

Lecture 6: Assembly Programs

- Today's topics:
 - Control instructions
 - Procedures
 - Examples

Procedures

- Each procedure (function, subroutine) maintains a scratchpad of register values – when another procedure is called (the callee), the new procedure takes over the scratchpad – values may have to be saved so we can safely return to the caller

- parameters (arguments) are placed where the callee can see them
- control is transferred to the callee
- acquire storage resources for callee
- execute the procedure
- place result value where caller can access it
- return control to caller

jal - jal

jal

\$sp, \$fp

\$v0, \$v1

jr \$ra

Jump-and-Link

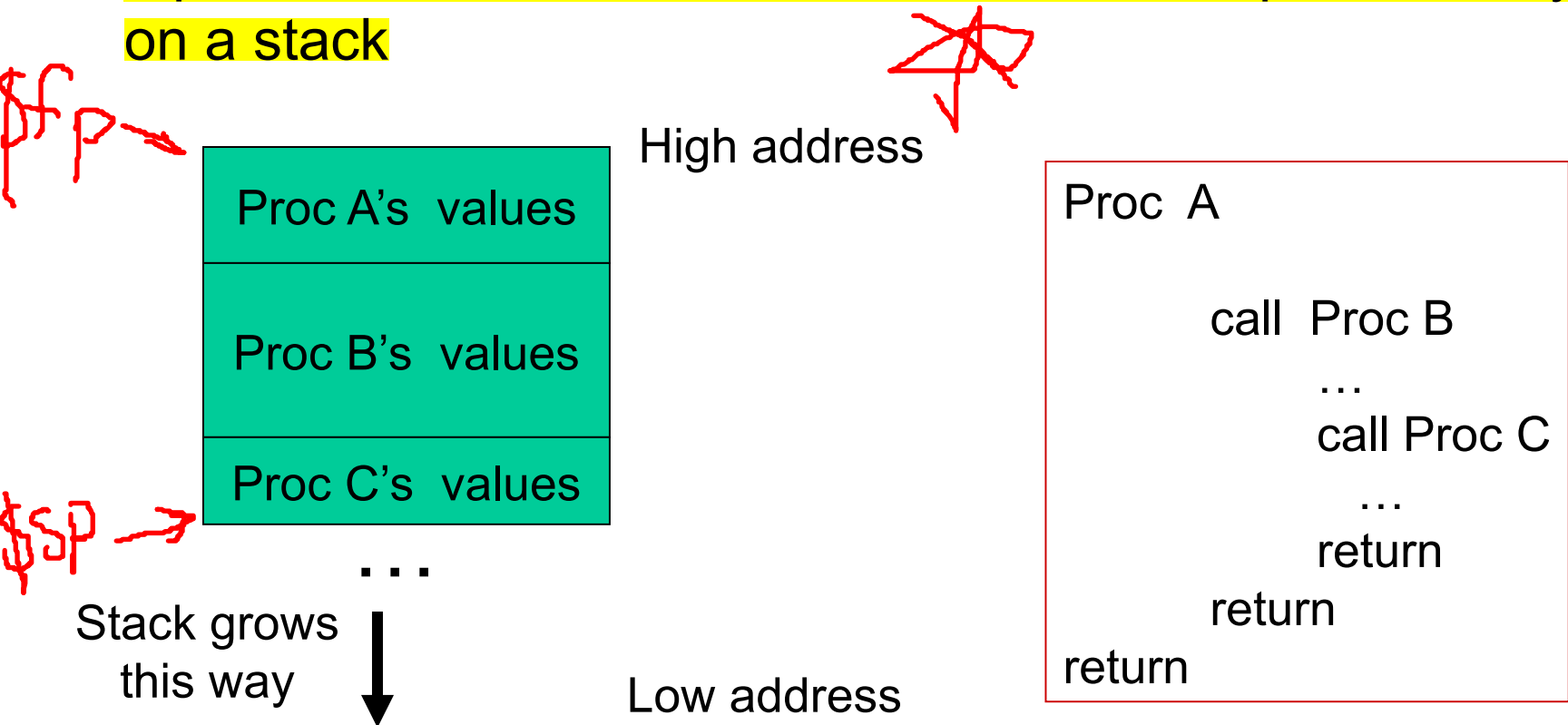
- A special register (storage not part of the register file) maintains the address of the instruction currently being executed – this is the *program counter* (PC)
- The procedure call is executed by invoking the jump-and-link (jal) instruction – the current PC (actually, PC+4) is saved in the register \$ra and we jump to the procedure's address (the PC is accordingly set to this address)

 **return address** `jal NewProcedureAddress`

- Since jal may over-write a relevant value in \$ra, it must be saved somewhere (in memory?) before invoking the jal instruction
- How do we return control back to the caller after completing the callee procedure? ถ้าทำฟังก์ชันซ้อนหลายๆชั้นมันจะ **return** ไม่ได้เพราะ ra ก่อนโดนเขียนทับ เลยต้องเซฟ ra ไว้ที่อื่นก่อน

The Stack

The register scratchpad for a procedure seems volatile – it seems to disappear every time we switch procedures – a procedure's values are therefore backed up in memory on a stack



Storage Management on a Call/Return

- A new procedure must create space for all its variables on the stack
- Before/after executing the jal, the caller/callee must save relevant values in \$s0-\$s7, \$a0-\$a3, \$ra, temps into the stack space
กระดาษทคองไครขงมัน หาก jal ควรเซฟตัวแปรไว้ที่อื่นด้วย
- Arguments are copied into \$a0-\$a3; the jal is executed
- After the callee creates stack space, it updates the value of \$sp
- Once the callee finishes, it copies the return value into \$v0, frees up stack space, and \$sp is incremented
เพิ่มขึ้นเพราะว่าเราต้องกลับไปด้านบน (เราเริ่มจากเลขมากมาน้อย)
- On return, the caller/callee brings in stack values, ra, temps into registers
- The responsibility for copies between stack and registers may fall upon either the caller or the callee

Example 1 (pg. 98)

```
int leaf_example (int g, int h, int i, int j)
{
    int f ;
    f = (g + h) - (i + j);
    return f;
}
```

Notes:

In this example, the callee took care of saving the registers it needs.

The caller took care of saving its \$ra and \$a0-\$a3.

เขียนเก่งจริงไม่ต้องใช้ Stack

Could have avoided using the stack altogether.

leaf_example:

```
addi    $sp, $sp, -12
sw      $t1, 8($sp)
sw      $t0, 4($sp)
sw      $s0, 0($sp)
add     $t0, $a0, $a1
add     $t1, $a2, $a3
sub     $s0, $t0, $t1
add     $v0, $s0, $zero
lw      $s0, 0($sp)
lw      $t0, 4($sp)
lw      $t1, 8($sp)
addi    $sp, $sp, 12
jr      $ra
```

save old registers
ten

Saving Conventions

- Caller saved: Temp registers \$t0-\$t9 (the callee won't bother saving these, so save them if you care), \$ra (it's about to get over-written), \$a0-\$a3 (so you can put in new arguments)
- Callee saved: \$s0-\$s7 (these typically contain “valuable” data)
- Read the Notes on the class webpage on this topic

Example 2 (pg. 101)

blt < slt
beq

```
int fact (int n)
{
    if (n < 1) return (1);
    else return (n * fact(n-1));
}
```

```
fact:
    slti    $t0, $a0, 1
    beq     $t0, $zero, L1
    addi    $v0, $zero, 1
    jr      $ra
```

L1:

```
    addi    $sp, $sp, -8
    sw      $ra, 4($sp)  save ra
    sw      $a0, 0($sp)  save n
    addi    $a0, $a0, -1
    jal     fact
    lw      $a0, 0($sp)
    lw      $ra, 4($sp)
    addi    $sp, $sp, 8
    mul     $v0, $a0, $v0
    jr      $ra
```

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Notes:

The caller saves \$a0 and \$ra
in its stack space.

Temp register \$t0 is never saved.