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TERM-PROJECT COMPILER

- 01: $\text{CODE} \rightarrow \text{VDECL CODE} \mid \text{FDECL CODE} \mid \text{CDECL CODE} \mid \epsilon$
- 02: $\text{VDECL} \rightarrow \text{vtype id semi} \mid \text{vtype ASSIGN semi}$
- 03: $\text{ASSIGN} \rightarrow \text{id assign RHS}$
- 04: $\text{RHS} \rightarrow \text{EXPR} \mid \text{literal} \mid \text{character} \mid \text{boolstr}$
- 05: $\text{EXPR} \rightarrow \text{EXPR addsub EXPR} \mid \text{EXPR multdiv EXPR}$
- 06: $\text{EXPR} \rightarrow \text{lparen EXPR rparen} \mid \text{id} \mid \text{num}$
- 07: $\text{FDECL} \rightarrow \text{vtype id lparen ARG rparen lbrace BLOCK RETURN rbrace}$
- 08: $\text{ARG} \rightarrow \text{vtype id MOREARGS} \mid \epsilon$
- 09: $\text{MOREARGS} \rightarrow \text{comma vtype id MOREARGS} \mid \epsilon$
- 10: $\text{BLOCK} \rightarrow \text{STMT BLOCK} \mid \epsilon$
- 11: $\text{STMT} \rightarrow \text{VDECL} \mid \text{ASSIGN semi}$
- 12: $\text{STMT} \rightarrow \text{if lparen COND rparen lbrace BLOCK rbrace ELSE}$
- 13: $\text{STMT} \rightarrow \text{while lparen COND rparen lbrace BLOCK rbrace}$
- 14: $\text{COND} \rightarrow \text{COND comp COND} \mid \text{boolstr}$
- 15: $\text{ELSE} \rightarrow \text{else lbrace BLOCK rbrace} \mid \epsilon$
- 16: $\text{RETURN} \rightarrow \text{return RHS semi}$
- 17: $\text{CDECL} \rightarrow \text{class id lbrace ODECL rbrace}$
- 18: $\text{ODECL} \rightarrow \text{VDECL ODECL} \mid \text{FDECL ODECL} \mid \epsilon$

1) Discard an ambiguity in the CFG:

In production 01 ($\text{CODE} \rightarrow \text{VDECL CODE} \mid \text{FDECL CODE} \mid \text{CDECL CODE} \mid \epsilon$), the non-terminal CODE can derive a sequence of VDECL, FDECL, and CDECL non-terminals followed by another CODE non-terminal. This allows for multiple derivations.

In production 05 ($\text{EXPR} \rightarrow \text{EXPR addsub EXPR} \mid \text{EXPR multdiv EXPR}$), the non-terminal EXPR on both sides allows for ambiguity in expressions involving addition/subtraction and multiplication/division.

In production 11 ($\text{STMT} \rightarrow \text{VDECL} \mid \text{ASSIGN semi}$), the non-terminal STMT can derive either a VDECL or an ASSIGN statement followed by a semi-colon.

To remove it we can rewrite the production like so:

- 01: $\text{CODE} \rightarrow \text{DECLS}$
- 02: $\text{DECLS} \rightarrow \text{DECL DECLS} \mid \text{epsilon}$
- 03: $\text{DECL} \rightarrow \text{VDECL} \mid \text{FDECL} \mid \text{CDECL}$
- 04: $\text{VDECL} \rightarrow \text{vtypeid semi} \mid \text{vtype ASSIGN semi}$
- 05: $\text{ASSIGN} \rightarrow \text{id assign RHS}$
- 06: $\text{RHS} \rightarrow \text{EXPR} \mid \text{literal} \mid \text{character} \mid \text{boolstr}$

07: `EXPR -> EXPR addsub TERM | TERM`
 08: `TERM -> TERM multdiv FACTOR | FACTOR`
 09: `FACTOR -> lparen EXPR rparen | id | num`
 10: `FDECL -> vtype id lparen ARGS rparen lbrace BLOCK RETURN rbrace`
 11: `ARGS -> vtype id MOREARGS | epsilon`
 12: `MOREARGS -> comma vtype id MOREARGS | epsilon`
 13: `BLOCK -> STMT BLOCK | epsilon`
 14: `STMT -> VDECL | ASSIGN semi | IFSTMT | WHILESTMT`
 15: `IFSTMT -> if lparen COND rparen lbrace BLOCK rbrace ELSE`
 16: `ELSE -> else lbrace BLOCK rbrace | epsilon`
 17: `WHILESTMT -> while lparen COND rparen lbrace BLOCK rbrace`
 18: `COND -> EXPR comp EXPR | boolstr`
 19: `RETURN -> return RHS semi`
 20: `CDECL -> class id lbrace ODECL rbrace`
 21: `ODECL -> VDECL ODECL | FDECL ODECL | epsilon`

2) Construct a SLR parsing table for the non-ambiguous CFG:

You can find in *SLRParserGenerator.pdf* file, the constructed SLR parsing table for the non-ambiguous CFG. To create it, we used the tool specified in the subject : <https://jsmachines.sourceforge.net/machines/slr.html> website.

3) Implement a SLR parsing program for the simplified Java programming language by using the constructed table:

We use python3 for this project. Then, to use our program you just need to execute this command in your Linux terminal:

```
python3 syntax_analyzer.py examples/your_desired_file.java
```

As you will see, we provide in total 3 input files in examples/ folder that you can use freely:

- Adder.java
- OctalDecimal.java
- LeastCommonMultiple.java

The output of our program is in two steps. First, we display the tokens' value and type obtained from the input file and secondly, the parsing tree is displayed.

```

-----TOKENS-----
class  (L1)  ---->  class
Adder  (L1)  ---->  id
{      (L1)  ---->  lbrace
public (L3)  ---->  id
static (L3)  ---->  id
void   (L3)  ---->  vtype
main   (L3)  ---->  id
(      (L3)  ---->  lparen
String (L3)  ---->  vtype
args   (L3)  ---->  id
)      (L3)  ---->  rparen
{      (L3)  ---->  lbrace
int    (L5)  ---->  vtype
first  (L5)  ---->  id
=      (L5)  ---->  assign

```

```

-----PARSE TREE-----
CODE
CDECL
  class
  Adder
  ODECL
    VDECL
      int
      first
      ASSIGN
        =
        EXPR
          TERM
            RHS
              10
    ODECL
      VDECL

```

4) Bonus:

- As you can see above, we can give in input a real java file and not just a list of tokens as suggested in the subject. This feature allows us to create a more complete syntax analyser, just like in the real world.
- Also this pre-parsing step is able to avoid java comment lines (“// Comment”).
- All our code is inline documented. You can read it and understand it step by step.
- Also, when we display an error, the line when it occurs is specified.

Exception: Parsing error (L3): Expected rbrace, found id