

0.1 Deletions Handler

0.1.1 Handler Algorithm

Algorithm 1: Handle	
Input: L, B, R, M	
1	$T_L \leftarrow \text{treeToText}(L);$
2	$T_B \leftarrow \text{treeToText}(B);$
3	$T_R \leftarrow \text{treeToText}(R);$
4	foreach $d_l \in D_L$ do
5	if $d_l.\text{children} \neq \emptyset$ then
6	$r \leftarrow \text{find}(r \in R \rightarrow r.\text{id} = d_l.\text{id});$
7	$m \leftarrow \text{find}(m \in M \rightarrow m.\text{id} = d_l.\text{id});$
8	if $\text{sameShape}(d_l, r) \wedge d_l.\text{body} = r.\text{body}$ then $\text{removeNode}(m, M);$
9	else if $\text{newReference}(d_l.\text{id}, T_B, T_R)$ then $m.\text{parent.addChild}(r, m.\text{index});$
10	else
11	$\text{removeNode}(m, M);$
12	$a_l \leftarrow \text{renamingMatch}(A_L, d_l, T_B, T_L);$
13	if $a_l \neq \text{null}$ then
14	$r.\text{id} \leftarrow a_l.\text{id};$
15	$\text{removeNode}(a_l, M);$
16	$m.\text{parent.addChild}(r, m.\text{index});$
17	else
18	$n.\text{id} \leftarrow r.\text{id};$
19	$n.\text{type} \leftarrow r.\text{type};$
20	$n.\text{body} \leftarrow \text{conflict}(\varepsilon, d_l, r);$
21	$m.\text{parent.addChild}(n, m.\text{index});$
22	end
23	end
24	end
25	end
26	foreach $d_r \in D_R$ do
27	if $d_r.\text{children} \neq \emptyset$ then
28	$l \leftarrow \text{find}(l \in L \rightarrow l.\text{id} = d_r.\text{id});$
29	$m \leftarrow \text{find}(m \in M \rightarrow m.\text{id} = d_r.\text{id});$
30	if $\text{sameShape}(d_r, l) \wedge d_r.\text{body} = l.\text{body}$ then $\text{removeNode}(m, M);$
31	else if $\text{newReference}(d_r.\text{id}, T_B, T_L)$ then $m.\text{parent.addChild}(l, m.\text{index});$
32	else
33	$\text{removeNode}(m, M);$
34	$a_r \leftarrow \text{renamingMatch}(A_R, d_r, T_B, T_R);$
35	if $a_r \neq \text{null}$ then
36	$l.\text{id} \leftarrow a_r.\text{id};$
37	$\text{removeNode}(a_r, M);$
38	$m.\text{parent.addChild}(l, m.\text{index});$
39	else
40	$n.\text{id} \leftarrow l.\text{id};$
41	$n.\text{type} \leftarrow l.\text{type};$
42	$n.\text{body} \leftarrow \text{conflict}(l, d_r, \varepsilon);$
43	$m.\text{parent.addChild}(n, m.\text{index});$
44	end
45	end
46	end
47	end

Algorithm 2: Same Shape**Input:** A, B **Output:** whether nodes A and B have same shape

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1 if  $A.children = \emptyset \wedge B.children = \emptyset$  then return  $A.type = B.type$ ;
2 if  $A.children = \emptyset \vee B.children = \emptyset$  then return  $FALSE$ ;
3 if  $|A.children| \neq |B.children|$  then return  $FALSE$ ;
4  $result \leftarrow TRUE$ ;
5 foreach  $(a, b) \in (A.children, B.children)$  do
6    $result \leftarrow result \wedge sameShape(a, b)$ ;
7 end
8 return  $result$ ;

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Algorithm 3: New Reference**Input:** id, T_B, T **Output:** whether there is a new reference to id in T

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1 return  $countReferences(id, T) > countReferences(id, T_B)$ ;

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Algorithm 4: Renaming Match**Input:** A, d, T_B, T **Output:** added node a in A with the same shape and similar body as deleted node d , such that there are no new references to a 's id in T

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1 return  $find(a \in A \rightarrow sameShape(a, d) \wedge a.body \approx d.body \wedge \neg newReference(a.id, T_B, T))$ ;

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