

〈色の塗り方〉

$\text{colorYB}(x, n, z) \stackrel{\text{def}}{=} \begin{cases} Y & (x \leq z \leq x+n \wedge \text{eqn}(z \bmod 2, 0) == \text{true}) \\ B & (x \leq z \leq x+n \wedge \text{eqn}(z \bmod 2, 1) == \text{true}) \\ Y & (\text{otherwise}) \end{cases}$

Lem
 $(\text{eqn}(z \bmod 2, 0) == \text{true}) \Leftrightarrow \text{even}(z)$
 $\text{even}(z) \text{ と略記 (定義?)}$

Lem YB1

$\forall x, n, \text{even}(x) \rightarrow$

$\forall i (0 \leq i \leq n \wedge \text{even}(i)) \rightarrow \text{colorYB}(x, n, x+i) = Y$

pf. even(x)となる x をとる.

$0 \leq i \leq n \Rightarrow \text{even}(i)$ となる n, i をとる.

$0 \leq i \leq n \wedge x \leq x+i \leq x+n$.

$\text{even}(x), \text{even}(i) \Rightarrow \text{even}(x+i)$.

よって, $\text{colorYB}(x, n, x+i) = Y$. //

不要?

Lem YB2

$\forall x, n, \text{even}(x) \rightarrow$

$\forall i (0 \leq i \leq n \wedge \text{odd}(i)) \rightarrow \text{colorYB}(x, n, x+i) = B$

pf. even(x)となる x をとる.

$0 \leq i \leq n \Rightarrow \text{odd}(i)$ となる n, i をとる.

$0 \leq i \leq n \wedge x \leq x+i \leq x+n$.

$\text{even}(x), \text{odd}(i) \Rightarrow \text{odd}(x+i)$.

よって, $\text{colorYB}(x, n, x+i) = B$. //

Lem YB3

$\forall x, n, \text{odd}(x) \rightarrow$

$\forall i (0 \leq i \leq n \wedge \text{even}(i) \rightarrow \text{colorYB}(x, n, x+i) = B)$

prf. $\text{odd}(x)$ となる x をとる.

$0 \leq i \leq n$ かつ $\text{even}(i)$ となる n, i をとる.

$0 \leq i \leq n$ 且 $x \leq x+i \leq x+n$.

$\text{odd}(x), \text{even}(i)$ 且 $\text{odd}(x+i)$.

よって, $\text{colorYB}(x, n, x+i) = B$. //

→ 不要?

Lem YB4

$\forall x, n, \text{odd}(x) \rightarrow$

$\forall i (0 \leq i \leq n \wedge \text{odd}(i) \rightarrow \text{colorYB}(x, n, x+i) = Y)$

prf. $\text{odd}(x)$ となる x をとる.

$0 \leq i \leq n$ かつ $\text{odd}(i)$ となる n, i をとる.

$0 \leq i \leq n$ 且 $x \leq x+i \leq x+n$.

$\text{odd}(x), \text{odd}(i)$ 且 $\text{even}(x+i)$.

よって, $\text{colorYB}(x, n, x+i) = Y$. //

Lem YB5

$\forall x, n, \text{even}(n) \rightarrow$

$(\text{colorYB}(x, n, x) = \text{colorYB}(x, n, x+n))$.

prf. $x \in \mathbb{Z}$.

$\text{even}(n)$ となる $n \in \mathbb{Z}$.

$0 \leq 0, 0 \leq n \Rightarrow 0 \leq 0 \leq n$.

$0 \leq n, n \leq n \Rightarrow 0 \leq n \leq n$.

• $\text{even}(x)$ のとき

Lem YB1 より $\text{colorYB}(x, n, x+0) = Y$

Lem YB1 より $\text{colorYB}(x, n, x+n) = Y$

よって, $\text{colorYB}(x, n, x) = \text{colorYB}(x, n, x+n)$.

• $\text{odd}(x)$ のとき

Lem YB3 より $\text{colorYB}(x, n, x+0) = B$.

Lem YB3 より $\text{colorYB}(x, n, x+n) = B$.

よって, $\text{colorYB}(x, n, x) = \text{colorYB}(x, n, x+n)$. //

EvenA

$\forall x, y, n \in \mathbb{N}$,

$\forall i (0 \leq i \leq n \rightarrow \text{Cpos}(x+i, y, \text{colorYB}(x, n, x+i))) \rightarrow \forall i (0 \leq i \leq n-1 \rightarrow \text{Cpos}(x+i, y+1, R))$

prf. x, y をとる.

$\text{even}(n)$ となる n をとる.

$0 \leq i \leq n-1$ となる i をとる.

$i \leq n-1 \leq n$ となるので $0 \leq i \leq n$

$\therefore \text{Cpos}(x+i, y, \text{colorYB}(x, n, x+i))$.

また, $i \leq n-1 \neq i+1 \leq n$

つまり, $0 \leq i \neq i+1 \leq n$ となるので, $0 \leq i+1 \leq n$

$\therefore \text{Cpos}(x+i, y, \text{colorYB}(x, n, x+i+1))$.

C-exists $\exists c' \text{ Cpos}(x+i, y+1, c) \text{ となる } c \text{ の存在}.$

$c = c' \wedge \exists c' \text{ Cpos}(x+i, y+1, c')$.

o even(x) のとき

- even(i) のとき

even(x), even(i) \neq even($x+i$).

$\therefore \text{colorYB}(x, n, x+i) = Y$.

すなわち, $\text{Cpos}(x+i, y, Y)$.

even($x+i$), odd(i) \neq odd($x+i+1$)

$\therefore \text{colorYB}(x, n, x+i+1) = B$

C-mix $\exists c' c' = \text{mix}(Y, B) = R$

$\therefore \text{Cpos}(x+i, y+1, R)$

- $\text{odd}(i) \wedge$

$\text{even}(x), \text{odd}(i) \Rightarrow \text{odd}(x+i)$.

$\vdash \tau, \text{colorYB}(x, n, x+i) = B$.

すなはち, $\text{Cpos}(x+i, y, B)$.

$\text{odd}(x+i), \text{odd}(i) \Rightarrow \text{even}(x+i+1)$

$\vdash \tau, \text{colorYB}(x, n, x+i+1) = Y$

C-mix $f' \quad c' = \text{mix}(B, Y) = R$

$\vdash \tau, \text{Cpos}(x+i, y+1, R)$

- $\text{odd}(x) \wedge$

- $\text{even}(i) \wedge$

$\text{odd}(x), \text{even}(i) \Rightarrow \text{odd}(x+i)$.

$\vdash \tau, \text{colorYB}(x, n, x+i) = B$.

すなはち, $\text{Cpos}(x+i, y, Y)$.

$\text{odd}(x+i), \text{odd}(i) \Rightarrow \text{even}(x+i+1)$

$\vdash \tau, \text{colorYB}(x, n, x+i+1) = Y$

C-mix $f' \quad c' = \text{mix}(B, Y) = R$

$\vdash \tau, \text{Cpos}(x+i, y+1, R)$

- $\text{odd}(i) \wedge$

$\text{odd}(x), \text{odd}(i) \Rightarrow \text{even}(x+i)$.

$\vdash \tau, \text{colorYB}(x, n, x+i) = Y$.

すなはち, $\text{Cpos}(x+i, y, B)$.

$\text{even}(x+i), \text{odd}(i) \Rightarrow \text{odd}(x+i+1)$

$\vdash \tau, \text{colorYB}(x, n, x+i+1) = B$

C-mix $f' \quad c' = \text{mix}(B, Y) = R$

$\vdash \tau, \text{Cpos}(x+i, y+1, R) //$

All Red (証明済み)

$\forall x, y, n$

$$\left(\forall i (0 \leq i \leq n \rightarrow C_{\text{pos}}(x+i, y, R)) \right) \rightarrow C_{\text{pos}}(x, y+n, R).$$

Even B

$\forall x, y, n \in \mathbb{N}$

$$\left(\begin{array}{l} \text{even}(n) \\ \forall i (0 \leq i \leq n \rightarrow C_{\text{pos}}(x+i, y, \text{color}^{\text{YB}}(x, n, x+i))) \end{array} \right) \rightarrow C_{\text{pos}}(x, (y+n), R).$$

prf. x, y をとる。

$\text{even}(n)$ をとる $n \in \mathbb{N}$.

$$\text{Even A} \models' \forall i (0 \leq i \leq n-1 \rightarrow C_{\text{pos}}(x+i, y+1, R))$$

$$\text{All Red} \models' C_{\text{pos}}(x, (y+1)+(n-1), R).$$

$$\models \top, C_{\text{pos}}(x, y+n, R).$$

false Color (証明済み)

$\forall x, y, \forall c_0, c_1$

$$(c_0 \neq c_1 \wedge C_{\text{pos}}(x, y, c_0) \wedge C_{\text{pos}}(x, y, c_1)) \rightarrow \text{false}.$$

EvenC

$\forall x, y, n \in \mathbb{N}$

$$\text{even}(n) \rightarrow \exists \left(\begin{array}{l} \forall C: \text{color}, \forall f: \text{nat} \rightarrow \text{Color} \\ \text{Triangle}(x, y, n, f(x), f(x+n), C) \end{array} \right)$$

prf. x, y をとる.

even(n) となる n をとる.

$\left(\begin{array}{l} \forall C: \text{color}, \forall f: \text{nat} \rightarrow \text{Color} \\ \text{Triangle}(x, y, n, f(x), f(x+n), C) \end{array} \right)$ を仮定して矛盾を示す.

$\forall i (0 \leq i \leq n \rightarrow \text{Cpos}(x+i, y, \text{colorYB}(x+i)))$ を示す.

$0 \leq i \leq n$ となる i をとり, C-paint が成立する.

○ EvenB より $\text{Cpos}(x, y+n, R)$.

○ $0 \leq o \leq n$ より $\text{Cpos}(x+o, y, \text{colorYB}(x, n, x))$.

$0 \leq n \leq n$ より $\text{Cpos}(x+n, y, \text{colorYB}(x, n, x+n))$.

C-exists より $\text{Cpos}(x, y+n, c)$ となる $c \in \text{Color}$ が存在する.

$c = c'$ とする $\text{Cpos}(x, y+n, c')$.

仮定 より $\text{Triangle}(x, y, n, \text{colorYB}(x, n, x), \text{colorYB}(x, n, x+n), c')$

よって, $c' = \text{mix}(\text{colorYB}(x, n, x), \text{colorYB}(x, n, x+n))$

さらに, LemYB5 より $\text{colorYB}(x, n, x) = \text{colorYB}(x, n, x+n)$ より

$c' = \text{mix}(\text{colorYB}(x, n, x), \text{colorYB}(x, n, x))$

- even(x) のとき LemYB1 より $\text{colorYB}(x, n, x) = Y$.

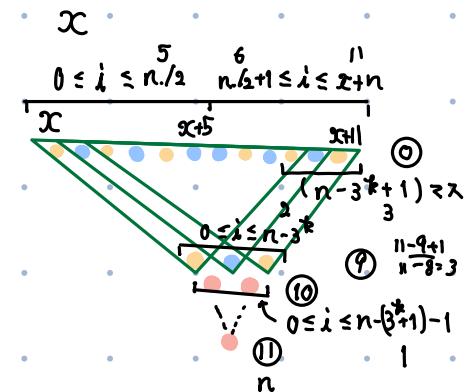
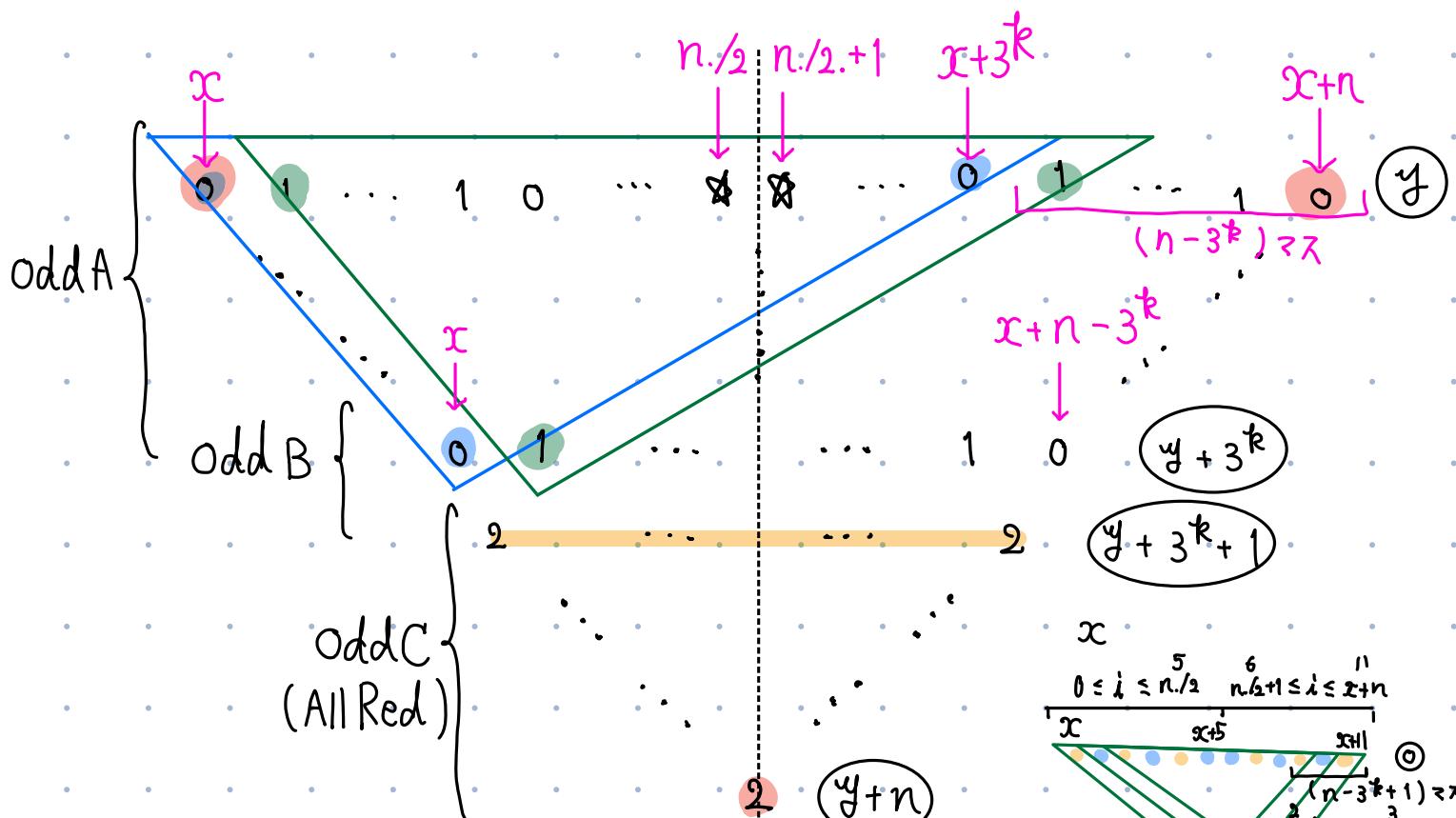
よって $c' = Y$ となるので, falsecolor より 矛盾

- odd(x) のとき LemYB3 より $\text{colorYB}(x, n, x) = B$.

よって $c' = B$ となるので, falsecolor より 矛盾 //

〈色の塗り方〉

$$\text{colorYBYB}(x, n, z) \stackrel{\text{def}}{=} \begin{cases} Y & \left(x \leq z \leq x+n/2 \text{ } \& \& \text{eqn}(z \bmod 2, 0) == \text{true} \right) \\ & \left(x+n/2+1 \leq z \leq x+n \text{ } \& \& \text{eqn}(z \bmod 2, 1) == \text{true} \right) \\ B & \left(x \leq z \leq x+n/2 \text{ } \& \& \text{eqn}(z \bmod 2, 1) == \text{true} \right) \\ & \left(x+n/2+1 \leq z \leq x+n \text{ } \& \& \text{eqn}(z \bmod 2, 0) == \text{true} \right) \\ Y & (\text{otherwise}) \end{cases}$$



Lem YBYB1

$\forall x, n, \text{even}(x) \rightarrow$

$\forall i ((0 \leq i \leq n/2 \wedge \text{even}(i)) \rightarrow \text{colorYBYB}(x, n, x+i) = Y).$

prf. $\text{even}(x)$ となる x をとる。

$0 \leq i \leq n/2$ かつ $\text{even}(i)$ となる n, i をとる。

$0 \leq i \leq n/2$ かつ $x \leq x+i \leq x+n/2$.

$\text{even}(x), \text{even}(i)$ かつ $\text{even}(x+i)$.

よって, $\text{colorYBYB}(x, n, x+i) = Y$. //

Lem YBYB2

$\forall x, n, \text{even}(x) \rightarrow$

$\forall i ((n/2+1 \leq i \leq n \wedge \text{odd}(i)) \rightarrow \text{colorYBYB}(x, n, x+i) = Y).$

prf. $\text{even}(x)$ となる x をとる。

$n/2+1 \leq i \leq n$ かつ $\text{odd}(i)$ となる n, i をとる。

$n/2+1 \leq i \leq n$ かつ $x+n/2+1 \leq x+i \leq x+n$.

$\text{even}(x), \text{odd}(i)$ かつ $\text{odd}(x+i)$.

よって, $\text{colorYBYB}(x, n, x+i) = Y$. //

Lem YBYB 3

$\forall x, n, \text{even}(x) \rightarrow$

$\forall i ((0 \leq i \leq n/2 \wedge \text{odd}(i)) \rightarrow \text{colorYBYB}(x, n, x+i) = B).$

prf. $\text{even}(x)$ となる x をとる。

$0 \leq i \leq n/2 \Leftrightarrow \text{odd}(i)$ となる n, i をとる。

$0 \leq i \leq n/2 \Leftrightarrow x \leq x+i \leq x+n/2.$

$\text{even}(x), \text{odd}(i) \Rightarrow \text{odd}(x+i).$

$\therefore \text{colorYBYB}(x, n, x+i) = B. //$

Lem YBYB 4

$\forall x, n, \text{even}(x) \rightarrow$

$\forall i ((n/2+1 \leq i \leq n \wedge \text{even}(i)) \rightarrow \text{colorYBYB}(x, n, x+i) = B).$

prf. $\text{even}(x)$ となる x をとる。

$n/2+1 \leq i \leq n \Leftrightarrow \text{even}(i)$ となる n, i をとる。

$n/2+1 \leq i \leq n \Leftrightarrow x+n/2+1 \leq x+i \leq x+n.$

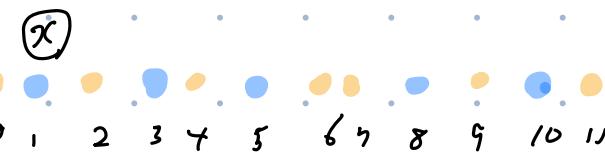
$\text{even}(x), \text{even}(i) \Rightarrow \text{even}(x+i).$

$\therefore \text{colorYBYB}(x, n, x+i) = B. //$

Lem YBYB5

$\forall x, n, \text{odd}(x) \rightarrow$

$\forall i ((0 \leq i \leq n/2 \wedge \text{even}(i)) \rightarrow \text{colorYBYB}(x, n, x+i) = B).$



Lem YBYB6

$\forall x, n, \text{odd}(x) \rightarrow$

$\forall i ((n/2 + 1 \leq i \leq n \wedge \text{odd}(i)) \rightarrow \text{colorYBYB}(x, n, x+i) = B).$

Lem YBYB7

$\forall x, n, \text{odd}(x) \rightarrow$

$\forall i ((0 \leq i \leq n/2 \wedge \text{odd}(i)) \rightarrow \text{colorYBYB}(x, n, x+i) = Y).$

Lem YBYB8

$\forall x, n, \text{odd}(x) \rightarrow$

$\forall i ((n/2 + 1 \leq i \leq n \wedge \text{even}(i)) \rightarrow \text{colorYBYB}(x, n, x+i) = Y).$

OddA

$\forall x, y, n, k,$

$$\left(\begin{array}{l} \bullet 3^k \leq n \leq 3^k * 2 \wedge \text{odd}(n) \rightarrow \\ \bullet \forall x', y', c_0, c_1, c_2. \text{Triangle}(x', y', c_0, c_1, c_2) \rightarrow \\ \bullet \forall i (0 \leq i \leq n \rightarrow \text{Cpos}(x+i, y, \text{colorYBYB}(x, n, x+i))) \end{array} \right) \rightarrow$$
$$\forall i (0 \leq i \leq n - 3^k \rightarrow \text{Cpos}(x+i, y+3^k, \text{colorYB}(x, n - 3^k, x+i)))$$

prf.

Odd B

$\forall x, y, n, k,$

$$\left(\begin{array}{l} \circ 3^k \leq n \leq 3^k * 2 \wedge \text{odd}(n) \rightarrow \\ \circ \forall x', y', c_0, c_1, c_2. \text{Triangle}(x', y', c_0, c_1, c_2) \rightarrow \\ \circ \forall i (0 \leq i \leq n \rightarrow \text{Cpos}(x+i, y, \text{colorYB}(x, n, x+i))) \end{array} \right) \rightarrow$$

$$\forall i (0 \leq i \leq (n - 3^k) - 1 \rightarrow \text{Cpos}(x+i, y + (3^k + 1), R))$$

prf. $x, y, n, k \in \mathbb{Z}$.

Odd A $\models' \forall i (0 \leq i \leq n - 3^k \rightarrow \text{Cpos}(x+i, y+3^k, \text{colorYB}(x, n-3^k, x+i)))$

Even A $\models' \forall i (0 \leq i \leq (n - 3^k) - 1 \rightarrow \text{Cpos}(x+i, (y+3^k)+1, R)) //$

All Red (証明済み)

$\forall x, y, n$

$$(\forall i (0 \leq i \leq n \rightarrow \text{Cpos}(x+i, y, R)) \rightarrow \text{Cpos}(x, y+n, R)).$$

Odd C

$\forall x, y, n, k,$

$$\left(\begin{array}{l} \circ 3^k \leq n \leq 3^k * 2 \wedge \text{odd}(n) \rightarrow \\ \circ \forall x', y', c_0, c_1, c_2. \text{Triangle}(x', y', 3^k, c_0, c_1, c_2) \rightarrow \\ \circ \forall i (0 \leq i \leq n \rightarrow \text{Cpos}(x+i, y, \text{colorYBYB}(x, n, x+i))) \end{array} \right) \rightarrow \text{Cpos}(x, y+n, R)$$

prf. x, y, n, k をとる。

$$\text{OddB} \models \forall i (0 \leq i \leq (n-3^k)-1 \rightarrow \text{Cpos}(x+i, y+(3^k+1), R))$$

$$\text{AllRed} \models \text{Cpos}(x, (y+3^k+1)+((n-3^k)-1), R)$$

$$d, \mathcal{T}, \text{Cpos}(x, y+n, R) //$$

false Color (証明済み)

$\forall x, y, \forall c_0, c_1$

$$(\text{c}_0 \neq \text{c}_1 \wedge \text{Cpos}(x, y, c_0) \wedge \text{Cpos}(x, y, c_1)) \rightarrow \text{false}.$$

Bottom_color_of_triangle (証明済み)

$\forall x, y, n, \forall c_0, c_1$

$\forall c'_0, c'_1, c'_2. \text{Triangle}(x, y, n, c'_0, c'_1, c'_2) \rightarrow$

$\text{Cpos}(x, y, c_0) \rightarrow \text{Cpos}(x+n, y, c_1) \rightarrow \text{Cpos}(x, y+n, \text{mix}(c_0, c_1)).$

OddD

$\forall x, y, n, R,$

$\left(\begin{array}{l} \text{odd}(n) \\ 3^k \leq n \leq 3^k * 2 \end{array} \right) \rightarrow \exists \left(\begin{array}{l} \forall C: \text{color}, \forall f: \text{nat} \rightarrow \text{Color} \\ \text{Triangle}(x, y, n, f(x), f(x+n), C) \end{array} \right)$

prf. x, y をとる。

$\text{odd}(n), 3^k \leq n \leq 3^k * 2$ となる n, R をとする。

$\left(\begin{array}{l} \forall C: \text{color}, \forall f: \text{nat} \rightarrow \text{Color} \\ \text{Triangle}(x, y, n, f(x), f(x+n), C) \end{array} \right)$ を仮定して 矛盾を示す。

十分性は示されていないので $\forall x', y', c_0, c_1, c_2. \text{Triangle}(x', y', 3^k, c_0, c_1, c_2)$

$C\text{-paint } f \upharpoonright i (0 \leq i \leq n \rightarrow \text{Cpos}(x+i, y, \text{colorYBYB}(x, n, x+i)))$

- Odd C 且 $\text{Cpos}(x, y+n, R).$

- $0 \leq o \leq n$ 且 $\text{Cpos}(x+o, y, \text{colorYBYB}(x, n, x+o)).$

- $0 \leq n \leq n$ 且 $\text{Cpos}(x+n, y, \text{colorYBYB}(x, n, x+n)).$

$C\text{-exist}\& \text{ } f \upharpoonright \text{Cpos}(x, y+n, c)$ となる c が存在する。 $(c = c' \text{ とする})$

仮定 $f \upharpoonright \text{Triangle}(x, y, n, \text{colorYBYB}(x, n, x+o), \text{colorYBYB}(x, n, x+n), c')$

$\text{Bottom_color_of_triangle } f \upharpoonright \text{Cpos}(x, y+n, \text{mix}(\text{colorYBYB}(x, n, x+o), c'))).$

○ even(x) のとき

$0 \leq o \leq n \Leftrightarrow \text{even}(o)$, LemYBYB1 より $\text{colorYBYB}(x, n, x+o) = Y$.

$n/2+1 \leq n \leq n \Leftrightarrow \text{odd}(n)$, LemYBYB2 より $\text{colorYBYB}(x, n, x+n) = Y$.

したがって, $\text{mix}(\bullet) = \text{mix}(Y, Y) = Y$ となり falsecolor より矛盾.

○ odd(x) のとき

$0 \leq o \leq n \Leftrightarrow \text{even}(o)$, LemYBYB5 より $\text{colorYBYB}(x, n, x+o) = B$.

$n/2+1 \leq n \leq n \Leftrightarrow \text{odd}(n)$, LemYBYB6 より $\text{colorYBYB}(x, n, x+n) = B$.

したがって, $\text{mix}(\bullet) = \text{mix}(B, B) = B$ となり falsecolor より矛盾.

//