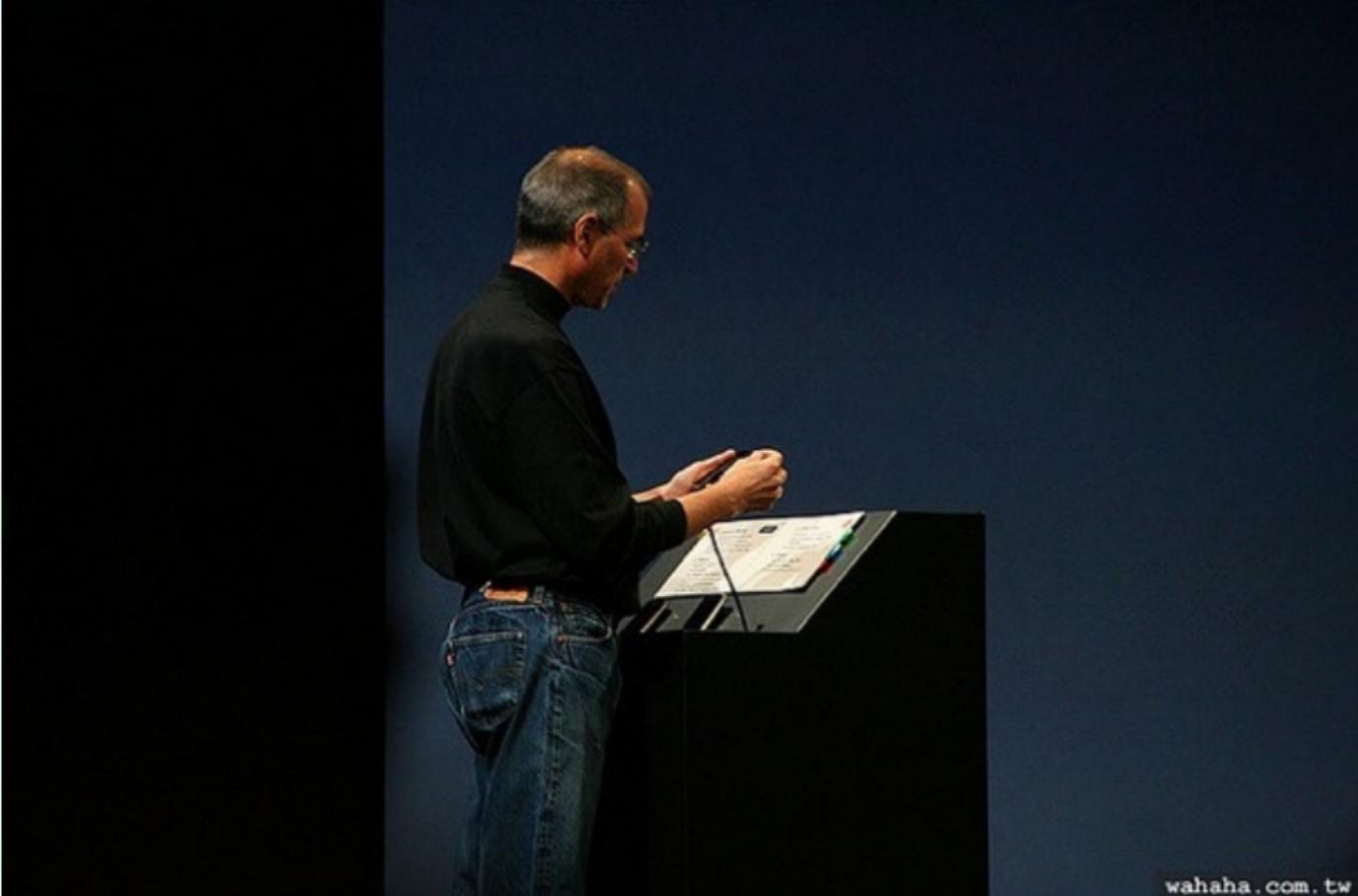


Channel Manipulation as a Coding Technique

Hsin-Po Wang (EECS, UC Berkeley)

What do the
following pictures
have in common?

(Images from the internet)



(Images from the internet)



(Images from the internet)



(Images from the internet)





<https://jobtalk.symbol.codes>



<https://jobtagrambol.codes>

Coding

= adding redundancies
in a smart way

Coding
= adding ~~redundancies~~
in a smart way



(Images from Wikipedia, NASA, IBM, Qualcomm)

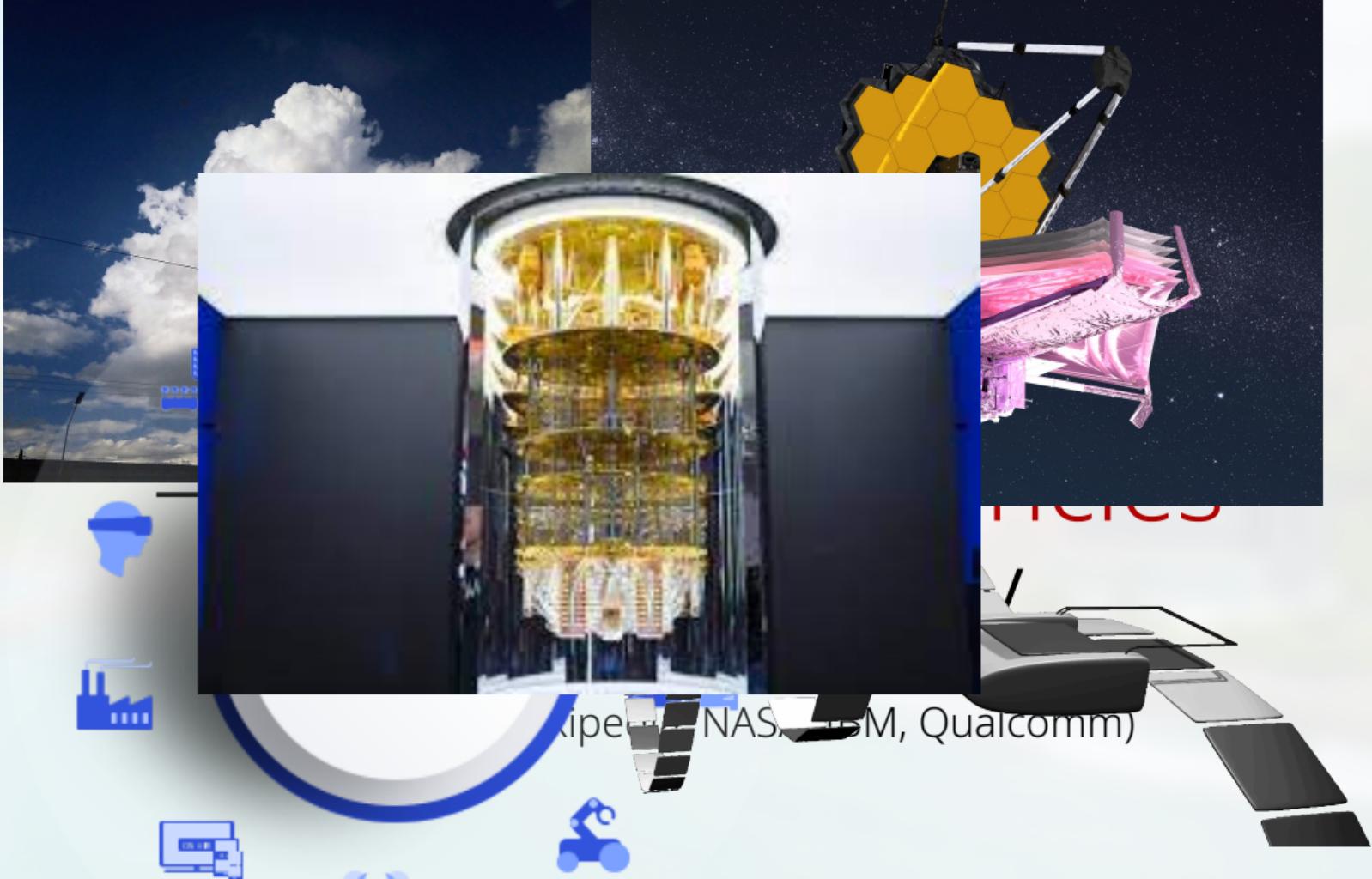


- during ~~redundancies~~
in a smart way

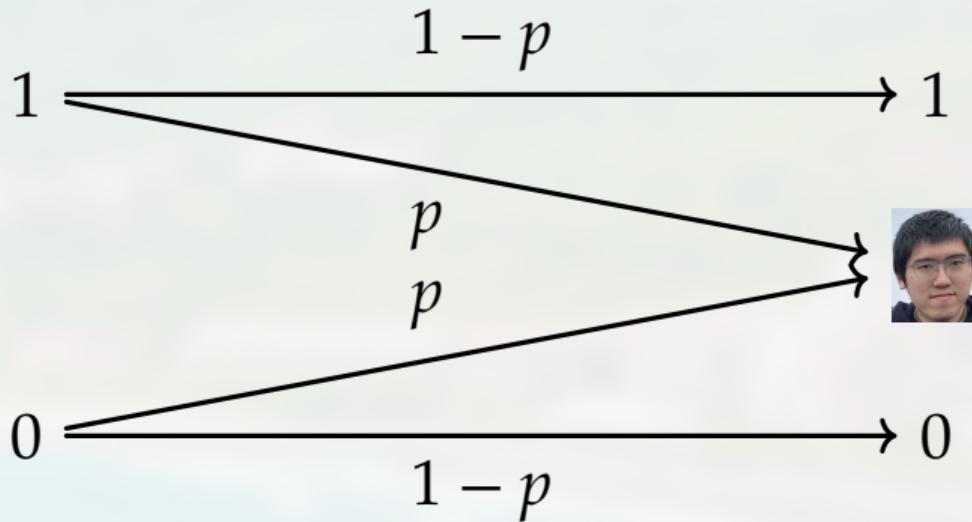
(Images from Wikipedia, NASA, IBM, Qualcomm)







Math Framework of Coding



Binary erasure channel with erasure probability p : BEC(p)

N = 78

S = 83

Y = 89

U = 85

The “NSYSU” polynomial:

$$f(x) = 78 + 83x + 89x^2 + 83x^3 + 85x^4$$

$$N = 78$$

$$S = 83$$

$$Y = 89$$

$$U = 85$$

The “NSYSU” polynomial:

$$f(x) = 78 + 83x + 89x^2 + 83x^3 + 85x^4$$

$$f(-3) = 274 \quad f(-2) = 964 \quad f(-1) = 86 \quad f(0) = 78$$

$$f(1) = 418 \quad f(2) = 624 \quad f(3) = 254 \quad f(4) = 906$$

$N = 78$

$S = 83$

$Y = 89$

$U = 85$

The “NSYSU” polynomial:

$$f(x) = \sum_{i=0}^4 a_i x^i$$


$$f(-3) = 274 \quad f(-2) = 964 \quad f(-1) = 86 \quad f(0) = ?$$



$$f(1) = 418 \quad f(2) = 624 \quad f(4) = 254 \quad f(4) = 906$$

$N = 78$

$S = 83$

$Y = 89$

$U = 85$

The “NSYSU” polynomial:

$$f(x) = \sum_{i=0}^4 a_i x^i$$


$$f(-3) = 274 \quad f(-2) = 964 \quad f(-1) = 86 \quad f(0) = 1$$



$$f(1) = 4 \quad f(2) = 624 \quad f(4) = 254 \quad f(4) = 906$$

$N = 78$

$S = 83$

$Y = 89$

$U = 85$

The “NSYSU” polynomial:

$$f(x) = \sum_{i=0}^4 a_i x^i$$


$$f(-3) = 274$$

$$f(-2) = 964$$

$$f(-1) = 8$$

$$f(0) = 1$$



$$f(1) = 4$$

$$f(2) = 624$$

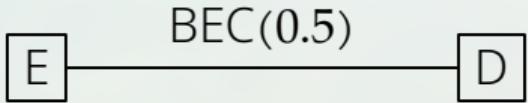
$$f(3) = 254$$

$$f(4) = 906$$

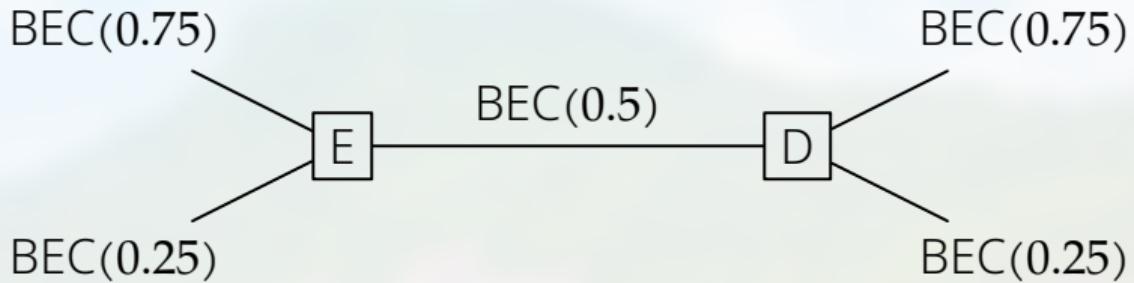
New Idea



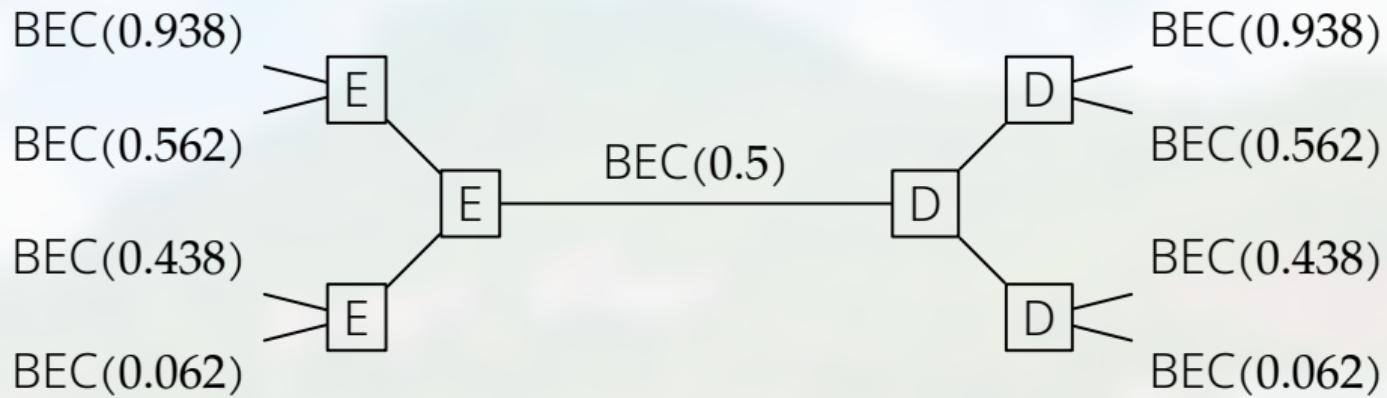
Polar Codes



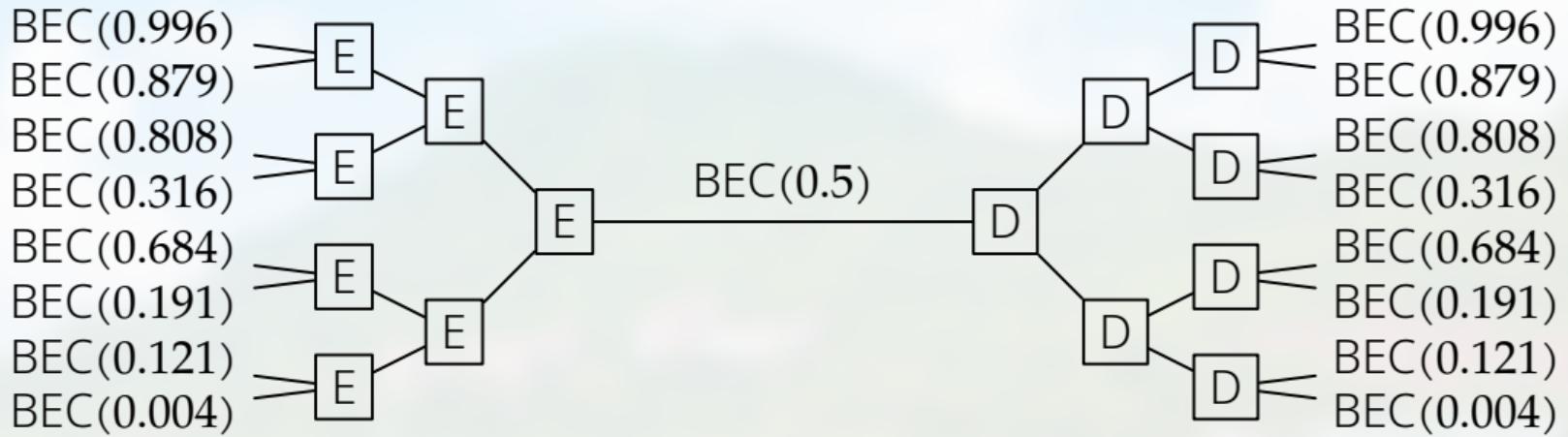
Suppose there are magic devices \boxed{E} and \boxed{D} that turns $\text{BEC}(x)$ into $\text{BEC}(x^2)$ and $\text{BEC}(2x - x^2)$.



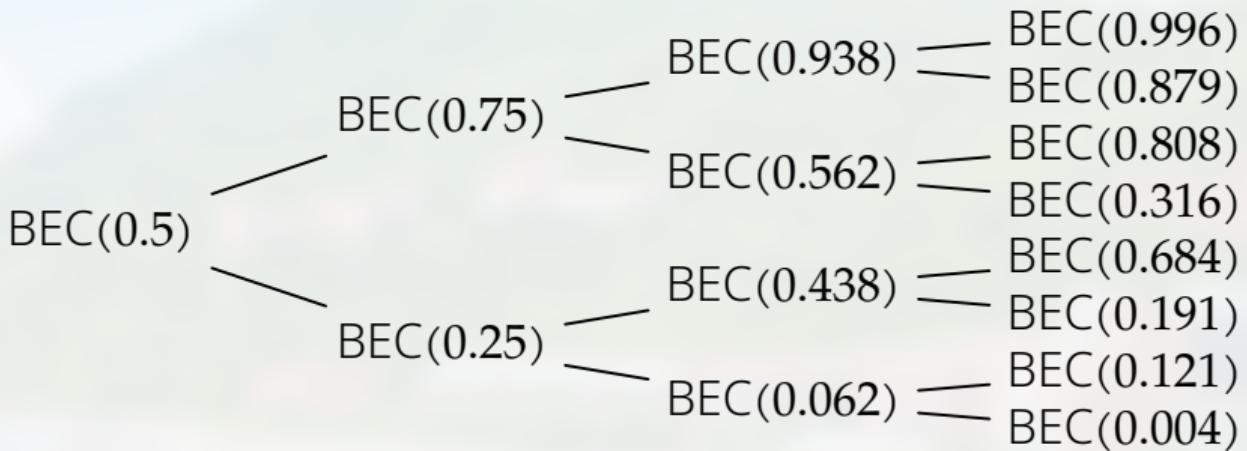
Suppose there are magic devices \boxed{E} and \boxed{D} that turns $\text{BEC}(x)$ into $\text{BEC}(x^2)$ and $\text{BEC}(2x - x^2)$.



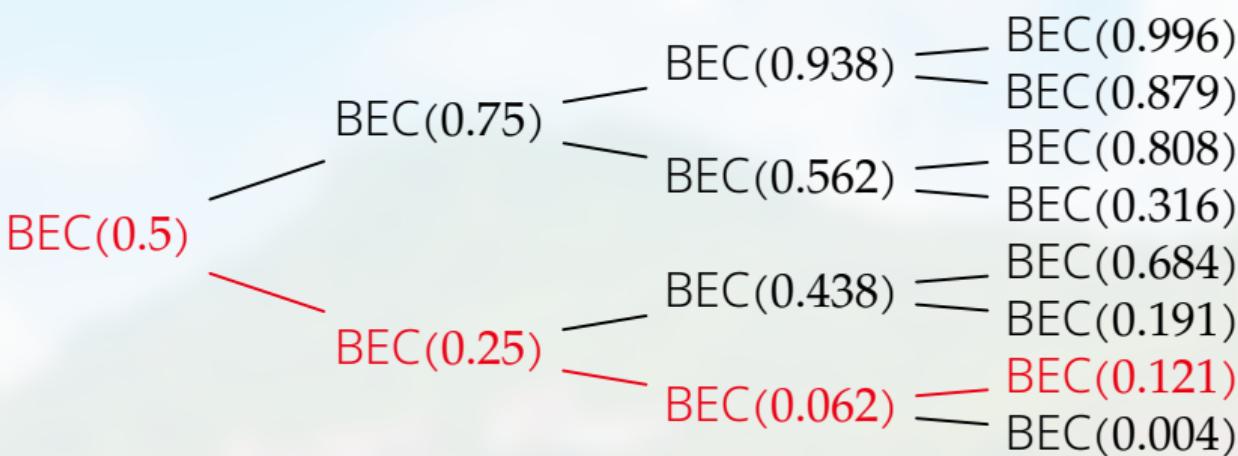
What if we apply more magic devices?



What if we apply more magic devices?
And more and more and more???

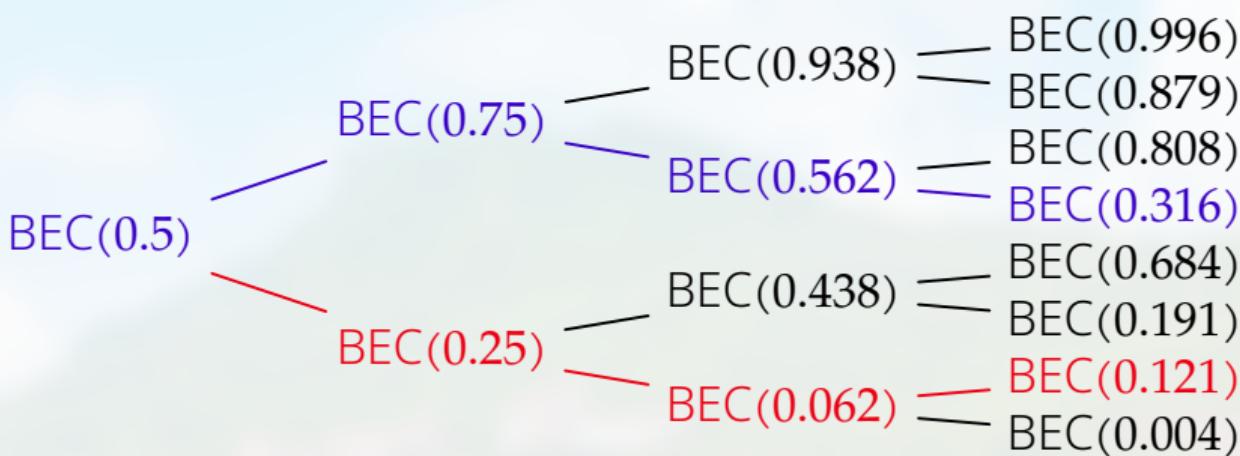


This is a tree



: This is a martingale

$$M_{n+1} := \begin{cases} 2M_n - M_n^2 & \text{with prob. } 1/2, \\ M_n^2 & \text{with prob. } 1/2. \end{cases}$$



: This is a martingale

$$M_{n+1} := \begin{cases} 2M_n - M_n^2 & \text{with prob. } 1/2, \\ M_n^2 & \text{with prob. } 1/2. \end{cases}$$



: Bounded martingale converges.

$$M_{n+1} := \begin{cases} 2M_n - M_n^2 & \text{with prob. } 1/2, \\ M_n^2 & \text{with prob. } 1/2. \end{cases}$$

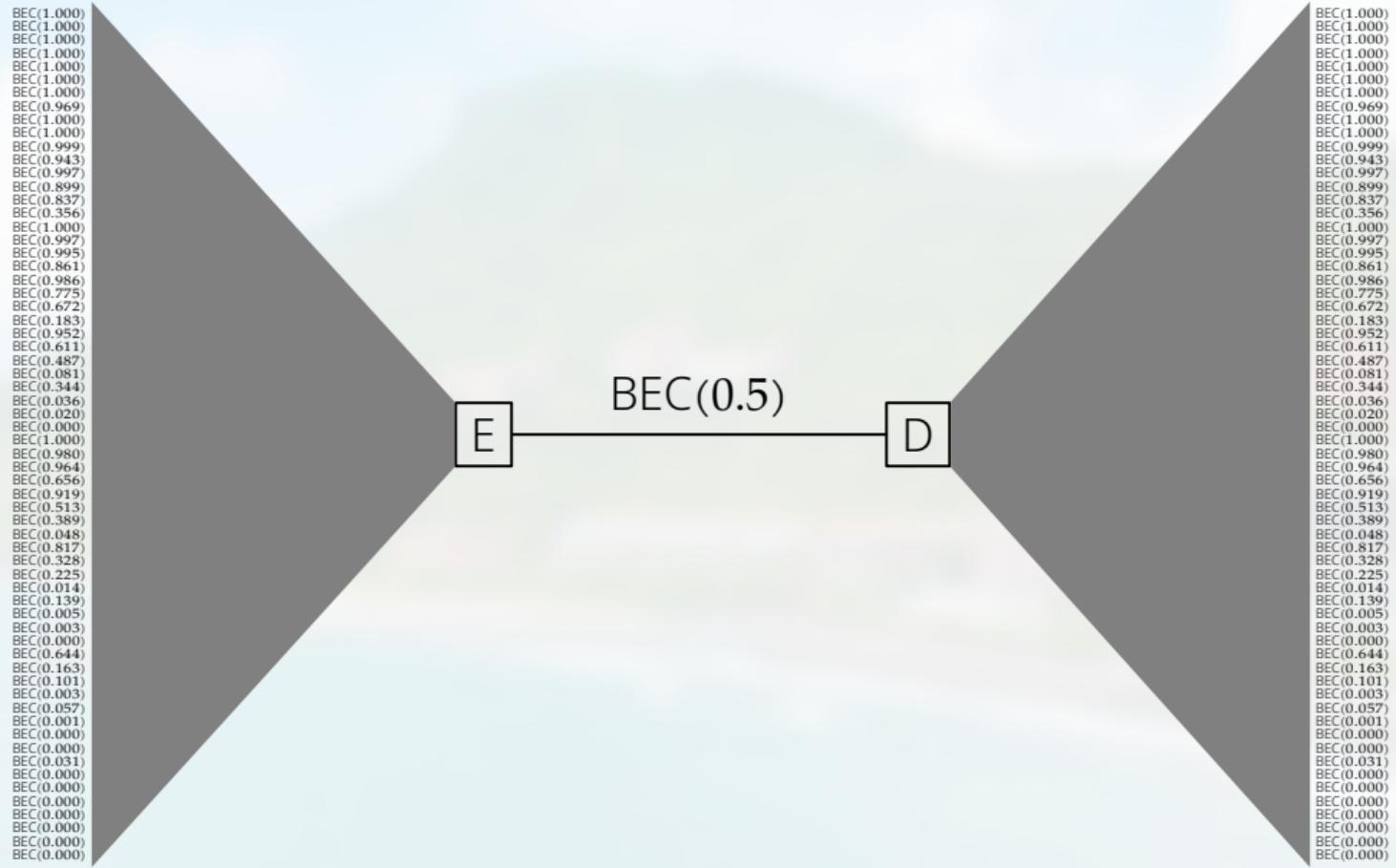


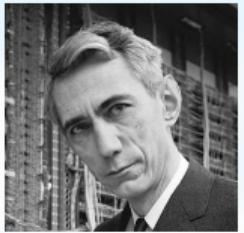
: Bounded martingale converges.



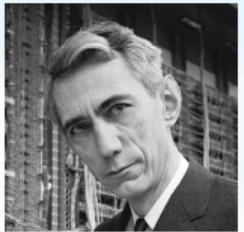
: BECs converges to 0 or 1.

$$M_{n+1} := \begin{cases} 2M_n - M_n^2 & \text{with prob. } 1/2, \\ M_n^2 & \text{with prob. } 1/2. \end{cases}$$





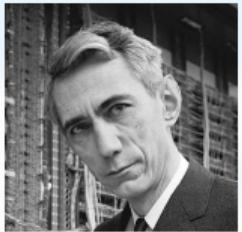
: 5000 bits / 10000 channel uses



: 5000 bits / 10000 channel uses



: Polar code 490bits/10000uses

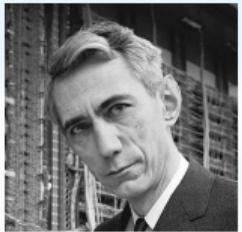


: 5000 bits / 10000 channel uses



: Polar code 490bits/10000uses

In general, $(1/2 - \varepsilon)N$ bits / N uses



: 5000 bits / 10000 channel uses



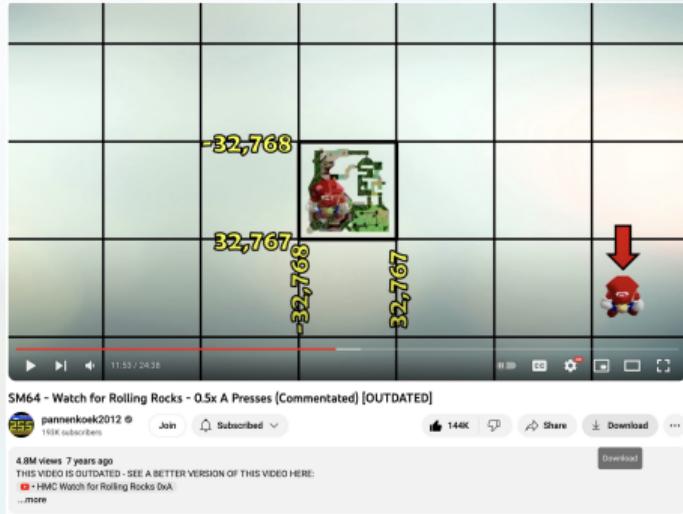
: Polar code 490bits/10000uses

In general, $(1/2 - \varepsilon)N$ bits / N uses

THE END?

5G is for

5G is for



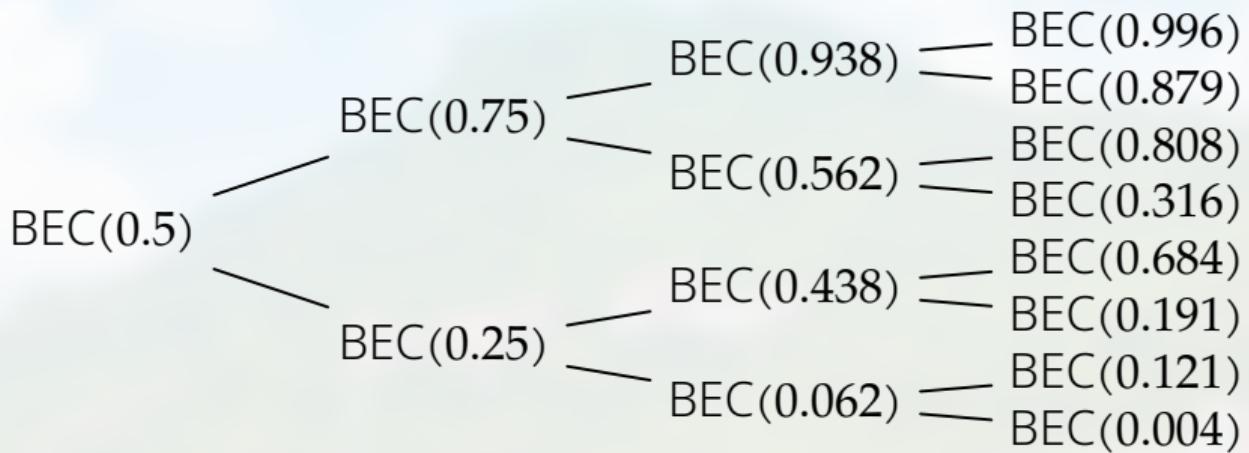
video streaming

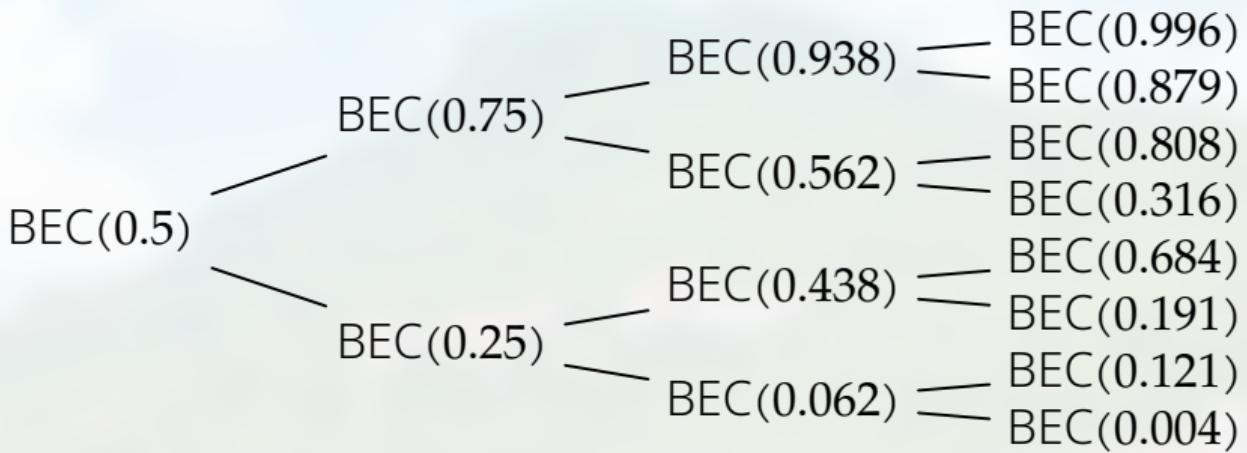
5G is for



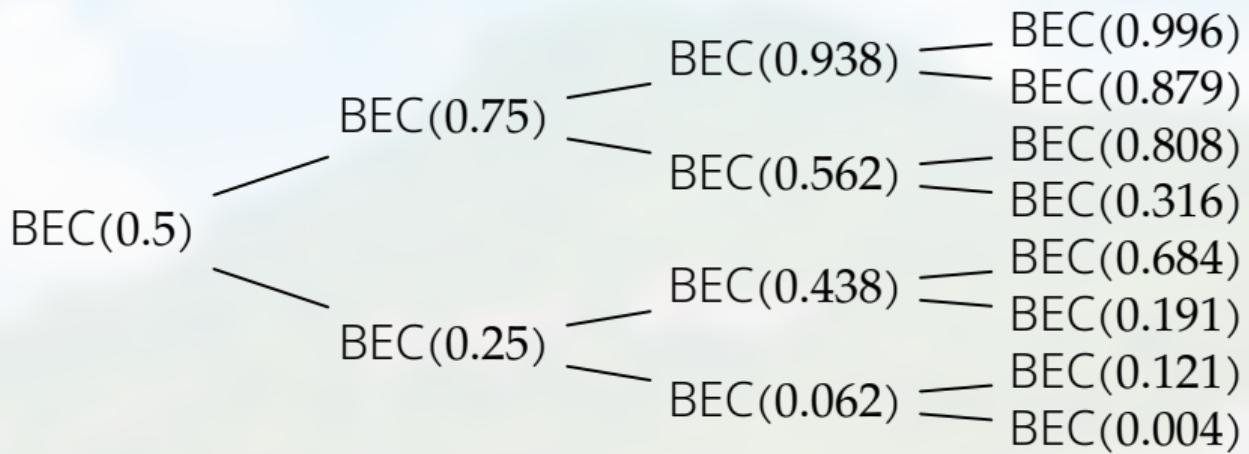
video streaming

flying drone





Early channels are not useful.



Early channels are not useful.
How to accelerate polarization?

$(1/2 - N^{-\rho})N$ bits / N channel uses

History of ρ over BMSC:0

1/2

2015 Guruswami-Xia \longleftrightarrow

2012 Goli-Hassani-Urbanke \longleftrightarrow

2014 Hassani-Alishahi-Urbanke \longleftrightarrow

2014 Goldin-Burshtein \longleftrightarrow

2016 Mondelli-Hassani-Urbanke \longleftrightarrow

$(1/2 - N^{-\rho})N$ bits / N channel uses

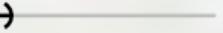
History of ρ over BMSC:0

1/2

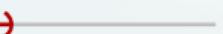
2015 Guruswami-Xia 

2012 Goli-Hassani-Urbanke 

2014 Hassani-Alishahi-Urbanke 

2014 Goldin-Burshtein 

2016 Mondelli-Hassani-Urbanke 

2022 W.-Lin-Vardy-Gabrys 

Improve ρ over BEC: 0

1/2

- 2010 Hassani–Alishahi–Urbanke 2×2 x
- 2010 Korada–Montanari–Telatar–Urbanke 2×2 •
- 2014 Fazeli–Vardy 8×8 •
- 2021 Trofimiuk–Trifonov 16×16 •
- 2021 Trofimiuk 24×24 •
- 2021 Yao–Fazeli–Vardy 32×32 •
- 2021 Yao–Fazeli–Vardy 64×64 •

Improve ρ over BEC: 0

1/2

2010 Hassani–Alishahi–Urbanke 2×2 x

2010 Korada–Montanari–Telatar–Urbanke 2×2 •

2014 Fazeli–Vardy 8×8 •

2021 Trofimiuk–Trifonov 16×16 •

2022 Duursma–Gabrys–Guruswami–Lin–W.2 $\times 2$ /GF4 •

2021 Trofimiuk 24×24 •

2021 Yao–Fazeli–Vardy 32×32 •

2021 Yao–Fazeli–Vardy 64×64 •

The optimal ρ : 0

1/2

2019 Pfister–Urbanke
 q -ary erasure channel, $q \rightarrow \infty$

2021 Fazeli–Hassani–Mondelli–Vardy
binary erasure channel

2022 Guruswami–Riazanov–Ye
binary symmetric memoryless channel

The optimal ρ : 0

1/2

2019 Pfister–Urbanke
 q -ary erasure channel, $q \rightarrow \infty$

2021 Fazeli–Hassani–Mondelli–Vardy
binary erasure channel

2022 Guruswami–Riazanov–Ye
binary symmetric memoryless channel

2021 W.–Duursma
discrete memoryless channel

2011 Alamdar-Yazdi-Kschischang:
Prune the tree to reduce complexity.

2017 El-Khamy-Mahdavifar-Feygin-Lee-Kang:
Pruning reduces complexity by a scalar; still $O(N \log N)$.

2021 Mondelli-Hashemi-Cioffi-Goldsmith,
2021 Hashemi-Mondelli-Fazeli-Vardy-Cioffi-Goldsmith:
Study parallelism vs latency.

2011 Alamdar-Yazdi-Kschischang:
Prune the tree to reduce complexity.

2017 El-Khamy-Mahdavifar-Feygin-Lee-Kang:
Pruning reduces complexity by a scalar; still $O(N \log N)$.

2021 W.-Duursma: $O(N \log \log N)$
if relax the performance requirement.

Trade-off: complexity $\approx O(N \log(-\log(\text{decode error})))$.

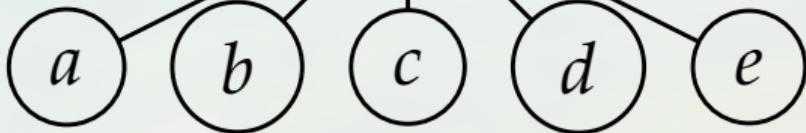
2021 Mondelli-Hashemi-Cioffi-Goldsmith,
2021 Hashemi-Mondelli-Fazeli-Vardy-Cioffi-Goldsmith:
Study parallelism vs latency.

5	3			7			
6			1	9	5		
	9	8				6	
8			6				3
4		8	3				1
7		2				6	
	6			2	8		
		4	1	9			5
			8		7	9	

Low-Density Parity-Check (LDPC) Codes

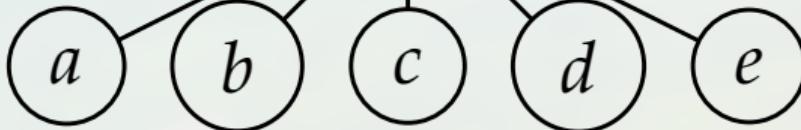
Rule: Every
Sum to an even number

check node



Rule: Every
Sum to an even number

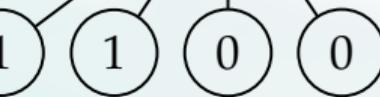
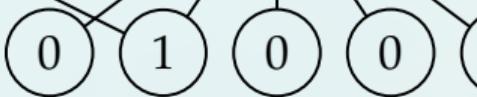
check node

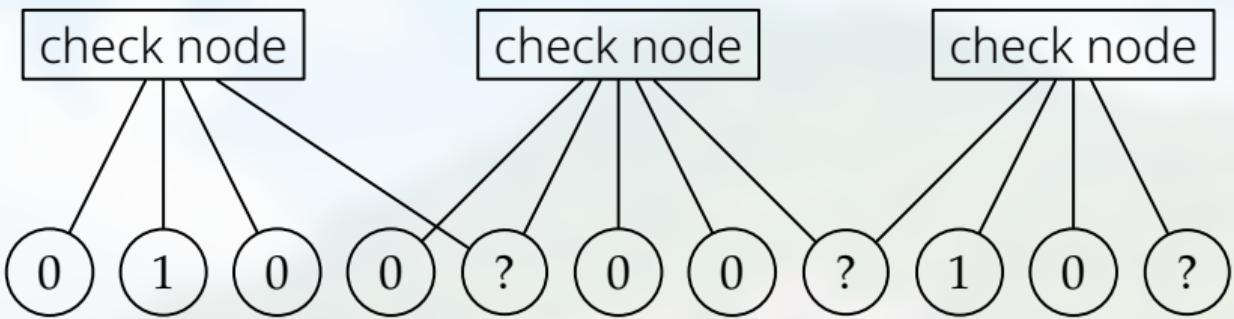


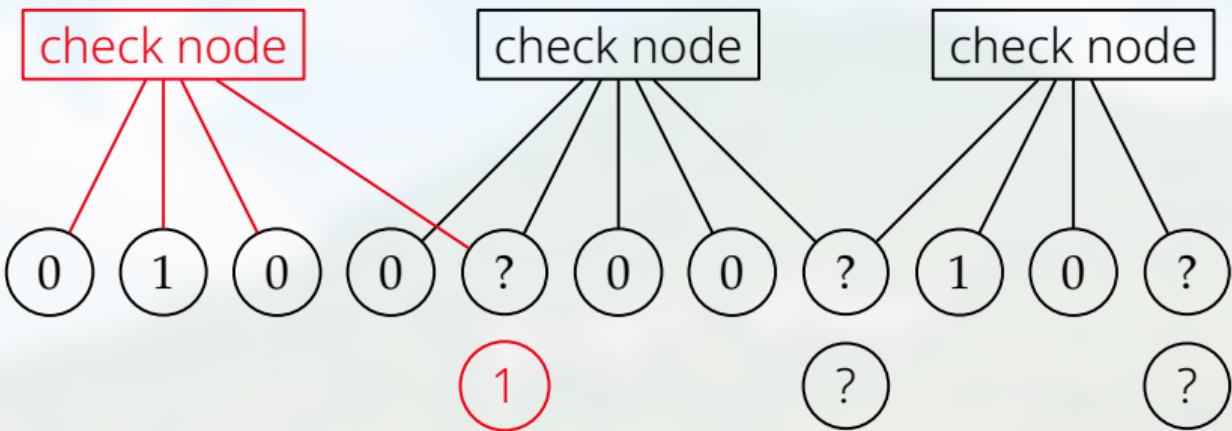
check node

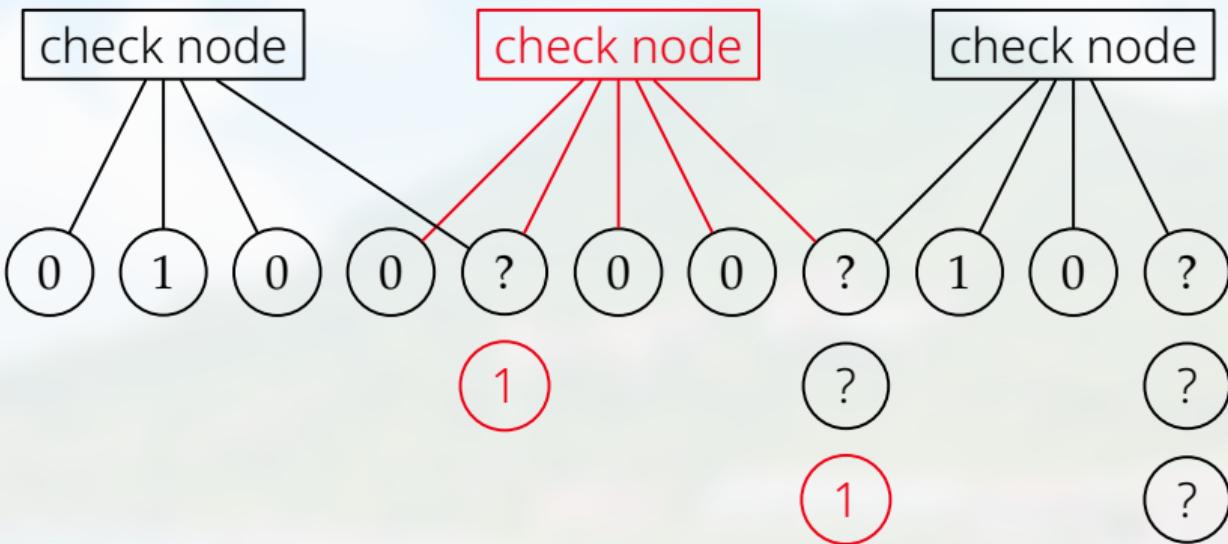
check node

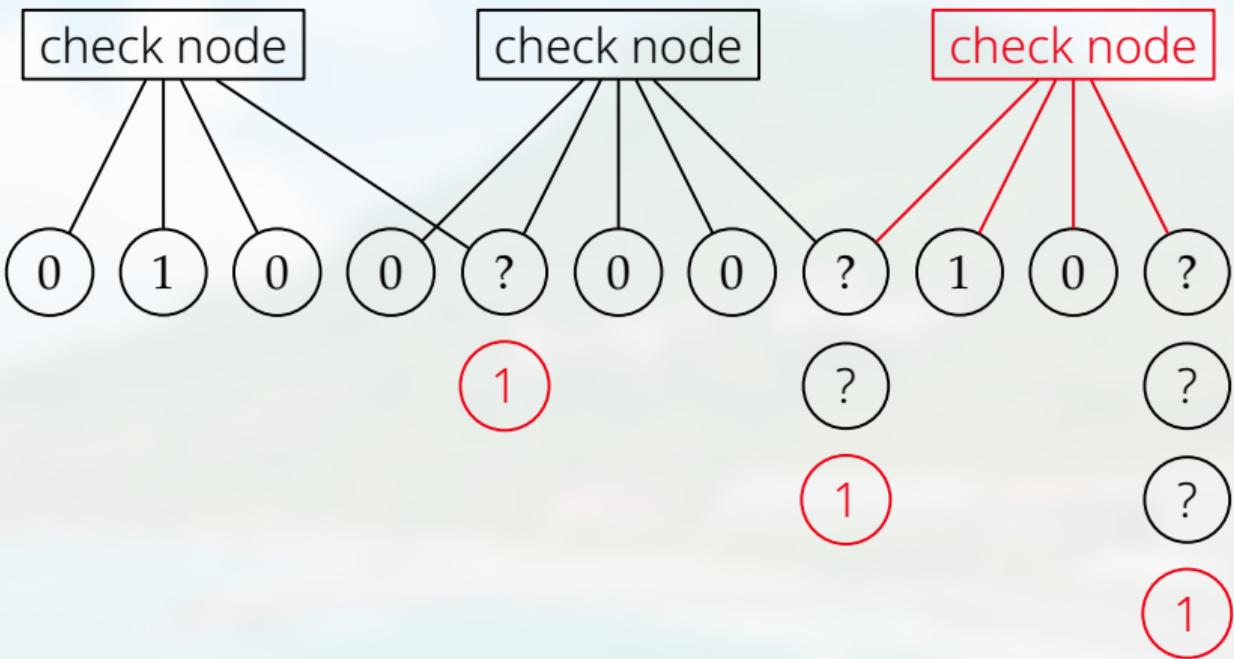
check node

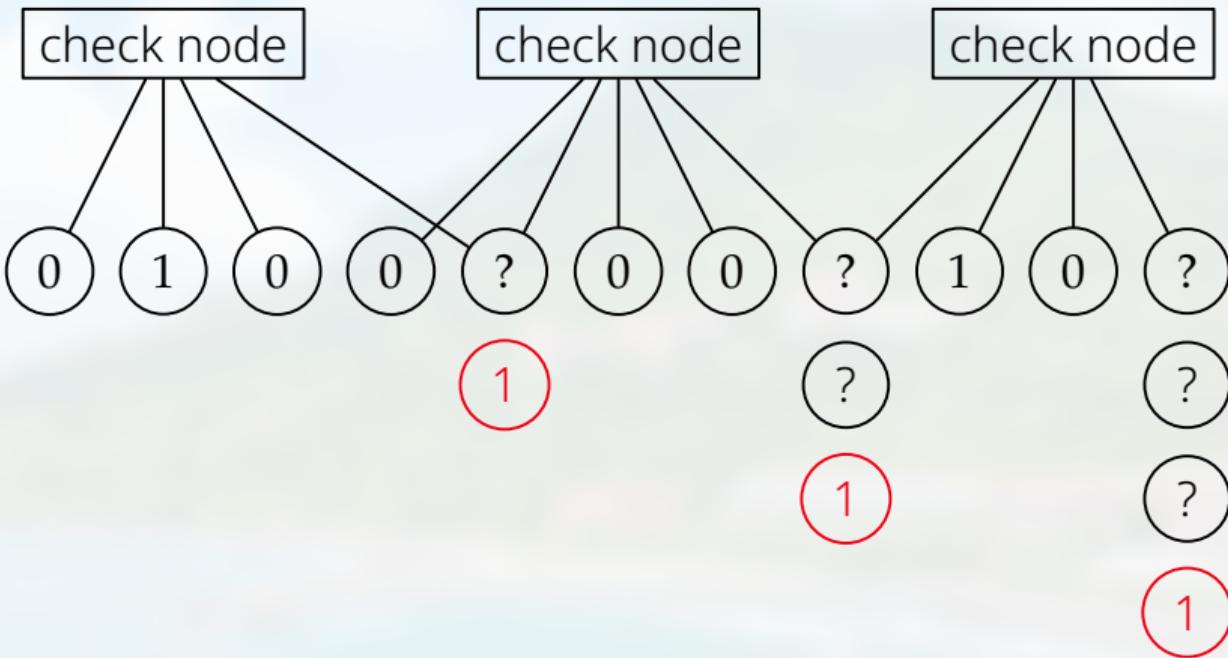




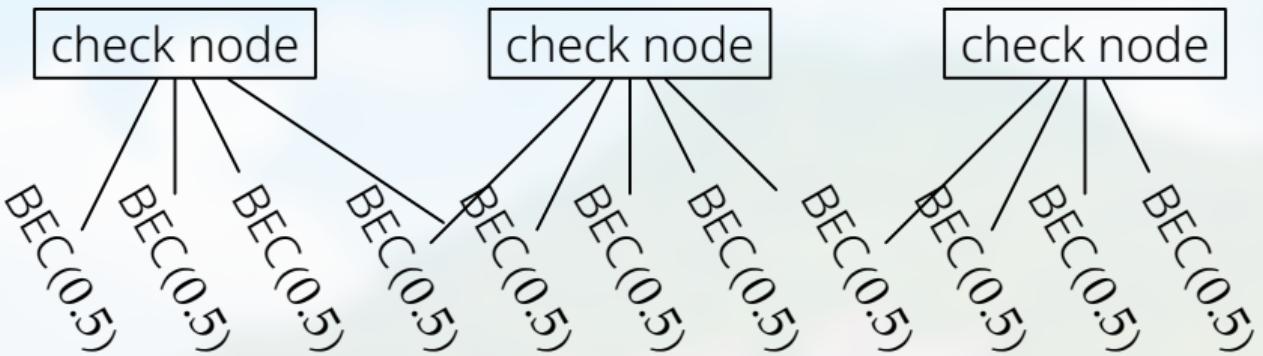


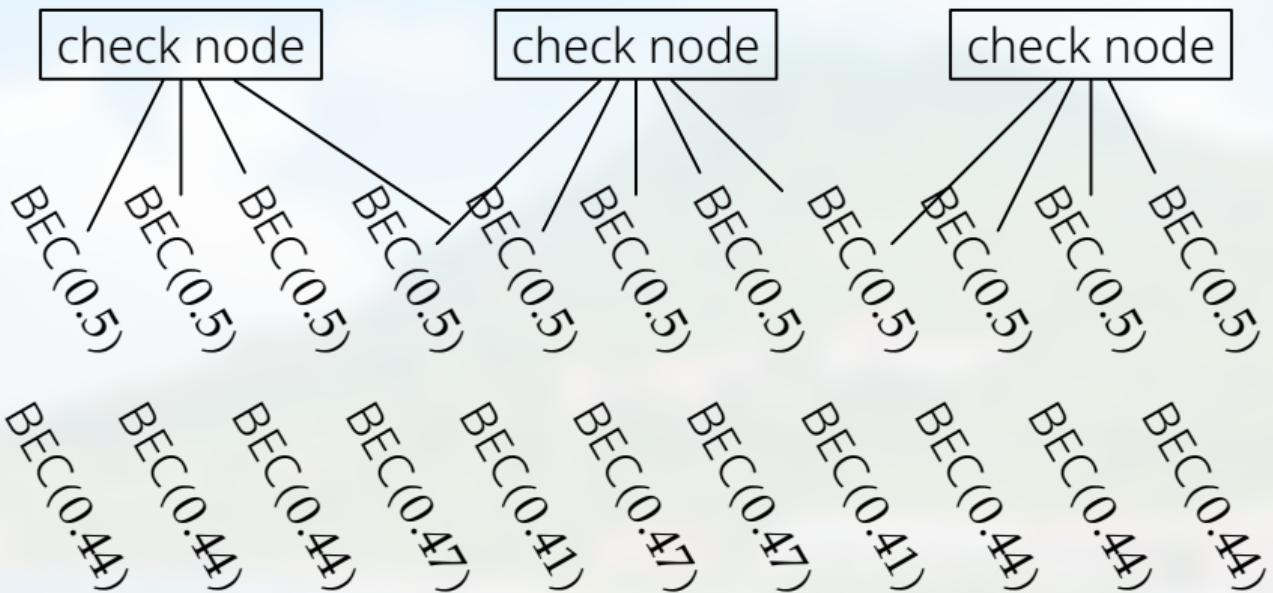


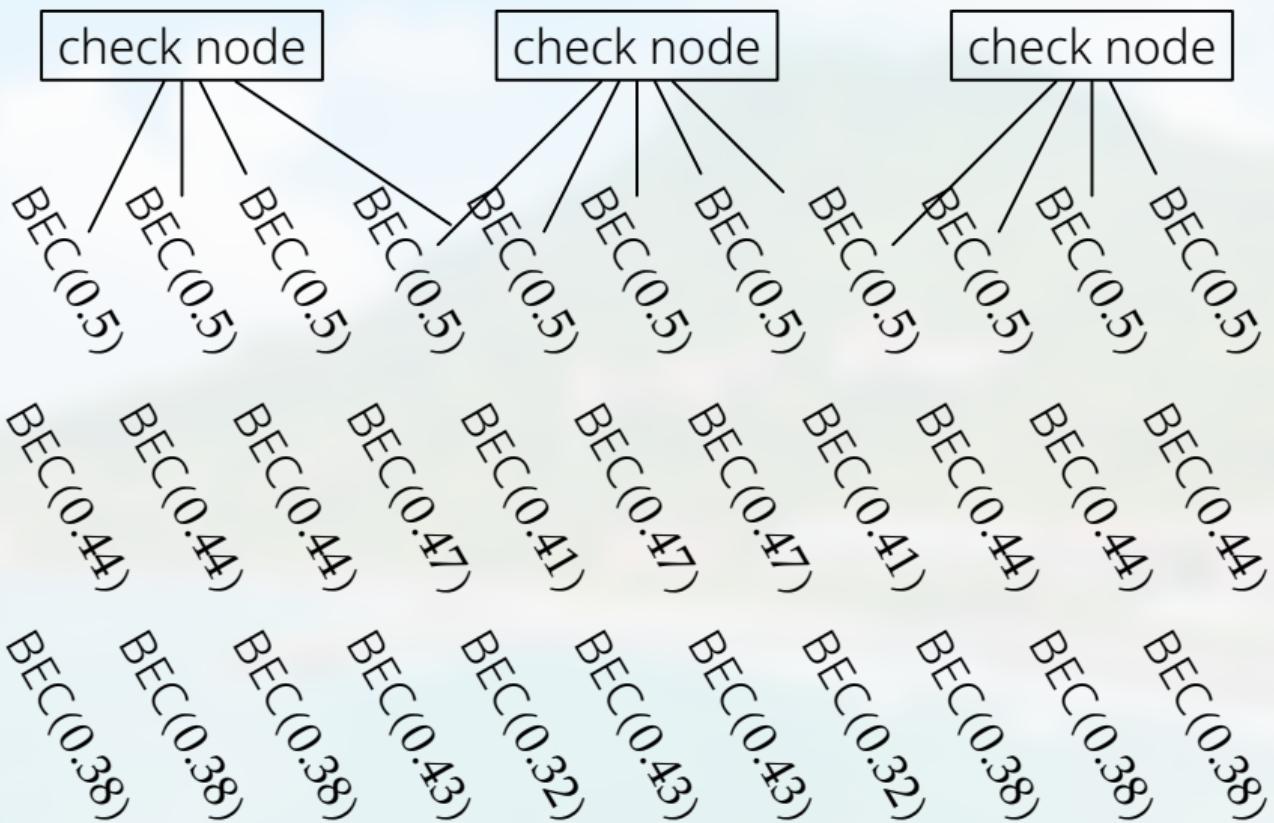




How to analyze, mathematically?









How To Make An Ice Bubble



Rosemary Daniels
692 subscribers

Subscribe



732



0



Share



Download

...



How To Make An Ice Bubble



Rosemary Daniels
690 subscribers

Subscribe



732



0



Share



Download



...

How To Make An Ice Bubble



Rosemary Daniels
690 subscribers

Subscribe



732



0



Share



Download



...

How To Make An Ice Bubble



Rosemary Daniels
690 subscribers

Subscribe



732



0



Share



Download



...



How To Make An Ice Bubble



Rosemary Daniels
690 subscribers

Subscribe



732



0



Share



Download



...



How To Make An Ice Bubble



Rosemary Daniels
690 subscribers

Subscribe



732



0



Share



Download



How To Make An Ice Bubble



Rosemary Daniels
690 subscribers

Subscribe



732



0



Share



Download





How To Make An Ice Bubble



Rosemary Daniels
690 subscribers

Subscribe

732 Share Download



How To Make An Ice Bubble



Rosemary Daniels
690 subscribers

Subscribe

732 Share Download



How To Make An Ice Bubble



Rosemary Daniels
690 subscribers

Subscribe

732 Share Download



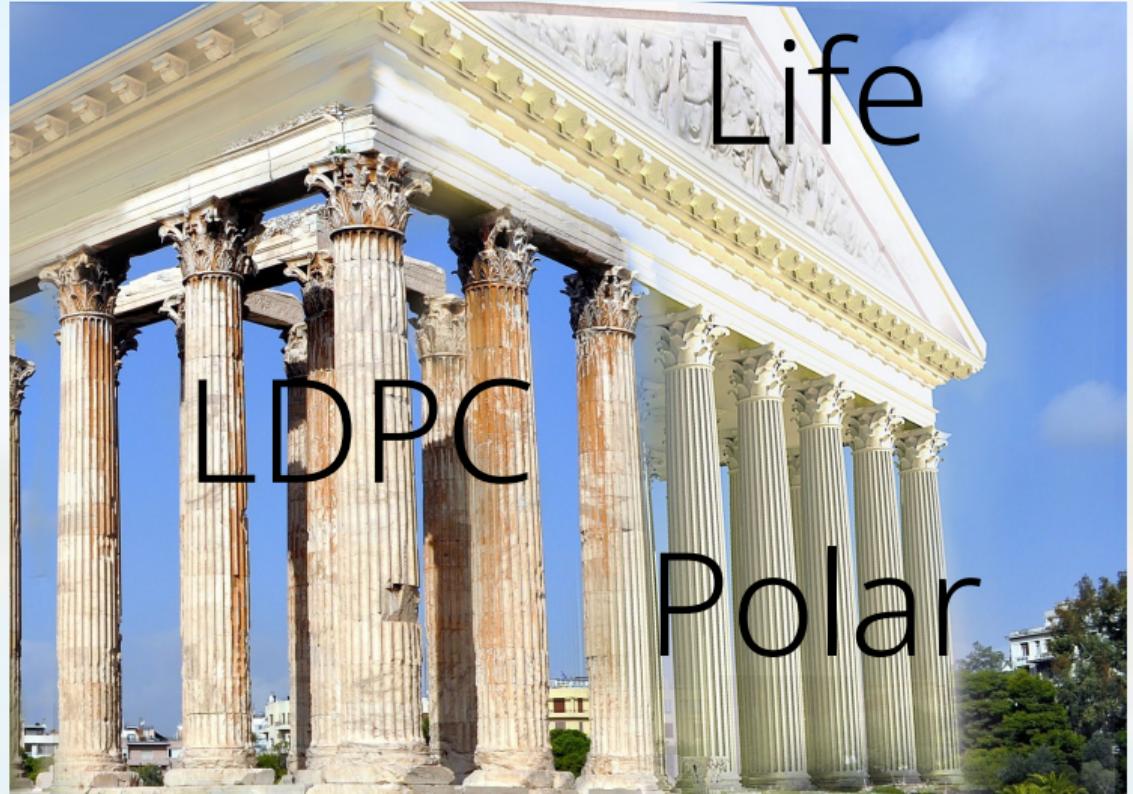
How To Make An Ice Bubble



Rosemary Daniels
690 subscribers

Subscribe

732 Share Download



Channel Evolution

Wikipedia



Game: Breath of the Wild

Challenges

Channel Manipulation is Coding

H B Wang







DNA coding

Permutation

Amplification

Substitution

Insertion

Deletion