Lecture S Relational A

Relational Al

- Operators are applied
 Operands
 - Operands are relations
 - Operators produce resuredrelations

Relational Al

State	Area	Population	
AL	52,419	4,661,900	
CA	163,696	36,756,666	
GA	59,425	9,685,744	
NC	53,819	9,222,414	
TX	268,820	24,326,974	

Relational Al

Five basic relational alge

– Selection σ

Projection

Union

Difference

Selection Ope

- σ_{condition}(R)
 - Selects tuples from relatio specified criteria
 - A condition may include

Selection Ope

citypop

Birmingham Huntsville Atlanta Charlotte Greensboro
Atlanta Charlotte
Charlotte
0.000
Croonahara
Greensboro
Huntsville

Projection Op

- $\Pi_{attribute1,...,attributek}$ (R)
 - Deletes attributes not spe
 - Must also remove dupli maintain the set propert

Projection Op

citypop

City	
Birmingham	
Huntsville	
Atlanta	
Charlotte	
Greensboro	
Huntsville	

Rename Ope

- $\rho_A(B)$
 - Renames relation B to A w relation

Rename Ope

City Stat

Set Operat

- Union, Intersection, an
 - Traditional set operations for relations defined for th types

Set Operat

govtemployees

UID	L
4232	
12408	
31023	W
007	
8938	

UID L

Union Oper

- R(X) U S(X)
 - Results in a relation that c attributes X from either R

Union Oper

presidents U postalworkers =

UID	Last Name	
4232	Roosevelt	
12408	Franklin	

Intersection Op

- $R(X) \cap S(X)$
 - Results in a relation that c attributes X from relations common to both R and S.

Intersection Op

presidents ∩ postalworkers = N

govtemployees ∩ postalworker

Difference Op

- R(X) S(X)
 - Results in a relation that c attributes X that belong to also belong to relation S.

Difference Op

govtemployees - presidents =

UID	
12408	

govtemployees - postalworker



Cartesian Pr

• RXS

Given arity k₁ tuple of relation S, the Ca pairing of all possible tupl tuples of S

Cartesian Pr

R

Α	В
a	4
b	1
С	5

S

D	Е
0	false
1	true

Joins

- Combination of a Cartes
 Selection operation
- Several variations
 - Theta Join

Theta Jo

- R ⋈ condition S
 - Given arity k₁ tuple of relation S, the Ca pairing of all possible tupl
 I f S h h

Theta Jo

	Α	В
D	1	4
R	3	1
	2	5
	D	Е
	ъ	-
S	2	4

Theta Jo

$$R \bowtie_{C < D} S = \sigma_{C < D} (R X S)$$

RXS =

Α	В	С
1	4	2
1	4	2
3	1	3
3	1	3

Equi-Joi

- $R\bowtie_{\text{equality}} S$
 - Given arity k₁ tuple of relation S, the Ca pairing of all possible tupl

- *R* ⋈ *S*
 - Compute R X S
 - For each attribute A_i com
 select all tuples that agree

	Α	В
D	1	4
R	3	1
	2	5
	Α	C
S	3	3

(Step #1) Compute R X S

R.A	R.B	R.C	S.A
1	4	2	3
1	4	2	5
1	4	2	2
1	4	2	1
3	1	3	3

(Step #2) Keep all tuples such th R.A = S.A AND R.C

R.A	R.B	R.C	S.A
1	4	2	1
3	1	3	3
2	5	6	2

Division

```
    R / S "Such That"
    -R / S = { x | ∀ y ∈ S : ∃<</li>
    "For Each" "Ther
```

Division

В	
pants	
shirt	
pants	
shirt	
shoes	
shirt	
h	

R

Laws and The

From these five basic operations one can derive the following oper

The operations $\{X, U, \bowtie, n\}$ are associative.

Laws and The

$$\sigma_{x \text{ AND } y}(R) = \sigma_{x}(\sigma_{y}(R))$$

$$\sigma_x (\sigma_y (R)) = \sigma_y (\sigma_x (R))$$

Tuple Relationa

Relational algebra is a procedu

TRC is a nonprocedural query I

A query as expressed in Tuple Re