1$. So there exists a monotone coupling of $\text{Binom}(n, p)$ and $\text{Binom}(n, q)$. By the proposition we have proven above, $\text{Binom}(n, p) \succeq \text{Binom}(n, q)$.

Necessity:

Prove by contradiction. If $p < q$, consider the case where $a = n$. $Pr(X \geq a) = Pr(X = n) = p^n < q^n = Pr(Y = n) = Pr(Y \geq a)$. This violates $\text{Binom}(n, p) \succeq \text{Binom}(n, q)$. So $p \geq q$.

2.3 Random Graph

Suppose $G \sim \mathcal{G}(n, p)$, $H \sim \mathcal{G}(n, q)$. Consider such a coupling ω of $\mathcal{G}(n, p)$ and $\mathcal{G}(n, q)$ where we generate G and H simultaneously. For each pair of vertices $\{i, j\}$ we independently pick a real r in $[0, 1]$ uniformly at random. Let G have edge $\{i, j\}$ iff $r \leq p$ and let H have edge $\{i, j\}$ iff $r \leq q$.

For any $p, q \in [0, 1]$ satisfying $p \geq q$, H is a subgraph of G , so ω is a monotone coupling. So by the proposition we have proven $Pr_{G \sim \mathcal{G}(n, p)}(G \text{ is connected}) \geq Pr_{H \sim \mathcal{G}(n, q)}(H \text{ is connected})$

3 Total Variation Distance is Non-Increasing

Let $X_0 \sim \mu_0$ and $Y_0 \sim \pi$. For any $t \geq 0$, we can couple the distributions of random variables X_t and Y_t such that $Pr(X_t \neq Y_t) = \Delta(t)$. This coupling is feasible because we have proven, in problem 1, that an optimal coupling exists.

Then we can construct a coupling of the distributions of X_{t+1} and Y_{t+1} with this coupling. We define

$$\begin{cases} X_{t+1} = Y_{t+1}, & \text{if } X_t = Y_t \\ X_{t+1} \sim \mu_{t+1}, Y_{t+1} \sim \pi & \text{if } X_t \neq Y_t \end{cases}$$

Then by the coupling lemma again we have

$$\Delta(t+1) \leq Pr(X_{t+1} \neq Y_{t+1}) \quad (22)$$

By our construction of coupling at t and $t+1$ we have

$$Pr(X_{t+1} \neq Y_{t+1}) \quad (23)$$

$$= Pr(X_{t+1} \neq Y_{t+1} | X_t = Y_t) Pr(X_t = Y_t) + Pr(X_{t+1} \neq Y_{t+1} | X_t \neq Y_t) Pr(X_t \neq Y_t) \quad (24)$$

$$\leq Pr(X_t \neq Y_t) \quad (25)$$

That inequality is true because all possibilities must be in $[0, 1]$.

So it follows That

$$\Delta(t+1) \leq Pr(X_{t+1} \neq Y_{t+1}) \leq Pr(X_t \neq Y_t) \leq \Delta(t), \forall t \geq 0$$

And that is true for any $t \geq 0$ by induction. The induction process is actually how the Markov Chain X_t and Y_t forms. Since any step of this induction process is the same as our construction of coupling and proof above, we do not write it again.

4 Acknowledgements

4.1

For Problem 1 the inspiration of constructing the optimal coupling is from the lecture notes of a MIT course 6.896 Probability and Computation. Actually my construction is not entirely like its demonstration, but the core ideas are the same.

4.2

For Problem 2 I actually learnt our teacher's lecture notes 5 from last year and then solved these questions. I think it will be extremely difficult to solve these questions without knowing the theorem about monotone coupling.

4.3

For Problem 3 I referred to the lecture note 7 of a UC Berkeley course CS294 Partition Functions(do not know why such a course contains materials about our course). I cannot justify that my works are different from that lecture note because I cannot figure out another easy approach to solve this problem.