

Lecture 3: September 12

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3.1 Variance bound for binary classification with constant noise and threshold functions

Let $\mathcal{X} = [0, 1]$, $\mathcal{Y} = \{0, 1\}$, $\mathcal{Z} = \mathcal{X} \times \mathcal{Y}$ and for $w \in [0, 1]$ let $\ell_w : \mathcal{Z} \rightarrow \{0, 1\}$ be defined by

$$\ell_w(x, y) = \mathbb{I}\{\mathbb{I}\{x \geq w\} \neq y\}.$$

Let

$$o_w(x, y) = \mathbb{I}\{x \notin [w - \varepsilon, w]\}.$$

Let P be a distribution over \mathcal{Z} . As usual, Pf stands for $\int f dP$. For some $w^* \in [0, 1]$, we let P_{w^*} be a distribution over \mathcal{Z} such that if $(X, Y) \sim P_{w^*}$ then X is uniformly distributed over $[0, 1]$ and $\mathbb{P}(Y = 1|X = x) = 1 - p$ if $x \geq w^*$ and $\mathbb{P}(Y = 1|X = x) = p$ otherwise.

We consider functions of the form $g_{w,w^*} = \ell_w o_w - \ell_{w^*}$. Note that from this choice, it follows that

$$P_{w^*} g_{w,w^*} \geq -p\varepsilon.$$

Indeed, writing $P = P_{w^*}$ to minimize clutter, on the one hand we have

$$P\ell_{w^*} = p,$$

and also

$$\min_{w \in [0, 1]} P\ell_w o_w \geq \min_{w' \in [0, 1]} \min_{w \in [0, 1]} P\ell_w o_{w'} = \min_{w' \in [0, 1]} P\ell_{w^*} o_{w'} = P\ell_{w^*} o_\varepsilon = p - p\varepsilon = p(1 - \varepsilon).$$

Hence,

$$Pg_{w,w^*} = P\ell_w o - P\ell_{w^*} \geq p(1 - \varepsilon) - p = -p\varepsilon.$$

Our next goal is to find the “smallest possible” $c_0, c_1 \geq 0$ such that for any w, w^*, ε, p ,

$$P_{w^*} g_{w,w^*}^2 \leq c_0^2 + c_1 P_{w^*} g_{w,w^*} \tag{3.1}$$

holds. In fact, we want to minimize first c_0 and keep c_1 bounded.

To minimize clutter, we let $P = P_{w^*}$, $o = o_w$, but these dependencies should still be kept in mind. First, algebra gives

$$\begin{aligned} Pg_{w,w^*}^2 &= P\ell_w o - 2P\ell_w \ell_{w^*} o + P\ell_{w^*} \\ &= P\ell_w o - 2P\ell_w \ell_{w^*} o + p. \end{aligned}$$

Furthermore,

$$Pg = P\ell_w o - P\ell_{w^*} = P\ell_w o - p.$$

Hence, Eq. (3.1) is equivalent to

$$P\ell_w o - 2P\ell_w \ell_{w^*} o + p \leq c_0^2 + c_1(P\ell_w o - p),$$

which is equivalent to

$$c_1 p + p \leq c_0^2 + (c_1 - 1)P\ell_w o + 2P\ell_w \ell_{w^*} o. \quad (3.2)$$

Now, we consider three cases based on the location of w^* relative to $w - \varepsilon$ and w and we will in each case lower bound the value of

$$A = (c_1 - 1)P\ell_w o + 2P\ell_w \ell_{w^*} o,$$

while enforcing that

$$c_1 \geq 1. \quad (3.3)$$

Case 1: $w^* \leq w - \varepsilon$. In this case we have

$$\begin{aligned} P\ell_w o &= (w - w^*)(1 - 2p) + p - \varepsilon(1 - p), \\ P\ell_w \ell_{w^*} o &= w^*(1 - p) - wp + p. \end{aligned}$$

Then,

$$A = (c_1 - 1) \{ (w - w^*)(1 - 2p) + p - \varepsilon(1 - p) \} + 2 \{ w^*(1 - p) - wp + p \}.$$

The coefficient of w^* in the above expression is

$$\alpha = (c_1 - 1)(2p - 1) + 2(1 - p).$$

If $\alpha \leq 0$, the value of A is minimized in w^* by choosing $w^* = w - \varepsilon$ and the minimum value is equal to

$$\begin{aligned} A' &= (c_1 - 1) \{ \varepsilon(1 - 2p) + p - \varepsilon(1 - p) \} + 2 \{ (w - \varepsilon)(1 - p) - wp + p \} \\ &= (c_1 - 1)p(1 - \varepsilon) + 2 \{ (w(1 - 2p) - \varepsilon(1 - p) + p \} \\ &\geq (c_1 - 1)p(1 - \varepsilon) + 2 \{ p - \varepsilon(1 - p) \} =: A'', \end{aligned}$$

where the last inequality follows since $2w(1 - 2p) \geq 0$. Now, Eq. (3.2) will follow if

$$c_1 p(1 - \varepsilon) + p + c_1 p \varepsilon = c_1 p + p \leq c_0^2 + A''.$$

Plugging in the definition of A'' , we get that the latter inequality is equivalent to

$$\begin{aligned} \cancel{c_1 p(1 - \varepsilon)} + p + c_1 \varepsilon &\leq c_0^2 + \cancel{(c_1 - 1)p(1 - \varepsilon)} + 2 \{ p - \varepsilon(1 - p) \} \\ &= c_0^2 - \cancel{p} + p \varepsilon + \cancel{2p} - 2\varepsilon(1 - p) \\ &= c_0^2 + 3\varepsilon p + p - 2\varepsilon \end{aligned}$$

Reordering gives

$$c_0^2 \geq \varepsilon(c_1 + 2 - 3p). \quad (3.4)$$

Since the $p < 1/2$ by assumption and, as we shall see $c_1 \geq 0$, this constraint can be satisfied.

Finally, simple algebra shows that $\alpha \leq 0$ is equivalent to

$$c_1 \geq 1 + \frac{1-p}{\frac{1}{2}-p}, \quad (3.5)$$

where we used that $p < 1/2$.

Case 2: $w - \varepsilon < w^* \leq w$. In this case we have

$$\begin{aligned} P\ell_w o &= p(1 - \varepsilon), \\ P\ell_w \ell_{w^*} o &= p(1 - \varepsilon). \end{aligned}$$

Then,

$$A = (c_1 - 1)p(1 - \varepsilon) + 2p(1 - \varepsilon) = c_1 p(1 - \varepsilon) + p(1 - \varepsilon). \quad (3.6)$$

Now, Eq. (3.2) will follow if $c_1(1 - \varepsilon)p + p(1 - \varepsilon) + (c_1 + 1)p\varepsilon \leq c_0^2 + A$. Plugging in the definition of A , we get that the latter inequality is equivalent to

$$\cancel{c_1(1-\varepsilon)p} + \cancel{(1-\varepsilon)p} + (c_1 + 1)p\varepsilon \leq c_0^2 + \cancel{c_1 p(1-\varepsilon)} + \cancel{p(1-\varepsilon)},$$

which gives

$$c_0^2 \geq (c_1 + 1)\varepsilon p. \quad (3.7)$$

Case 3: $w < w^*$. In this case we have

$$\begin{aligned} P\ell_w o &= (w - w^*)(2p - 1) + p(1 - \varepsilon), \\ P\ell_w \ell_{w^*} o &= (w - w^*)p + p(1 - \varepsilon). \end{aligned}$$

Then,

$$A = (c_1 - 1) \{ (w - w^*)(2p - 1) + p(1 - \varepsilon) \} + 2 \{ (w - w^*)p + p(1 - \varepsilon) \}.$$

Note that here $w - w^* \leq 0$. The coefficient of $w - w^*$ in this expression is

$$\alpha = (c_1 - 1)(2p - 1) + 2p.$$

When $\alpha \leq 0$, the value of A is minimized by choosing the largest possible value for $w - w^*$, which is zero. In this case the value of A is

$$A' = (c_1 - 1)p(1 - \varepsilon) + 2p(1 - \varepsilon) = c_1 p(1 - \varepsilon) + p(1 - \varepsilon),$$

which is the value we got in the previous case (cf. Eq. (3.6)). As such, as long as Eq. (3.7) holds, Eq. (3.1) will be satisfied.

To guarantee that $\alpha \leq 0$, we need to have

$$(c_1 - 1)(2p - 1) + 2p \leq 0.$$

Thanks to $2p - 1 < 0$, this holds if and only if

$$c_1 \geq 1 + \frac{p}{\frac{1}{2}-p}. \quad (3.8)$$

Thus, we derived the following:

Proposition 3.1. *Let $p < 1/2$ and let*

$$c_1 = 1 + \frac{1-p}{\frac{1}{2}-p}, \quad c_0^2 = \varepsilon(c_1 + 2).$$

Then, for any w, w^ it holds that*

$$P_{w^*} g_{w,w^*}^2 \leq c_0^2 + c_1 P_{w^*} g_{w,w^*}.$$

Proof. Choosing c_1 as above satisfies both Eqs. (3.5) and (3.8). Furthermore, because $\varepsilon(c_1 + 2) \geq \varepsilon(c_1 + 2 - 3p)$ and $\varepsilon(c_1 + 2) \geq \varepsilon p(c_1 + 1)$, choosing $c_0^2 = \varepsilon(c_1 + 2)$ will satisfy both Eqs. (3.4) and (3.7). \square