

Digital Electronics and Microprocessor Systems (ELEC211)

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Week 4 – Lecture 10

Microprocessor Systems



Question

What values are held in r4, r5 and r6 after the execution of the following assuming that the value in r0 is 0xE0300A0B?

r0:=1110 0000 0011 0000 0000 1010 0000 1011₂

LSLS r4, r0, #8 ; 2 hex digits left

ASRS r5, r0, #16 ; 4 hex digits right

LSRS r6, r0, #12 ; 3 hex digits right

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Answer

LSLS r4, r0, #8

r0 held:

0xE0300A0B =

1110 0000 0011 0000 0000 1010 0000 1011₂

after the instruction is executed r4 would hold:

0011 0000 0000 1010 0000 1011 0000 0000₂



So the value in r4 is 0x300A0B00.

Answer

ASRS r5, r0, #16


r0 held:

0xE0300A0B =

1110 0000 0011 0000 0000 1010 0000 1011₂

after the instruction is executed r5 would hold:

1111 1111 1111 1111 1110 0000 0011 0000₂



So the value in r5 is 0xFFFFE030.

Answer

LSRS r6, r0, #12

r0 held:

0xE0300A0B =

1110 0000 0011 0000 0000 1010 0000 1011₂

after the instruction is executed r6 would hold:

0000 0000 0000 1110 0000 0011 0000 0000₂



So the value in r6 is 0x000E0300.



Question

What values are held in r3 and r4 after the execution of the following assuming that the value in r0 is 0x00000005 ($= 5_{10}$)?

```
LSLS r3, r0, #1  
LSLS r4, r0, #6
```

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Answer

LSLS r3, r0, #1

r0 held 0x00000005 which is:

0000 0000 0000 0000 0000 0000 0000 0101₂

shift left by 1 bit so that the value in r3 is 0x0000000A

0000 0000 0000 0000 0000 0000 0000 1010₂

$$= 10_{10} = (2 \times 5), \text{ where } 2 = 2^1$$

LSLS r4, r0, #6

shift left by 6 bits so that the value in r4 is 0x00000140

0000 0000 0000 0000 0000 0001 0100 0000₂

$$= 320_{10} = (64 \times 5), \text{ where } 64 = 2^6$$



Question

? Assuming that the values in r5 are:

(i) $0x0000001C$ ($= 28_{10}$)

0000 0000 0000 0000 0000 0000 0001 1100₂

and (ii) $0xFFFFFE4$ ($= -28_{10}$)

1111 1111 1111 1111 1111 1111 1110 0100₂

What values are held in r7 after the execution of the following instruction?

ASRS r7, r5, #2

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Answer

ASRS r7, r5, #2

If value in r5 is 0x0000001C:

0000 0000 0000 0000 0000 0000 0001 1100₂

then value in r7 is

0000 0000 0000 0000 0000 0000 0000 0111₂ =

0x00000007 = (28₁₀ ÷ 4), where 4 = 2²

If value in r5 is 0xFFFFFFE4:

1111 1111 1111 1111 1111 1111 1110 0100₂

then value in r7 is

1111 1111 1111 1111 1111 1111 1111 1001₂ =

0xFFFFFFFF9 = -7₁₀ = (-28₁₀ ÷ 4), where 4 = 2²

Week 4 – Lecture 11

Microprocessor Systems



Question

? What value is stored in the link register, r14, when the ADDS instruction is executed?

Address Instruction		Subroutine	
0x00000FC2	MOV r1, #0	Address	Instruction
0x00000FC4	BL subX	subX	BICS r6, r1
0x00000FC8	MOV r0, #1	0x00008108	ANDS r2, r1
0x00000FCA	ADDS r1, r0, #1	0x0000810A	EORS r5, r6
0x00000FCC	BL subX	0x0000810C	BX lr
0x00000FD0	SUBS r3, r2, #2		

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Answer

Address	Instruction	Subroutine	
0x00000FC2	MOV r1, #0	Address	Instruction
0x00000FC4	BL subX	subX	BICS r6, r1
0x00000FC8	MOV r0, #1	0x00008108	ANDS r2, r1
0x00000FCA	ADDS r1, r0, #1	0x0000810A	EORS r5, r6
0x00000FCC	BL subX	0x0000810C	BX lr
0x00000FD0	SUBS r3, r2, #2		

- The 'subX' subroutine is called twice (**BL subX**)
- The **ADDS** instruction is after the first pass through the subroutine but before the second pass.
- So the **link register**, r14, will hold the return address for the first pass, which is **0x00000FC8**.



Question

? What value is stored in the link register, r14, when the SUBS instruction is executed?

Address	Instruction	Subroutine	
0x00000FC2	MOV r1, #0	Address	Instruction
0x00000FC4	BL subX	subX	BICS r6, r1
0x00000FC8	MOV r0, #1	0x00008108	ANDS r2, r1
0x00000FCA	ADDS r1, r0, #1	0x0000810A	EORS r5, r6
0x00000FCC	BL subX	0x0000810C	BX lr
0x00000FD0	SUBS r3, r2, #2		

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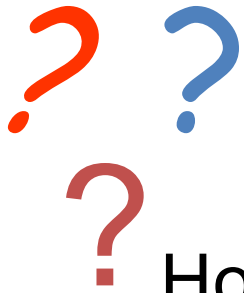
Answer

Address	Instruction	Subroutine	
0x00000FC2	MOV r1, #0	Address	Instruction
0x00000FC4	BL subX	subX	BICS r6, r1
0x00000FC8	MOV r0, #1	0x00008108	ANDS r2, r1
0x00000FCA	ADDS r1, r0, #1	0x0000810A	EORS r5, r6
0x00000FCC	BL subX	0x0000810C	BX lr
0x00000FD0	SUBS r3, r2, #2		

- The 'subX' subroutine is called twice (**BL subX**)
- The **SUBS** instruction is after the second pass through the subroutine.
- So the **link register**, r14, will hold the return address for the second pass, which is **0x00000FD0**.

Week 4 – Lecture 12

Microprocessor Systems



Question

How many clock cycles do the following instructions take to execute if the value in r0 is (a) not equal to 0 and (b) equal to 0? What is the CPI?

```
        MOVS r0, r0
        BNE cont          ; branch if not zero
        ADDS r3, r4, r5
        EORS r2, r1, r4
cont     SUBS r6, r7, r2
```

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Answer

(a) r0:= value not equal to 0

- The branch is taken so that needs 3 clock cycles.
- The MOVS and SUBS instructions take 1 each
- so 5 clock cycles are required in total.
- 5 clock cycles and 3 instructions execute = 1.66 CPI

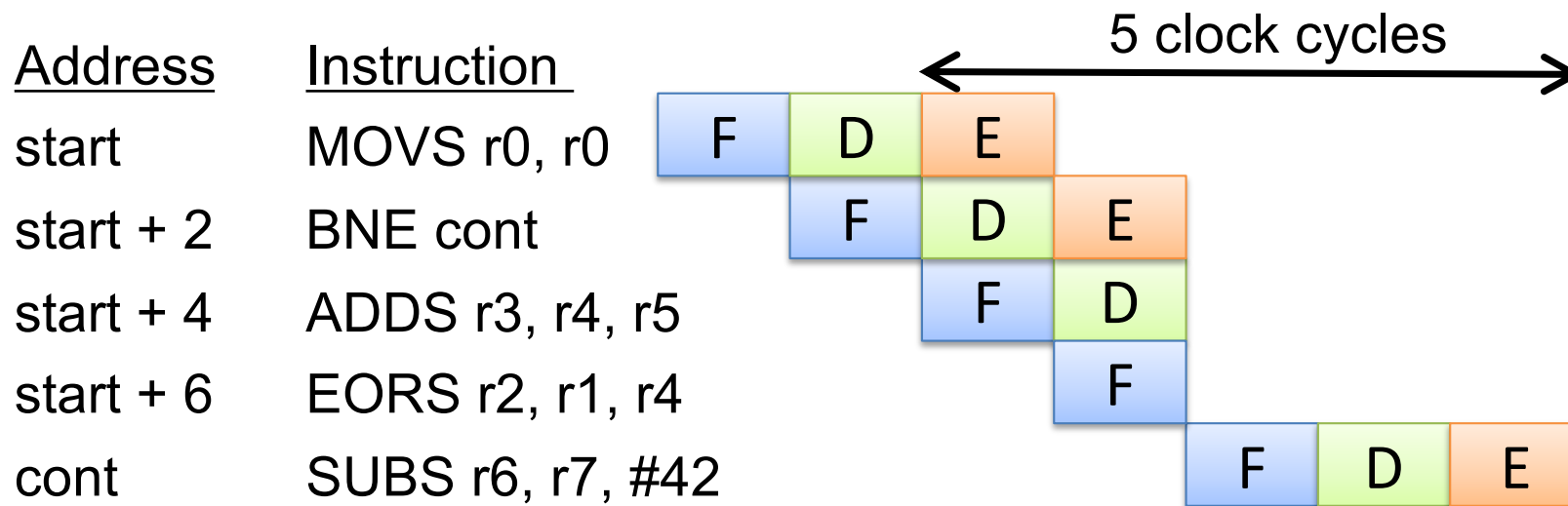
(b) r0:= value equal to 0

- The branch is not taken so that only needs one clock cycle. The other four instructions (MOVS, ADDS, EORS and SUBS) each take 1 so again 5 clock cycles are needed.
- 5 clock cycles and 5 instructions execute = 1 CPI

Answer

(a) $r0 := \text{value not equal to } 0$

The first instruction clears the zero flag ($Z=0$) so that the condition, NE, for the branch is met. The program counter is reloaded with the address of the SUBS instruction. The ADDS and EORS instructions are not executed.

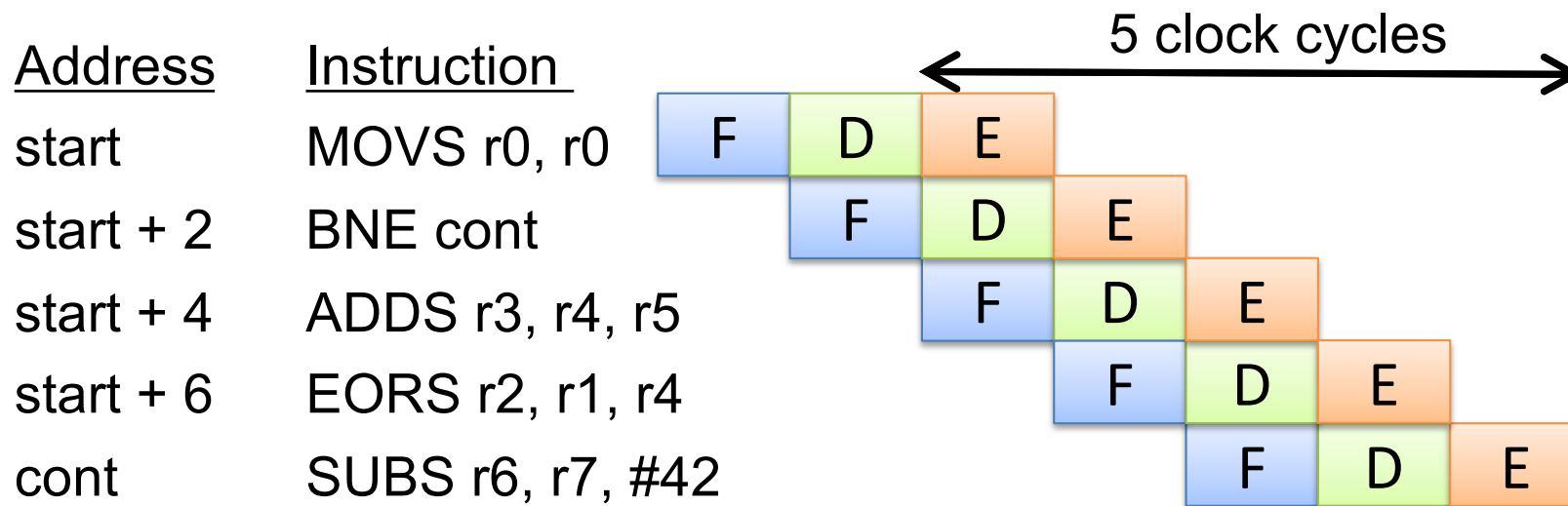


$$5/3 = 1.66 \text{ CPI}$$

Answer

(b) $r0 := \text{value equal to } 0$

The first instruction sets the zero flag ($Z=1$) so that the condition, NE, for the branch is not met. The program counter continues to increment by 2 so that the ADDS, EORS and SUBS instructions are executed in that order.



$$5/5 = 1 \text{ CPI (Optimum)}$$