

1 Renaming Handler

1.1 Early Concepts

1.1.1 Possibly renamed without body changes nodes

$$R_{wobc}(T, B) = \{b \in B \mid (\neg \exists t \in T)(t.id = b.id) \wedge (\exists t \in T)(t.body = b.body)\}$$

1.1.2 Possibly deleted or renamed with body changes nodes

$$DR_{wbc}(T, B) = \{b \in B \mid (\neg \exists t \in T)(t.id = b.id \vee t.body = b.body)\}$$

1.1.3 Nodes IDs similarity

$$a.id \approx b.id \leftrightarrow a.id.name = b.id.name \vee a.id.params = b.id.params$$

1.2 Match Algorithm

Algorithm 1: Match Algorithm

Input: L, B, R, M

Output: Set of quadruples (l, b, r, m) consisting of the base node b and its corresponding left node l , right node r and merge node m

```
1 matches  $\leftarrow \emptyset$ ;  
2 foreach  $b \in DR_{wbc}(L, B) \cup DR_{wbc}(R, B) \cup R_{wobc}(L, B) \cup R_{wobc}(R, B)$  do  
3    $l \leftarrow \text{correspondentNode}(b, L)$ ;  
4    $r \leftarrow \text{correspondentNode}(b, R)$ ;  
5    $m \leftarrow \text{mergeNode}(l, r, M)$ ;  
6    $matches \leftarrow matches \cup (l, b, r, m)$ ;  
7 end  
8 return matches
```

Algorithm 2: Correspondent Node

Input: b, T

Output: b's correspondent node on tree T

```
1  $t \leftarrow \text{findFirst}(t \in T \rightarrow t.id = b.id)$ ;  
2 if  $t = \text{null}$  then  
3    $t \leftarrow \text{findFirst}(t \in T \rightarrow t.body = b.body)$ ;  
4 end  
5 if  $t = \text{null}$  then  
6    $t \leftarrow \text{findFirst}(t \in T \rightarrow t.body \approx b.body \wedge t.id \approx b.id)$ ;  
7 end  
8 if  $t = \text{null}$  then  
9    $t \leftarrow \text{findFirst}(t \in T \rightarrow t.body = \text{substring}(b.body) \vee b.body = \text{substring}(t.body))$ ;  
10 end  
11 return t;
```

Algorithm 3: Merge Node**Input:** l, r, M **Output:** l and r 's merge node on tree M

```

1 if  $l \neq null$  then
2   | return  $find(m \in M \rightarrow m.id = l.id)$ ;
3 end
4 if  $r \neq null$  then
5   | return  $find(m \in M \rightarrow m.id = r.id)$ ;
6 end
7 return  $null$ ;

```

1.3 Handle Algorithms**Algorithm 4: Check References and Merge Methods Variant****Input:** $(l, b, r, m), M$

```

1 if  $l.id = b.id \vee r.id = b.id$  then
2   |  $m.body = textualMerge(l, b, r)$ ;
3   |  $removeUnmatchedNode(l, r, m, M)$ ;
4 else if  $l.id \neq r.id$  then
5   |  $m.body = conflict(l, b, r)$ ;
6   |  $removeUnmatchedNode(l, r, m, M)$ ;
7 else if  $l.body \neq r.body$  then
8   | if  $newReferenceTo(l) \vee newReferenceTo(r)$  then
9     |  $m.body = conflict(l, b, r)$ ;
10  | else
11  |  $m.body = textualMerge(l, b, r)$ ;
12  | end
13  |  $removeUnmatchedNode(l, r, m, M)$ ;
14 end

```

Algorithm 5: Merge Methods Variant**Input:** $(l, b, r, m), M$

```

1  $m.body = textualMerge(l, b, r)$ ;
2  $removeUnmatchedNode(l, r, m, M)$ ;

```

Algorithm 6: Check Textual and Keep Both Methods Variant**Input:** $(l, b, r, m), M$

```

1 if  $l.id = b.id \vee r.id = b.id$  then
2   | if  $textualMergeHasConflictInvolvingSignature(b)$  then
3     |  $m.body = conflict(l, b, r)$ ;
4     |  $removeUnmatchedNode(l, r, m, M)$ ;
5   | end
6 else if  $l.id \neq r.id \wedge l.body = r.body$  then
7   |  $m.body = conflict(l, b, r)$ ;
8   |  $removeUnmatchedNode(l, r, m, M)$ ;
9 end

```

Algorithm 7: Keep Both Methods Variant**Input:** $(l, b, r, m), M$

```
1 if  $(l.id = b.id \vee r.id = b.id) \wedge \text{hasConflict}(m)$  then  
2   |  $\text{removeConflict}(m);$   
3 end
```

Algorithm 8: Remove Unmatched Node**Input:** l, r, m, M

```
1 if  $l.id = m.id \wedge r.id \neq m.id$  then  
2   |  $\text{removeNode}(r, M);$   
3 end
```