

36. 0s: (A) → Packet 1 → (R)
 1s: (R) → Packet 1 → (B) / (A) → Packet 2 → (R)
 2s: (A) → P3 (R) → P2 Ack1 ← (B)
 3s: (A) → P4 Ack1 ← (R) → P3 Ack2 ← (B)
 4s: (A) → P5 Ack2 ← (R) → P4 Ack3 ← (B)
 5s: (A) → P6 Ack3 ← (R) → P5 Ack4 ← (B)

b) 0s: (A) → P1, P2, P3, P4 (R) (B)
 1s: (A) (R) → P1, P2, P3, P4 (B)
 2s: (A) (R) Ack1,2,3,4 ← (B)
 3s: (A) Ack1,2,3,4 ← (R) (B)
 4s: (A) → P5, P6, P7, P8 (R) (B)
 5s: (A) (R) → P5, P6, P7, P8 (B)

37. 0s: (A) → P1-4 (R) (B)
 1s: (A) P2-4 → (R) → P1 (B)
 2s: (A) P3,4 → (R) → P2 A1 ← (B)
 3s: (A) A1 ← P4 → (R) → P3 A2 ← (R)
 4s: (A) → P5 A2 ← P5 → (R) → P4 A3 ← (B)

Queue @ R

0

3

2

1

1

42. a) Throughput: 100 Mbps, Transfer Time: 46.4 μs.
 $100,000,000 \text{ bps} = X / 0.0000464 \text{ s}$
 $X = 4,640$

Minimum Packet Size: 8,192 bits

b) Higher cost of retransmission, more wasted utilization for data that doesn't align to 8,192 packets.

53. If a node 'D' outside of node 'B's' range begins communication with 'C' at the same time + channel as 'B' transmits to 'A', then collisions would occur if 'C' is within 'B's' range.

54. An ethernet node sees all traffic, and can simultaneously Tx/Rx. A wireless node can only transmit or receive at any time on a channel, and it may not be aware of all traffic on other nodes.

55. CSMA/CA has an ACK from receiver to sender which will force retransmission if a hidden terminal collides with the packet. There is also CTS-RTS which alerts hidden terminals to not transmit for a while.