EOPSY Lab 4

Memory Management

Performing Simulation

In order to perform the task we must run the simulation that is constructed with the use of Java language. Unpack the file called task 4 and in directory ftp execute ./make setup. This script creates directory work where you have to type commands: make compile and then make run to see the simulation window.

Mapping virtual pages to physical pages

In order to configure our mapping we have to edit file called *memory.conf*. Our task is to map any 8 pages of physical memory to the first 8 pages of virtual memory.

```
| memset | virt | page # | physical | page # | R | (read | from) | M | (modified) | inMemTime | (ns) | lastTouchTime | (ns) | memset | 0 7 0 0 0 0 0 | memset | 2 5 0 0 0 0 0 | memset | 2 5 0 0 0 0 0 | memset | 3 4 0 0 0 0 0 | memset | 4 3 0 0 0 0 0 | memset | 5 2 0 0 0 0 0 | memset | 5 2 0 0 0 0 0 | memset | 7 0 0 0 0 0 0 | memset | 7 0 0 0 0 0 0 | memset | 3 4 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 0 | memset | 3 0 0 0 0 | mem
```

As presented above we assign our virtual pages to the physical pages according to the scheme below:

Affair al Davis	District Design
Virtual Page	Physical Page
0	7
1	6
2	5
3	4
4	3
5	2
6	1
7	0

Reading virtual pages

Next we have to configure file <u>commands</u> which will specify read operation for every 64 of virtual pages. As written in *memory.conf* file we've established that pagesize is 16384 addresses. Thus to ensure that address form each distinct virtual page. I've decided to use 1st address form ith page.

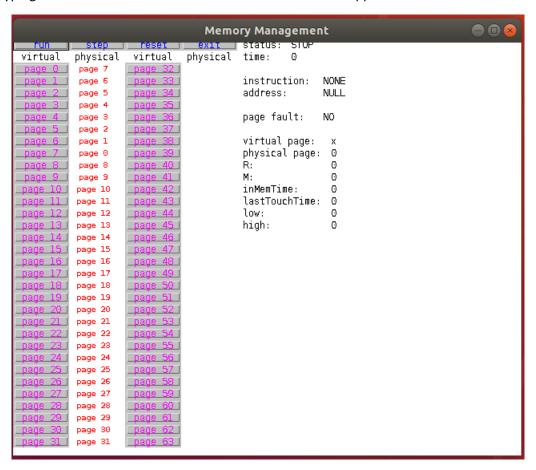
	commands
READ 0	
READ 16384	
READ 32768	
READ 49152	
READ 65536	
READ 81920	
READ 98304	
READ 114688	
READ 131072	
READ 147456	
READ 163840	
READ 180224	
READ 196608	
READ 212992	
READ 229376	
READ 245760	
READ 262144	
READ 278528	
READ 294912	
READ 311296	
READ 327680	
READ 344064	
READ 360448	
READ 376832	
READ 393216	
READ 409600	
READ 425984	
READ 442368	
READ 458752	
READ 475136	
READ 491520	
READ 507904 READ 524288	
READ 544268 READ 540672	
READ 557056	
READ 573440	
READ 589824	
READ 606208	
READ 622592	
READ 638976	
READ 655360	
READ 671744	
READ 688128	
READ 704512	
READ 720896	
READ 737280	
READ 753664	
READ 770048	
READ 786432	
READ 802816	
READ 819200	
READ 835584	
READ 851968	

READ	868352			
READ	884736			
READ	901120			
READ	917504			
READ	933888			
READ	950272			
READ	966656			
READ	983040			
READ	999424			
READ	1015808			
READ	1032192			

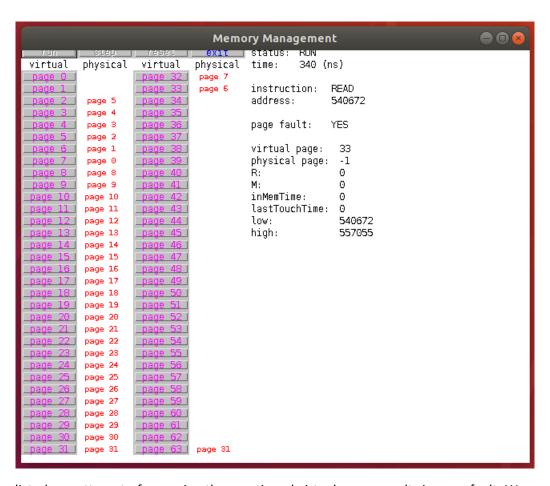
Page faults predictions

Now we will run our simulation and try to predict which READ operation will trigger page fault. Page fault is a situation where virtual page is not mapped to any physical page and we try to read address from this virtual page. In such situation proper physical page will be mapped to virtual page that informed about page fault.

After typing make run command we can see simulation window appear:



As visible on the screenshot first 8 virtual pages have been have been assigned physical pages. As for the rest mapping have been performed by default. Also we can observe that for pages 32-63 there is no physical page assigned. Therefore I predict that any attempt of accessing this pages will result in a page fault. To verify this assumption I will start this simulation and see if my assumption is correct.



As predicted any attempt of accessing the mentioned virtual pages results in page fault. We can also notice that physical page assigned to those values is equal to -1.

Another indication that this assumption is correct may be found in *tracefile* file. Let's see it's contents.

```
tracefile
   READ 0 ... okay
1)
   READ 16384 ... okay
   READ 32768 ... okay
   READ 49152 ... okay
5)
   READ 65536 ... okay
   READ 81920 ... okay
7)
   READ 98304 ... okay
   READ 114688 ... okay
8)
9)
   READ 131072 ... okay
10) READ 147456 ... okay
11) READ 163840 ... okay
12) READ 180224 ... okay
13) READ 196608 ... okay
14) READ 212992 ... okay
15) READ 229376 ... okay
16) READ 245760 ... okay
17) READ 262144 ... okay
18) READ 278528 ... okay
19) READ 294912 ... okay
20) READ 311296 ... okay
21) READ 327680 ... okay
22) READ 344064 ... okay
23) READ 360448 ... okay
24) READ 376832 ... okay
25) READ 393216 ... okay
26) READ 409600 ... okay
```

```
27) READ 425984 ... okay
28) READ 442368 ... okay
29) READ 458752 ... okay
30) READ 475136 ... okay
31) READ 491520 ... okay
32) READ 507904 ... okay
33) READ 524288 ... page fault
34) READ 540672 ... page fault
35) READ 557056 ... page fault
36) READ 573440 ... page fault
37) READ 589824 ... page fault
38) READ 606208 ... page fault
39) READ 622592 ... page fault
40) READ 638976 ... page fault
41) READ 655360 ... page fault
42) READ 671744 ... page fault
43) READ 688128 ... page fault
44) READ 704512 ... page fault
45) READ 720896 ... page fault
46) READ 737280 ... page fault
47) READ 753664 ... page fault
48) READ 770048 ... page fault
49) READ 786432 ... page fault
50) READ 802816 ... page fault
51) READ 819200 ... page fault
52) READ 835584 ... page fault
53) READ 851968 ... page fault
54) READ 868352 ... page fault
55) READ 884736 ... page fault
56) READ 901120 ... page fault
57) READ 917504 ... page fault
58) READ 933888 ... page fault
59) READ 950272 ... page fault
60) READ 966656 ... page fault
61) READ 983040 ... page fault
62) READ 999424 ... page fault
63) READ 1015808 ... page fault
64) READ 1032192 ... page fault
```

As we can see for last 32 pages (32-63) there is a page fault error. This concludes our prediction.

Page Replacement Algorithm

Our page replacement algorithm can be found in ./tast4/work/PageFault.java file. To be exact the specific function is called:

```
public static void replacePage(Vector mem, int virtPageNum, int replacePageNum, ControlPanel controlPanel)
```

In the documentation we read that implemented algorithm is FIFO(First-In First-Out).

Description: Operating system holds a list of all pages currently in memory, with the most recent arrival at the tail and the least recent arrival at the head. On a page fault, the page at the head is removed and the new page is added to the tail of the list.

If we take a look at the end of the simulation we can observe that the first mapping removed was the first one that occurred.

