

## 0.1 Multiple Initialization Blocks Handler

### 0.1.1 Handler Algorithm

Algorithm 1: Handle	
<b>Input:</b> L, B, R, M	
1	$IB_L \leftarrow \{n \in A_L \mid n.type = INITBLOCK\};$
2	$IB_R \leftarrow \{n \in A_R \mid n.type = INITBLOCK\};$
3	$IB_B \leftarrow \{n \in D_B \mid n.type = INITBLOCK\};$
4	$eIB_L \leftarrow \text{editedNodes}(IB_L, IB_B);$
5	$eIB_R \leftarrow \text{editedNodes}(IB_R, IB_B);$
6	$dIB_L \leftarrow \text{deletedNodes}(IB_L, IB_B, eIB_L);$
7	$dIB_R \leftarrow \text{deletedNodes}(IB_R, IB_B, eIB_R);$
8	<b>foreach</b> $b \in IB_B$ <b>do</b>
9	$l \leftarrow eIB_L[b];$
10	$r \leftarrow eIB_R[b];$
11	<b>if</b> $l \neq null \wedge r \neq null$ <b>then</b>
12	$\text{updateMergeTree}(l, b, r, M);$
13	<b>else if</b> $l \neq null \vee r \neq null$ <b>then</b>
14	<b>if</b> $l \neq null$ <b>then</b>
15	$r \leftarrow \text{find}(r \in dIB_R \rightarrow r.body = b.body);$
16	<b>if</b> $r \neq null$ <b>then</b> $\text{removeNode}(b, M);$
17	<b>else</b>
18	$l \leftarrow \text{find}(l \in dIB_L \rightarrow l.body = b.body);$
19	<b>if</b> $l \neq null$ <b>then</b> $\text{removeNode}(b, M);$
20	<b>end</b>
21	$\text{updateMergeTree}(l, b, r, M);$
22	<b>else</b>
23	$m \leftarrow \text{find}(m \in M \rightarrow m.body = b.body);$
24	$\text{removeNode}(m, M);$
25	<b>end</b>
26	<b>end</b>
27	$aIB_L \leftarrow \text{addedNodes}(IB_L, IB_B, eIB_L);$
28	$aIB_R \leftarrow \text{addedNodes}(IB_R, IB_B, eIB_R);$
29	$DEP \leftarrow \text{dependentNodes}(aIB_L, aIB_R);$
30	<b>foreach</b> $(l, rs) \in DEP$ <b>do</b>
31	$s \leftarrow \varepsilon;$
32	<b>foreach</b> $r \in rs$ <b>do</b>
33	$s \leftarrow s + r.body;$
34	$m \leftarrow \text{find}(m \in M \rightarrow m.body = r.body);$
35	$\text{removeNode}(r, M);$
36	<b>end</b>
37	$m \leftarrow \text{find}(m \in M \rightarrow m.body = l.body);$
38	$m.body \leftarrow \text{conflict}(l.body, \varepsilon, s);$
39	<b>end</b>
40	<b>foreach</b> $l \in aIB_L$ <b>do</b>
41	<b>foreach</b> $r \in aIB_R$ <b>do</b>
42	<b>if</b> $l.body = r.body$ <b>then</b>
43	$m \leftarrow \text{find}(m \in M \rightarrow m.body = r.body);$
44	$\text{removeNode}(m, M);$
45	<b>end</b>
46	<b>end</b>
47	<b>end</b>

**Algorithm 2: Edited Nodes****Input:**  $IB, IB_B$ **Output:** map associating a deleted base node  $b$  in  $IB_B$  and its correspondent added branch node  $a$  in  $IB$ 

```

1  $D \leftarrow \{d \in IB_B \mid (\neg \exists a \in IB)(d.body = a.body)\};$ 
2  $A \leftarrow \{a \in IB \mid (\neg \exists d \in IB_B)(a.body = d.body)\};$ 
3  $matches \leftarrow \emptyset;$ 
4 foreach  $a \in A$  do
5    $S \leftarrow \{d \in D \mid a.body \approx d.body\};$ 
6    $b \leftarrow \underset{s \in S}{\operatorname{argmax}} (\operatorname{similarity}(s.body, a.body));$ 
7   if  $b \neq null$  then  $matches \leftarrow matches \cup \{b : a\};$ 
8 end
9 return  $matches$ 

```

**Algorithm 3: Added Nodes****Input:**  $IB, IB_B, eIB$ **Output:** set of initialization block nodes added by branch

```

1  $A \leftarrow \{n \in IB \mid (\neg \exists b \in IB_B)(n.body = b.body)\};$ 
2  $A \leftarrow \{n \in A \mid (\neg \exists e \in eIB)(n.body = e.value.body)\};$ 
3 return  $A;$ 

```

**Algorithm 4: Deleted Nodes****Input:**  $IB, IB_B, eIB$ **Output:** set of initialization block nodes deleted by branch

```

1  $D \leftarrow \{b \in IB_B \mid (\neg \exists n \in IB)(b.body = n.body)\};$ 
2  $D \leftarrow \{n \in D \mid (\neg \exists e \in eIB)(n.body = e.key.body)\};$ 
3 return  $D;$ 

```

**Algorithm 5: Update Merge Tree****Input:**  $l, b, r, M$ 

```

1  $m \leftarrow \operatorname{find}(m \in M \rightarrow m.body = l.body);$ 
2  $m.body \leftarrow \operatorname{textualMerge}(l.body, b.body, r.body);$ 
3  $m \leftarrow \operatorname{find}(m \in M \rightarrow m.body = r.body);$ 
4 removeNode $(m, M);$ 

```

**Algorithm 6: Dependent Nodes****Input:**  $aIB_L, aIB_R$ **Output:** map associating an added left node  $l$  in  $aIB_L$  and all added right nodes  $r$  in  $aIB_R$  with common global variables

```

1  $DEP \leftarrow \emptyset;$ 
2 foreach  $l \in aIB_L$  do
3    $DEP \leftarrow DEP \cup \{l : \emptyset\};$ 
4    $V_L \leftarrow \operatorname{globalVariables}(l);$ 
5   foreach  $r \in aIB_R$  do
6      $V_R \leftarrow \operatorname{globalVariables}(r);$ 
7     if  $V_L \cap V_R \neq \emptyset$  then  $DEP[l] \leftarrow DEP[l] \cup r;$ 
8   end
9 end
10 return  $DEP;$ 

```