

Static Analysis of Heap-Manipulating Programs using Separation Logic

Author

Tomáš Brablec

Supervisor

Ing. Tomáš Dacík

Advisor

prof. Ing. Tomáš Vojnar, Ph.D.



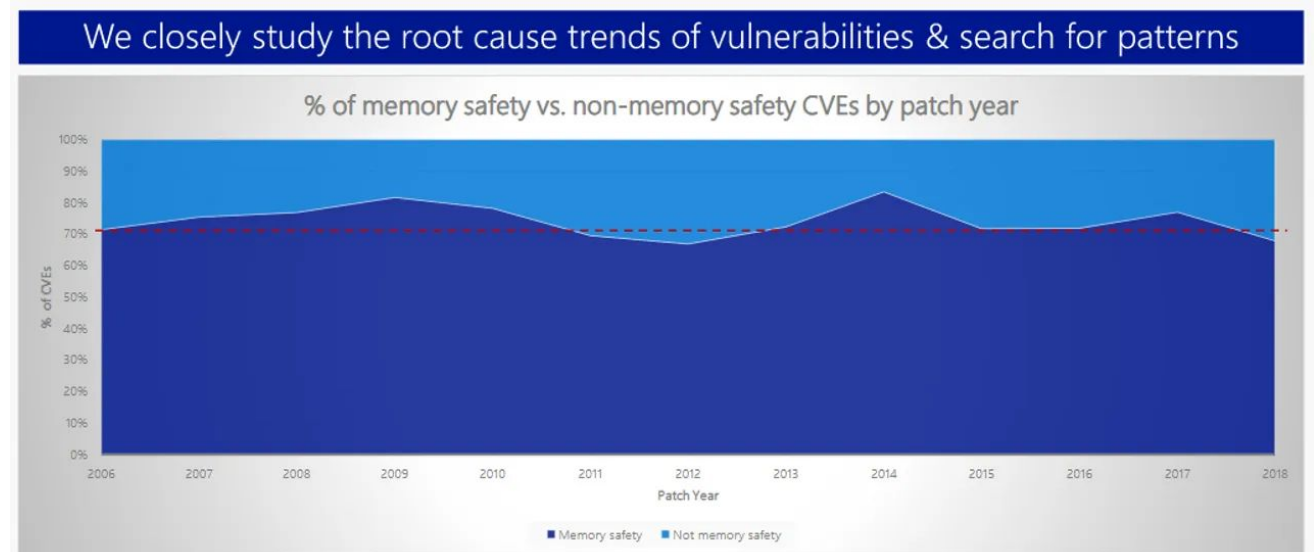
- **Manual memory management** creates a whole class of bugs
- Use-after-free, double-free, memory leaks, ...
- Memory safety errors are a common source of security vulnerabilities
- **Linked lists** are a common data structure in low level software

CVE-2024-12382 Detail

Description

Use after free in Translate in Google Chrome prior to 131.0.6778.139 allowed a remote attacker to potentially exploit heap corruption via a crafted HTML page. (Chromium security severity: High)

Example of a recent memory safety bug



Microsoft: Around 70% of CVEs are memory related

- This work introduces the **KTSN** verification tool
- Goals: verification of memory safety and bug detection
- Focused on programs with variants of **linked lists**
- Abstract memory states are encoded in **separation logic (SL)**
- Compared to existing approaches, a dedicated solver **Astral** is used for checking SL formulae:
 - Better modularity
 - Higher precision than syntactic comparison
- Implemented in the **Frama-C** framework using dataflow analysis



Astral



Software Analyzers

- Analyzed program's AST is **preprocessed** into simple instructions
- These instructions are interpreted as changes to the SL formulae
- Invariants for loops are found using **abstraction** and checking entailments of formulae using the solver
- Analyses of function calls are cached using function **summaries**

Code	State formulae
1 Node *x = malloc(size);	$x \mapsto f'_1$
2 Node *y = x;	$x \mapsto f'_1 * y = x$
3 while (rand()) {	$x \mapsto f'_1 * y = x$ $x \mapsto y * y \mapsto f'_2$ $ls_{2+}(x, y) * y \mapsto f'_3$
4 y->next = malloc(size);	$x \mapsto f'_1 * f'_1 \mapsto f'_2 * y = x$ $x \mapsto y * y \mapsto f'_2 * f'_2 \mapsto f'_3$ $ls_{2+}(x, y) * y \mapsto f'_3 * f'_3 \mapsto f'_4$
5 y = y->next;	$x \mapsto y * y \mapsto f'_2$ $x \mapsto f'_4 * f'_4 \mapsto y * y \mapsto f'_3$ $ls_{2+}(x, f'_5) * f'_5 \mapsto y * y \mapsto f'_4$
6 }	$x \mapsto f'_1 * y = x$ $ls_{1+}(x, y) * y \mapsto f'_2$
7 y->next = NULL;	$x \mapsto nil * y = x$ $ls_{1+}(x, y) * y \mapsto nil$

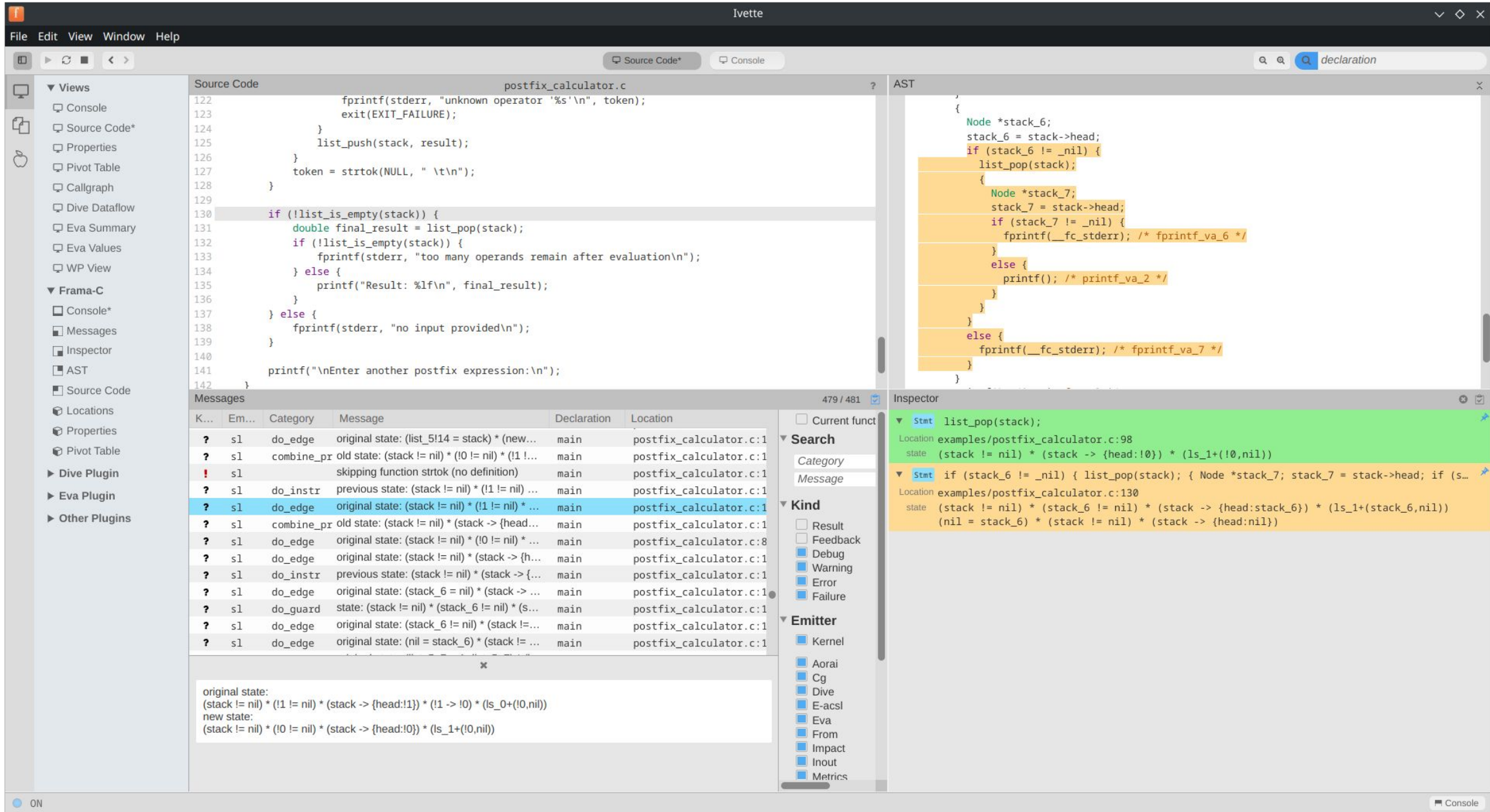
simplified analysis progress

- Tested on crafted programs:
 - Verification of programs with singly/doubly linked lists and nested lists
 - Can find use-after-free and double-free bugs, null-pointer dereferences, and memory leaks
- Tested on [SV-COMP](#) benchmarks, surpasses all but two other tools on linked lists
- Outperforms the [EVA](#) analyzer from Frama-C on most benchmarks in this category
- Presented the results at [Excel@FIT](#)

Analyser	Score	Safe	Unsafe
PredatorHP	220	96	28
CPAchecker	208	90	28
KTSN	154	74	6
symbiotic	130	51	28
CBMC	127	50	28
Bubaak	114	48	18
Mopsa	96	48	0
ESBMC	96	50	28
2LS	90	46	14
DIVINE	84	42	0
UAutomizer	22	10	2
sv-sanitizers	17	0	17
UTaipan	14	6	2

Comparison with
SV-COMP'25 participants

Results in Ivette (Frama-C GUI)



The screenshot displays the Ivette GUI interface, which is used for analyzing C code with Frama-C. The interface is divided into several panes:

- Views:** A sidebar on the left containing icons for Console, Source Code*, Properties, Pivot Table, Callgraph, Dive Dataflow, Eva Summary, Eva Values, WP View, Frama-C, Console*, Messages, Inspector, AST, Source Code, Locations, Properties, Pivot Table, Dive Plugin, Eva Plugin, and Other Plugins.
- Source Code:** The main pane showing the source code of `postfix_calculator.c`. The code includes a function `list_pop` and a `main` function. The current line of execution is highlighted at line 130.
- AST:** The right pane shows the Abstract Syntax Tree (AST) for the selected code. It displays the structure of the `list_pop` function and the `if` statement.
- Messages:** The bottom-left pane shows a list of messages generated during the analysis. The messages are organized by category (e.g., `do_edge`, `combine_pr`, `do_instr`) and location. The current message is "original state: (stack != nil) * (!1 != nil) * ...".
- Inspector:** The bottom-right pane shows the state of the program at the current location. It displays the state of variables `stack_6` and `stack_7`, and the state of the `if` statement.

The bottom status bar shows "ON" and "Console".

- Implemented a static analyzer able to verify the memory safety of programs working with linked lists
- The method outperforms most other analyzers in this category, while being more flexible

Future work

- Compete with other analyzers in SV-COMP 2026
- Extend the analysis to other types of lists
- Integrate our method into the EVA analyzer to improve its analysis of programs with linked lists

“Co znamenají řádky "Correct (true)", "Correct (false)", "Incorrect (true)" a "Incorrect (false)" v tabulce 1?”

Table 1: The results of evaluating different analyzers on the SV-COMP dataset

<u>Tests (134 total)</u>	<u>KTSN</u>	<u>PredatorHP</u>	<u>EVA</u>
Correct	80	124	56
Correct (true)	74	96	50
Correct (false)	6	28	6
Incorrect (true)	0	0	6
Incorrect (false)	0	0	48
Timeout	11	10	4
Unknown	43	0	20

- Correct (true) - Analyzer correctly verified 74 programs
- Correct (false) - Analyzer correctly identified 6 programs with a bug
- Incorrect (true) - Analyzer did not report any false positives
- Incorrect (false) - Analyzer did not report any false negatives

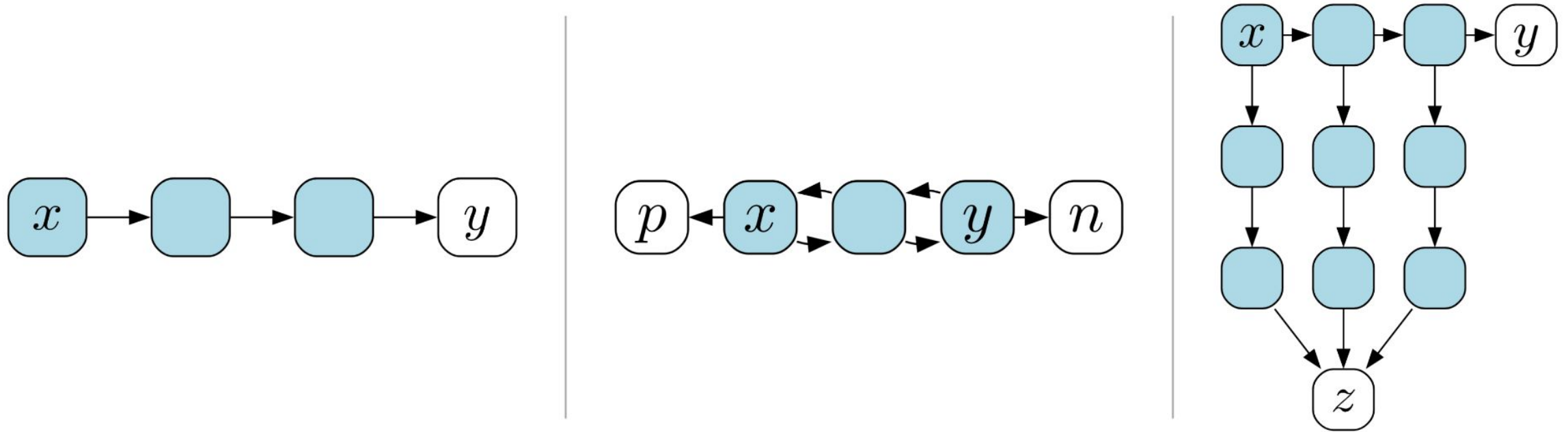


Figure 1: Supported types of lists with the allocated nodes colored blue:
singly-linked list, doubly-linked list, nested list

Comparison of analysis times

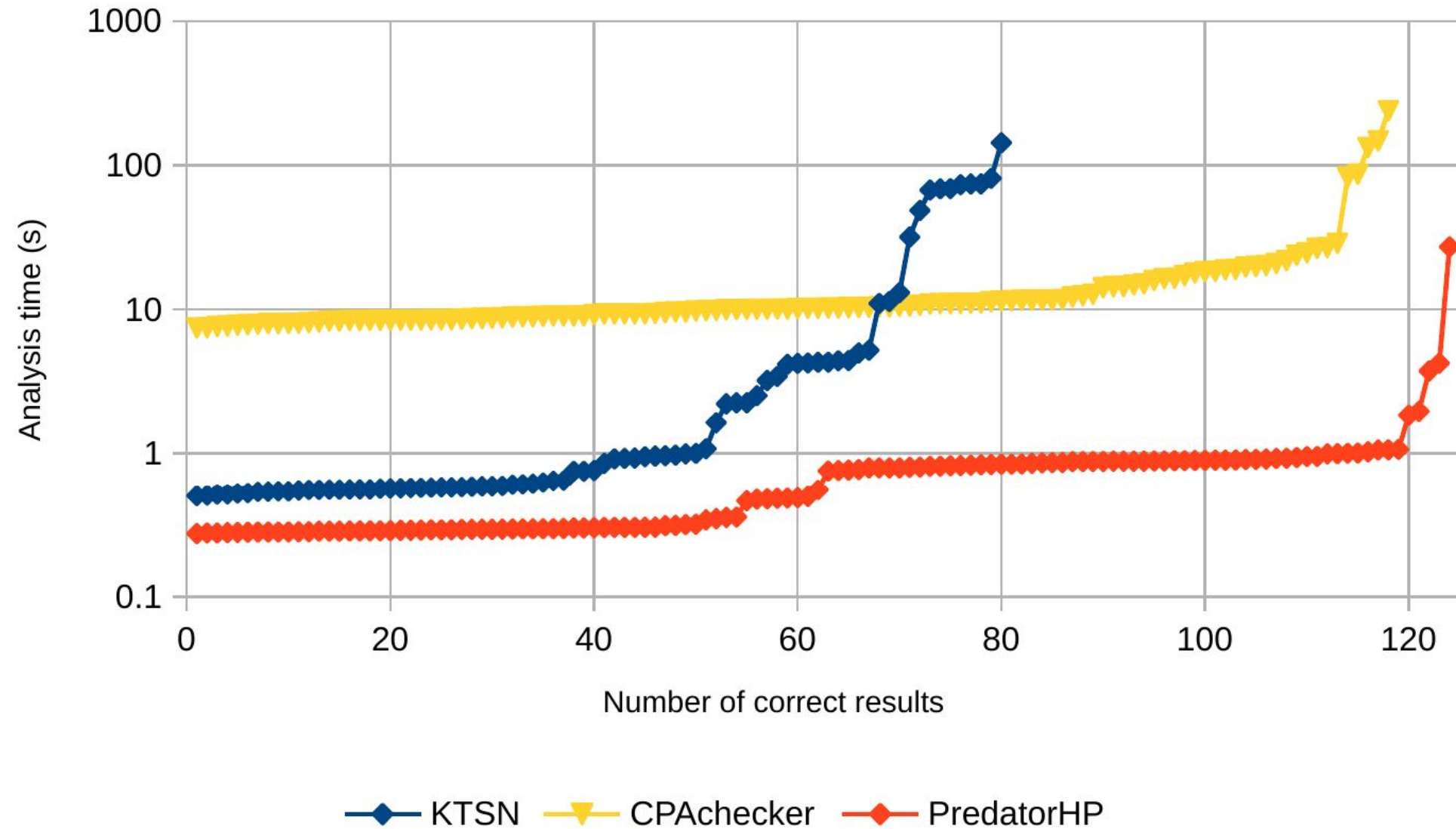


Table 3: Examples of the fastest and slowest benchmarks and the time spent in the solver versus in the analyzer itself

Benchmark	Total time (s)	Astral time (s)	Analyzer time (s)	Number of queries
sll2c_append_unequal.c	0.50	0.00	0.50	2
sll_shallow_copy-1.c	0.51	0.00	0.51	2
dll2c_prepend_equal.c	0.51	0.02	0.49	4
...				
dll-rb-cnstr_1-2.c	16.29	15.64	0.65	141
sll-sorted-2.c	31.79	31.65	0.14	1472
dll-simple-white-blue-2.c	48.61	48.13	0.48	816