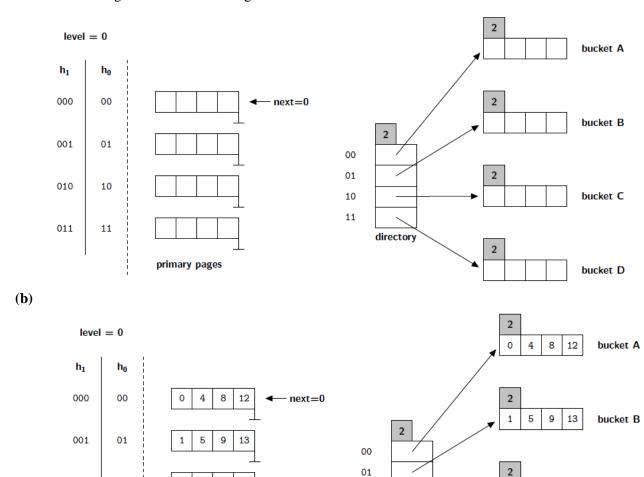
EECS 495

Introduction to Database Systems Fall 2018

Instructor: Mas-ud Hussain Solution: Homework Assignment No. 4

Problem No. 1:

(a) This is a trick question. Even before entering any data, extendible hashing has more pages than linear hashing as extendible hashing needs "bucket address table".



10

11

directory

bucket C

bucket D

10

11 15

3

(c)

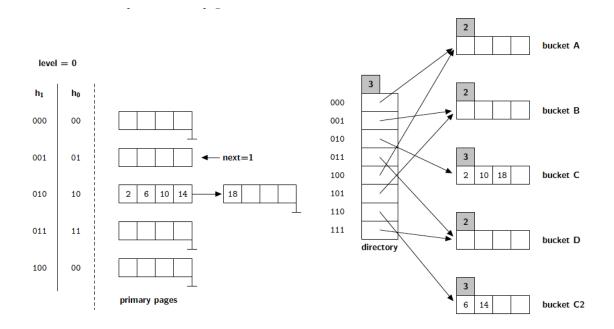
010

011

11

10

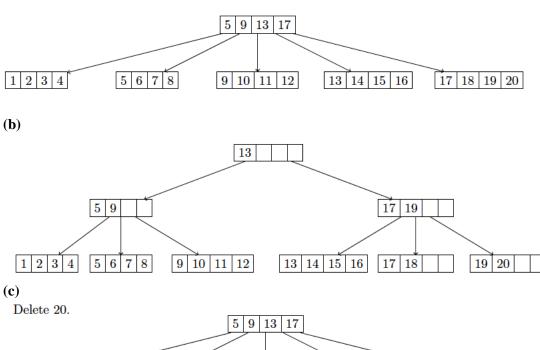
primary pages



Problem No. 2:

(a)

Insert and delete 21. The tree is full, so the height would have to increase to handle the insertion. Then, when deleting 21, the right-most node would redistribute with its neighbor to maintain the 50 percent occupancy requirement without decreasing the height.



9 10 11 12

13 14 15 16

17 18 19

Problem No. 3:

1 2 3 4

(a) <u>Sorted Runs of A:</u> A1 = (8,19,20,20)

5 6 7 8

Step-9:

 $P_{A1} = 19$, $P_{A2} = 13$; Lowest A = 13

 $P_{B1} = 13$, $P_{B2} = 10$, $P_{B3} = 15$; Lowest B = 10

As Lowest B < Lowest A; move P_{B2} to next.

Step-10:

 $P_{A1} = 19$, $P_{A2} = 13$; Lowest A = 13

 $P_{B1} = 13$, $P_{B2} = 13$, $P_{B3} = 15$; Lowest B = 13

As Lowest B = Lowest A for both (P_{A2}, P_{B2}) and (P_{A2}, P_{B1}) pairs; **output** (13,13) and (13, 13). Move all of P_{B1} , P_{B2} and P_{A2} to next.

Step-11:

 $P_{A1} = 19$, $P_{A2} = 18$; Lowest A = 18

 $P_{B1} = 15$, $P_{B2} = 19$, $P_{B3} = 15$; Lowest B = 15

As Lowest B < Lowest A; move both P_{B1} and P_{B3} to next. Run B1 is finished processing.

Step-12:

 $P_{A1} = 19$, $P_{A2} = 18$; Lowest A = 18

 $P_{B2} = 19$, $P_{B3} = 16$; Lowest B = 16

As Lowest B < Lowest A; move P_{B3} to next. Run B3 is finished processing.

Step-13:

 $P_{A1} = 19$, $P_{A2} = 18$; Lowest A = 18

 $P_{B2} = 19$; Lowest B = 19

As Lowest A < Lowest B; move P_{A2} to next.

Step-14:

 $P_{A1} = 19$, $P_{A2} = 20$; Lowest A = 19

 $P_{B2} = 19$; Lowest B = 19

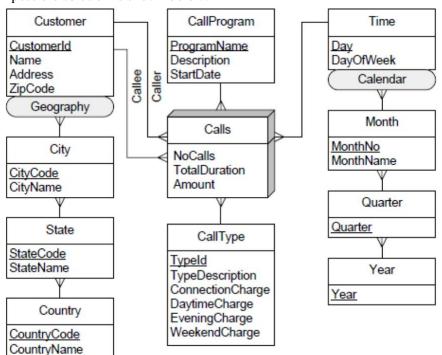
As Lowest B = Lowest A; output (19,19). Move both P_{B2} and P_{A1} . Run B2 is finished processing.

As all runs of B is finished processing, there cannot be any more matches. So, the algorithm also terminates.

Problem No. 4:

(a)

A possible solution is shown below:



```
(b)
1.
Calls1 ← _DICE(Calls, Caller.City = 'Brussels')
Calls2 ← _DICE(Calls1, Time.Year = 2016)
Result ← _ROLLUP*(Calls3, SUM(TotalDuration))
Calls1 ← _DICE(Calls, Time.Year = 2016)
Calls2 ← _ROLLUP(Calls1, Time → Year, SUM(Amount))
Calls3 ← _ROLLUP(Calls2, Caller → _All, SUM(Amount))
Calls4 ← _ROLLUP(Calls3, Callee → _All, SUM(Amount))
Result \leftarrow ROLLUP(Calls4, CallType \rightarrow All, SUM(Amount))
Calls1 DICE(Calls, Time.DayOfWeek = 'Saturday' OR
Time.DayOfWeek = 'Sunday')
Calls2 ← _DICE(Calls1, Caller.City = 'Brussels')
Calls3 ← _DICE(Calls2, Callee.City = 'Antwerp')
Calls4← _DICE(Calls3, Time.Year = 2012)
Result← ROLLUP*(Calls4, SUM(NoCalls))
4.
Calls1 ← _DICE(Calls, Caller.Country = 'Belgium')
Calls2 ← DICE(Calls1, Callee.Country <> 'Belgium')
Calls3 ← _DICE(Calls2, Time. Year = 2012)
Result ← ROLLUP*(Calls3, SUM(TotalDuration))
```