

Block Cipher Operations

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November 2nd, 2021
Intro to Cryptology
Hands On Exercise 9

1)

CBC Mode

$$\begin{aligned} P_i &= C_{i-1} \text{ xor } D_k(C_i) & C_i &= E_k(P_i \text{ xor } C_{i-1}) \\ &= C_{i-1} \text{ xor } D_k(E_k(P_i \text{ xor } C_{i-1})) \\ &= C_{i-1} \text{ xor } P_i \text{ xor } C_{i-1} \\ &= P_i \end{aligned}$$

CFB Mode

$$\begin{aligned} P_i &= C_i \text{ xor } L_s(E_k(X_i)) & C_i &= P_i \text{ xor } L_s(E_k(X_i)) \\ &= P_i \text{ xor } L_s(E_k(X_i)) \text{ xor } L_s(E_k(X_i)) \\ &= P_i \end{aligned}$$

2)

a)

CFB-32

$$P_i = C_i \text{ xor } L_{32}(E_k(X_i)) \qquad X_{i+1} = R_{32}(X_i) \parallel C_i$$

b)

$$\begin{aligned} C &= C_1 \parallel C_2 \parallel C_3 \parallel C_4 \dots C_n \\ P_1' &= C_1' \text{ xor } L_{32}(E_k(X_1)) & X_2' &= R_{32}(X_1) \parallel C_1' \\ P_2' &= C_2 \text{ xor } L_{32}(E_k(X_2')) & X_3' &= C_1' \parallel C_2 \\ P_3' &= C_3 \text{ xor } L_{32}(E_k(X_3')) & X_4' &= C_2 \parallel C_3 \\ P_4' &= C_4 \text{ xor } L_{32}(E_k(X_4)) & X_5' &= C_3 \parallel C_4 \end{aligned}$$

The only affected blocks are P_1' , P_2' , and P_3' , therefore this concludes that $P_i' = P_i$ for all $i \leq 4$.

3)

$$\begin{aligned} C &= C_1 \parallel C_2 \parallel C_3 \parallel C_4 \dots C_n \\ P_1' &= C_0 \text{ xor } D_k(C_1) \\ P_2' &= C_1 \text{ xor } D_k(C_2) \\ P_3' &= C_2 \text{ xor } D_k(C_3) \\ P_4' &= C_3 \text{ xor } D_k(C_4) \end{aligned}$$

Setting the affected ciphertext block to C_1 , we can see in our example that P_1 and P_2 are the only blocks affected in the plaintext.

Therefore, the two blocks which are affected are block i and $i + 1$.