DAA PRACTICAL 6 SUCCESSFUL AND UNSUCCESSFUL SEARCH

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Aim: Construction of OBST

Problem Statement: Smart Library Search Optimization

TASK 1 CODE -

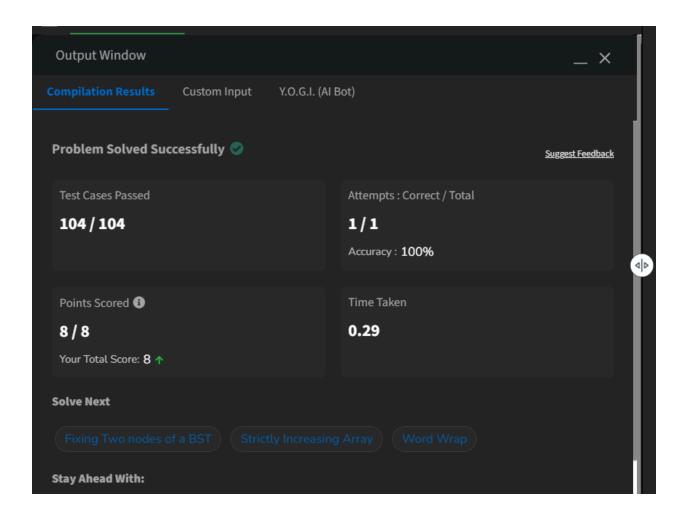
```
import java.util.*;
public class Main {
  public static void optimalBST(double[] p, double[] q, int n) {
     double[][] E = new double[n + 1][n + 1];
     double[][] W = new double[n + 1][n + 1];
     int[][] R = new int[n + 1][n + 1];
     // Step 1: Initialization
     for (int i = 0; i \le n; i++) {
        E[i][i] = q[i];
        W[i][i] = q[i];
        R[i][i] = 0;
     }
     // Step 2: Compute optimal costs
     for (int d = 1; d \le n; d++) {
        for (int i = 0; i \le n - d; i++) {
           int j = i + d;
           E[i][j] = Double.MAX VALUE;
           W[i][j] = W[i][j - 1] + p[j - 1] + q[j];
           for (int k = i + 1; k \le j; k++) {
              double cost = E[i][k - 1] + E[k][j] + W[i][j];
              if (cost < E[i][j]) {
                 E[i][i] = cost;
                 R[i][j] = k;
              }
           }
```

```
}
   }
   // Step 3: Print results
   System.out.println("E Matrix:");
   for (int i = 0; i \le n; i++) {
     for (int j = 0; j \le n; j++) {
        System.out.printf("%.2f ", E[i][j]);
      System.out.println();
   }
   System.out.println("\nW Matrix:");
   for (int i = 0; i \le n; i++) {
     for (int j = 0; j \le n; j++) {
        System.out.printf("%.2f", W[i][j]);
      System.out.println();
   }
   System.out.println("\nR Matrix:");
   for (int i = 0; i \le n; i++) {
     for (int j = 0; j \le n; j++) {
        System.out.print(R[i][j] + " ");
     System.out.println();
}
public static void main(String[] args) {
   int n = 3;
   double[] p = \{0.15, 0.10, 0.05\};
   double[] q = \{0.05, 0.10, 0.05, 0.05\};
   optimalBST(p, q, n);
}
```

}

```
Output
E Matrix:
0.05 0.45 0.90 1.25
0.00 0.10 0.40 0.70
0.00 0.00 0.05 0.25
0.00 0.00 0.00 0.05
W Matrix:
0.05 0.30 0.45 0.55
0.00 0.10 0.25 0.35
0.00 0.00 0.05 0.15
0.00 0.00 0.00 0.05
R Matrix:
0 1 1 2
0 0 2 2
0 0 0 3
0 0 0 0
```

```
{
           sum[i][j] = sum[i][j - 1] + freq[j]; // Sum of frequencies from i to j
     }
     // Step 2: Fill the dp array for the base case when there's only one key
     for (int i = 0; i < n; i++)
     {
        dp[i][i] = freq[i]; // Cost when there's only one key is just its frequency
     }
     // Step 3: Solve for larger subarrays
     for (int len = 2; len \leq n; len++)
     { // len is the subarray length
        for (int i = 0; i \le n - len; i++)
           int j = i + len - 1; // j is the end index of the subarray
           // Initialize dp[i][j] to a large value
           dp[i][j] = Integer.MAX VALUE;
           // Try all possible roots for the current subarray keys[i...j]
           for (int r = i; r <= j; r++)
           {
              // Cost if r is the root
              int cost = (r > i ? dp[i][r - 1] : 0) + (r < j ? dp[r + 1][j] : 0) + sum[i][j];
             // Update dp[i][j] with the minimum cost
              dp[i][j] = Math.min(dp[i][j], cost);
           }
        }
     }
     // The answer is the minimum cost of the BST for the whole array
     return dp[0][n - 1];
  }
}
```



Compilation Completed					
• Case 1					
Input: 🗓					
2 10 12 34 50					
Your Output:					
118					
Expected Outp					
118					